

Anhang A

Typentheorie: Syntax, Semantik, Inferenzregeln

A.1 Parametertypen

variable : Variablennamen (ML Datentyp `var`)

Zulässige Elemente sind Zeichenketten der Form $[a-zA-Z0-9_-\%]^+$

natural : Natürliche Zahlen einschließlich der Null (ML Datentyp `int`).

Zulässige Elemente sind Zeichenketten der Form $0 + [1-9][0-9]^*$.

token : Zeichenketten für Namen (ML Datentyp `tok`).

Zulässige Elemente bestehen aus allen Symbolen des NuPRL Zeichensatzes mit Ausnahme der Kontrollsymbole.

string : Zeichenketten für Texte (ML Datentyp `string`).

Zulässige Elemente bestehen aus allen Symbolen des NuPRL Zeichensatzes mit Ausnahme der Kontrollsymbole.

level-expression : Ausdrücke für das Level eines Typuniversums (ML Datentyp `level_exp`)

Die genaue Syntax wird in Abschnitt ?? bei der Diskussion der Universenhierarchie besprochen.

Die Namen für Parametertypen dürfen durch ihre ersten Buchstaben abgekürzt werden.

A.2 Operatorentabelle

	kanonisch (Typen)	(Elemente)	nichtkanonisch
	$\mathbf{var}\{x:v\}()$	x	
$\mathbf{int}\{()\}$ \mathbb{Z}	$\mathbf{natnum}\{n:n\}()$ $\mathbf{minus}\{(\mathbf{natnum}\{n:n\}())\}$	n $-n$	$\mathbf{ind}\{(\overline{u}; x, f_x.s; \mathit{base}; y, f_y.t)\}$ $\mathbf{minus}\{(\overline{u})\}$, $\mathbf{add}\{(\overline{u}; \overline{v})\}$, $\mathbf{sub}\{(\overline{u}; \overline{v})\}$ $\mathbf{mul}\{(\overline{u}; \overline{v})\}$, $\mathbf{div}\{(\overline{u}; \overline{v})\}$, $\mathbf{rem}\{(\overline{u}; \overline{v})\}$ $\mathbf{int_eq}\{(\overline{u}; \overline{v}; s; t)\}$, $\mathbf{less}\{(\overline{u}; \overline{v}; s; t)\}$ $\mathbf{ind}\{(\overline{u}; x, f_x.s; \mathit{base}; y, f_y.t)\}$ $-\overline{u}$, $\overline{u} + \overline{v}$, $\overline{u} - \overline{v}$ $\overline{u} * \overline{v}$, $\overline{u} \div \overline{v}$, $\overline{u} \mathbf{rem} \overline{v}$ $\mathbf{if} \overline{u} = \overline{v} \mathbf{then} s \mathbf{else} t$, $\mathbf{if} \overline{u} < \overline{v} \mathbf{then} s \mathbf{else} t$
$\mathbf{lt}\{(u;v)\}$ $u < v$	$\mathbf{Axiom}\{()\}$ Axiom		
$\mathbf{void}\{()\}$ void			$\mathbf{any}\{(e)\}$ $\mathbf{any}(e)$
$\mathbf{Atom}\{()\}$ Atom	$\mathbf{token}\{string:t\}()$ "string"		$\mathbf{atom_eq}\{(\overline{u}; \overline{v}; s; t)\}$ $\mathbf{if} \overline{u} = \overline{v} \mathbf{then} s \mathbf{else} t$
$\mathbf{U}\{j:1\}()$ U_j	(alle kanonischen Typen)		
$\mathbf{fun}\{(S; x.T)\}$ $x:S \rightarrow T$	$\mathbf{lam}\{(x.t)\}$ $\lambda x.t$		$\mathbf{apply}\{(\overline{f}; t)\}$ $\overline{f} t$
$\mathbf{prod}\{(S; x.T)\}$ $x:S \times T$	$\mathbf{pair}\{(s;t)\}$ $\langle s, t \rangle$		$\mathbf{spread}\{(\overline{e}; x, y.u)\}$ $\mathbf{let} \langle x, y \rangle = \overline{e} \mathbf{in} u$
$\mathbf{union}\{(S;T)\}$ $S + T$	$\mathbf{inl}\{(s)\}$, $\mathbf{inr}\{(t)\}$ $\mathbf{inl}(s)$, $\mathbf{inr}(t)$		$\mathbf{decide}\{(\overline{e}; x.u; y.v)\}$ $\mathbf{case} \overline{e} \mathbf{of} \mathbf{inl}(x) \mapsto u \mid \mathbf{inr}(y) \mapsto v$
$\mathbf{equal}\{(s;t;T)\}$ $s = t \in T$	$\mathbf{Axiom}\{()\}$ Axiom		
$\mathbf{list}\{(T)\}$ $T \mathit{list}$	$\mathbf{nil}\{()\}$, $\mathbf{cons}\{(t;l)\}$ $[], t.l$		$\mathbf{list_ind}\{(\overline{s}; \mathit{base}; x, l, f_{xl}.t)\}$ $\mathbf{list_ind}(\overline{s}; \mathit{base}; x, l, f_{xl}.t)$
$\mathbf{set}\{(S; x.T)\}$ $\{x:S \mid T\}$	(Elemente von S)		
$\mathbf{quotient}\{(T; x, y.E)\}$ $x, y : T // E$	(Elemente von T)		
$\mathbf{rec}\{(X.T_X)\}$ $\mathit{rectype} X = T_X$	(Elemente gemäß T_X)		$\mathbf{rec_ind}\{(\overline{e}; f, x.t)\}$ $\mathbf{let}^* f(x) = t \mathbf{in} f(\overline{e})$
$\mathbf{pfun}\{(S; T)\}$ $S \not\rightarrow T$	$\mathbf{fix}\{(f, x.t)\}$ $\mathbf{letrec} f(x) = t$		$\mathbf{dom}\{(\overline{f})\}$, $\mathbf{apply_p}\{(\overline{f}; t)\}$ $\mathbf{dom}(\overline{f})$, $\overline{f}(t)$

Man beachte, daß in der Implementierung von NuPRL 4.0 die konkreten Namen und Displayformen zum Teil etwas anders sind als hier angegeben. An der Anpassung wird gearbeitet.

