

CS 683 Lecture 29

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Scribes: Seth Marvel, Shaomei Wu

ARROW'S CLAIM AND PROOF

We define a *global ranking* (GR) of the set $S = \{a, b, c, \dots\}$ as a function that maps sets of N orderings on S (e.g. $b > a = c > d > \dots$) onto single orderings on S . Note that the orderings, or *voter rankings*, may include both greater than ($>$) and equality ($=$) relations.

By definition, a GR must be

- (i) total, i.e. for all a and b in S , either $a > b$, $b > a$, or $a = b$ in GR
- (ii) transitive, i.e. if $a > b$ and $b > c$ in GR, then $a > c$ in GR

For many applications, we would like the GR to reflect “the general opinion” of the voter rankings. Hence, we may seek GRs that satisfy the following 3 axioms:

- (1) Non-Dictator Axiom: None of the voter rankings is a dictator of GR.
For all a and b in S , let an *ab-dictator* of GR be a voter ranking where (1) the order of a and b in GR is the same as that in the *ab-dictator*, and (2) no changes to the order of a and b in other voter rankings can change the order of a and b in the GR. A *dictator* of GR is an *ab-dictator* of GR for all a and b in S .
- (2) Unanimity Axiom: For all a and b in S , if $a > b$ in all voter rankings, then $a > b$ in GR.
- (3) Axiom of Independence of Irrelevant Alternatives (AIIA): For all a and b in S , the relative order of a and b in GR depends only on the relative order of a and b in the voter rankings.

Arrow proved that no GR satisfies all 3 of the above axioms:

Arrow's Theorem: If a GR satisfies (2) and (3), then it violates (1).

Lemma 1: Consider a set of voter rankings in which each voter ranks b first or last. Then b is either first or last in GR.

Proof: For suppose not. Then there exists some a and c in S such that $a > b > c$ in GR, so $a > c$ in GR (by transitivity of GR). Suppose that we reorder a and c such that $c > a$ in each voter ranking while keeping b in its original position as first or last. Then, by the Unanimity Axiom, $c > a$ in GR. On the other hand, the relative orders of a and b , and b and c are unaffected by re-ranking a and c since b keeps its extreme position. So, by AIIA, we still have $a > b$ and $b > c$ in GR, implying $a > c$ in GR (again, by transitivity). The contradiction proves Lemma 1.

Now consider a set of voter rankings in which each voter assigns b last place. By the Unanimity Axiom, b must also be last in GR. We now select an order for the voters and sequentially move

the b 's from last to first in each ranking. When all b 's have been moved to first, the Unanimity Axiom guarantees b is first in the GR. By Lemma 1, then, there exists a voter ranking v_b such that moving b from last to first in v_b moves b from last to first in GR.

Let **state I** denote the set of voter rankings just before b is moved from last to first in v_b , and **state II** denote the set of voter rankings just after b is moved from last to first in v_b .

Lemma 2: v_b is an ac -dictator of GR in state II.

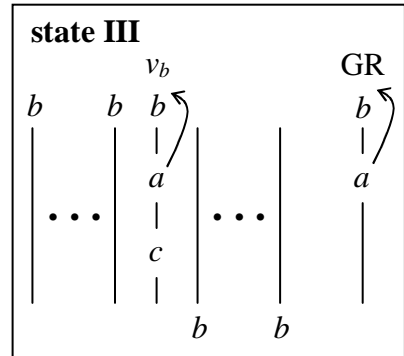
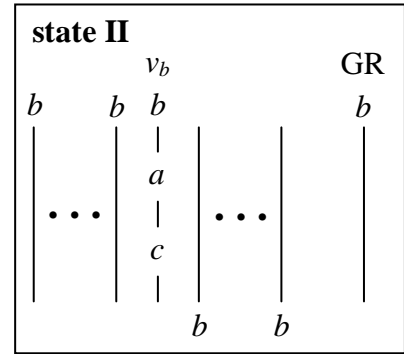
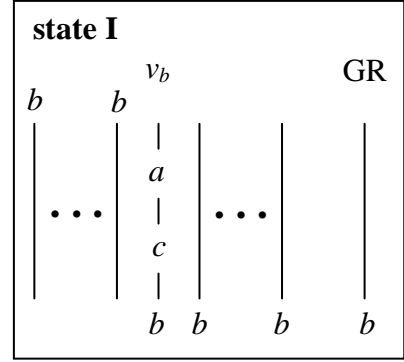
Proof: Without loss of generality, suppose $a > c$ in v_b . We first prove that $a > c$ in GR in state II, the first condition necessary for v_b to be an ac -dictator. Suppose we move a above b in v_b so that $a > b$ in v_b , thereby creating **state III**. In state III, a and b have the same order in all voter rankings as they did in state I, when b was last in GR and $a > b$. So, by AIIA, $a > b$ in GR in state III. Furthermore, in state III, b and c have the same order in all voter rankings as they did in state II, so by AIIA, $b > c$ in GR in state III, implying $a > c$ in GR in state III. Finally, the transition from state II to state III does not change the order of a and c in any voter rankings, so by AIIA, $a > c$ in GR in state II.

