Efficient On-Demand Operations in Distributed Infrastructures

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We need to design responsive datacenters that handle massive data.

- Will use monitoring (querying) as a main example
The Need

- **Cloud computing**

**Apps in the Cloud**
- Social networking sites: facebook, LiveJournal, etc
- In-Cloud desktop software: Google docs, Photoshop express, Apple’s MobileMe, etc
- Data-intensive apps: log analysis, scientific computing, etc

**Massive data in the Cloud**
- Arise in many domains, e.g., scientific computing, social networking sites, web search engines, etc
- TB to PB (e.g., Yahoo 12TB/day)

**Cloud initiatives**
- Industry: HP, IBM, Amazon, Google, Yahoo, MS, ...
- Academia: Illinois CCT, Google-IBM cluster, etc
The Requirement

Responsiveness while handling massive data

- Web 2.0 apps: users expect desktop-quality responsiveness.
- Data-intensive processing: long, long time to finish

Importance

- Amazon: every 100ms of latency cost them 1% in sales.
- Google: an extra .5 seconds in search page generation time dropped traffic by 20%.
- “Users really respond to speed” – Google VP Marissa Mayer
- Coadd took 70 days to complete with over 30 sites and 4,500 CPUs.
- Coadd? a spatial processing application with 44,000 tasks accessing 588,900 files in total
Not Easy to Do Both

Datacenter monitoring example

- Monitoring responsiveness: quick query results regarding the state of the datacenter

Case of HP OpenView

- 144 attributes, store them in a database
- **Data scale**: # of attributes * # of machines
- First-cut attempt in 2006: “6 ½ hours for 6,200 machines, and 12 hours for 11,500 machines”
- Not possible to query the most up-to-date state

More than merely scale

- Static vs. dynamic attributes
Challenges - Scale & Dynamism

Monitoring example

(Data) Scale

Scaling centralized (e.g., Replicated DB)

Centralized (e.g., DB)

Static Attributes (e.g., CPU type)

Centralized (e.g., DB)

Dynamic Attributes (e.g., CPU util.)

(Attribute) Dynamism
On-Demand Operations

On-demand operations

- Operations that are responsive in spite of massive data
- Challenges: scale & dynamism

Examples

- On-demand monitoring: quick response even when querying the most up-to-date state of the infrastructure
- On-demand scheduling for data-intensive workload: quick scheduling decisions reflecting the current resource availability
On-Demand Group Querying

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Problem

Group monitoring (querying) in large-scale infrastructures
Necessity

Groups are formed naturally

• A set of machines running a service in a datacenter, PlanetLab Slices, Grid sites, etc

Users want to query attributes from groups

• Avg. CPU utilization of a slice
• Top-3 loaded machines running Linux and executing a MapReduce task
• List of machines in rack R with either CPU util. < 10% or Mem util. < 5% running less than 2 VMs on either Xen or VMware
Requirements

Efficient group support with expressive queries

- Single groups & multiple groups
  - Queries targeting unions and intersections of groups
  - List of machines in rack R with CPU util. < 10% or Mem util. < 5% running less than 2 VMs on Xen or VMware
  - $A \text{ and } (B \text{ or } C) \text{ and } D \text{ and } (E \text{ or } F)$
- Static groups & dynamic groups
  - Static groups: Linux machines, web servers, etc
  - Dynamic groups: CPU > 90%, # of VMs < 2, etc
- Query: (Aggregation function, Query attribute, Group predicate)

Responsiveness while handling massive data

- Quick query results even for the most up-to-date state
- Massive data (many attributes over many machines)
Previous Solutions

Centralized DB-based systems

• Collect data to a DB (HP OpenView, IBM Tivoli, ...)
• Advantage: very expressive (SQL)
• Disadvantage: can’t deliver freshness & scalability at the same time (e.g., every 5 min for CoMon on PL) – gathering all attributes every time