A.3 Redizes und Kontrakta

Redex		Kontraktum
$(\lambda x. u) t$	$\xrightarrow{\beta}$	$u[t/x]$
$\text{let } \langle x, y \rangle = \langle s, t \rangle \text{ in } u$	$\xrightarrow{\beta}$	$u[s, t / x, y]$
$\text{case inl}(s) \text{ of inl}(x) \mapsto u \mid \text{inr}(y) \mapsto v$	$\xrightarrow{\beta}$	$u[s/x]$
$\text{case inr}(t) \text{ of inl}(x) \mapsto u \mid \text{inr}(y) \mapsto v$	$\xrightarrow{\beta}$	$v[t/y]$
$\text{ind}(0; x, f_x.s; \text{base}; y, f_y.t)$	$\xrightarrow{\beta}$	base
$\text{ind}(n; x, f_x.s; \text{base}; y, f_y.t)$	$\xrightarrow{\beta}$	$t[n, \text{ind}(n-1; x, f_x.s; \text{base}; y, f_y.t) / x, f_x] \quad , \quad (n > 0)$
$\text{ind}(-n; x, f_x.s; \text{base}; y, f_y.t)$	$\xrightarrow{\beta}$	$s[-n, \text{ind}(-n+1; x, f_x.s; \text{base}; y, f_y.t) / y, f_y], \quad (n > 0)$
$-i$	$\xrightarrow{\beta}$	<i>Die Negation von i (als Zahl)</i>
$i+j$	$\xrightarrow{\beta}$	<i>Die Summe von i und j</i>
$i-j$	$\xrightarrow{\beta}$	<i>Die Differenz von i und j</i>
$i*j$	$\xrightarrow{\beta}$	<i>Das Produkt von i und j</i>
$i \div j$	$\xrightarrow{\beta}$	0 , falls $j=0$; ansonsten die Integer-Division von i und j
$i \text{ rem } j$	$\xrightarrow{\beta}$	0 , falls $j=0$; ansonsten der Rest der Division von i und j
$\text{if } i=j \text{ then } s \text{ else } t$	$\xrightarrow{\beta}$	s , falls $i = j$; ansonsten t
$\text{if } i < j \text{ then } s \text{ else } t$	$\xrightarrow{\beta}$	s , falls $i < j$; ansonsten t
$\text{list_ind}([], \text{base}; x, l, f_{xl}.t)$	$\xrightarrow{\beta}$	base
$\text{list_ind}(s.u; \text{base}; x, l, f_{xl}.t)$	$\xrightarrow{\beta}$	$t[s, u, \text{list_ind}(u; \text{base}; x, l, f_{xl}.t) / x, l, f_{xl}]$
$\text{if } u=v \text{ then } s \text{ else } t$	$\xrightarrow{\beta}$	s , falls $u = v$; ansonsten t
$\text{let}^* f(x) = t \text{ in } f(e)$	$\xrightarrow{\beta}$	$t[\lambda y. \text{let}^* f(x) = t \text{ in } f(y), e / f, x]$
$(\text{letrec } f(x) = t) (u)$	$\xrightarrow{\beta}$	$t[\text{letrec } f(x) = t, u / f, x]$
$\text{dom}(\text{letrec } f(x) = t)$	$\xrightarrow{\beta}$	$\lambda x. \text{rectype } F = \mathcal{E}[[t]]$

A.4 Urteile

A.4.1 Typsemantik

$\mathbb{Z} = \mathbb{Z}$	
$i_1 < j_1 = i_2 < j_2$	falls $i_1 = i_2 \in \mathbb{Z}$ und $j_1 = j_2 \in \mathbb{Z}$
void = void	
Atom = Atom	
$U_{j_1} = U_{j_2}$	falls $j_1 = j_2$ (als natürliche Zahl)
$x_1 : S_1 \rightarrow T_1 = x_2 : S_2 \rightarrow T_2$	falls $S_1 = S_2$ und $T_1[s_1/x_1] = T_2[s_2/x_2]$ für alle Terme s_1, s_2 mit $s_1 = s_2 \in S_1$.
$T = S_2 \rightarrow T_2$	falls $T = x_2 : S_2 \rightarrow T_2$ für ein beliebiges $x_2 \in \mathcal{V}$.
$S_1 \rightarrow T_1 = T$	falls $x_1 : S_1 \rightarrow T_1 = T$ für ein beliebiges $x_1 \in \mathcal{V}$.
$x_1 : S_1 \times T_1 = x_2 : S_2 \times T_2$	falls $S_1 = S_2$ und $T_1[s_1/x_1] = T_2[s_2/x_2]$ für alle Terme s_1, s_2 mit $s_1 = s_2 \in S_1$.
$T = S_2 \times T_2$	falls $T = x_2 : S_2 \times T_2$ für ein beliebiges $x_2 \in \mathcal{V}$.
$S_1 \times T_1 = T$	falls $x_1 : S_1 \times T_1 = T$ für ein beliebiges $x_1 \in \mathcal{V}$.
$S_1 + T_1 = S_2 + T_2$	falls $S_1 = S_2$ und $T_1 = T_2$.
$s_1 = t_1 \in T_1 = s_2 = t_2 \in T_2$	falls $T_1 = T_2$ und $s_1 = s_2 \in T_1$ und $t_1 = t_2 \in T_1$.
$T_1 \text{ list} = T_2 \text{ list}$	falls $T_1 = T_2$
$\{x_1 : S_1 \mid T_1\} = \{x_2 : S_2 \mid T_2\}$	falls $S_1 = S_2$ und es gibt eine Variable x , die weder in T_1 noch in T_2 vorkommt, und zwei Terme p_1 und p_2 mit der Eigenschaft $p_1 \in \forall x : S_1. T_1[x/x_1] \Rightarrow T_2[x/x_2]$ und $p_2 \in \forall x : S_1. T_2[x/x_2] \Rightarrow T_1[x/x_1]$.
$T = \{S_2 \mid T_2\}$	falls $T = \{x_2 : S_2 \mid T_2\}$ für ein beliebiges $x_2 \in \mathcal{V}$.
$\{S_1 \mid T_1\} = T$	falls $\{x_1 : S_1 \mid T_1\} = T$ für ein beliebiges $x_1 \in \mathcal{V}$.
$x_1, y_1 : T_1 // E_1 = x_2, y_2 : T_2 // E_2$	falls $T_1 = T_2$ und es gibt (verschiedene) Variablen x, y, z , die weder in E_1 noch in E_2 vorkommen, und Terme p_1, p_2, r, s und t mit der Eigenschaft $p_1 \in \forall x : T_1. \forall y : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_2[x, y/x_2, y_2]$ und $p_2 \in \forall x : T_1. \forall y : T_1. E_2[x, y/x_2, y_2] \Rightarrow E_1[x, y/x_1, y_1]$ und $r \in \forall x : T_1. E_1[x, x/x_1, y_1]$ und $s \in \forall x : T_1. \forall y : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_1[y, x/x_1, y_1]$ und $t \in \forall x : T_1. \forall y : T_1. \forall z : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_1[y, z/x_1, y_1] \Rightarrow E_1[x, z/x_1, y_1]$
rectype $X_1 = T_{X_1} = \text{rectype } X_2 = T_{X_2}$	falls $T_{X_1}[X/X_1] = T_{X_2}[X/X_2]$ für alle Typen X
$S_1 \not\rightarrow T_1 = S_2 \not\rightarrow T_2$	falls $S_1 = S_2$ und $T_1 = T_2$

