

# Principled Programming

Introduction to Coding in Any Imperative Language

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*Department of Computer Science*

*Cornell University*

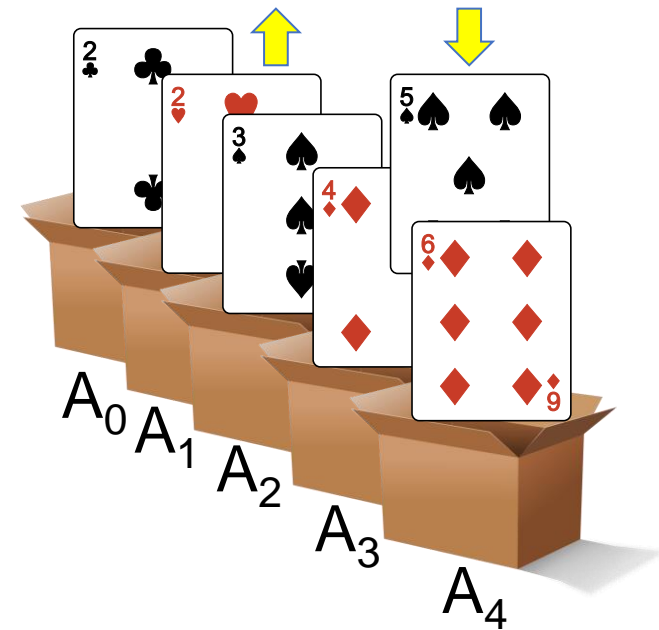
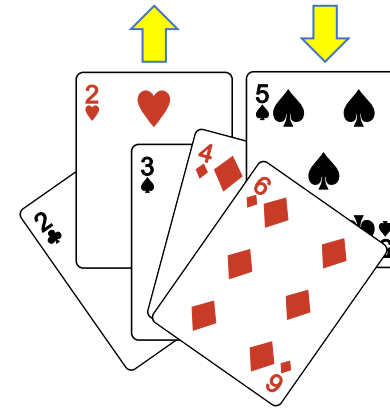
## One-Dimensional Array Rearrangements

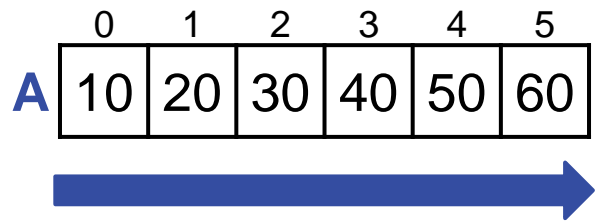
The need to rearrange values in an array is commonplace, and facility in doing so is important.

Everyday experience is helpful, e.g., manipulating a hand of playing cards. However, beware that when cards are deleted or inserted, others move over automatically. A better analogy is cards in boxes, but even this is flawed because values are *copied* from variables, not *pulled*, like cards.

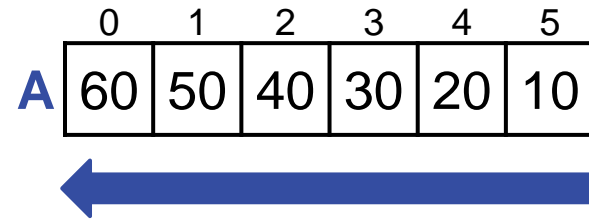
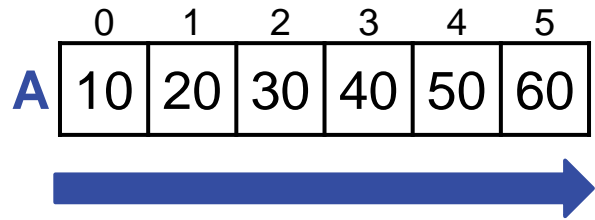
We consider:

- Reverse
- LeftShift
- LeftRotate
- Partitioning
- Collation





**Application:** Reverse the order of an array.

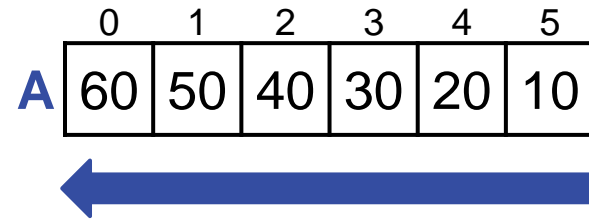
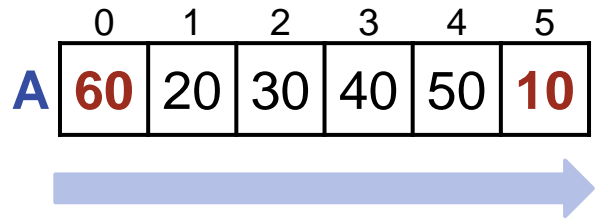


**Application:** Reverse the order of an array.

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 **There is no shame in reasoning with concrete examples.**

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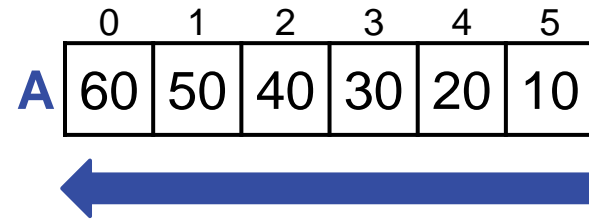
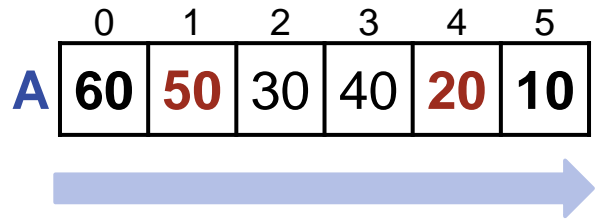


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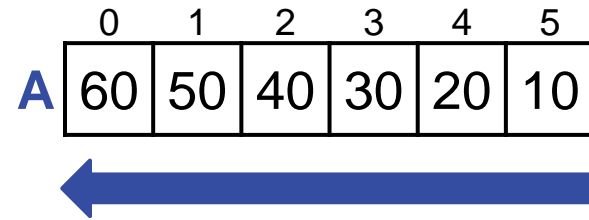
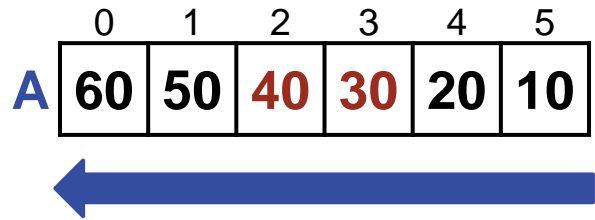


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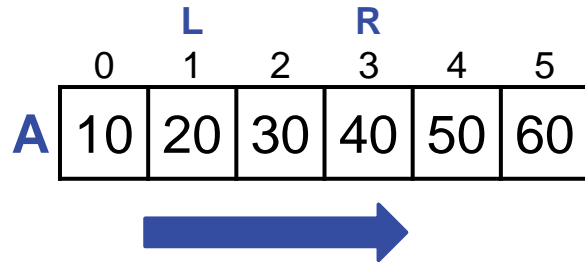


**Application:** Reverse the order of an array.

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**Application:** Reverse the order of (a subsequence of) an array.

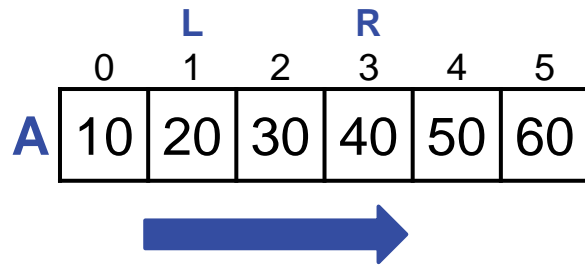
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   in situ without affecting the rest of A. */
static void Reverse( int A[], int L, int R ) {
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---

👉 **A header-comment says exactly what a method must accomplish, not how it does so.**

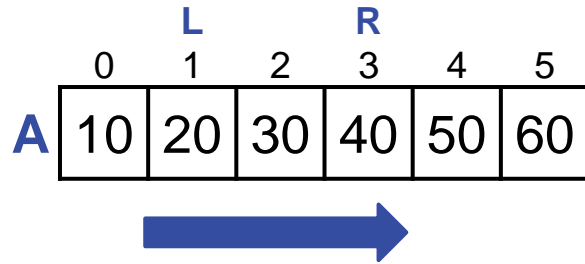
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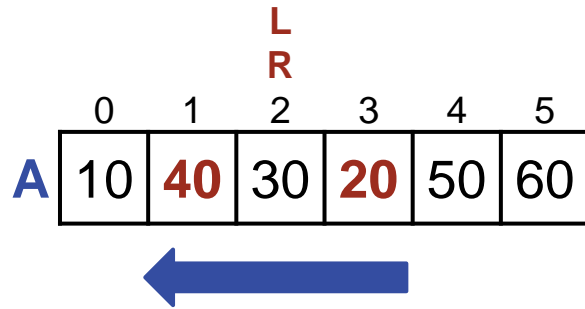
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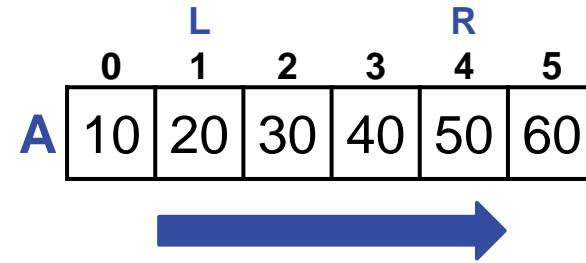
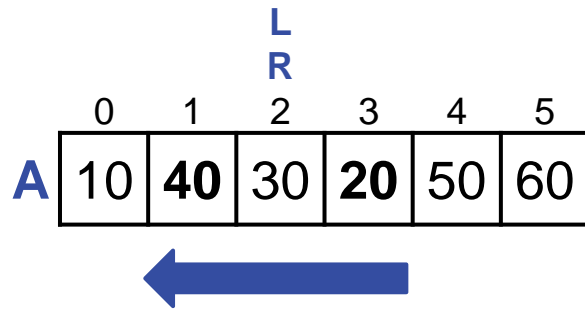
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/* Given int array A[0..n-1], reverse the order of the subsequence A[L..R]
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static void Reverse( int A[], int L, int R ) {
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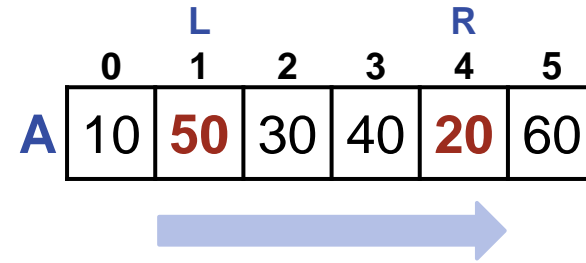
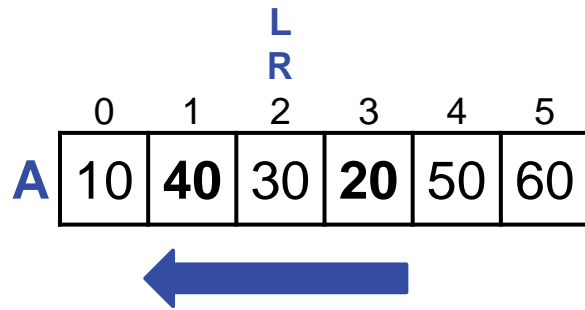
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---

 **Validate output thoroughly.**

---



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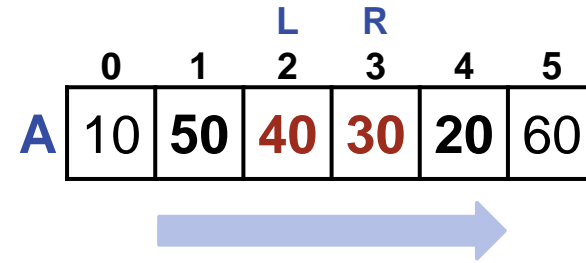
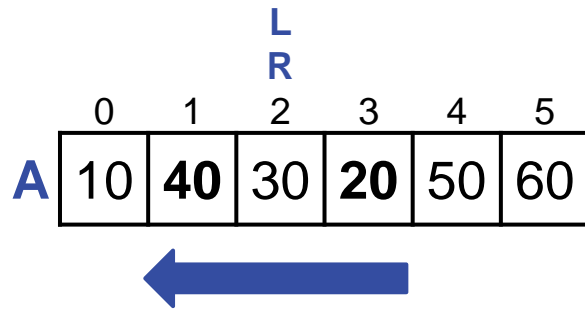
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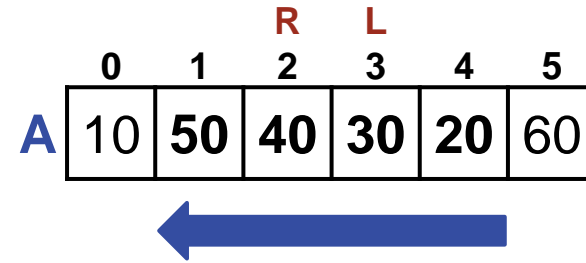
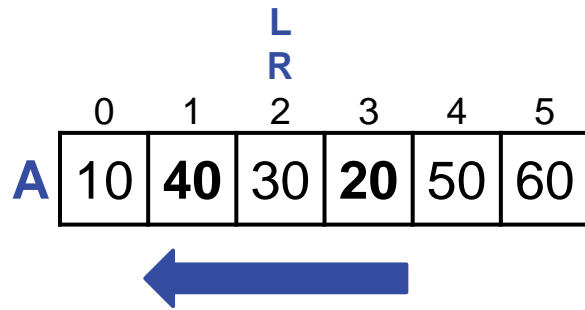
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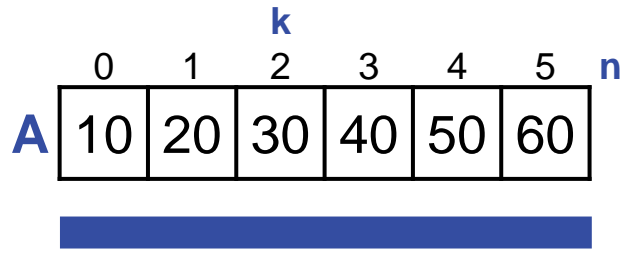
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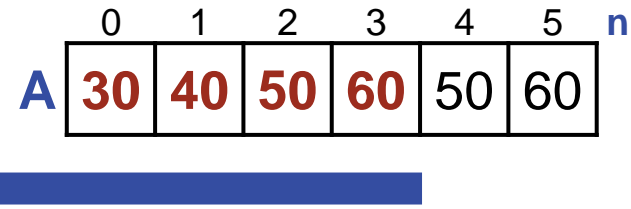
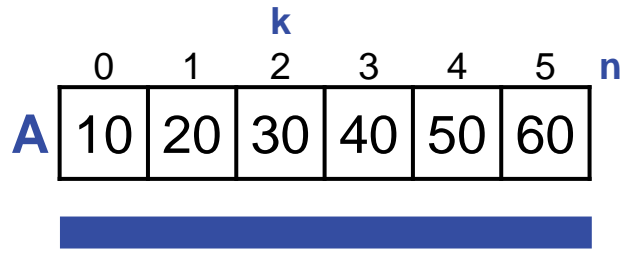
It would have been an error for the *condition* to have been  $L \neq R$ .

