

### Staged Simulation

Improving Scale and Performance of Wireless Network Simulations

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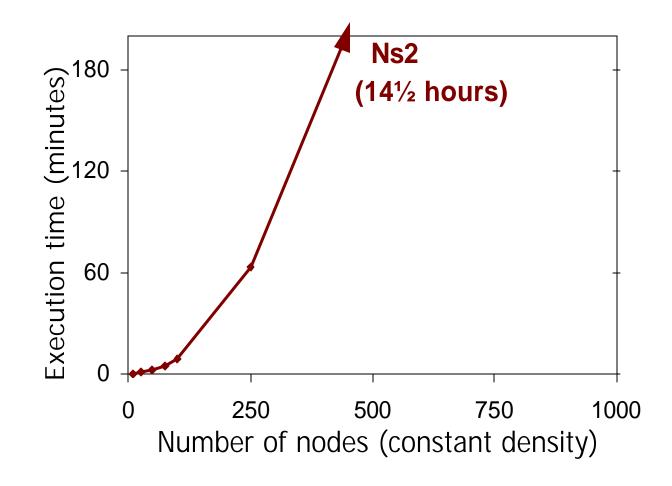


### Wireless Simulation

- Simulation for networking research
  - Ad hoc routing protocols
  - Mobile applications
  - Sensor networks
- Wireless simulation is often slow and does not scale
  - Takes too long
  - Works only for small networks

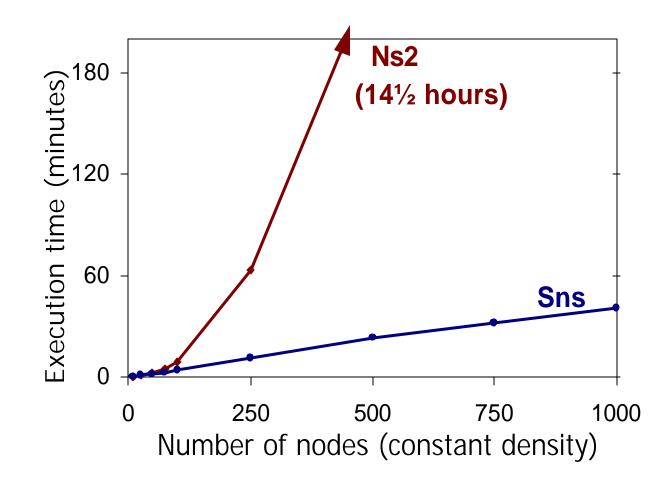


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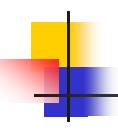


- Redundant computations, due to conservative structure of simulator
  - Can be highly dynamic
  - Recompute state whenever inputs may have changed
- State recomputed more often than strictly necessary



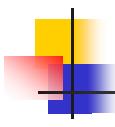
# Redundancy Across Simulator Runs

- Simulator often used for large batch runs
  - Nearly identical scenarios/parameters
- Leads to redundancy across runs
  - Users can identify or predict



### Contributions

- Identify problem
  - Redundant computations
- Propose approach: Staging
  - Reduce time spent in redundant computations
  - Intra- & inter-simulation staging
- Demonstrate speedup with practical simulator
  - No loss of accuracy or change in interface



### Staging: Approach

- Basic approach: Result caching and reuse
  - Limited use, due to real valued and continuously varying inputs: e.g. current simulation time
- Restructure events within a discrete event simulator to make them amenable to caching
  - Expose redundancy for caching and reuse
  - Leads to small, time-independent computations
- Use time-shifting where possible
  - Reorder events for maximum efficiency



### Simulator Model

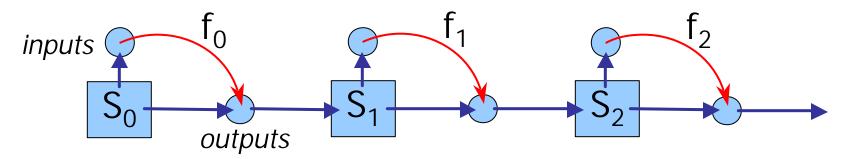
- Progression of states S<sub>i</sub>
- An event  $e_i$  acts on  $S_i$  and produces  $S_{i+1}$ 
  - Ex. Packet transmission, topology computation





### Simulator Model

- Progression of states S<sub>i</sub>
- An event  $e_i$  acts on  $S_i$  and produces  $S_{i+1}$ 
  - Ex. Packet transmission, topology computation
- Can view events as functions
  - Inputs taken from current state
  - Output describes modifications to get next state



- Use currying
  - Decompose event into multiple events
  - Group part of computation that depends on slowly varying or discrete inputs
  - Example: node position nearly static

time

$$r = f(t, p_1, p_2, p_3, ...)$$

$$r' = f(t', p_1, p_2, p_3, ...)$$



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 $r' = f(t', p_1, p_2, p_3, ...)$ 

$$r = f'(t, g(p_1, p_2, p_3, ...))$$

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 $r' = f(t', p_1, p_2, p_3, ...)$   $r' = f'(t', g(p_1, p_2, p_3, ...))$ 