A.4.2 Elementsemantik – Universen

Elementsemantik – Universen	
$\mathbf{Z} = \mathbf{Z} \in U_j$ $i_1 < j_1 = i_2 < j_2 \in U_j$	falls $i_1 = i_2 \in \mathbf{Z}$ und $j_1 = j_2 \in \mathbf{Z}$
$\text{void} = \text{void} \in U_j$ $\text{Atom} = \text{Atom} \in U_j$	
$U_{j_1} = U_{j_2} \in U_j$	falls $j_1 = j_2 < j$ (als natürliche Zahl)
$x_1 : S_1 \rightarrow T_1 = x_2 : S_2 \rightarrow T_2 \in U_j$	falls $S_1 = S_2 \in U_j$ und $T_1[s_1/x_1] = T_2[s_2/x_2] \in U_j$ für alle Terme s_1, s_2 mit $s_1 = s_2 \in S_1$.
$T = S_2 \rightarrow T_2 \in U_j$	falls $T = x_2 : S_2 \rightarrow T_2 \in U_j$ für ein beliebiges $x_2 \in \mathcal{V}$.
$S_1 \rightarrow T_1 = T \in U_j$	falls $x_1 : S_1 \rightarrow T_1 = T \in U_j$ für ein beliebiges $x_1 \in \mathcal{V}$.
$x_1 : S_1 \times T_1 = x_2 : S_2 \times T_2 \in U_j$	falls $S_1 = S_2 \in U_j$ und $T_1[s_1/x_1] = T_2[s_2/x_2] \in U_j$ für alle Terme s_1, s_2 mit $s_1 = s_2 \in S_1$.
$T = S_2 \times T_2 \in U_j$	falls $T = x_2 : S_2 \times T_2 \in U_j$ für ein beliebiges $x_2 \in \mathcal{V}$.
$S_1 \times T_1 = T \in U_j$	falls $x_1 : S_1 \times T_1 = T \in U_j$ für ein beliebiges $x_1 \in \mathcal{V}$.
$S_1 + T_1 = S_2 + T_2 \in U_j$	falls $S_1 = S_2 \in U_j$ und $T_1 = T_2 \in U_j$.
$s_1 = t_1 \in T_1 = s_2 = t_2 \in T_2 \in U_j$	falls $T_1 = T_2 \in U_j$ und $s_1 = s_2 \in T_1$ und $t_1 = t_2 \in T_1$.
$T_1 \text{ list} = T_2 \text{ list} \in U_j$	falls $T_1 = T_2 \in U_j$
$\{x_1 : S_1 \mid T_1\} = \{x_2 : S_2 \mid T_2\} \in U_j$	falls $S_1 = S_2 \in U_j$ und $T_1[s/x_1] \in U_j$ sowie $T_2[s/x_2] \in U_j$ gilt für alle Terme s mit $s \in S_1$ und es gibt eine Variable x , die weder in T_1 noch in T_2 vorkommt, und zwei Terme p_1 und p_2 mit der Eigenschaft $p_1 \in \forall x : S_1. T_1[x/x_1] \Rightarrow T_2[x/x_2]$ und $p_2 \in \forall x : S_1. T_2[x/x_2] \Rightarrow T_1[x/x_1]$
$T = \{S_2 \mid T_2\} \in U_j$	falls $T = \{x_2 : S_2 \mid T_2\} \in U_j$ für ein beliebiges $x_2 \in \mathcal{V}$.
$\{S_1 \mid T_1\} = T \in U_j$	falls $\{x_1 : S_1 \mid T_1\} = T \in U_j$ für ein beliebiges $x_1 \in \mathcal{V}$.
$x_1, y_1 : T_1 // E_1 = x_2, y_2 : T_2 // E_2 \in U_j$	falls $T_1 = T_2 \in U_j$ und für alle Terme s, t mit $s \in T_1$ und $t \in T_1$ gilt $E_1[s, t/x_1, y_1] \in U_j$ sowie $E_2[s, t/x_2, y_2] \in U_j$ und es gibt (verschiedene) Variablen x, y, z , die weder in E_1 noch in E_2 vorkommen, und Terme p_1, p_2, r, s und t mit der Eigenschaft $p_1 \in \forall x : T_1. \forall y : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_2[x, y/x_2, y_2]$ und $p_2 \in \forall x : T_1. \forall y : T_1. E_2[x, y/x_2, y_2] \Rightarrow E_1[x, y/x_1, y_1]$ und $r \in \forall x : T_1. E_1[x, x/x_1, y_1]$ und $s \in \forall x : T_1. \forall y : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_1[y, x/x_1, y_1]$ und $t \in \forall x : T_1. \forall y : T_1. \forall z : T_1. E_1[x, y/x_1, y_1] \Rightarrow E_1[y, z/x_1, y_1]$
$\text{rectype } X_1 = T_{X_1} = \text{rectype } X_2 = T_{X_2} \in U_j$	falls $T_{X_1}[X/X_1] = T_{X_2}[X/X_2] \in U_j$ für alle Terme X mit $X \in U_j$
$S_1 \not\rightarrow T_1 = S_2 \not\rightarrow T_2 \in U_j$	falls $S_1 = S_2 \in U_j$ und $T_1 = T_2 \in U_j$

A.4.3 Elementsemantik

$i = i \in \mathbb{Z}$	
Axiom = Axiom $\in s < t$	falls es ganze Zahlen i und j gibt, wobei i kleiner als j ist, für die gilt $s \xrightarrow{l} i$ und $t \xrightarrow{l} j$
$s = t \in \text{void}$	<i>gilt niemals !</i>
„string“ = „string“ $\in \text{Atom}$	
$\lambda x_1. t_1 = \lambda x_2. t_2 \in x : S \rightarrow T$	falls $x : S \rightarrow T$ Typ und $t_1[s_1/x_1] = t_2[s_2/x_2] \in T[s_1/x]$ für alle Terme s_1, s_2 mit $s_1 = s_2 \in S$.
$\langle s_1, t_1 \rangle = \langle s_2, t_2 \rangle \in x : S \times T$	falls $x : S \times T$ Typ und $s_1 = s_2 \in S$ und $t_1 = t_2 \in T[s_1/x]$.
$\text{inl}(s_1) = \text{inl}(s_2) \in S + T$	falls $S + T$ Typ und $s_1 = s_2 \in S$.
$\text{inr}(t_1) = \text{inr}(t_2) \in S + T$	falls $S + T$ Typ und $t_1 = t_2 \in T$.
Axiom = Axiom $\in s = t \in T$	falls $s = t \in T$
$[] = [] \in T \text{ list}$	falls T Typ
$t_1.l_1 = t_2.l_2 \in T \text{ list}$	falls T Typ und $t_1 = t_2 \in T$ und $l_1 = l_2 \in T \text{ list}$
$s = t \in \{x : S \mid T\}$	falls $\{x : S \mid T\}$ Typ und $s = t \in S$ und es gibt einen Term p mit der Eigenschaft $p \in T[s/x]$
$s = t \in x, y : T // E$	falls $x, y : T // E$ Typ und $s \in T$ und $t \in T$ und es gibt einen Term p mit der Eigenschaft $p \in E[s, t/x, y]$
$s = t \in \text{rectype } X = T_X$	falls $\text{rectype } X = T_X$ Typ und $s = t \in T_X[\text{rectype } X = T_X / X]$
$\text{letrec } f_1(x_1) = t_1 = \text{letrec } f_2(x_2) = t_2 \in S \not\rightarrow T$	falls $S \not\rightarrow T$ Typ und für alle Terme s_1 und s_2 mit $s_1 = s_2 \in S$ $\{x : S \mid \text{dom}(\text{letrec } f_1(x_1) = t_1)(x)\}$ $= \{x : S \mid \text{dom}(\text{letrec } f_2(x_2) = t_2)(x)\}$ und $t_1[\text{letrec } f_1(x_1) = t_1, s_1 / f_1, x_1]$ $= t_2[\text{letrec } f_2(x_2) = t_2, s_2 / f_2, x_2] \in T$

A.5 Inferenzregeln

A.5.1 Ganze Zahlen

$\Gamma \vdash \mathbf{Z} \in \mathbf{U}_j$ $_{[Ax]}$
by intEq

$\Gamma \vdash n \in \mathbf{Z}$ $_{[Ax]}$
by natnumEq

$\Gamma \vdash \text{ind}(u_1; x_1, f_{x_1}.s_1; \text{base}_1; y_1, f_{y_1}.t_1)$
 $= \text{ind}(u_2; x_2, f_{x_2}.s_2; \text{base}_2; y_2, f_{y_2}.t_2)$
 $\in T[u_1/z]$ $_{[Ax]}$
by indEq $z T$
 $\Gamma \vdash u_1 = u_2$ $_{[Ax]}$
 $\Gamma, x: \mathbf{Z}, v: x < 0, f_x: T[(x+1)/z]$
 $\vdash s_1[x, f_x/x_1, f_{x_1}] = s_2[x, f_x/x_2, f_{x_2}] \in T[x/z]$
 $_{[Ax]}$
 $\Gamma \vdash \text{base}_1 = \text{base}_2 \in T[0/z]$ $_{[Ax]}$
 $\Gamma, x: \mathbf{Z}, v: 0 < x, f_x: T[(x-1)/z]$
 $\vdash t_1[x, f_x/y_1, f_{y_1}] = t_2[x, f_x/y_2, f_{y_2}] \in T[x/z]$ $_{[Ax]}$