 **Be alert to high-risk coding steps associated with binary choices.**



**New Application:** Shift an array left  $k$  places.



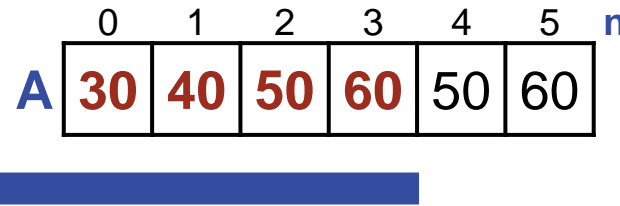
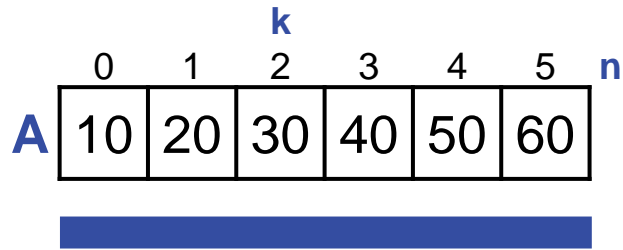


**Application:** Shift an array left k places.

---

 There is no shame in reasoning with concrete examples.

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**Application:** Shift an array left k places.

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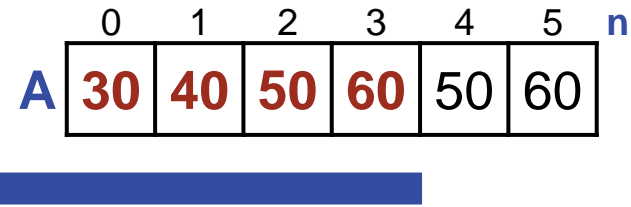
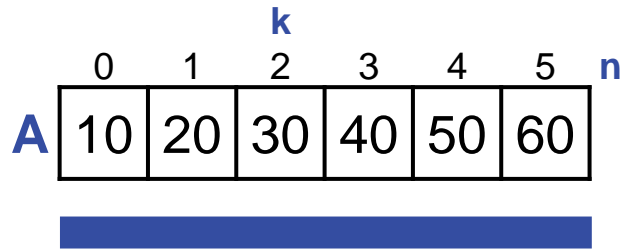
/* Given array A[0..n-1], and  $0 \leq k$ , shift elements of A left k places.
   Values shifted off the left end of the array are lost. Values not
   overwritten remain as they were originally. */
static void LeftShiftK( int[] A, int n, int k ) {
    } /* LeftShiftK */

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 **A header-comment says exactly what a method must accomplish, not how it does so.**

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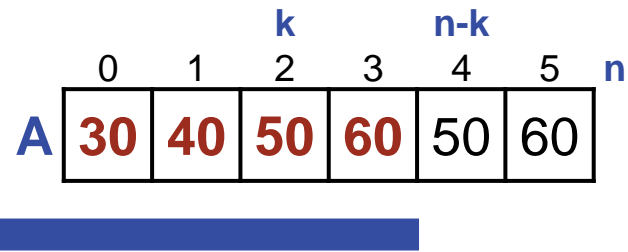
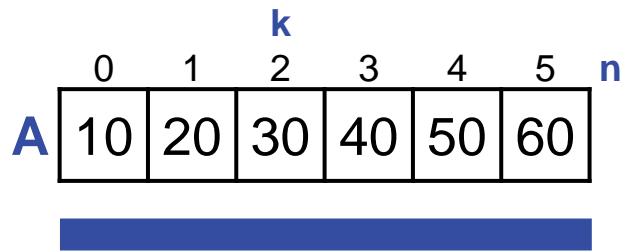


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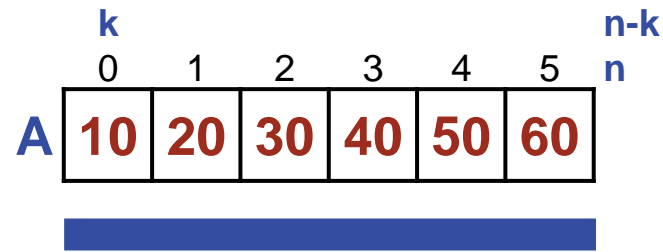
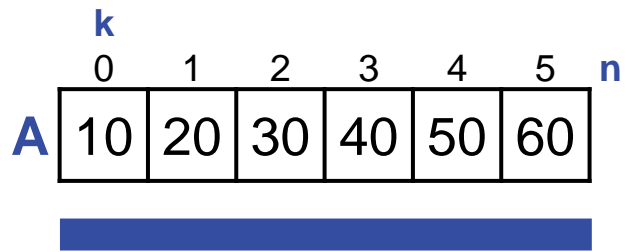


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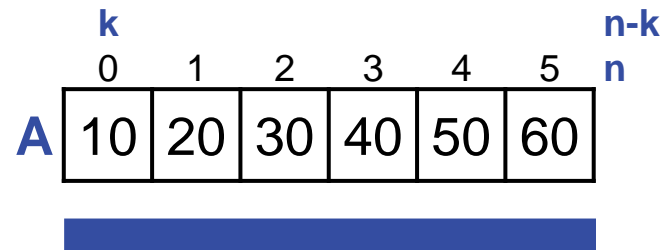
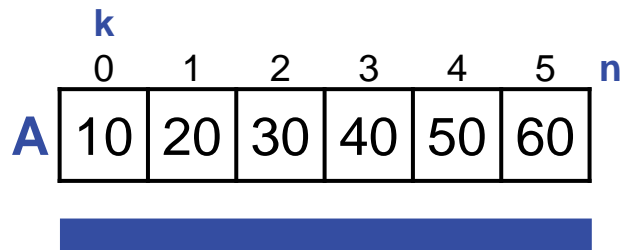
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 **Boundary conditions. Dead last, but don't forget them.**

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    if ( k > 0 )
        for (int j=0; j<n-k; j++) A[j] = A[j+k];
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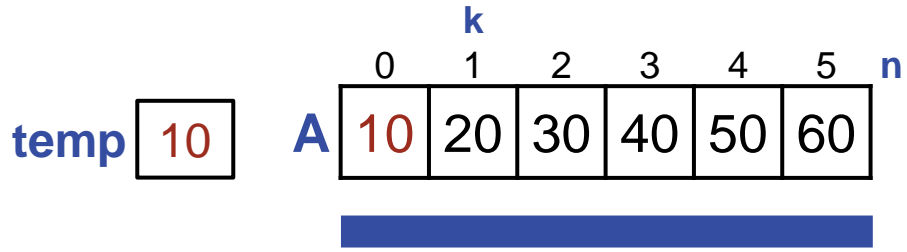
```

It would have been correct without this test, but offensive that for  $k=0$  we would do the most work.

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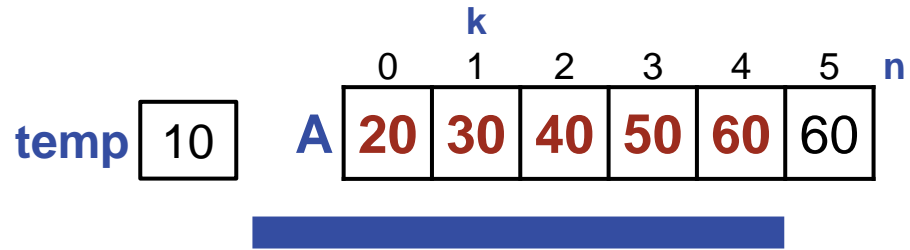
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**New application: Rotate an array left 1 place.**

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/* Given int array A[0..n-1], shift A[1..n-1] 1 place left, with the value
   originally in A[0] reentering at right in A[n-1]. */
static void LeftRotateOne(int A[], int n) {
    int temp = A[0];
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    A[n-1] = temp;
} /* LeftRotateOne */
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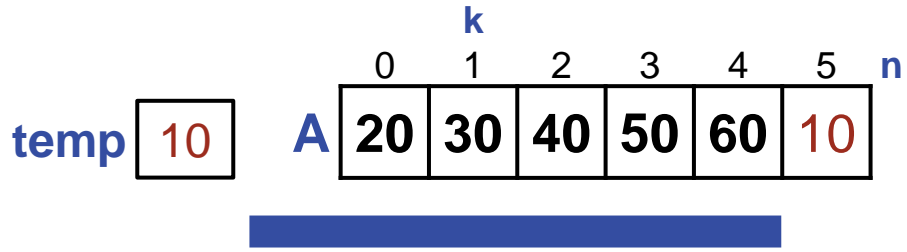
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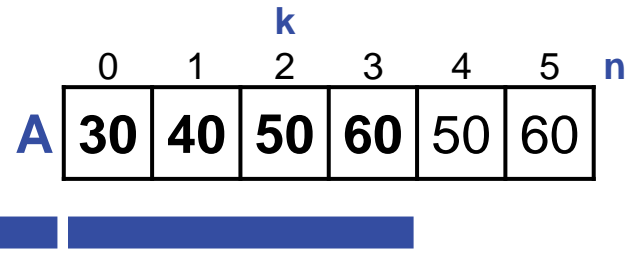
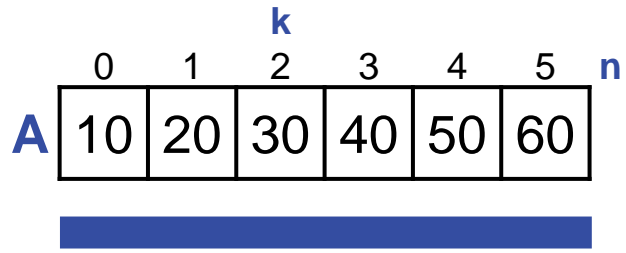


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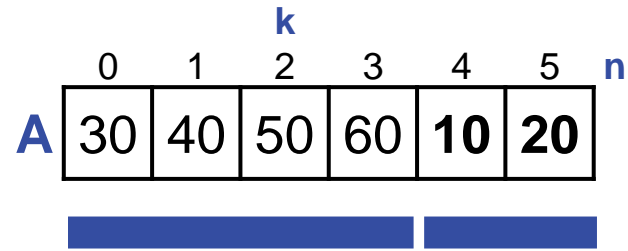
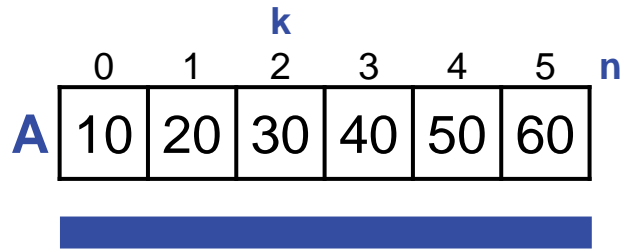
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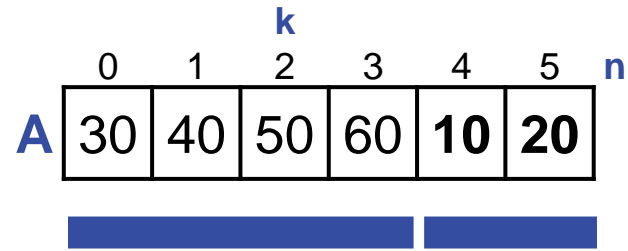
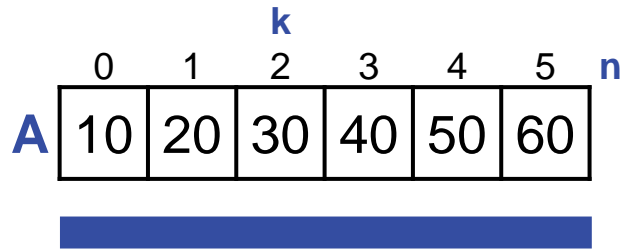
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```

We shall consider four distinct approaches:

- Repeated Left-Rotate-1
- Swap Generalization
- Three Flips
- Juggle in Cycles

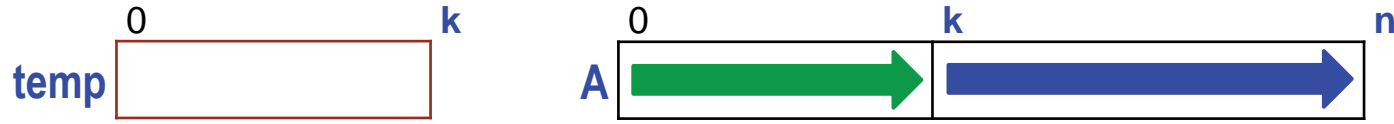
**Approach 1: Repeated left rotation 1 place.**

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/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
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for (int j=0; j<k; j++) LeftRotateOne(A, n);
```



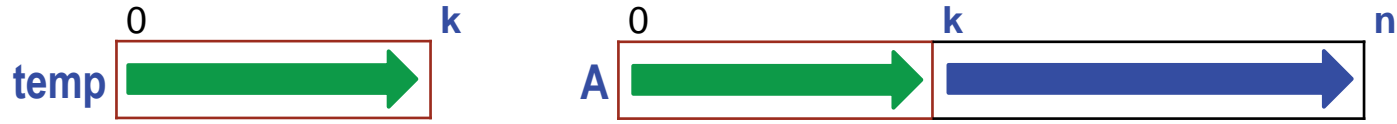
**Approach 2:** k-wide generalization of swap.