Distributed aggregation systems

• Build end-host aggregation trees (Astrolabe, SDIMS, Seaweed, MON, etc)
• Advantage: scalability with responsiveness
• Disadvantages (next slide)
Previous Solutions

**No support for groups**

- One group = the entire system
- They all collect data from the entire system

**E.g.,**

- Just to monitor problematic 3 machines in your PlanetLab slice with 400 machines
- => Need to get data from all 400 machines
Moara Overview

A group-based querying system

- Allows users to target specific groups of machines
- Fast query resolution & expressive group predicates

Group trees

- Moara builds per-group trees & uses these for aggregation
- Different groups can share a tree for scalability.
- Leverages SDIMS (a Pastry-based aggregation framework)

Two optimization techniques

- Multi-group optimization
- Single-group optimization for dynamic groups
Multi-Group Aggregation

Goal

- Quick results with minimum B/W when querying union (or) & intersection (and) of groups
- E.g., \{(A or B) and C\} or \{(D and E)\}?
- Naïve, but correct approach: send a query to every group and collect the data

A = \{Apache\}  B = \{CPU > 90\%\}
Multi-Group Aggregation

Optimization opportunity for latency and bandwidth

- Querying \((A \text{ and } B)\): need to send queries to only one (smaller) of the two
- What about complex expressions? e.g., \(\{(A \text{ or } B) \text{ and } C\} \text{ or } (D \text{ and } E)\)?
- Solution: CNF transformation & a few more tricks

\[ \begin{align*}
A &= \{\text{Apache}\} \\
B &= \{\text{CPU} > 90\%\}
\end{align*} \]
Efficient Dynamic Group Management

Group tree management

T: CPU-Util > 50%
F: CPU-Util <= 50%
Efficient Dynamic Group Management

Group tree management

- Two choices: keep or cut

T: CPU-Util > 50%
F: CPU-Util <= 50%
Efficient Dynamic Group Management

**Tradeoff**

- Aggressive group tree management: saves per-query cost, but increases management cost
- Lazy group tree management: increases per-query cost, but saves management cost

**Adaptive group tree management**

- Moara strikes the balance between aggressive and lazy tree management (best of both worlds)
- Adaptively adjust aggressiveness of management by continuously tracking query cost and management cost
Evaluation

PlanetLab: 200 nodes
- Represents wide-area infrastructures

Emulab: 500 instances on 50 machines
- Represents medium-size datacenters

Simulation: up to 10K nodes
- Don’t care about latency
- Only measure # of msgs
Emulab

Latency & Bandwidth (static group, 500 instances)
Emulab

Latency & Bandwidth (static group, 500 instances)

Latency (ms)  Msg/query
Simulation

Adaptive tree mechanism (10K nodes)

- No tree management, one global tree
- Moara with aggressive tree management
- Moara with adaptive tree management

Graph showing messages per node (Msgs/node) and group churn over queries. The graph compares different management strategies.
Beyond On-Demand Operations
Beyond On-Demand Operations

Data-intensive apps in the Cloud

- Apps access & generate massive amount of data (TB to PB)

Need new solutions to old problems

- E.g., scheduling for data-intensive apps (Hadoop scheduler, worker-centric scheduler – Middleware ‘07, etc)

New problems

- How to coordinate data management and scheduling
- How to handle intermediate data
  - Data that exists only during the lifetime of an app
  - E.g., data from the Map phase
Beyond On-Demand Operations

Web 2.0 apps in the Cloud
- Problem: responsiveness
- Users expect desktop-quality responsiveness

Best practices
- Make each component as fast as possible
- E.g., multi-layer caching: DB in-memory cache, in-datacenter distributed cache (e.g., memcached), CDN, etc

What’s missing? Coordination
- Possibility: Design components that are “aware” of one another.
- To achieve maximum possible performance
Summary

We need to design responsive datacenters that handle massive data.

- Web 2.0 apps in the Cloud
- Data-intensive apps in the Cloud

On-demand operations

- Operations that achieve the goal
- Scale & dynamism are the challenges.