$\Gamma \vdash -s_1 = -s_2$ $_{[Ax]}$
by minusEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$

$\Gamma \vdash s_1 + t_1 = s_2 + t_2$ $_{[Ax]}$
by addEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 = t_2$ $_{[Ax]}$

$\Gamma \vdash s_1 - t_1 = s_2 - t_2$ $_{[Ax]}$
by subEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 = t_2$ $_{[Ax]}$

$\Gamma \vdash s_1 * t_1 = s_2 * t_2$ $_{[Ax]}$
by mulEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 = t_2$ $_{[Ax]}$

$\Gamma \vdash s_1 \div t_1 = s_2 \div t_2$ $_{[Ax]}$
by divEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 = t_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 \neq 0$ $_{[Ax]}$

$\Gamma \vdash s_1 \text{ rem } t_1 = s_2 \text{ rem } t_2$ $_{[Ax]}$
by remEq
 $\Gamma \vdash s_1 = s_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 = t_2$ $_{[Ax]}$
 $\Gamma \vdash t_1 \neq 0$ $_{[Ax]}$

$\Gamma \vdash \text{if } u_1 = v_1 \text{ then } s_1 \text{ else } t_1$
 $= \text{if } u_2 = v_2 \text{ then } s_2 \text{ else } t_2 \in T$ $_{[Ax]}$
by int_eqEq
 $\Gamma \vdash u_1 = u_2$ $_{[Ax]}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } n_j}$
by natnumI n

$\Gamma, z: \mathbf{Z}, \Delta \vdash C$
 $_{\text{ext ind}(z; x, f_x.s[Axiom/v]; \text{base}; x, f_x.t[Axiom/v])}$
by intE i
 $\Gamma, z: \mathbf{Z}, \Delta, x: \mathbf{Z}, v: x < 0, f_x: C[(x+1)/z]$
 $\vdash C[x/z]$ $_{\text{ext } s_j}$
 $\Gamma, z: \mathbf{Z}, \Delta \vdash C[0/z]$ $_{\text{ext base}_j}$
 $\Gamma, z: \mathbf{Z}, \Delta, x: \mathbf{Z}, v: 0 < x, f_x: C[(x-1)/z]$
 $\vdash C[x/z]$ $_{\text{ext } t_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } -s_j}$
by minusI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s+t_j}$
by addI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } t_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s-t_j}$
by subI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } t_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s*t_j}$
by mulI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } t_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s \div t_j}$
by divI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } t_j}$

$\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s \text{ rem } t_j}$
by remI
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } s_j}$
 $\Gamma \vdash \mathbf{Z}$ $_{\text{ext } t_j}$

$\Gamma \vdash \text{if } u_1 < v_1 \text{ then } s \text{ else } t$
 $= \text{if } u_2 < v_2 \text{ then } s_2 \text{ else } t_2 \in T$ $_{[Ax]}$
by lessEq
 $\Gamma \vdash u_1 = u_2$ $_{[Ax]}$

$$\Gamma \vdash \text{ind}(i; x, f_x.s; \text{base}; y, f_y.t) = t_2 \in T_{[A_x]}$$

by indRedDown

$$\Gamma \vdash t[i, \text{ind}(i+1; x, f_x.s; \text{base}; y, f_y.t) / x, f_x]$$

$$= t_2 \in T_{[A_x]}$$

$$\Gamma \vdash i < 0_{[A_x]}$$

$$\Gamma \vdash \text{ind}(i; x, f_x.s; \text{base}; y, f_y.t) = t_2 \in T_{[A_x]}$$

by indRedUp

$$\Gamma \vdash t[i, \text{ind}(i-1; x, f_x.s; \text{base}; y, f_y.t) / y, f_y]$$

$$= t_2 \in T_{[A_x]}$$

$$\Gamma \vdash 0 < i_{[A_x]}$$

$$\Gamma \vdash \text{ind}(i; x, f_x.s; \text{base}; y, f_y.t) = t_2 \in T_{[A_x]}$$

by indRedBase

$$\Gamma \vdash \text{base} = t_2 \in T_{[A_x]}$$

$$\Gamma \vdash i = 0_{[A_x]}$$

$$\Gamma \vdash \text{if } u=v \text{ then } s \text{ else } t = t_2 \in T_{[A_x]}$$

by int_eqRedT

$$\Gamma \vdash s = t_2 \in T_{[A_x]}$$

$$\Gamma \vdash u=v_{[A_x]}$$

$$\Gamma \vdash \text{if } u=v \text{ then } s \text{ else } t = t_2 \in T_{[A_x]}$$

by int_eqRedF

$$\Gamma \vdash t = t_2 \in T_{[A_x]}$$

$$\Gamma \vdash u \neq v_{[A_x]}$$

$$\Gamma \vdash \text{if } u < v \text{ then } s \text{ else } t = t_2 \in T_{[A_x]}$$

by lessRedT

$$\Gamma \vdash s = t_2 \in T_{[A_x]}$$

$$\Gamma \vdash u < v_{[A_x]}$$

$$\Gamma \vdash \text{if } u < v \text{ then } s \text{ else } t = t_2 \in T_{[A_x]}$$