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/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
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**Approach 2:** k-wide generalization of swap.

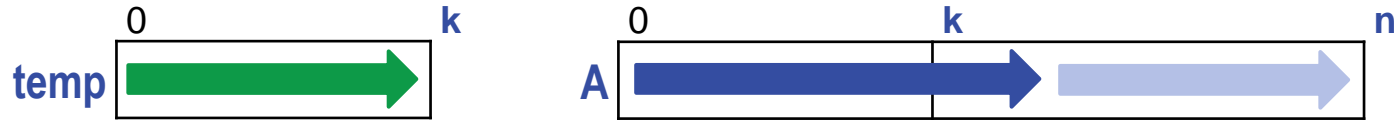
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int temp[] = new int[k];
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**Approach 2:** k-wide generalization of swap.

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/* temp[0..k-1] = A[0..k-1]; */
```



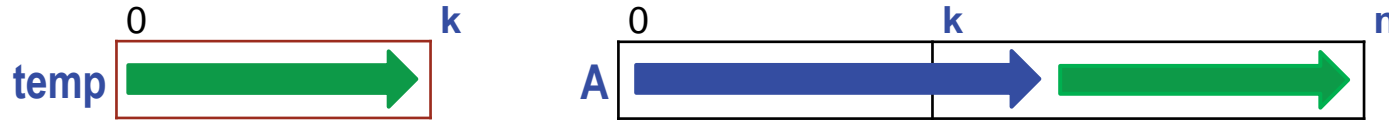


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LeftShiftK(A, n, k);

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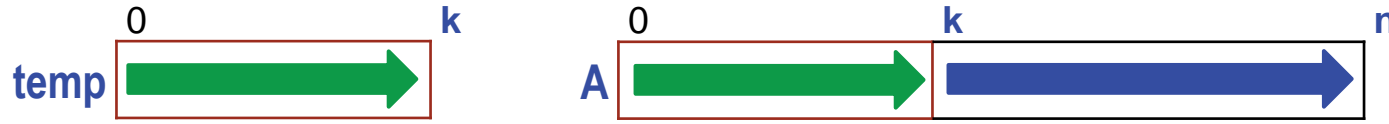
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```

---

 **Defer challenging code for later; do the easy parts first.**

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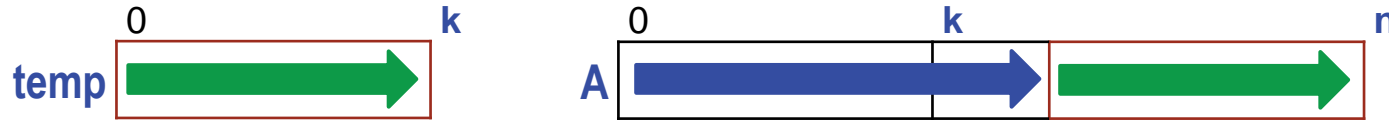


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/* A[___..n-1] = temp[0..k-1]; */

```



**Approach 2:** k-wide generalization of swap.

```

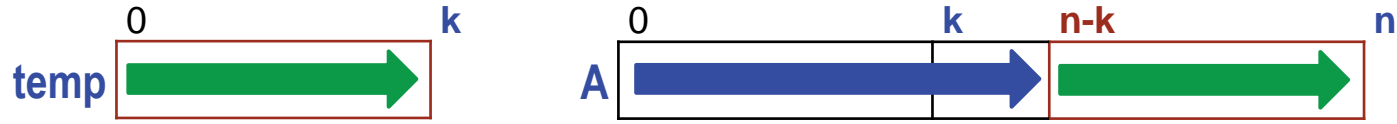
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int temp[] = new int[k];
/* temp[0..k-1] = A[0..k-1]; */
for (int j=0; j<k; j++) temp[j] = A[j];
LeftShiftK(A, n, k);
/* A[___..n-1] = temp[0..k-1]; */
for (int j=0; j<k; j++) A[___] = temp[j];

```

---

 **Avoid gratuitous differences in code. Reuse code patterns, if possible.**

---

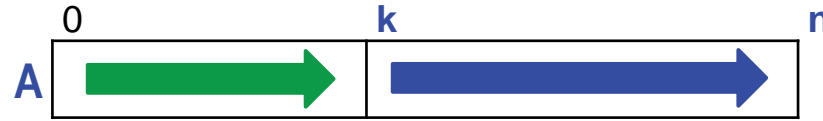


**Approach 2:** k-wide generalization of swap.

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int temp[] = new int[k];
/* temp[0..k-1] = A[0..k-1]; */
for (int j=0; j<k; j++) temp[j] = A[j];
LeftShiftK(A, n, k);
/* A[n-k..n-1] = temp[0..k-1]; */
for (int j=0; j<k; j++) A[n-k+j] = temp[j];

```



**Approach 3: Three Flips.** Consider the two parts of the array.

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
k places, with values originally in A[0..k-1] reentering at right. */
```



**Approach 3:** Represent the values in those parts as **green** and **blue** arrows.

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
k places, with values originally in A[0..k-1] reentering at right. */
```



### Approach 3: Reverse first k

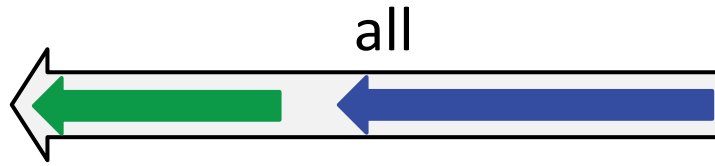
```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
k places, with values originally in A[0..k-1] reentering at right. */  
Reverse(A, 0, k-1);
```





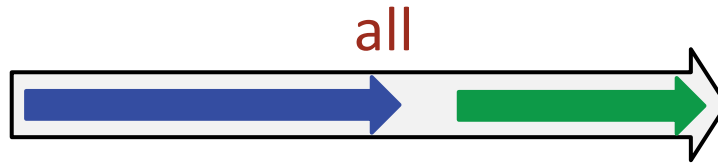
**Approach 3:** Reverse first k, then rest of elements

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
Reverse(A, 0, k-1);
Reverse(A, k, n-1);
```



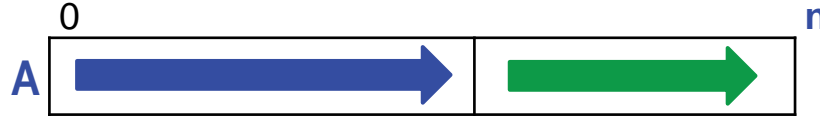
**Approach 3:** Reverse first k, then rest of elements

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]  
k places, with values originally in A[0..k-1] reentering at right. */  
Reverse(A, 0, k-1);  
Reverse(A, k, n-1);
```



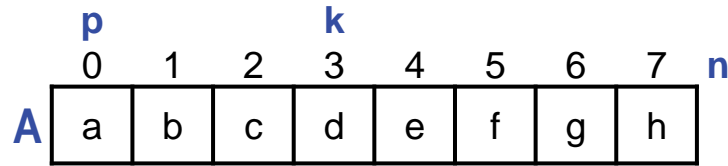
**Approach 3:** Reverse first k, then rest of elements, **then all elements.**

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
Reverse(A, 0, k-1);
Reverse(A, k, n-1);
Reverse(A, 0, n-1);
```



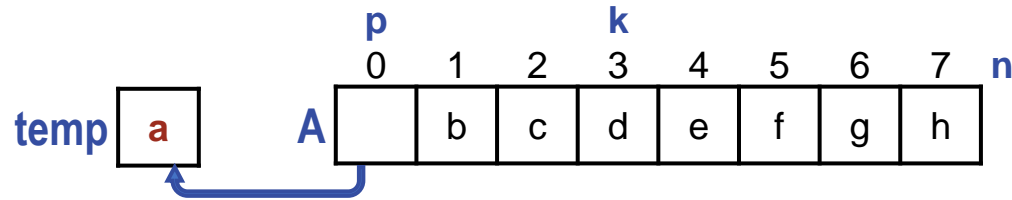
**Approach 3:** Reverse first k, then rest of elements, then all elements.

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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Reverse(A, 0, k-1);
Reverse(A, k, n-1);
Reverse(A, 0, n-1);
```



**Approach 4: Juggle elements in a stride of k.**

```
/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
   k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
```

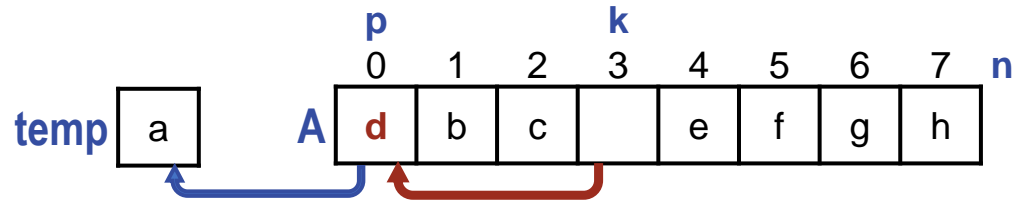


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
int temp = A[0];     // and make a hole there.

```

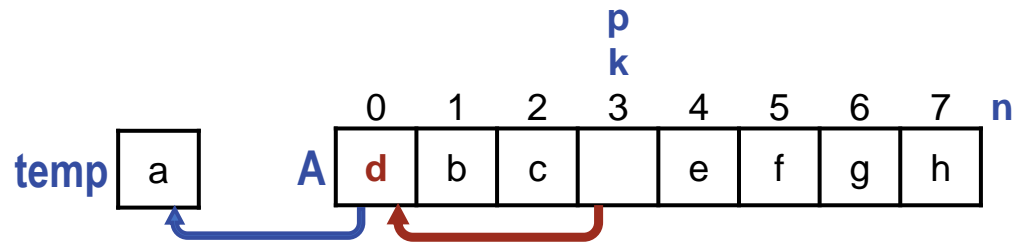


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
int temp = A[0];    // and make a hole there.
while ( _____ ) {
    A[p] = A[ p+k ]; // Fill hole at p, making a new hole.
}

```



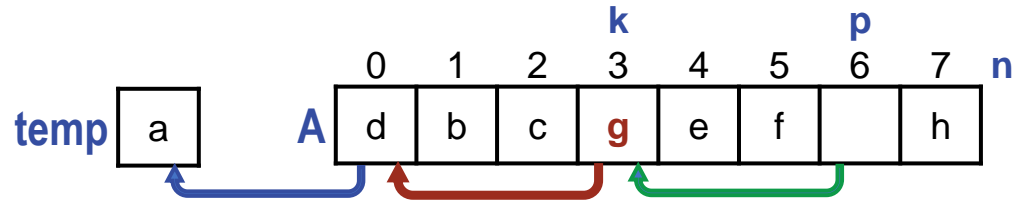
**Approach 4:** Juggle elements in a stride of  $k$ .

```

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int temp = A[0];    // and make a hole there.
while ( _____ ) {
    A[p] = A[ p+k ]; // Fill hole at p, making a new hole.
    p = p+k;        // Advance to the new hole.
}

```



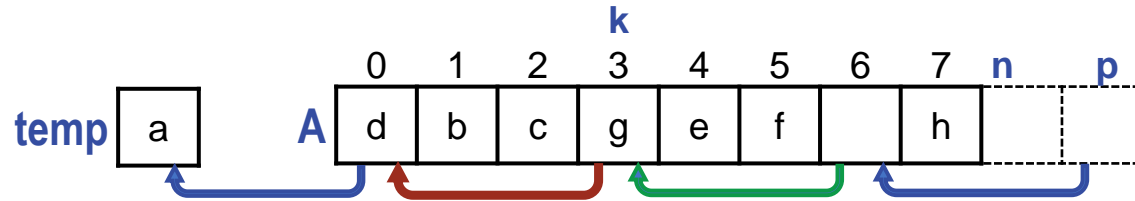


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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while ( _____ ) {
    A[p] = A[ p+k ]; // Fill hole at p, making a new hole.
    p = p+k;        // Advance to the new hole.
}

```

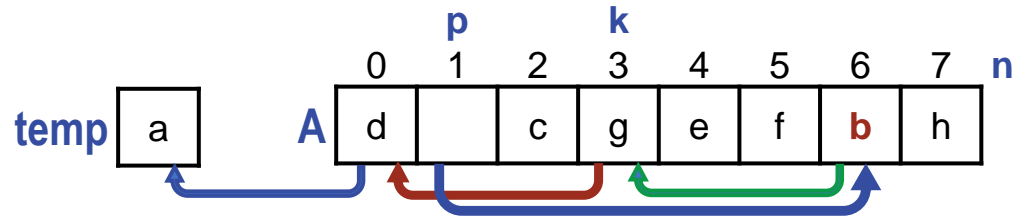


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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while ( _____ ) {
    A[p] = A[ p+k ]; // Fill hole at p, making a new hole.
    p = p+k;        // Advance to the new hole.
}

```

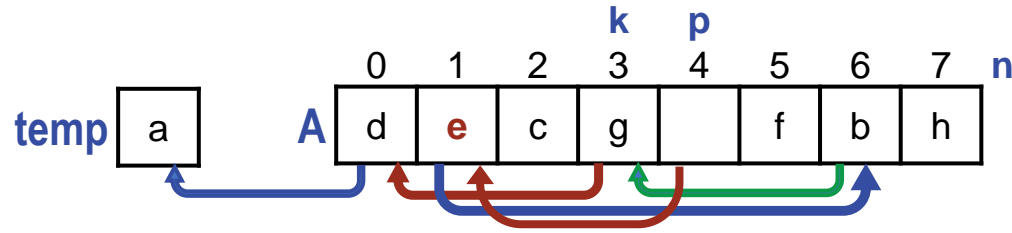


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
int temp = A[0];     // and make a hole there.
while ( _____ ) {
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}