To satisfy the second condition for v_b to be an ac -dictator, consider 2 cases:

Case 1: Starting with state II, suppose we change the order of a and c in an arbitrary number of voter rankings other than v_b , while leaving b in its extreme position in each of these rankings. Since the relative order of a and b , and b and c , does not change in any ranking by this rearrangement, the above proof of the first condition goes through and $a > c$ in GR regardless of the rearrangement.

Case 2: Suppose we do another rearrangement similar to that in Case 1, but this time do not require that b remain in its extreme position in all voter rankings. If this case were to change $a > c$ to $a = c$ or $a < c$ in GR, then AIIA would be violated since the order of a and c would depend on their order relative to b . Hence, $a > c$ in GR regardless of these rearrangements.

Lemma 2 asserts that in state II, v_b is an ac -dictator of GR for all a and c that are not b . If we repeat the proof of Lemma 2 with c playing the role of b , we obtain an ab -dictator v_c . As the final step in proving Arrow's Theorem, we prove the following lemma:



Lemma 3: $v_b = v_c$

Proof: Suppose $v_b \neq v_c$ and, without loss of generality, suppose we encounter v_b before v_c in our selected order of voters. Consider a set of voter rankings in which each voter initially ranks a in last place. We now sequentially move the a 's from last to first in each ranking. If the voter rankings all agree that $c > b$, then by the Unanimity Axiom, $c > b$ in GR. Furthermore, just after a is moved from last to first in the ac -dictator v_b , v_b ranks $a > c$, so it must be that $a > c$ in GR. By transitivity of GR, this implies $a > b$ in GR. However, v_c is an ab -dictator and ranks $b > a$ before a is moved from last to first in its ranking, so by contradiction, $v_b = v_c$.

HARE VOTING SYSTEM

(used in faculty rankings at Cornell)

We describe this system via an example. Suppose:

4 individuals rank: $a > b > c$

3 individuals rank: $b > c > a$

2 individuals rank: $c > b > a$

Now, examine the first column. Element c has the fewest first-place votes, so we eliminate c from all columns. Next, we tally the individuals who rank b first among the remaining elements (5) and those who instead rank a first (4). Since the former is greater than the latter, a is eliminated, leaving b the winner.

Unfortunately, the Hare Voting System can be sabotaged, and the mechanism of sabotage sometimes takes a bizarre form. Suppose:

7 individuals rank: $a > b > c > d$

6 individuals rank: $b > a > c > d$

5 individuals rank: $c > b > a > d$

3 individuals rank: $d > c > b > a$

By the algorithm described above, d is eliminated first, then b and finally c , leaving a . However, if the 3 individuals who voted $d > c > b > a$ would like to prevent a from winning and know how the other individuals are voting, they can sabotage a 's victory by moving a 's rank from last to first: $a > d > c > b$. If the voting process is then repeated, d is again eliminated first, then c , and finally a , to leave b the winner.

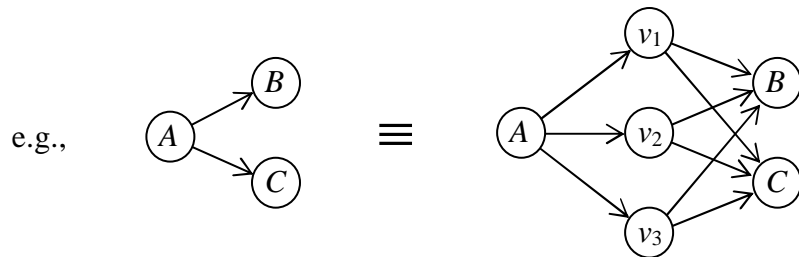
PAGERANK

Suppose our goal is to design an algorithm α that ranks vertices on a directed, strongly connected, unweighted graph. Altman and Tennenholtz[1] proposed 5 intuitively desirable axioms for α and proved only PageRank satisfies all 5 axioms.

Axiom 1: α assigns equal rank to isomorphic vertices.

Axiom 2: Adding a self-loop to vertex v does not change relative ranking of any other vertices.

Axiom 3: Ideal voting by committee ranks the vertices in the same order as direct voting.



On Monday: Axioms 4 and 5, plus a demonstration that PageRank is satisfied by all 5 axioms.

[1] A. Altman, M. Tennenholtz. Ranking Systems: The PageRank Axioms. In *Proceedings of the 6th ACM Conference on Electronic Commerce (EC '05)*.