by lessRedF

$$\Gamma \vdash t = t_2 \in T_{[A_x]}$$

$$\Gamma \vdash u \geq v_{[A_x]}$$

$$\Gamma \vdash 0 \leq s \text{ rem } t \wedge s \text{ rem } t < t_{[A_x]}$$

by remBounds1

$$\Gamma \vdash 0 \leq s_{[A_x]}$$

$$\Gamma \vdash 0 < t_{[A_x]}$$

$$\Gamma \vdash s \text{ rem } t \leq 0 \wedge s \text{ rem } t > -t_{[A_x]}$$

by remBounds3

$$\Gamma \vdash s \leq 0_{[A_x]}$$

$$\Gamma \vdash 0 < t_{[A_x]}$$

$$\Gamma \vdash s = (s \div t) * t + (s \text{ rem } t)_{[A_x]}$$

by divremSum

$$\Gamma \vdash s \in \mathbb{Z}_{[A_x]}$$

$$\Gamma \vdash t \neq 0_{[A_x]}$$

$$\Gamma \vdash s_1 < t_1 = s_2 < t_2 \in U_j_{[A_x]}$$

by ltEq

$$\Gamma \vdash s_1 = s_2_{[A_x]}$$

$$\Gamma \vdash t_1 = t_2_{[A_x]}$$

$$\Gamma \vdash 0 \leq s \text{ rem } t \wedge s \text{ rem } t < -t_{[A_x]}$$

by remBounds2

$$\Gamma \vdash 0 \leq s_{[A_x]}$$

$$\Gamma \vdash t < 0_{[A_x]}$$

$$\Gamma \vdash s \text{ rem } t \leq 0 \wedge s \text{ rem } t > t_{[A_x]}$$

by remBounds4

$$\Gamma \vdash s \leq 0_{[A_x]}$$

$$\Gamma \vdash t < 0_{[A_x]}$$

$$\Gamma \vdash \text{Axiom} \in s < t_{[A_x]}$$

by axiomEq_lt

$$\Gamma \vdash s < t_{[A_x]}$$

A.5.2 Void

$$\Gamma \vdash \text{void} \in U_j_{[A_x]}$$

by voidEq

$$\Gamma \vdash \text{any}(s) = \text{any}(t) \in T_{[A_x]}$$

by anyEq

$$\Gamma \vdash s = t \in \text{void}_{[A_x]}$$

$$\Gamma, z: \text{void}, \Delta \vdash C \text{ [ext any}(z)]$$

by voidE i

A.5.3 Universen

$$\Gamma \vdash U_j \in U_k_{[A_x]}$$

by univEq*

$$\Gamma \vdash T \in U_k_{[A_x]}$$

by cumulativity j*

$$\Gamma \vdash T \in U_{[A_x]}$$

A.5.4 Atom

$$\Gamma \vdash \text{Atom} = \text{Atom} \in \mathbf{U}_j \text{ [Ax]}$$

by atomEq

$$\Gamma \vdash \text{"string"} \in \text{Atom} \text{ [Ax]}$$

by tokEq

$$\Gamma \vdash \text{if } u_1=v_1 \text{ then } s_1 \text{ else } t_1$$

$$= \text{if } u_2=v_2 \text{ then } s_2 \text{ else } t_2 \in T \text{ [Ax]}$$

by atom_eqEq

$$\Gamma \vdash u_1=u_2 \in \text{Atom} \text{ [Ax]}$$

$$\Gamma \vdash v_1=v_2 \in \text{Atom} \text{ [Ax]}$$

$$\Gamma, v: u_1=v_1 \in \text{Atom} \vdash s_1 = s_2 \in T \text{ [Ax]}$$

$$\Gamma, v: \neg(u_1=v_1 \in \text{Atom}) \vdash t_1 = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash \text{if } u=v \text{ then } s \text{ else } t = t_2 \in T \text{ [Ax]}$$

by atom_eqRedT

$$\Gamma \vdash s = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash u=v \in \text{Atom} \text{ [Ax]}$$

$$\Gamma \vdash \text{Atom} \text{ [ext "string"]}$$

by tokI "string"

$$\Gamma \vdash \text{if } u=v \text{ then } s \text{ else } t = t_2 \in T \text{ [Ax]}$$

by atom_eqRedF

$$\Gamma \vdash t = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash \neg(u=v \in \text{Atom}) \text{ [Ax]}$$

A.5.5 Funktionenraum

$$\Gamma \vdash x_1:S_1 \rightarrow T_1 = x_2:S_2 \rightarrow T_2 \in \mathbf{U}_j \text{ [Ax]}$$

by functionEq

$$\Gamma \vdash S_1 = S_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, x:S_1 \vdash T_1[x/x_1] = T_2[x/x_2] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash \lambda x_1.t_1 = \lambda x_2.t_2 \in x:S \rightarrow T \text{ [Ax]}$$

by lambdaEq j

$$\Gamma, x':S \vdash t_1[x'/x_1] = t_2[x'/x_2] \in T[x'/x] \text{ [Ax]}$$

$$\Gamma \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash f_1 t_1 = f_2 t_2 \in T[t_1/x] \text{ [Ax]}$$

by applyEq x:S \rightarrow T

$$\Gamma \vdash f_1 = f_2 \in x:S \rightarrow T \text{ [Ax]}$$

$$\Gamma \vdash t_1 = t_2 \in S \text{ [Ax]}$$

$$\Gamma \vdash (\lambda x.t) s = t_2 \in T \text{ [Ax]}$$

by applyRed

$$\Gamma \vdash t[s/x] = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash f_1 = f_2 \in x:S \rightarrow T \text{ [ext } t_j]$$

by functionExt j x_1:S_1 \rightarrow T_1 x_2:S_2 \rightarrow T_2

$$\Gamma, x':S \vdash f_1 x' = f_2 x' \in T[x'/x] \text{ [ext } t_j]$$

$$\Gamma \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash f_1 \in x_1:S_1 \rightarrow T_1 \text{ [Ax]}$$

$$\Gamma \vdash f_2 \in x_2:S_2 \rightarrow T_2 \text{ [Ax]}$$

$$\Gamma \vdash x:S \rightarrow T \text{ [ext } \lambda x'.t_j]$$

by lambdaI j

$$\Gamma, x':S \vdash T[x'/x] \text{ [ext } t_j]$$

$$\Gamma \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, f: x:S \rightarrow T, \Delta \vdash C \text{ [ext } t[f s, \text{Axiom}/y, v]]$$

by functionE i s

$$\Gamma, f: x:S \rightarrow T, \Delta \vdash s \in S \text{ [Ax]}$$

$$\Gamma, f: x:S \rightarrow T, y:T[s/x],$$

$$v: y=f s \in T[s/x], \Delta \vdash C \text{ [ext } t_j]$$

$$\Gamma, f: S \rightarrow T, \Delta \vdash C \text{ [ext } t[f s, y]]$$

by functionE_indep i

$$\Gamma, f: S \rightarrow T, \Delta \vdash \bar{S} \text{ [ext } s_j]$$

$$\Gamma, f: S \rightarrow T, y:T, \Delta \vdash C \text{ [ext } t_j]$$

A.5.6 Produktraum

$$\Gamma \vdash x_1:S_1 \times T_1 = x_2:S_2 \times T_2 \in \mathbf{U}_j \text{ }_{[A_x]}$$

by productEq

$$\Gamma \vdash S_1 = S_2 \in \mathbf{U}_j \text{ }_{[A_x]}$$

$$\Gamma, x':S_1 \vdash T_1[x'/x_1] = T_2[x'/x_2] \in \mathbf{U}_j \text{ }_{[A_x]}$$

$$\Gamma \vdash \langle s_1, t_1 \rangle = \langle s_2, t_2 \rangle \in x:S \times T \text{ }_{[A_x]}$$

by pairEq j

$$\Gamma \vdash s_1 = s_2 \in S \text{ }_{[A_x]}$$

$$\Gamma \vdash t_1 = t_2 \in T[s_1/x] \text{ }_{[A_x]}$$

$$\Gamma, x':S \vdash T[x'/x] \in \mathbf{U}_j \text{ }_{[A_x]}$$

$$\Gamma \vdash \text{let } \langle x_1, y_1 \rangle = e_1 \text{ in } t_1$$

$$= \text{let } \langle x_2, y_2 \rangle = e_2 \text{ in } t_2 \in C[e_1/z] \text{ }_{[A_x]}$$

by spreadEq z C x:S \times T

$$\Gamma \vdash e_1 = e_2 \in x:S \times T \text{ }_{[A_x]}$$

$$\Gamma, s:S, t:T[s/x], y: e_1 = \langle s, t \rangle \in x:S \times T$$

$$\vdash t_1[s, t/x_1, y_1] = t_2[s, t/x_2, y_2] \in C[\langle s, t \rangle/z] \text{ }_{[A_x]}$$

$$\Gamma \vdash \text{let } \langle x, y \rangle = \langle s, t \rangle \text{ in } u = t_2 \in T \text{ }_{[A_x]}$$

by spreadRed

$$\Gamma \vdash u[s, t/x, y] = t_2 \in T \text{ }_{[A_x]}$$

$$\Gamma \vdash \langle s_1, t_1 \rangle = \langle s_2, t_2 \rangle \in S \times T \text{ }_{[A_x]}$$

by pairEq_indep

$$\Gamma \vdash s_1 = s_2 \in S \text{ }_{[A_x]}$$

$$\Gamma \vdash t_1 = t_2 \in T \text{ }_{[A_x]}$$

$$\Gamma \vdash x:S \times T \text{ }_{[\mathbf{ext} \langle s, t \rangle]}$$

by pairI j s

$$\Gamma \vdash s \in S \text{ }_{[A_x]}$$

$$\Gamma \vdash T[s/x] \text{ }_{[\mathbf{ext} t]}$$

$$\Gamma, x':S \vdash T[x'/x] \in \mathbf{U}_j \text{ }_{[A_x]}$$

$$\Gamma, z: x:S \times T, \Delta \vdash C \text{ }_{[\mathbf{ext} \text{let } \langle s, t \rangle = z \text{ in } u]}$$