```

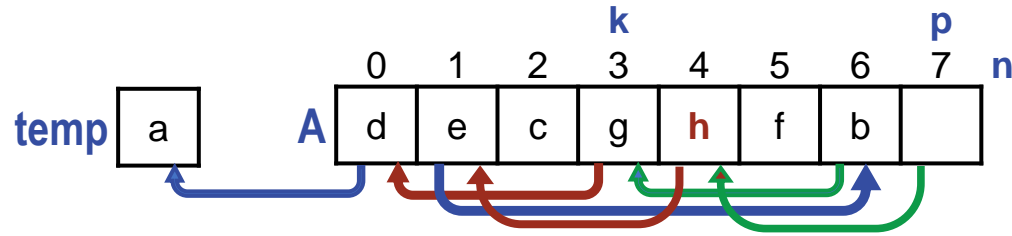


**Approach 4:** Juggle elements in a stride of  $k$ .

```

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    p = (p+k)%n;      // Advance to the new hole.
}

```

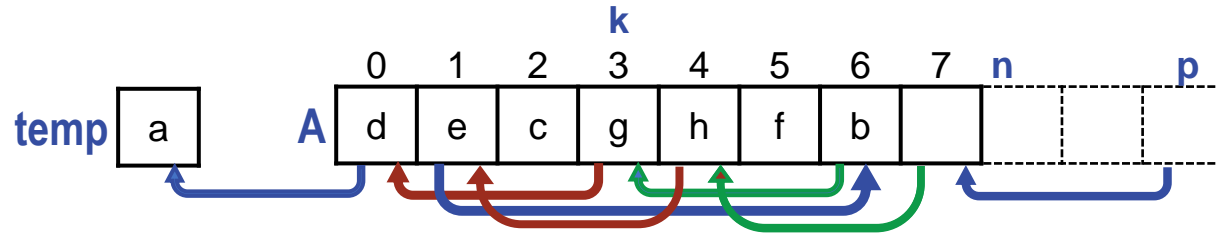


**Approach 4:** Juggle elements in a stride of  $k$ .

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    p = (p+k)%n;      // Advance to the new hole.
}

```

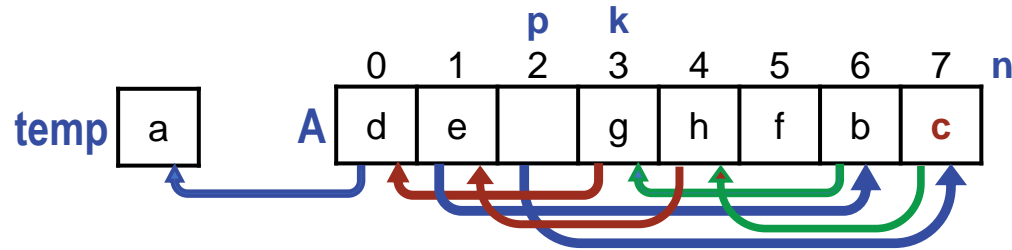


**Approach 4:** Juggle elements in a stride of  $k$ .

```

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    p = (p+k)%n;      // Advance to the new hole.
}

```

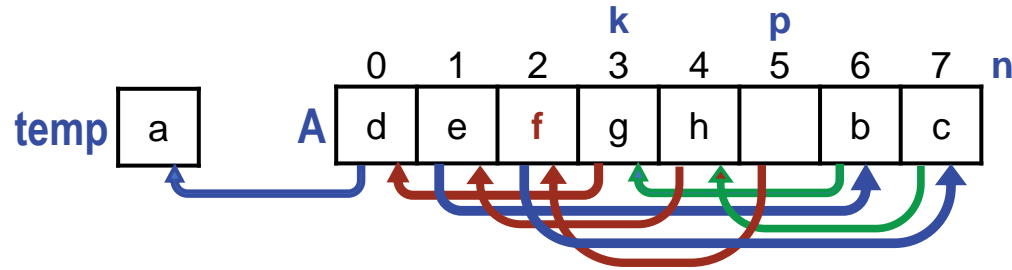


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
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while ( _____ ) {
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}

```



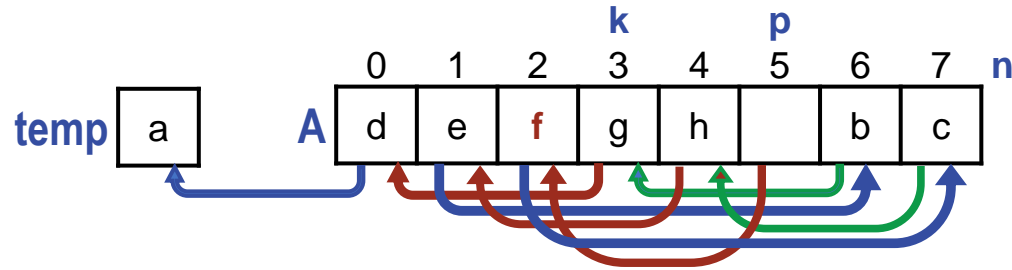
**Approach 4:** Juggle elements in a stride of  $k$ .

```

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    p = (p+k)%n;      // Advance to the new hole.
}

```



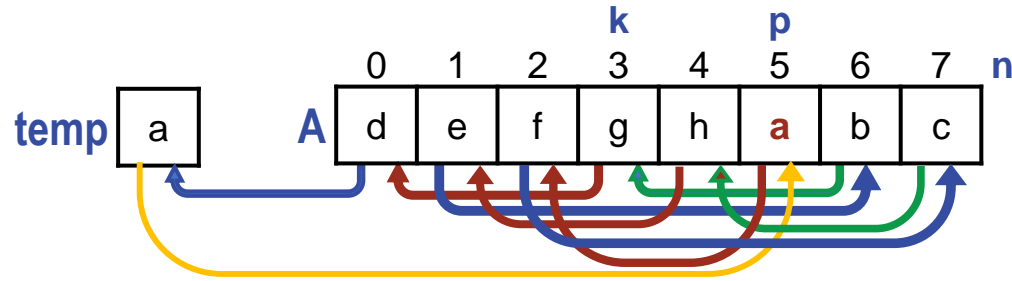


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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int temp = A[0];    // and make a hole there.
while ( (p+k)%n!=0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}

```

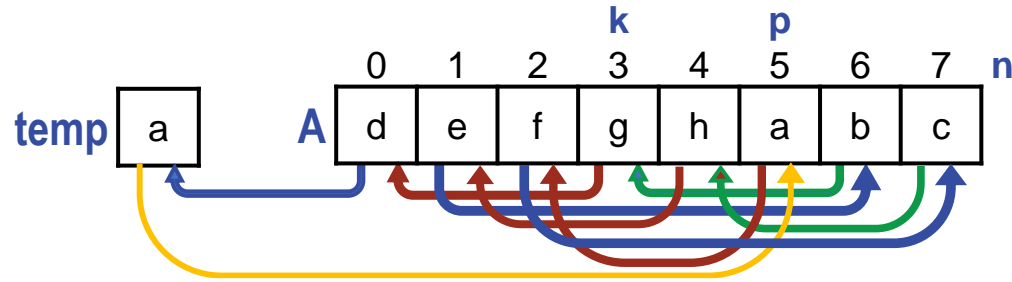


**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
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while ( (p+k)%n != 0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}
A[p] = temp;        // Fill the last hole from temp.

```



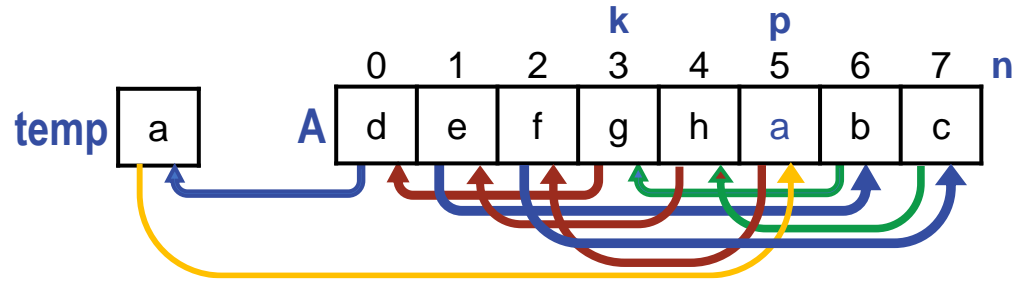
**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
int temp = A[0];    // and make a hole there.
while ( (p+k)%n != 0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}
A[p] = temp;        // Fill the last hole from temp.

```

Are we done?



**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
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while ( (p+k)%n != 0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}
A[p] = temp;        // Fill the last hole from temp.

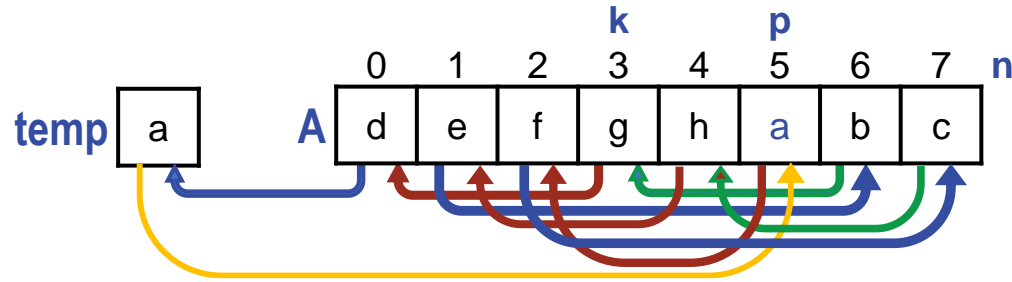
```

---

 **Beware of premature self-satisfaction.**

---

Are we done?



**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
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}
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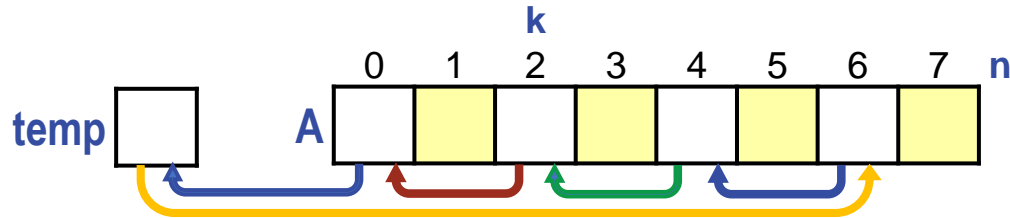
```

---

👉 **Validate output thoroughly.**

---

Are we done?



**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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int temp = A[0];     // and make a hole there.
while ( (p+k)%n != 0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
    p = (p+k)%n;      // Advance to the new hole.
}
A[p] = temp;        // Fill the last hole from temp.

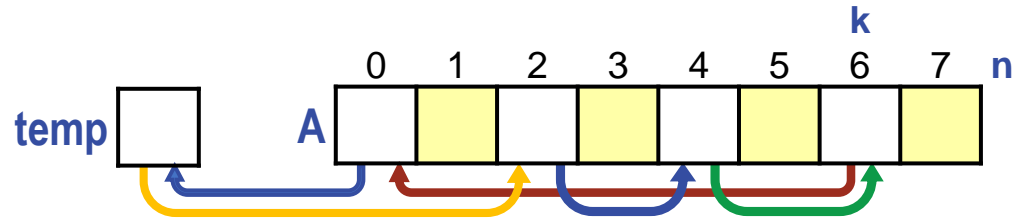
```

---

👉 **Validate output thoroughly.**

---

Are we done? Hardly.



**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int p = 0;           // Start at A[0]
int temp = A[0];     // and make a hole there.
while ( (p+k)%n != 0 ) { // Stop if p is about to be 0 again.
    A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
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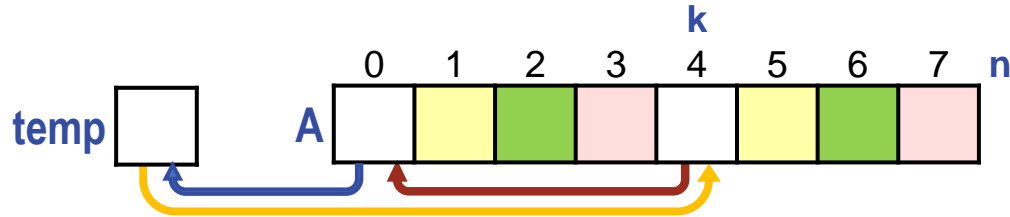
```

---

👉 **Validate output thoroughly.**

---

Are we done? Hardly.



**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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```

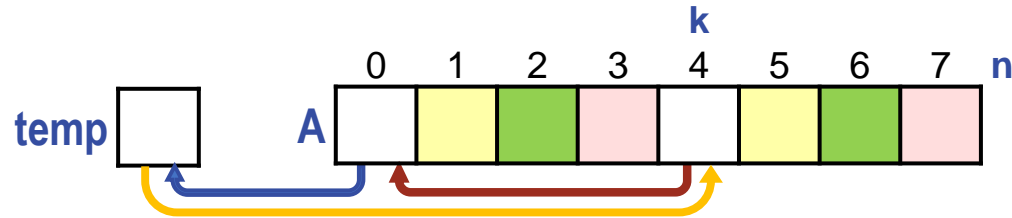
---

👉 **Validate output thoroughly.**

---

Are we done? Hardly. It only works if  $k$  and  $n$  are relatively prime!





**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
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    p = (p+k)%n;      // Advance to the new hole.
}
A[p] = temp;        // Fill the last hole from temp.

```

---

👉 **Validate output thoroughly.**

---

Are we done? Hardly. It only works if  $k$  and  $n$  are relatively prime! Now what?

Abandon, or learn that each of  $A[0..gcd(k, n) - 1]$  begins a disjoint cycle.

**Approach 4:** Juggle elements in a stride of  $k$ .

```

/* Given int array A[0..n-1], and integer k,  $0 \leq k < n$ , left shift A[k..n-1]
k places, with values originally in A[0..k-1] reentering at right. */
int g = gcd(n,k);
for (int j=0; j<g; j++) {
    int p = j;           // Start at A[j]
    int temp = A[p];    // and make a hole there.
    while ( (p+k)%n!=j ) { // Stop if p is about to be j again.
        A[p] = A[(p+k)%n]; // Fill hole at p, making a new hole.
        p = (p+k)%n;      // Advance to the new hole.
    }
    A[p] = temp; // Fill the last hole from temp.
}


```

## Assessment:

	Version of Left-Rotate-k	#moves	Explanation
☞	Repeated Left-Rotate-1	$k \cdot n$	Each Left-Rotate-1 moves all $n$ elements. Done $k$ times.
	Swap Generalization	$n+k$	The copies into and out from temp do $2 \cdot k$ moves, and the shift does $n-k$ moves.
	Three Flips	$2 \cdot n$	Each element moves once during the 1st two reverses, and then again for the 3rd reverse.
	Juggle in Cycles	$n + \text{gcd}(n, k)$	Each element moves once, plus the first element of each of the $\text{gcd}(n, k)$ cycles must first be saved in temp.