- Use incremental computation
  - Reuse results of similar computations
  - Often relies on continuity w.r.t. an input
  - Example: topology similar at nearby times

time

$$r = f(t, p_1, p_2, p_3, ...)$$

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- Use auxiliary results
  - Compute and save additional information
  - Example: bounded node speed

## time $r = f(t, p_1, ...)$ $r' = f(t', p'_1, ...)$

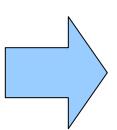


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### time

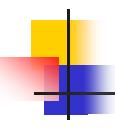
$$r = f(t, p_1, ...)$$

$$r' = f(t', p'_1, ...)$$



$$\{r, a, b\} = f'(t, p_1, ...)$$

$$\{r', a', b'\} = f'(t', p'_1, ..., a, b)$$



### Time-shifting

- Restructuring and caching provide opportunities for changing the time at which computations are performed
  - Smaller, time-independent events
- Architectural benefits
  - Better working set and cache performance using precomputation and event reordering
- Algorithmic benefits
  - More efficient algorithms using batch processing



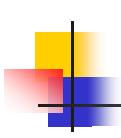
### Staging in Practice: Sns

- Based on ns-2 wireless network simulator
  - Ns-2 wireless is slow; scales poorly
  - Implementation is typical
- Inter- and intra-simulation staging
- No change in accuracy or interface



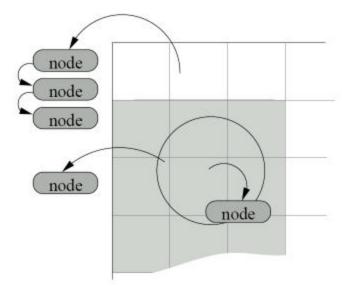
### Neighborhood Computation

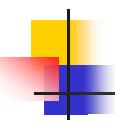
- During each packet send, compute nodes in neighborhood of sender
- Expensive in ns-2
  - Full scan of network on every packet
  - Leads to redundant computations for typical networks
- Staging applied to neighborhood computation



# 1: Grid-based Neighborhood Computation

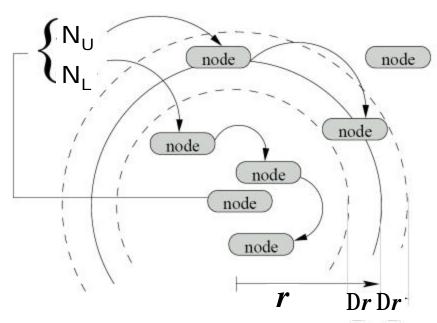
- Use a grid to compute nodes within range
  - Reuse results for nearby nodes
  - Share grid maintenance across calls
- Reduces number of nodes examined
- Elementary form of staging
  - Currying and incremental computation

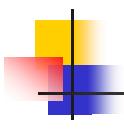




## 2: Neighborhood Caching

- Compute upper/lower bounds on neighborhood set, with expiration time? t
- Refine bounds into exact result
- Many computations share same bounds
- Reduces number of nodes examined
- Staging by auxiliary results





### 3: Time-shifting

- Precompute all neighborhood cache entries together
- Compute on schedule, every ? t epoch
- Reduces total work (batching)
- Improves memory locality
- May introduce new work
  - Perform only under heavy load

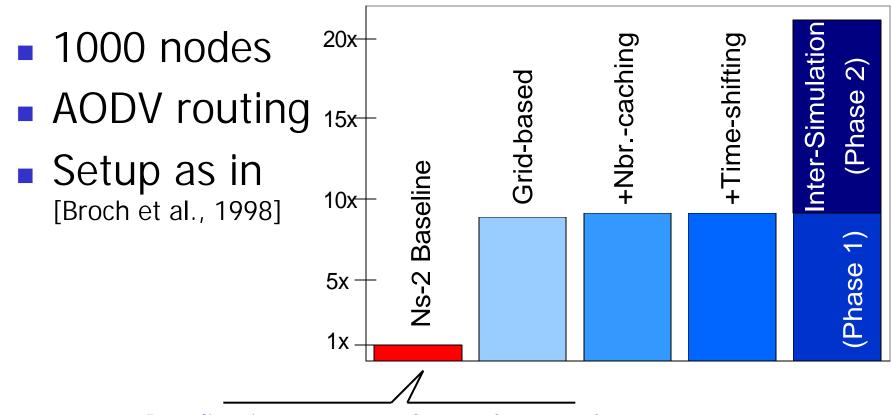


## 4: Inter-simulation Staging

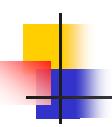
- Reuse neighborhood sets across simulation runs
  - All runs have same mobility scenario
  - Other simulation parameters may differ
- Generate phase: write neighborhood sets to disk
  - First run in batch
- Use phase: read sets from disk
  - No grid or topology needed if cache is complete
  - Need only the previously computed result cache