by productE i

$$\Gamma, z: x:S \times T, s:S, t:T[s/x], \Delta[\langle s, t \rangle/z]$$

$$\vdash C[\langle s, t \rangle/z] \text{ }_{[\mathbf{ext} u]}$$

$$\Gamma \vdash S \times T \text{ }_{[\mathbf{ext} \langle s, t \rangle]}$$

by pairI_indep

$$\Gamma \vdash S \text{ }_{[\mathbf{ext} s]}$$

$$\Gamma \vdash T \text{ }_{[\mathbf{ext} t]}$$

A.5.7 Disjunkte Vereinigung (Summe)

$$\Gamma \vdash S_1 + T_1 = S_2 + T_2 \in \mathbf{U}_j \text{ [Ax]}$$

by unionEq

$$\Gamma \vdash S_1 = S_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash T_1 = T_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash \text{inl}(s_1) = \text{inl}(s_2) \in S + T \text{ [Ax]}$$

by inlEq j

$$\Gamma \vdash s_1 = s_2 \in S \text{ [Ax]}$$

$$\Gamma \vdash T \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash \text{inr}(t_1) = \text{inr}(t_2) \in S + T \text{ [Ax]}$$

by inrEq j

$$\Gamma \vdash t_1 = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

$$\begin{aligned} \Gamma \vdash \text{case } e_1 \text{ of } \text{inl}(x_1) \mapsto u_1 \mid \text{inr}(y_1) \mapsto v_1 \\ = \text{case } e_2 \text{ of } \text{inl}(x_2) \mapsto u_2 \mid \text{inr}(y_2) \mapsto v_2 \\ \in C[e_1/z] \text{ [Ax]} \end{aligned}$$

by decideEq z C $S+T$

$$\Gamma \vdash e_1 = e_2 \in S + T \text{ [Ax]}$$

$$\Gamma, s:S, y: e_1 = \text{inl}(s) \in S + T$$

$$\vdash u_1[s/x_1] = u_2[s/x_2] \in C[\text{inl}(s)/z] \text{ [Ax]}$$

$$\Gamma, t:T, y: e_1 = \text{inr}(t) \in S + T$$

$$\vdash v_1[t/y_1] = v_2[t/y_2] \in C[\text{inr}(t)/z] \text{ [Ax]}$$

$$\begin{aligned} \Gamma \vdash \text{case } \text{inl}(s) \text{ of } \text{inl}(x) \mapsto u \mid \text{inr}(y) \mapsto v \\ = t_2 \in T \text{ [Ax]} \end{aligned}$$

by decideRedL

$$\Gamma \vdash u[s/x] = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash S + T \text{ [ext inl}(s)\text{]}$$

by inlI j

$$\Gamma \vdash S \text{ [ext } s\text{]}$$

$$\Gamma \vdash T \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash S + T \text{ [ext inr}(t)\text{]}$$

by inrI j

$$\Gamma \vdash T \text{ [ext } t\text{]}$$

$$\Gamma \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, z:S+T, \Delta \vdash C$$

$$\text{[ext case } z \text{ of } \text{inl}(x) \mapsto u \mid \text{inr}(y) \mapsto v \text{]}$$

by unionE i

$$\Gamma, z:S+T, x:S, \Delta[\text{inl}(x)/z]$$

$$\vdash C[\text{inl}(x)/z] \text{ [ext } u\text{]}$$

$$\Gamma, z:S+T, y:T, \Delta[\text{inr}(y)/z]$$

$$\vdash C[\text{inr}(y)/z] \text{ [ext } v\text{]}$$

$$\begin{aligned} \Gamma \vdash \text{case } \text{inr}(t) \text{ of } \text{inl}(x) \mapsto u \mid \text{inr}(y) \mapsto v \\ = t_2 \in T \text{ [Ax]} \end{aligned}$$

by decideRedR

$$\Gamma \vdash v[t/y] = t_2 \in T \text{ [Ax]}$$

A.5.8 Gleichheit

$$\Gamma \vdash s_1 = t_1 \in T_1 = s_2 = t_2 \in T_2 \in \bigcup_j \text{[Ax]}$$

by equalityEq

$$\Gamma \vdash T_1 = T_2 \in \bigcup_j \text{[Ax]}$$

$$\Gamma \vdash s_1 = s_2 \in T_1 \text{ [Ax]}$$

$$\Gamma \vdash t_1 = t_2 \in T_1 \text{ [Ax]}$$

$$\Gamma, x:T, \Delta \vdash x \in T \text{ [Ax]}$$

by hypEq *i*

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

by transitivity *t'*

$$\Gamma \vdash s = t' \in T \text{ [Ax]}$$

$$\Gamma \vdash t' = t \in T \text{ [Ax]}$$

A.5.9 Listen

$$\Gamma \vdash T_1 \text{ list} = T_2 \text{ list} \in \bigcup_j \text{[Ax]}$$

by listEq

$$\Gamma \vdash T_1 = T_2 \in \bigcup_j \text{[Ax]}$$

$$\Gamma \vdash [] \in T \text{ list [Ax]}$$

by nilEq *j*

$$\Gamma \vdash T \in \bigcup_j \text{[Ax]}$$

$$\Gamma \vdash t_1.l_1 = t_2.l_2 \in T \text{ list [Ax]}$$

by consEq

$$\Gamma \vdash t_1 = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash l_1 = l_2 \in T \text{ list [Ax]}$$

$$\Gamma \vdash \text{list_ind}(s_1; \text{base}_1; x_1, l_1, f_{xl1}.t_1) = \text{list_ind}(s_2; \text{base}_2; x_2, l_2, f_{xl2}.t_2) \in T[s_1/z] \text{ [Ax]}$$

by list_indEq *z T S list*

$$\Gamma \vdash s_1 = s_2 \in S \text{ list [Ax]}$$

$$\Gamma \vdash \text{base}_1 = \text{base}_2 \in T[[]/z] \text{ [Ax]}$$

$$\Gamma, x:S, l:S \text{ list}, f_{xl}:T[l/z] \vdash$$

$$t_1[x, l, f_{xl}/x_1, l_1, f_{xl1}] = t_1[x, l, f_{xl}/x_2, l_2, f_{xl2}] \in T[x.l/z] \text{ [Ax]}$$

$$\Gamma \vdash \text{list_ind}([], \text{base}; x, l, f_{xl}.t) = t_2 \in T \text{ [Ax]}$$

by list_indRedBase

$$\Gamma \vdash \text{base} = t_2 \in T \text{ [Ax]}$$

$$\Gamma \vdash \text{Axiom} \in s = t \in T \text{ [Ax]}$$

by axiomEq

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

$$\Gamma, z: s = t \in T, \Delta \vdash C \text{ [ext } u_j]$$

by equalityE *i*

$$\Gamma, z: s = t \in T, \Delta[\text{Axiom}/z]$$

$$\vdash C[\text{Axiom}/z] \text{ [ext } u_j]$$

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

by symmetry

$$\Gamma \vdash t = s \in T \text{ [Ax]}$$

$$\Gamma \vdash C[s/x] \text{ [ext } u_j]$$

by subst *j s = t \in T x C*

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

$$\Gamma \vdash C[t/x] \text{ [ext } u_j]$$

$$\Gamma, x:T \vdash C \in \bigcup_j \text{[Ax]}$$

$$\Gamma \vdash T \text{ list [ext } []_j]$$

by nilI *j*

$$\Gamma \vdash T \in \bigcup_j \text{[Ax]}$$

$$\Gamma \vdash T \text{ list [ext } t.l_j]$$

by consI

$$\Gamma \vdash T \text{ [ext } t_j]$$

$$\Gamma \vdash T \text{ list [ext } l_j]$$

$$\Gamma, z:T \text{ list}, \Delta \vdash C$$

$$\text{[ext list_ind}(z; \text{base}; x, l, f_{xl}.t)]$$

by listE *i*

$$\Gamma, z:T \text{ list}, \Delta \vdash C[[]/z] \text{ [ext base}_j]$$

$$\Gamma, z:T \text{ list}, \Delta, x:T, l:T \text{ list}, f_{xl}:C[l/z]$$

$$\vdash C[x.l/z] \text{ [ext } t_j]$$

$$\Gamma \vdash \text{list_ind}(s.u; \text{base}; x, l, f_{xl}.t) = t_2 \in T \text{ [Ax]}$$

by list_indRedUp

$$\Gamma \vdash t[s, u, \text{list_ind}(u; \text{base}; x, l, f_{xl}.t) / x, l, f_{xl}] = t_2 \in T \text{ [Ax]}$$