Worst #moves, by far. Easiest to understand.

**Assessment:**

Version of Left-Rotate-k	#moves	Explanation
Repeated Left-Rotate-1	$k \cdot n$	Each Left-Rotate-1 moves all $n$ elements. Done $k$ times.
 Swap Generalization	$n+k$	The copies into and out from temp do $2 \cdot k$ moves, and the shift does $n-k$ moves.
Three Flips	$2 \cdot n$	Each element moves once during the 1st two reverses, and then again for the 3rd reverse.
Juggle in Cycles	$n + \text{gcd}(n, k)$	Each element moves once, plus the first element of each of the $\text{gcd}(n, k)$ cycles must first be saved in temp.


Reasonable #moves but not *in situ*, i.e., requires extra space for temp.

**Assessment:**

Version of Left-Rotate-k	#moves	Explanation
Repeated Left-Rotate-1	$k \cdot n$	Each Left-Rotate-1 moves all $n$ elements. Done $k$ times.
Swap Generalization	$n+k$	The copies into and out from temp do $2 \cdot k$ moves, and the shift does $n-k$ moves.
 Three Flips	$2 \cdot n$	Each element moves once during the 1st two reverses, and then again for the 3rd reverse.
Juggle in Cycles	$n + \text{gcd}(n, k)$	Each element moves once, plus the first element of each of the $\text{gcd}(n, k)$ cycles must first be saved in temp.

Reasonable #moves. Good locality.

**Assessment:**

Version of Left-Rotate-k	#moves	Explanation
Repeated Left-Rotate-1	$k \cdot n$	Each Left-Rotate-1 moves all $n$ elements. Done $k$ times.
Swap Generalization	$n+k$	The copies into and out from temp do $2 \cdot k$ moves, and the shift does $n-k$ moves.
Three Flips	$2 \cdot n$	Each element moves once during the 1st two reverses, and then again for the 3rd reverse.
 Juggle in Cycles	$n + \text{gcd}(n, k)$	Each element moves once, plus the first element of each of the $\text{gcd}(n, k)$ cycles must first be saved in temp.

Hardest to understand. Poor locality.

**Assessment:**

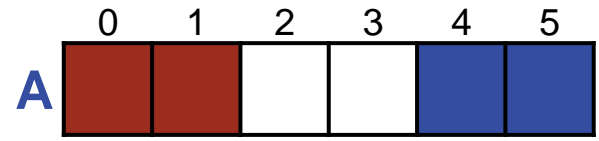
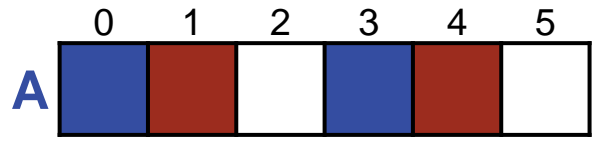
Version of Left-Rotate-k	#moves	Explanation
Repeated Left-Rotate-1	$k \cdot n$	Each Left-Rotate-1 moves all $n$ elements. Done $k$ times.
Swap Generalization	$n+k$	The copies into and out from temp do $2 \cdot k$ moves, and the shift does $n-k$ moves.
✓ Three Flips	$2 \cdot n$	Each element moves once during the 1st two reverses, and then again for the 3rd reverse.
Juggle in Cycles	$n + \text{gcd}(n, k)$	Each element moves once, plus the first element of each of the $\text{gcd}(n, k)$ cycles must first be saved in temp.

Personal favorite, and really elegant!

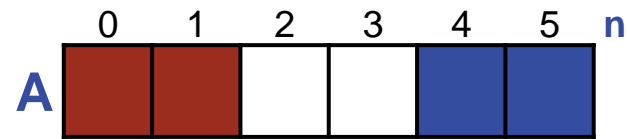
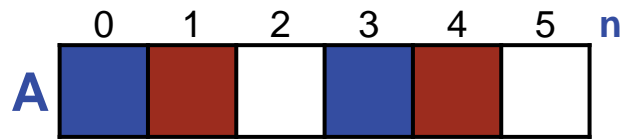
**New Application:** The Dutch National Flag problem.







**Application:** Rearrange an array into all red, then all white, then all blue.



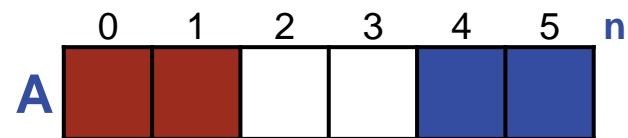
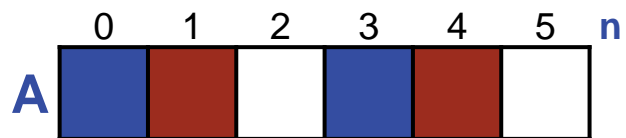
**Application:** Rearrange an array into all red, then all white, then all blue.

```
/* Given array A[0..n-1] consisting of only three values (red, white,  
and blue), rearrange A into all red, then white, then blue. */
```

---

 **A statement-comment says exactly what code must accomplish, not how it does so.**

---



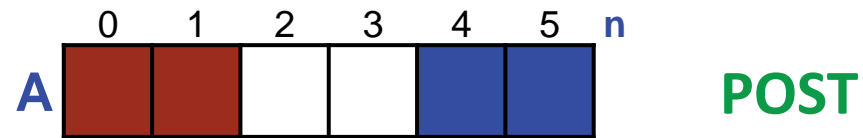
**Application:** Rearrange an array into all red, then all white, then all blue.

```
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
while ( _____ ) _____
```

---

👉 If you “smell a loop”, write it down.

---



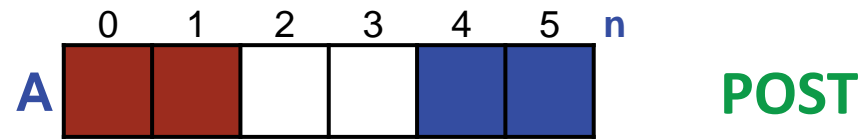
**Application:** Rearrange an array into all red, then all white, then all blue.

```
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
while ( _____ ) _____
```

---

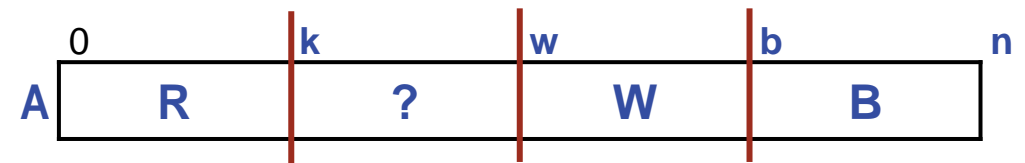
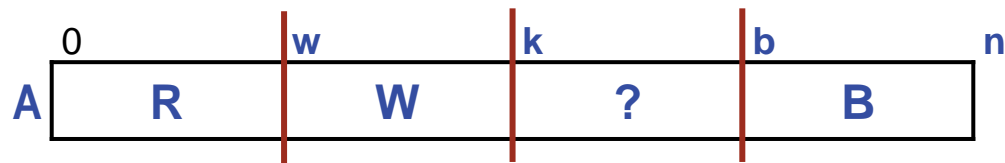
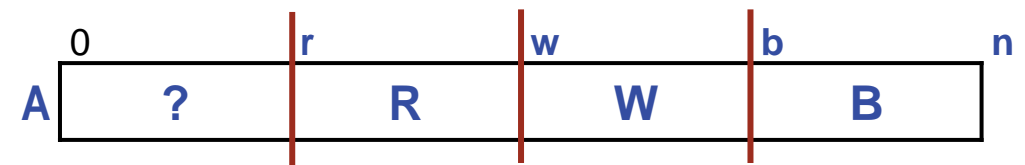
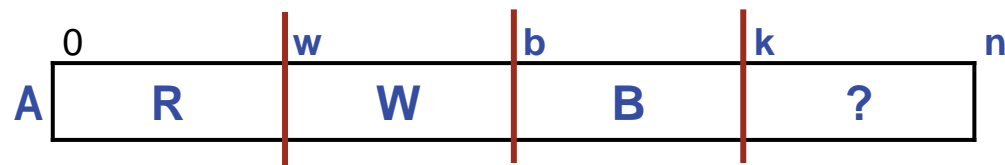
👉 To get to **POST** iteratively, choose a **weakened POST** as **INVARIANT**.

---



**Application:** Rearrange an array into all red, then all white, then all blue.

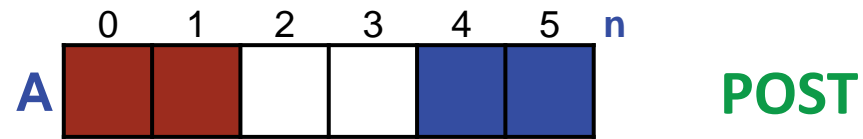
Here are four choices:




---

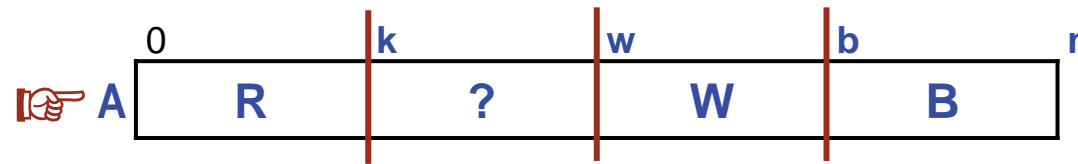
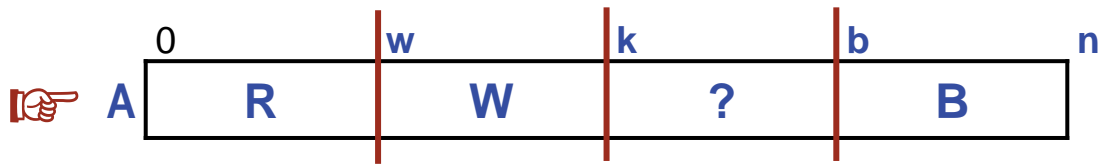
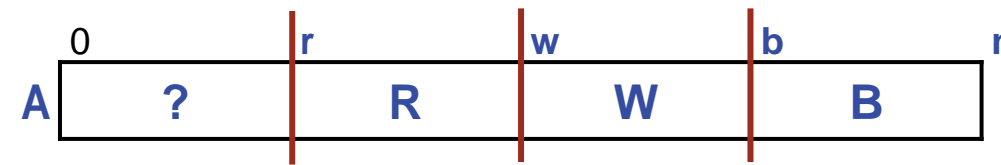
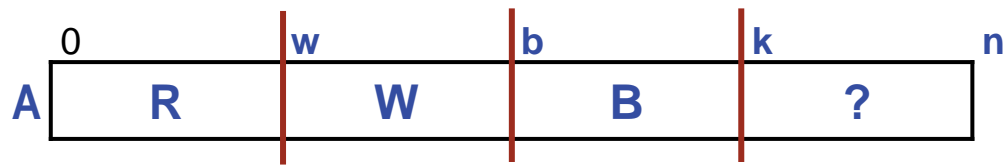
👉 To get to **POST** iteratively, choose a **weakened POST** as **INVARIANT**.

---

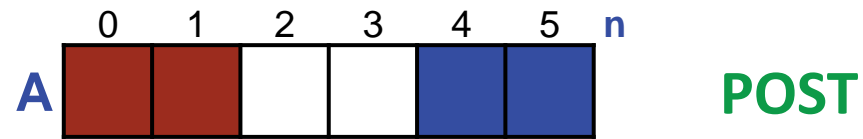


**Application:** Rearrange an array into all red, then all white, then all blue.

Here are four choices: How shall we choose?

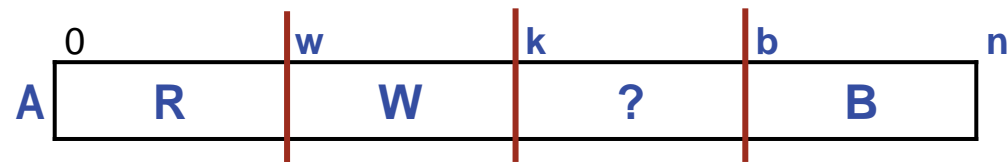
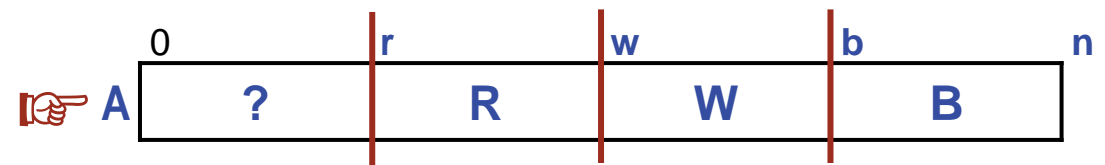
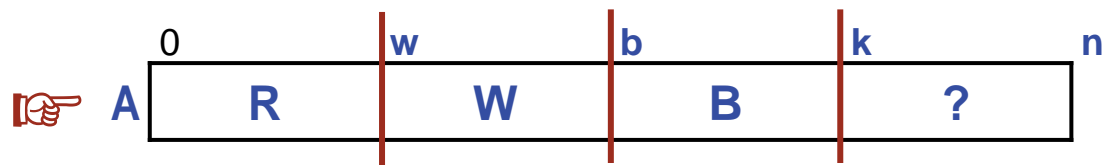


Symmetric, so discard one arbitrarily.

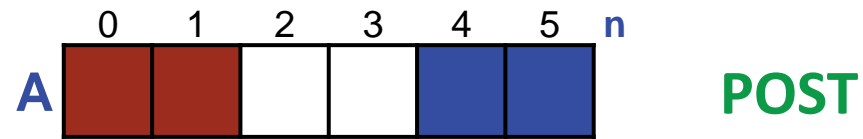


**Application:** Rearrange an array into all red, then all white, then all blue.

Here are four choices: How shall we choose?