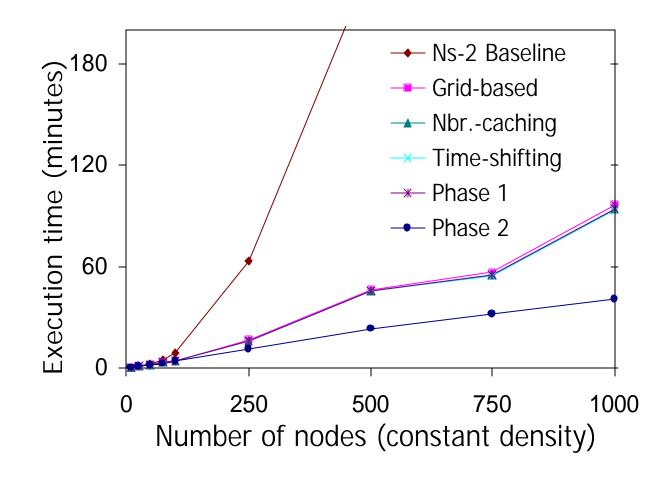
### **Execution Time Speedup**



Baseline is approx. 2x faster than stock ns-2 (using standard optimization techniques)



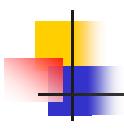
### Effect of Network Size





### Related & Prior Work

- Existing instances of intra-simulation staging
  - NixVectors for wired networks [Riley et al. 2000]
  - Selective packet transmission [Wu & Bonnet 2002]
- Instances of inter-simulation staging
  - Splitting [Glasserman et al. 1996]
  - Cloning [Hybinette & Fujimoto 1997]
  - Updateable Simulations [Ferenci et al. 2002]
- Staging in other domains
  - Compilation [Chambers 2002], iterative programming, memoization



### Conclusions

- Insight: Redundant computations are main bottleneck for wireless simulation
- Staging improves speed & scale by eliminating redundant computation
- No loss in accuracy
- Applicable to a variety of simulation engines
- $O(n^2)$  to O(n) speedup for ns-2

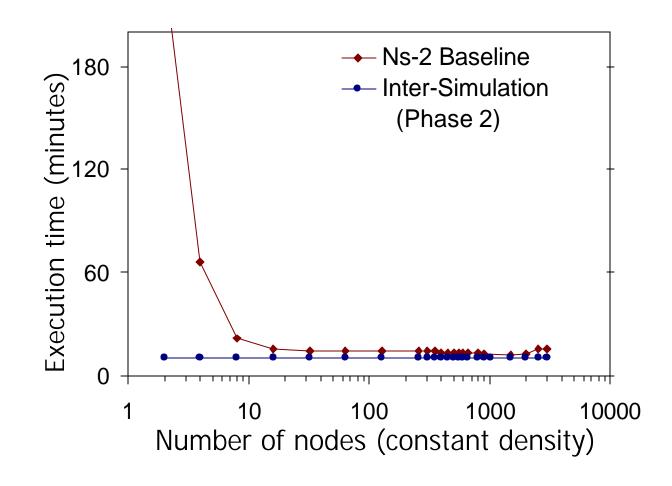
### http://www.cs.cornell.edu/People/egs/sns/

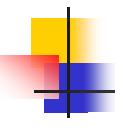
Staged Simulation for Improving the Scale and Performance of Wireless Network Simulations. Kevin Walsh and Emin Gün Sirer. In *Procedings of the Winter Simulation Conference*. December 2003.

Staged Simulation: A General Technique for Improving Simulation Scale and Performance. Kevin Walsh and Emin Gün Sirer. In ACM Transactions on Modeling and Computer Simulation (TOMACS). April 2004.



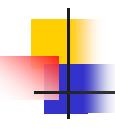
### **Grid Performance**





### Sensitivity to Parameters

- Initial application of staging (grid):
  - Very sensitive to granularity (both in memory and CPU)
  - Poorly tuned grid much worse than baseline
- Successive applications of staging:
  - Reduce sensitivity to granularity
  - Less sensitive to other parameters
  - 10% variation in execution time over wide range of parameters



### Memory Use and Performance

- Ns2 severely memory constrained
  - Artifact of simulator implementation
- Grid-based staging adds:
  - Typical: 1-10 KB, many new events
  - Poorly-tuned: 1-100+ MB, many new events
- Other intra-simulation staging adds:
  - Typical: 20-200 KB, few to no new events
  - Uses grid, but avoids worst-case scenarios
- Inter-simulation staging:
  - Eliminates grid entirely