A.5.10 Teilmengen

$$\Gamma \vdash \{x_1:S_1 \mid T_1\} = \{x_2:S_2 \mid T_2\} \in \mathbf{U}_j \text{ [Ax]}$$

by setEq

$$\Gamma \vdash S_1 = S_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, x:S_1 \vdash T_1[x/x_1] = T_2[x/x_2] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash s = t \in \{x:S \mid T\} \text{ [Ax]}$$

by elementEq j

$$\Gamma \vdash s = t \in S \text{ [Ax]}$$

$$\Gamma \vdash T[s/x] \text{ [Ax]}$$

$$\Gamma, x':S \vdash T[x'/x] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash \{x:S \mid T\} \text{ [ext } s_j]$$

by elementI j s

$$\Gamma \vdash s \in S \text{ [Ax]}$$

$$\Gamma \vdash T[s/x] \text{ [Ax]}$$

$$\Gamma, x':S \vdash T[x'/x] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, z: \{x:S \mid T\}, \Delta \vdash C \text{ [ext } (\lambda y.t) z]$$

by setE i

$$\Gamma, z: \{x:S \mid T\}, y:S, \llbracket v \rrbracket : T[y/x], \Delta[y/z]$$

$$\vdash C[y/z] \text{ [ext } t]$$

$$\Gamma \vdash s = t \in \{S \mid T\} \text{ [Ax]}$$

by elementEq_indep

$$\Gamma \vdash s = t \in S \text{ [Ax]}$$

$$\Gamma \vdash T \text{ [Ax]}$$

$$\Gamma \vdash \{S \mid T\} \text{ [ext } s_j]$$

by elementI_indep

$$\Gamma \vdash S \text{ [ext } s_j]$$

$$\Gamma \vdash T \text{ [Ax]}$$

A.5.11 Quotienten

$$\Gamma \vdash x_1, y_1 : T_1 // E_1 = x_2, y_2 : T_2 // E_2 \in \mathbf{U}_j \text{ [Ax]}$$

by quotientEq_weak

$$\Gamma \vdash T_1 = T_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, x:T_1, y:T_1$$

$$\vdash E_1[x, y/x_1, y_1] = E_2[x, y/x_2, y_2] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, x:T_1, y:T_1 \vdash E_1[x, x/x_1, y_1] \text{ [Ax]}$$

$$\Gamma, x:T_1, y:T_1, v: E_1[x, y/x_1, y_1]$$

$$\vdash E_1[y, x/x_1, y_1] \text{ [Ax]}$$

$$\Gamma, x:T_1, y:T_1, z:T_1, v: E_1[x, y/x_1, y_1],$$

$$v': E_1[y, z/x_1, y_1] \vdash E_1[x, z/x_1, y_1] \text{ [Ax]}$$

$$\Gamma \vdash x_1, y_1 : T_1 // E_1 = x_2, y_2 : T_2 // E_2 \in \mathbf{U}_j \text{ [Ax]}$$

by quotientEq

$$\Gamma \vdash x_1, y_1 : T_1 // E_1 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash x_2, y_2 : T_2 // E_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash T_1 = T_2 \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, v: T_1 = T_2 \in \mathbf{U}_j, x:T_1, y:T_1$$

$$\vdash E_1[x, y/x_1, y_1] \Rightarrow E_2[x, y/x_2, y_2] \text{ [Ax]}$$

$$\Gamma, v: T_1 = T_2 \in \mathbf{U}_j, x:T_1, y:T_1$$

$$\vdash E_2[x, y/x_2, y_2] \Rightarrow E_1[x, y/x_1, y_1] \text{ [Ax]}$$

$$\Gamma \vdash s = t \in x, y : T // E \text{ [Ax]}$$

by memberEq_weak j

$$\Gamma \vdash x, y : T // E \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

$$\Gamma \vdash s = t \in x, y : T // E \text{ [Ax]}$$

by memberEq j

$$\Gamma \vdash x, y : T // E \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash s \in T \text{ [Ax]}$$

$$\Gamma \vdash t \in T \text{ [Ax]}$$

$$\Gamma \vdash E[s, t/x, y] \text{ [Ax]}$$

$$\Gamma \vdash x, y : T // E \text{ [ext } t_j]$$

by memberI j

$$\Gamma \vdash x, y : T // E \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma \vdash T \text{ [ext } t_j]$$

$$\Gamma, v: s=t \in x, y : T // E, \Delta \vdash C \text{ [ext } u_j]$$

by quotient_eqE i j

$$\Gamma, v: s=t \in x, y : T // E, \llbracket v' \rrbracket : E[s, t/x, y], \Delta$$

$$\vdash C \text{ [ext } u_j]$$

$$\Gamma, v: s=t \in x, y : T // E, \Delta, x':T, y':T$$

$$\Gamma, z: x, y : T // E, \Delta \vdash s=t \in S \text{ [Ax]}$$

by quotientE i j

$$\Gamma, z: x, y : T // E, \Delta, x':T, y':T$$

$$\vdash E[x', y'/x, y] \in \mathbf{U}_j \text{ [Ax]}$$

$$\Gamma, z: x, y : T // E, \Delta \vdash S \in \mathbf{U}_j \text{ [Ax]}$$

A.5.12 Induktive Typen

$\Gamma \vdash \text{rectype } X_1 = T_{X_1} = \text{rectype } X_2 = T_{X_2} \in \mathbb{U}_j \text{ [Ax]}$
by recEq
 $\Gamma, X : \mathbb{U}_j \vdash T_{X_1}[X/X_1] = T_{X_2}[X/X_2] \in \mathbb{U}_j \text{ [Ax]}$

$\Gamma \vdash s = t \in \text{rectype } X = T_X \text{ [Ax]}$
by rec_memEq j
 $\Gamma \vdash s = t \in T_X[\text{rectype } X = T_X / X] \text{ [Ax]}$
 $\Gamma \vdash \text{rectype } X = T_X \in \mathbb{U}_j \text{ [Ax]}$

$\Gamma \vdash \text{let}^* f_1(x_1) = t_1 \text{ in } f_1(e_1)$
 $= \text{let}^* f_2(x_2) = t_2 \text{ in } f_2(e_2) \in T[e_1/z] \text{ [Ax]}$
by rec_indEq $z T \text{ rectype } X = T_X j$
 $\Gamma \vdash e_1 = e_2 \in \text{rectype } X = T_X \text{ [Ax]}$
 $\Gamma \vdash \text{rectype } X = T_X \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma, P : (\text{rectype } X = T_X) \rightarrow \mathbb{P}_j,$
 $f : (y : \{x : \text{rectype } X = T_X \mid P(x)\} \rightarrow T[y/z]),$
 $x : T_X[\{x : \text{rectype } X = T_X \mid P(x)\} / X]$
 $\vdash t_1[f, x / f_1, x_1] = t_2[f, x / f_2, x_2] \in T[x/z] \text{ [Ax]}$