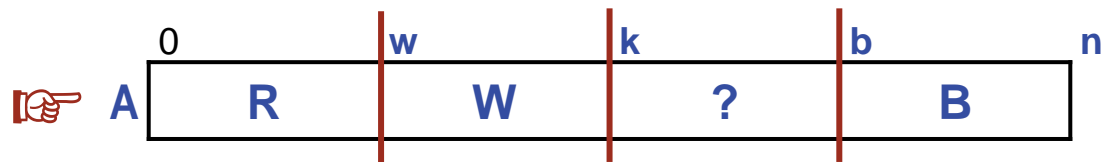
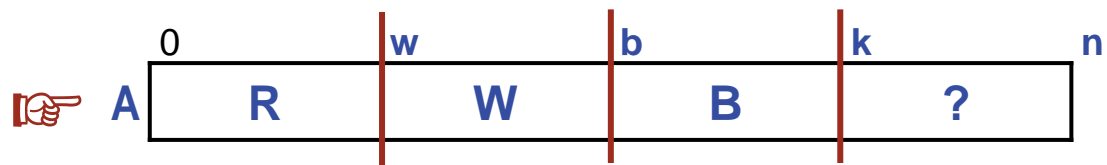


Left more intuitive, because the ? region seems more familiar, so discard right.



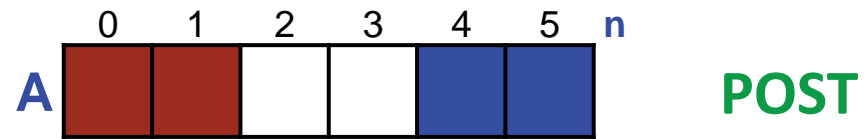
**Application:** Rearrange an array into all red, then all white, then all blue.

Here are four choices: How shall we choose?



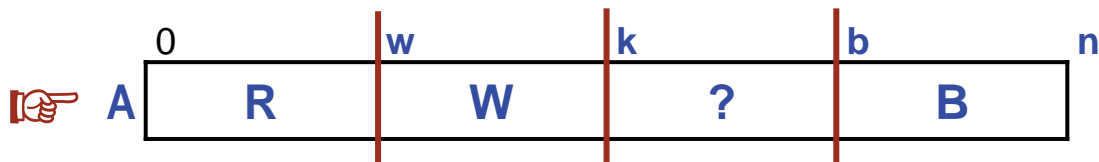
? region of top has only one degree of freedom, but bottom has two. Discard top.



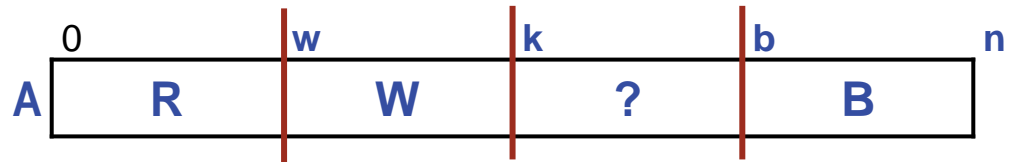


**Application:** Rearrange an array into all red, then all white, then all blue.

Here are four choices: How shall we choose?



This will be our **INVARIANT**.



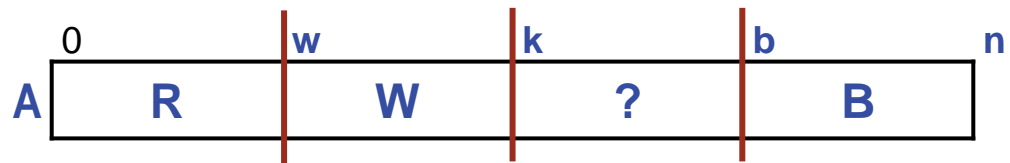
VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( _____ ) _____

```



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

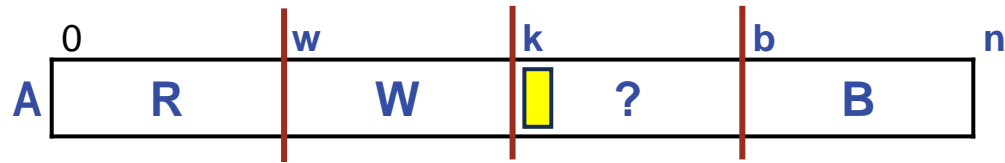
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( _____ ) _____

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

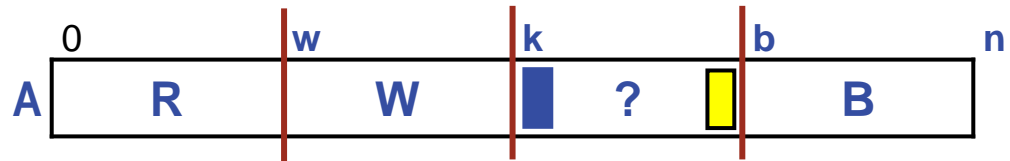
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( ____ )
   if ( A[k] == ____ ) ____
   else if ( A[k] == ____ ) ____
   else ____

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

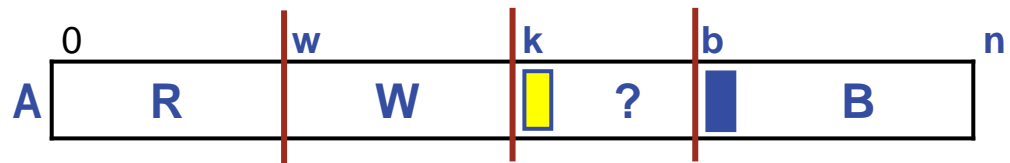
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( ____ )
   if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */  b--; }
   else if ( A[k] == ____ ) ____
   else ____

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

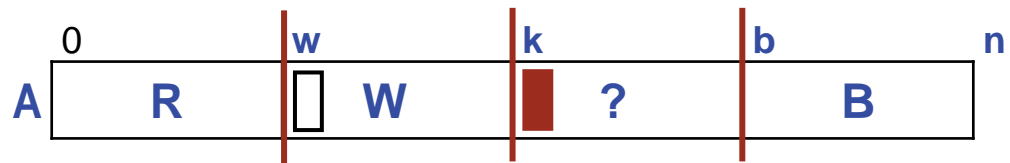
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( ____ )
   if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */ b--; }
   else if ( A[k] == ____ ) ____
   else ____

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

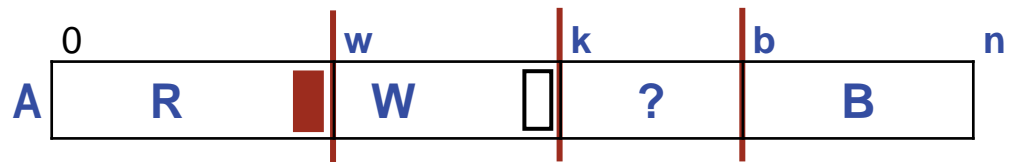
/* Given array A[0..n-1] consisting of only three values (red, white,
   and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
   int k = ____; int w = ____; int b = ____;
while ( ____ )
   if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */ b--; }
   else if ( A[k]==R ) { /* Swap A[w] and A[k]. */ w++; k++; }
   else ____

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
int k = ____; int w = ____; int b = ____;
while ( _____ )
    if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */ b--; }
    else if ( A[k]==R ) { /* Swap A[w] and A[k]. */ w++; k++; }
    else _____

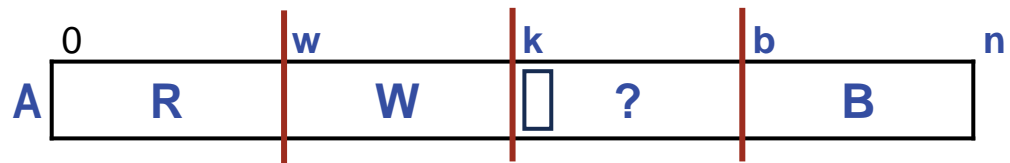
```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---





VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

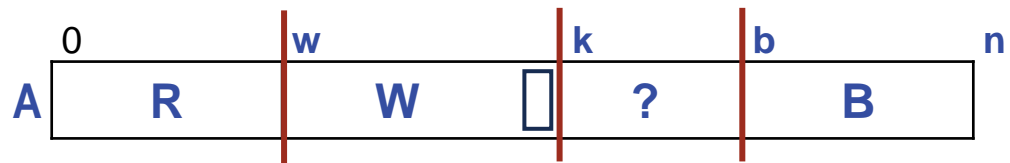
/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
int k = ____; int w = ____; int b = ____;
while ( _____ )
    if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */ b--; }
    else if ( A[k]==R ) { /* Swap A[w] and A[k]. */ w++; k++; }
    else /* A[k]==W */ k++;

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

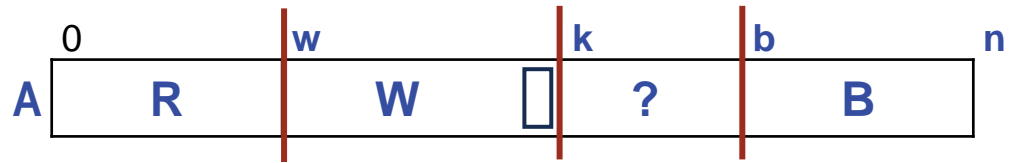
/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
  int k = ____; int w = ____; int b = ____;
while ( _____ )
  if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */  b--; }
  else if ( A[k]==R ) { /* Swap A[w] and A[k]. */  w++; k++; }
  else /* A[k]==W */ k++;

```

---

👉 **A Case Analysis in the loop body is often needed for characterizing different ways in which to decrease the loop variant while maintaining the loop invariant.**

---



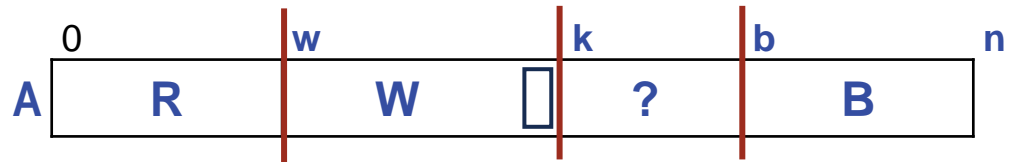
VARIANT: b-k

**Application:** Rearrange an array into all red, then all white, then all blue.

```

/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
  int k = ____; int w = ____; int b = ____;
while ( k!=b )
  if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */  b--; }
  else if ( A[k]==R ) { /* Swap A[w] and A[k]. */  w++; k++; }
  else /* A[k]==W */ k++;

```



VARIANT:  $b-k$

**Application:** Rearrange an array into all red, then all white, then all blue.

```

/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
int k = 0; int w = 0; int b = n;
while ( k!=b )
    if ( A[k]==B ) { /* Swap A[b-1] and A[k]. */ b--; }
    else if ( A[k]==R ) { /* Swap A[w] and A[k]. */ w++; k++; }
    else /* A[k]==W */ k++;

```

**Application:** Rearrange an array into all red, then all white, then all blue.

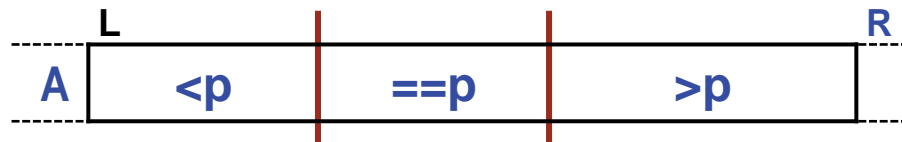
```
/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
  int k = 0; int w = 0; int b = n;
while ( k!=b )
  if ( A[k]==B ) {
    /* Swap A[b-1] and A[k]. */
    b--;
  }
  else if ( A[k]==R ) {
    /* Swap A[w] and A[k]. */
    w++; k++;
  }
  else /* A[k]==W */ k++;
```

**Application:** Rearrange an array into all red, then all white, then all blue.

```
/* Given array A[0..n-1] consisting of only three values (red, white,
and blue), rearrange A into all red, then white, then blue. */
/* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for  $0 \leq w \leq k \leq b \leq n$ . */
  int k = 0; int w = 0; int b = n;
while ( k!=b )
  if ( A[k]==B ) {
    /* Swap A[b-1] and A[k]. */
    int temp = A[b-1]; A[b-1] = A[k]; A[k] = temp;
    b--;
  }
  else if ( A[k]==R ) {
    /* Swap A[w] and A[k]. */
    int temp = A[k]; A[k] = A[w]; A[w] = temp;
    w++; k++;
  }
  else /* A[k]==W */ k++;
```

## Performance:

- Constant work per iteration.
- Variant reduced by 1 on each iteration.
- Thus, **running time linear in  $n$ .**



**New Application:** Rearrange (a segment of) an array into <p, ==p, and >p.

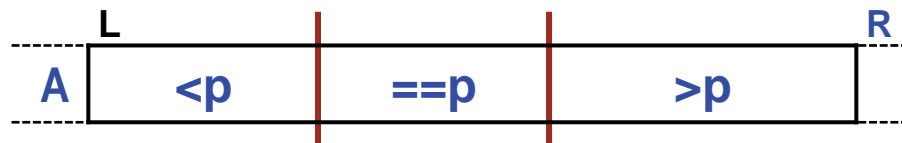
```
/* Rearrange A[L..R-1] into all <p, then all ==p, then all >p. */  
static void Partition( int A[], int L, int R, int p ) {  
    } /* Partition */
```

---

 A header-comment says exactly what a method must accomplish, not how it does so.

---





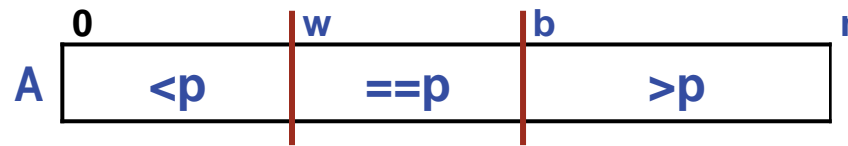
**Application:** Rearrange (a segment of) an array into  $<p$ ,  $=p$ , and  $>p$ .

```
/* Rearrange A[L..R-1] into all <p, then all ==p, then all >p. */  
static void Partition( int A[], int L, int R, int p ) {  
    <body of Dutch National Flag problem>  
} /* Partition */
```

---

 **Don't type if you can avoid it; clone. Cut and paste, then adapt.**

---



**Application:** Rearrange (a segment of) an array into <p, ==p, and >p.

```

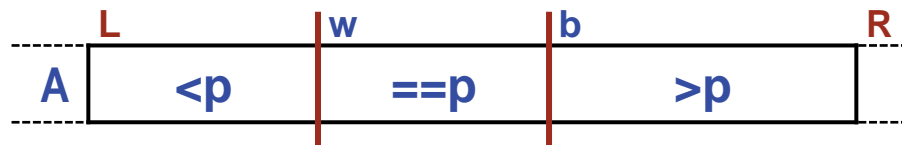
/* Rearrange A[L..R-1] into all <p, then all ==p, then all >p. */
static void Partition( int A[], int L, int R, int p ) {
    /* A[0..w-1] red, A[w..k-1] white, A[b..n-1] blue, for 0≤w≤k≤b≤n. */
    int k = 0; int w = 0; int b = n;
    while ( k!=b )
        if ( A[k]==B ) {
            /* Swap A[b-1] and A[k]. */
            int temp = A[b-1]; A[b-1] = A[k]; A[k] = temp;
            b--;
        }
        else if ( A[k]==R ) {
            /* Swap A[w] and A[k]. */
            int temp = A[k]; A[k] = A[w]; A[w] = temp;
            w++; k++;
        }
        else /* A[k]==W */ k++;
    } /* Partition */

```

Dutch National Flag problem

**Application:** Rearrange (a segment of) an array into  $<p$ ,  $=p$ , and  $>p$ .