$\Gamma \vdash \text{rectype } X = T_X \text{ [ext } t_j]$
by rec_memI j
 $\Gamma \vdash T_X[\text{rectype } X = T_X / X] \text{ [ext } t_j]$
 $\Gamma \vdash \text{rectype } X = T_X \in \mathbb{U}_j \text{ [Ax]}$

$\Gamma, z : \text{rectype } X = T_X, \Delta \vdash C$
 $\text{[ext let}^* f(x) = t[\lambda y. \Lambda / P] \text{ in } f(z)]$
by recE $i j$
 $\Gamma, z : \text{rectype } X = T_X, \Delta$
 $\vdash \text{rectype } X = T_X \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma, z : \text{rectype } X = T_X, \Delta$
 $P : (\text{rectype } X = T_X) \rightarrow \mathbb{P}_j,$
 $f : (y : \{x : \text{rectype } X = T_X \mid P(x)\} \rightarrow C[y/z]),$
 $x : T_X[\{x : \text{rectype } X = T_X \mid P(x)\} / X]$
 $\vdash C[x/z] \text{ [ext } t_j]$

$\Gamma, z : \text{rectype } X = T_X, \Delta \vdash C \text{ [ext } t[z/x]]$
by recE_unroll i
 $\Gamma, z : \text{rectype } X = T_X, \Delta,$
 $x : T_X[\text{rectype } X = T_X / X],$
 $v : z = x \in T_X[\text{rectype } X = T_X / X] \vdash C[x/z] \text{ [ext } t_j]$

A.5.13 Rekursive Funktionen

$\Gamma \vdash S_1 \not\rightarrow T_1 = S_2 \not\rightarrow T_2 \in \mathbb{U}_j \text{ [Ax]}$
by pfunEq
 $\Gamma \vdash S_1 = S_2 \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma \vdash T_1 = T_2 \in \mathbb{U}_j \text{ [Ax]}$

$\Gamma \vdash (\text{letrec } f_1(x_1) = t_1)$
 $= (\text{letrec } f_2(x_2) = t_2) \in S \not\rightarrow T \text{ [Ax]}$
by fixEq j
 $\Gamma \vdash \text{letrec } f_1(x_1) = t_1 \in S \not\rightarrow T \text{ [Ax]}$
 $\Gamma \vdash \text{letrec } f_2(x_2) = t_2 \in S \not\rightarrow T \text{ [Ax]}$
 $\Gamma \vdash \{x : S \mid \text{dom}(\text{letrec } f_1(x_1) = t_1)(x)\}$
 $= \{x : S \mid \text{dom}(\text{letrec } f_2(x_2) = t_2)(x)\} \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma, y : \{x : S \mid \text{dom}(\text{letrec } f_1(x_1) = t_1)(x)\}$
 $\vdash (\text{letrec } f_1(x_1) = t_1)(y)$
 $= (\text{letrec } f_2(x_2) = t_2)(y) \in T \text{ [Ax]}$

$\Gamma \vdash f_1(t_1) = f_2(t_2) \in T \text{ [Ax]}$
by apply_pEq $S \not\rightarrow T$
 $\Gamma \vdash f_1 = f_2 \in S \not\rightarrow T \text{ [Ax]}$
 $\Gamma \vdash t_1 = t_2 \in \{x : S \mid \text{dom}(f_1)(x)\} \text{ [Ax]}$

$\Gamma \vdash (\text{letrec } f(x) = t)(u) = t_2 \in T \text{ [Ax]}$
by apply_pBed

$\Gamma \vdash \text{letrec } f(x) = t \in S \not\rightarrow T \text{ [Ax]}$
by fixMem j
 $\Gamma \vdash S \not\rightarrow T \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma, f' : S \not\rightarrow T, x' : S \vdash \mathcal{E}[[t[f', x' / f, x]]] \in \mathbb{U}_j \text{ [Ax]}$
 $\Gamma, f' : S \not\rightarrow T, x' : S, \mathcal{E}[[t[f', x' / f, x]]]$
 $\vdash t[f', x' / f, x] \in T \text{ [Ax]}$

$\Gamma, f : S \not\rightarrow T, \Delta \vdash C \text{ [ext } t[f(s), \text{Axiom} / y, v]]$
by pfunE $i s$
 $\Gamma, f : S \not\rightarrow T, \Delta \vdash s \in \{x : S \mid \text{dom}(f_1)(x)\} \text{ [Ax]}$
 $\Gamma, f : S \not\rightarrow T, \Delta, y : T, v : y = f(s) \in T \vdash C \text{ [ext } t_j]$

$\Gamma \vdash \text{dom}(f_1) = \text{dom}(f_2) \in S \rightarrow \mathbb{P}_j \text{ [Ax]}$
by domEq $S \not\rightarrow T$

A.5.14 Strukturelle Regeln und sonstige Regeln

$$\Gamma, x:T, \Delta \vdash T \text{ [ext } x]$$

by hypothesis i

$$\Gamma \vdash T \text{ [ext } t]$$

by intro t

$$\Gamma \vdash t \in T \text{ [Ax]}$$

$$\Gamma, \Delta \vdash C \text{ [ext } (\lambda x.t) s]$$

by cut $i T$

$$\Gamma, \Delta \vdash T \text{ [ext } s]$$

$$\Gamma, x:T, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma, x:T, \Delta \vdash C \text{ [ext } t]$$

by thin i

$$\Gamma, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma \vdash C \text{ [ext } t]$$

by compute $tagC$

$$\Gamma \vdash C \downarrow_{tagC} \text{ [ext } t]$$

$$\Gamma \vdash C \text{ [ext } t]$$

by rev_compute $tagC$

$$\Gamma, \vdash C \uparrow_{tagC} \text{ [ext } t]$$

$$\Gamma, z:T, \Delta \vdash C \text{ [ext } t]$$

by computeHyp $i tagT$

$$\Gamma, z:T \downarrow_{tagT}, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma, z:T, \Delta \vdash C \text{ [ext } t]$$

by rev_computeHyp $i tagT$

$$\Gamma, z:T \uparrow_{tagT}, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma \vdash C \text{ [ext } t]$$

by unfold $def-name$

$$\Gamma \vdash C \downarrow \text{ [ext } t]$$

$$\Gamma \vdash C \downarrow \text{ [ext } t]$$

by fold $def-name$

$$\Gamma \vdash C \text{ [ext } t]$$

$$\Gamma, z:T, \Delta \vdash C \text{ [ext } t]$$

by unfoldHyp $i def-name$

$$\Gamma, z:T \downarrow, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma, z:T \downarrow, \Delta \vdash C \text{ [ext } t]$$

by foldHyp $i def-name$

$$\Gamma, z:T, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma, z:T, \Delta \vdash C \text{ [ext } t]$$

by replaceHyp $i S j$

$$\Gamma, z:S, \Delta \vdash C \text{ [ext } t]$$

$$\Gamma, z:T, \Delta \vdash T = S \in U_j \text{ [Ax]}$$

$$\Gamma \vdash C \text{ [ext } t]$$

by lemma $theorem-name$

$$\Gamma \vdash t \in T \text{ [Ax]}$$

by extract $theorem-name$

$$\Gamma \vdash C \text{ [ext } t[\sigma]]$$

by instantiate $\Gamma' C' \sigma$

$$\Gamma' \vdash C' \text{ [ext } t]$$

$$\Gamma \vdash C \text{ [ext } t[y/x]]$$

by rename $y x$

$$\Gamma[x/y] \vdash C[x/y] \text{ [ext } t]$$

$$\Gamma \vdash s = t \in T \text{ [Ax]}$$

by equality

$$\Gamma \vdash C \text{ [ext } t]$$

by arith j

$$\Gamma \vdash s_1 \in \mathbf{Z} \text{ [Ax]}$$

$$\vdots$$