```
/* Rearrange A[L..R-1] into all <p, then all ==p, then all >p. */
static void Partition( int A[], int L, int R, int p ) {
    /* A[L..w-1]<p, A[w..k-1]==p, A[b..R-1]>p, for L≤w≤k≤b≤R. */
    int k = L; int w = L; int b = R;
    while ( k!=b )
        if ( A[k]>p ) {
            /* Swap A[b-1] and A[k]. */
            int temp = A[b-1]; A[b-1] = A[k]; A[k] = temp;
            b--;
        }
        else if ( A[k]<p ) {
            /* Swap A[w] and A[k]. */
            int temp = A[k]; A[k] = A[w]; A[w] = temp;
            w++; k++;
        }
        else /* A[k]==p */ k++;
    } /* Partition */
```



**Application:** Rearrange (a segment of) an array into  $<p$ ,  $=p$ , and  $>p$ .

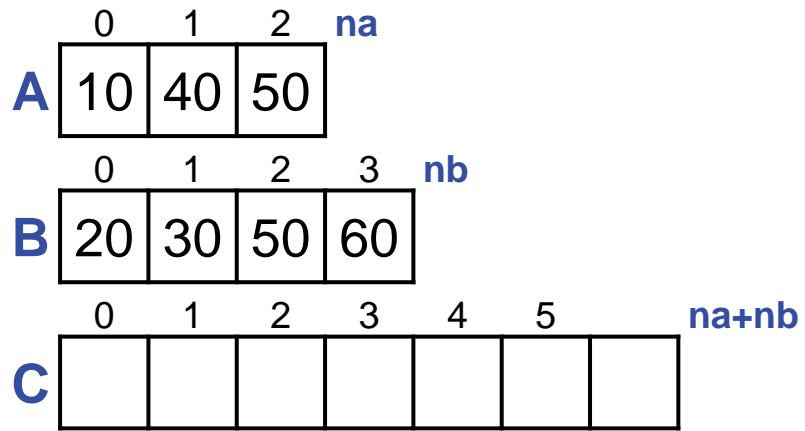
What value of  $p$  would tend to create “ $<p$ ” and “ $>p$ ” regions of near equal size?

If  $A[L..R-1]$  are in random order, any of those values is equally good for  $p$ .

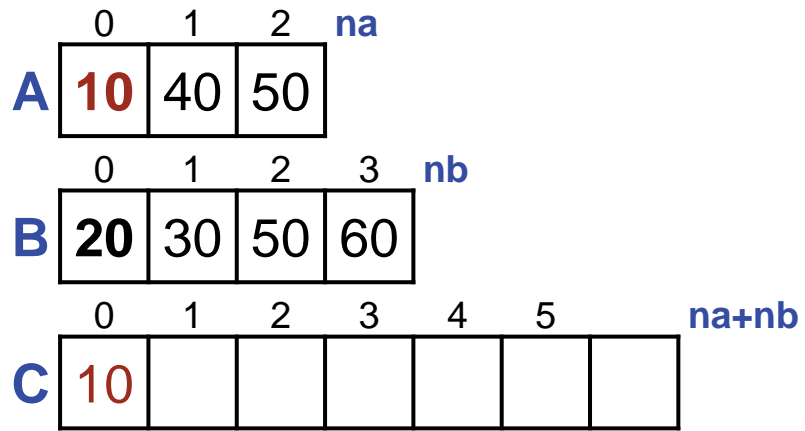
Choosing a  $p$  of  $(A[L]+A[R-1])/2$  guards against poor performance when  $A[L..R-1]$  is already ordered.

This is central to a Divide and Conquer approach to sorting called QuickSort.

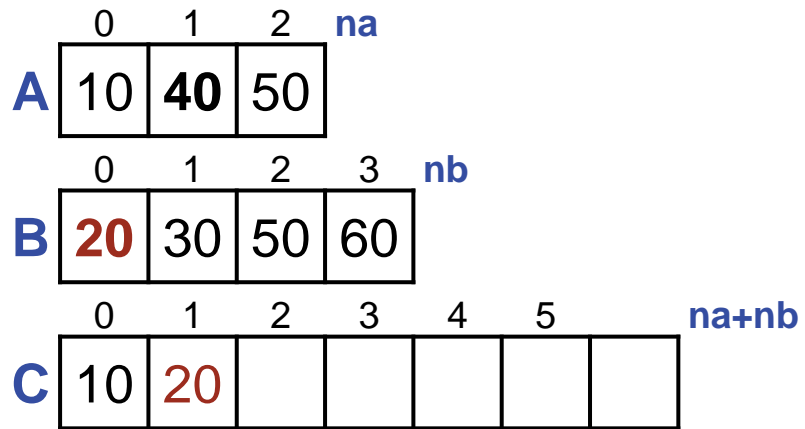
**New Application:** Collate ordered arrays A and B into array C.



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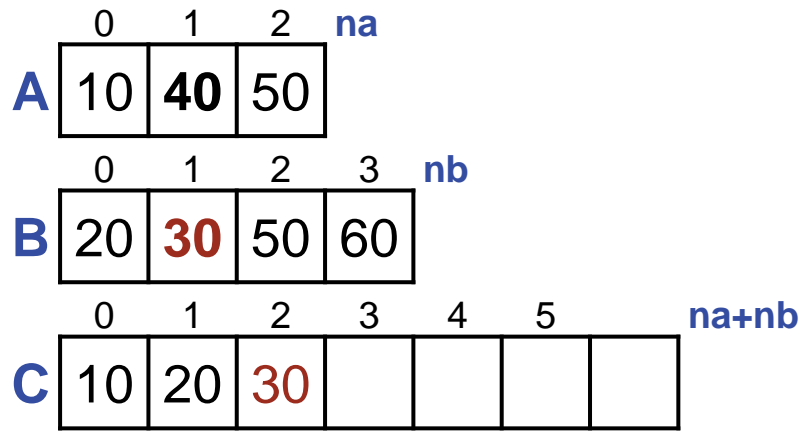


**Application:** Collate ordered arrays A and B into array C.

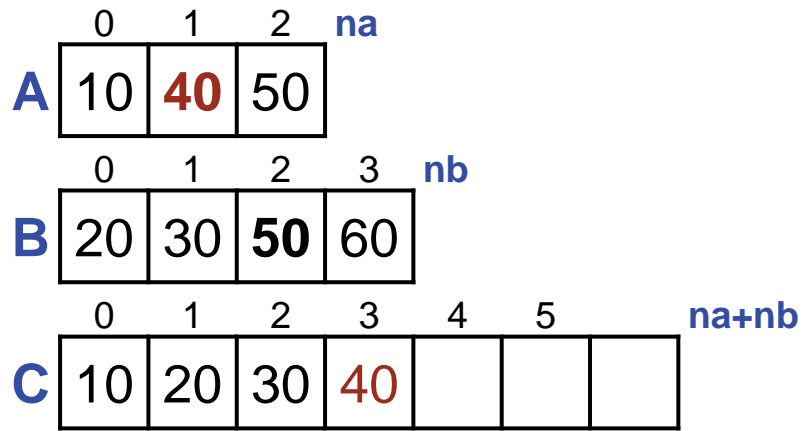


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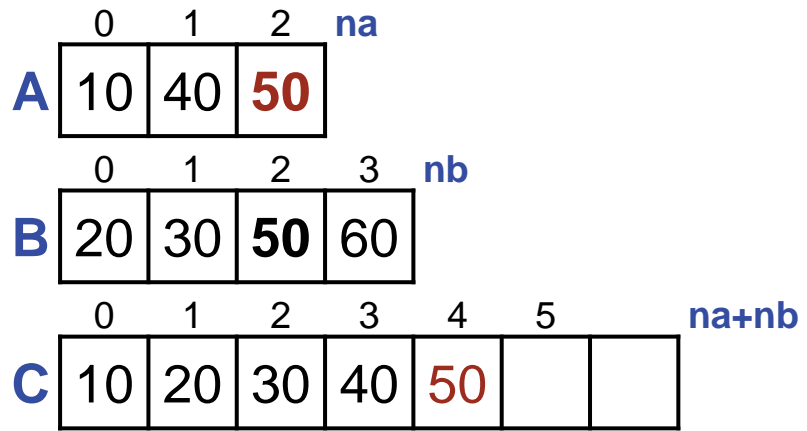




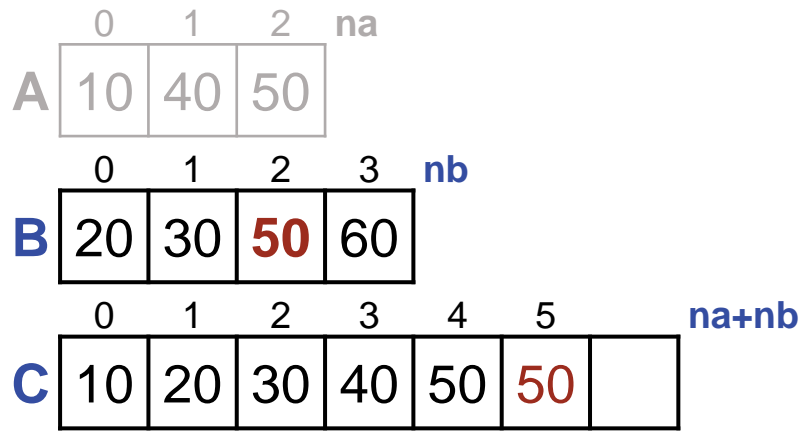
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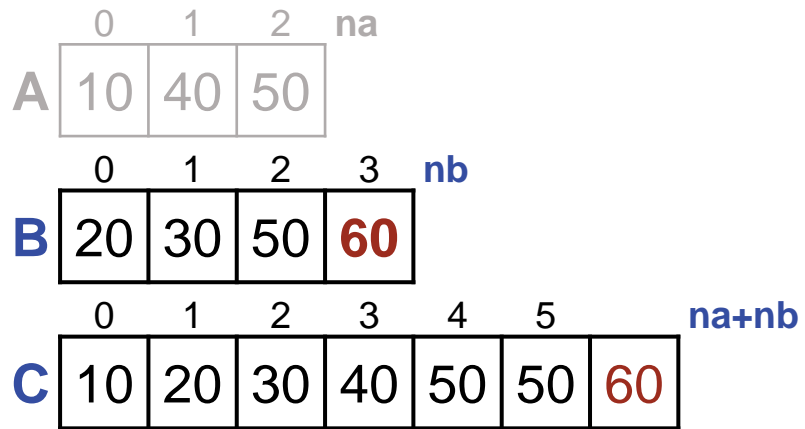
**Application:** Collate ordered arrays A and B into array C.



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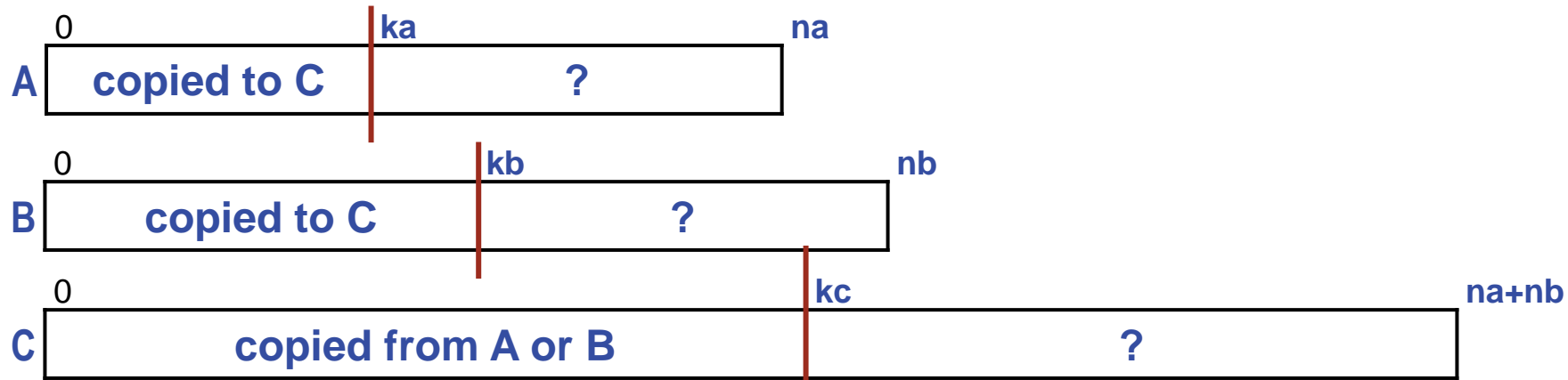


**Application:** Collate ordered arrays A and B into array C.



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Collation is central to a Divide and Conquer approach to sorting called MergeSort.



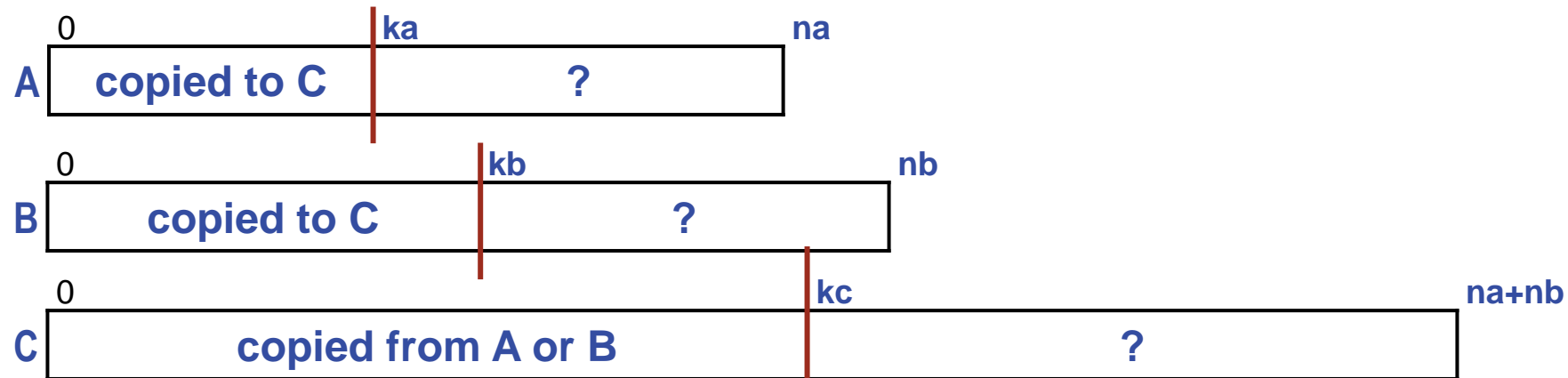
**Application:** Collate ordered arrays A and B into array C.

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/* Given ordered arrays A and B of lengths na and nb, create ordered
   array C of length na+nb consisting of those values. */
   int C[] = new int[na+nb];           // C[0..kc-1] is collation of
                                       // A[0..ka-1] and B[0..kb-1].

   int ka = __; int kb = __; int kc = __; // Indices in A, B, and C.
/* Copy values from A or B into C until one array is exhausted. */
/* Copy remaining values into C from the unexhausted array, A or B. */

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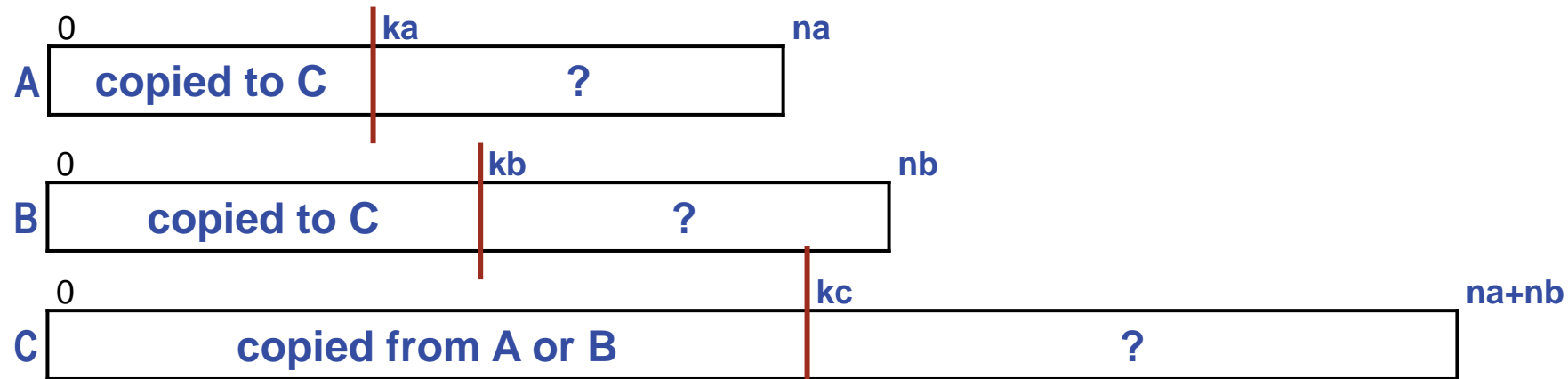
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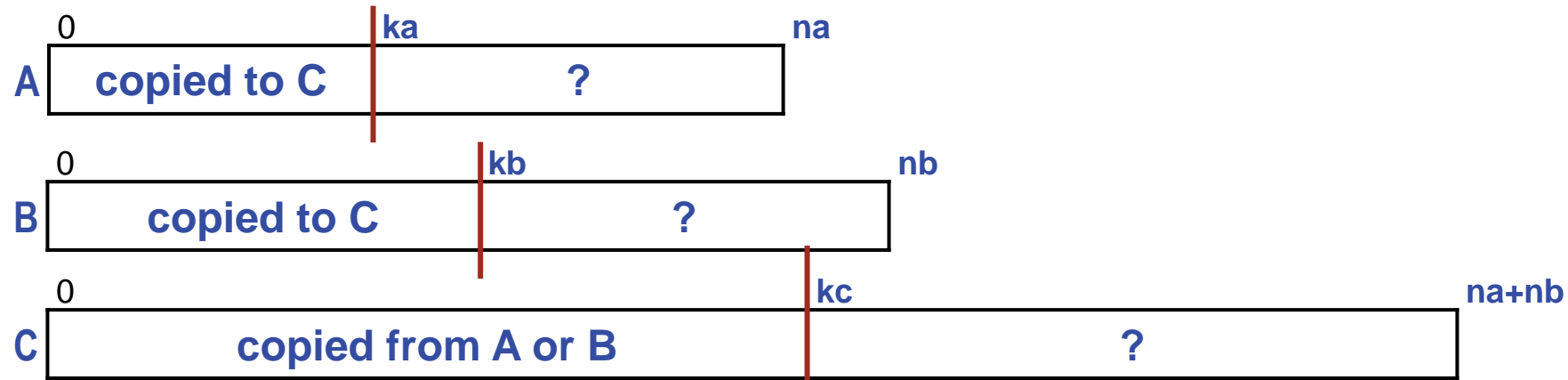
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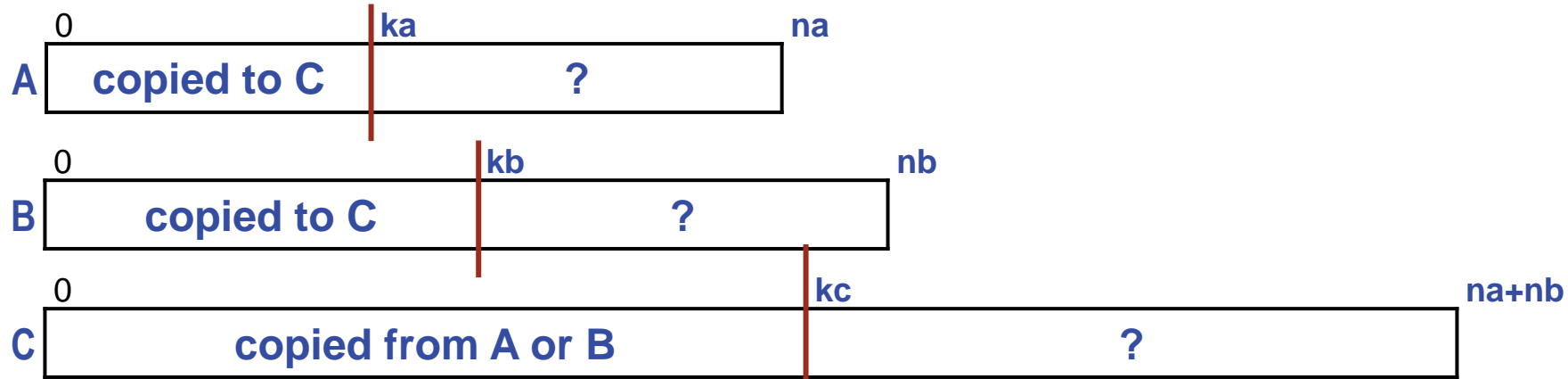
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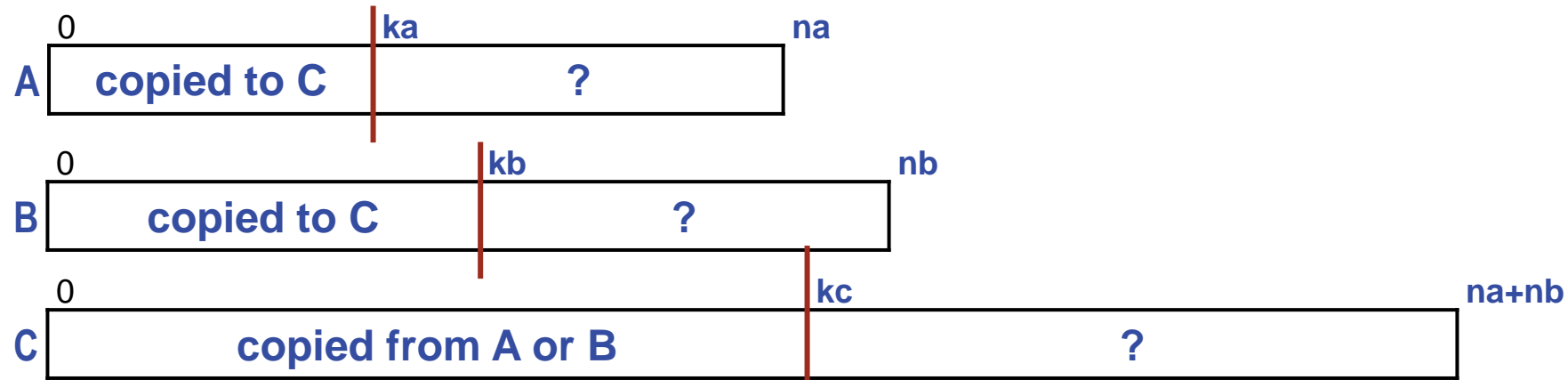
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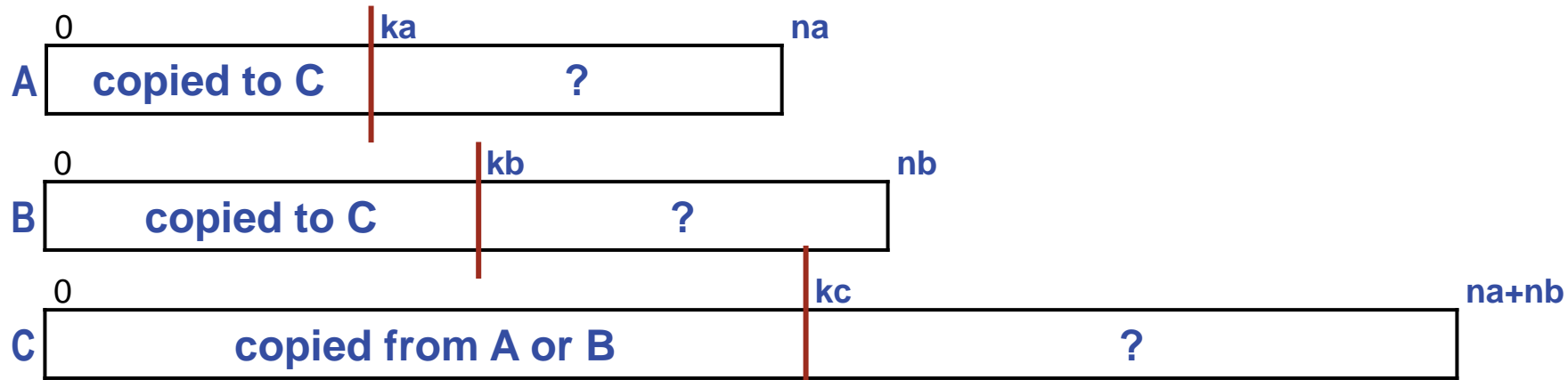
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while ( ka < na && kb < nb )
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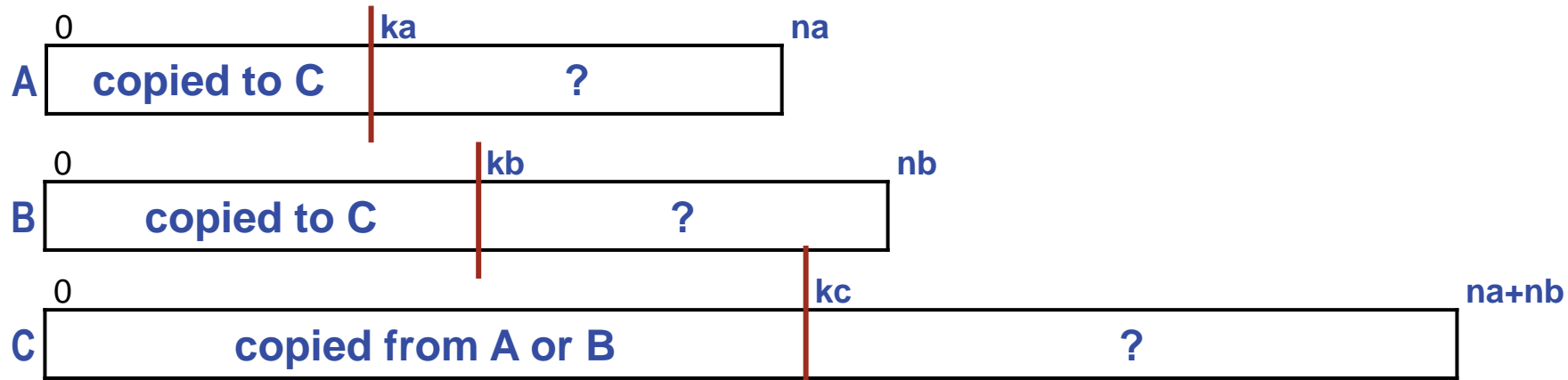
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/* Copy remaining values into C from the unexhausted array, A or B. */
while ( ka<na ) { C[kc] = A[ka]; ka++; kc++; }
while ( kb<nb ) { C[kc] = B[kb]; kb++; kc++; }

```