

Recall our question: how do you tell if two ~~things~~ "things" are different?
(statistical)

What is ~~language~~ are language models for? For the purposes of this class:
 → giving a ~~short~~ compact representation of what the language is 'like'.
 → comparing two language sources. (ex: no country)

We define a language model as a distribution over all strings w that we wish to consider
 "ok": "legal".

~~Some distributions aren't model-based per se, e.g.~~
 ex: HMMs (~~for now, no a lot's~~)

[a good counterexample to key in mind:
 the Poisson over strings of form "dude"
 - finite # of params, no equivalent PCFs.
 (see Booth & Thompson '73)]

• a finite non-empty set of states $q_1, \dots, q_M, q_b, q_e$
 distinguished "start" begin distinguished "end"

- a finite vocabulary $V \cup \{\langle \text{begin} \rangle, \langle \text{end} \rangle\}$, not in V .

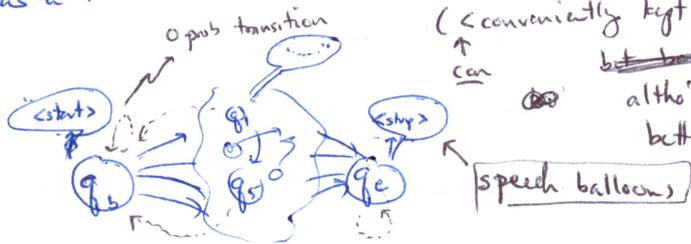
- each state ~~def~~ has an associated emission distribution over V^* (all sequences of items from V , including the empty string)

q_b emits only $\langle \text{begin} \rangle$
 q_e $\langle \text{end} \rangle$
 q_i does not emit any sequence containing those special items

this might be itself a (state-specific) L.M.

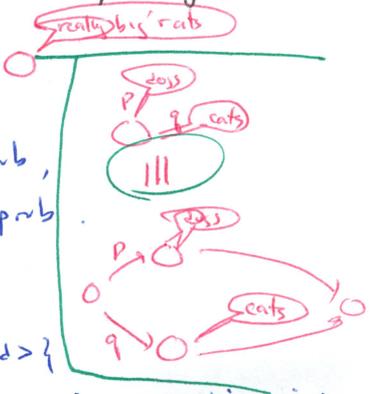
- each state has a transition distribution over all other states, s.t:
 (conveniently kept in a matrix, all rows sum to 1)
 can be ~~bad~~ but ~~bad~~ if altho' maybe a sparse rg would be better.

draw 1st
 - drawing really uncalibrated a lot



will further assume every q_i reachable from q_b w/ non-zero prob,
 q_i can reach q_e w/ non-zero prob
 → these definitely seem like necessary

~~$p(s) = \dots$~~
 $p(s) = \text{sum of probs of all paths generating } s, s \in \{\langle \text{begin} \rangle\} V^* \{\langle \text{end} \rangle\}$
 product of all the transition probs x prob of corresponding emissions.



nb: should make clear diff. b/w termination and progress (i.e. ~~the~~ and the l.m. itself being proper.

(so you can think of the emissions as also ~~def~~ part of what defines a path.

What does this tell you?

- broken ^{we have} & distribution over a potentially ∞ # of sequences to sth definable
- by a finite set of parameters (indeed, you have a model θ , so you're managing the complexity.

- structure of this model gives you insight into what is "preferred".

- learning algorithms for HMMs, either from labelled or unlabelled data, (altho' you have to set the # of states beforehand)

learnable two training-data paradigms: state-labelled data, unlabelled data

need to specify M ahead of time. Baum-welch / EM for max-likelihood estimation (tends to find saddle points or local maxima, b/c there's a lot) - REMEMBER: need constraints (ow., just set all params to ∞ !)

- why are n-gram models so much more common than HMMs?

- here, the data basically has state labels! Nothing is really "hidden".

$$P(\langle \text{cat} \rangle \text{ dog, dog } \langle \text{stop} \rangle)$$

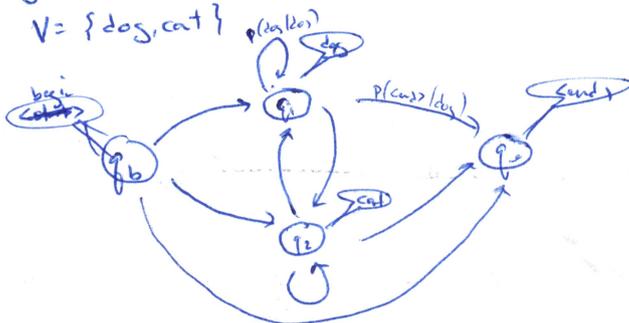
$$P(\rightarrow) \cdot 1 \cdot P(\rightarrow) \dots$$

vs. unigram model:



ex: a "bigram lm" as an HMM:

$V = \{\text{dog, cat}\}$



$q_i = \text{"just said } v_i \text{"}$

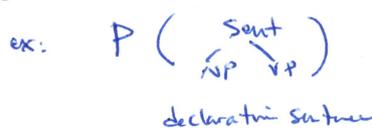
- linear dependence of what is to be next generated

~~we can build more co~~

- ~~but it's not just finite state~~ - we aren't just restricted to finite state models

ex: probabilistic context-free grammars (PCFGs) - analogous defns for dependency grammars. aka stochastic CFGs (SCFGs)

- each category type has a distribution over its decompositions, a finite set:



~~what~~

$P(s) = \sum_{\theta} \text{sum of probs of each trees generating } s$

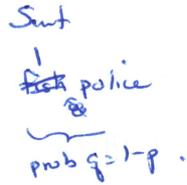
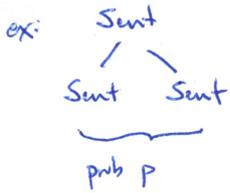
↳ product of each decomposition used in a tree

$$P \left(\begin{array}{c} \text{sent} \\ \swarrow \quad \searrow \\ \text{NP} \quad \text{VP} \\ | \quad | \\ \text{they} \quad \text{left} \end{array} \right) = P \left(\begin{array}{c} \text{sent} \\ \swarrow \quad \searrow \\ \text{NP} \quad \text{VP} \\ | \quad | \\ \text{they} \quad \text{left} \end{array} \right) P(\text{NP} \mid \text{they}) P(\text{VP} \mid \text{left})$$

q: ~~does this~~ does every choice of valid 'decomposition' distribution. 'india' a proper prob on all strings? Turns out it's tricky.

- [Booth; Thompson '73] conditions under which you would get a proper distribution over all strings.

[Chi & Gema '98]



For this $\theta = (p)$ (that's all we need)

$$P_0(\text{"police"}) = q$$

$$P_0(\text{"police police"}) = p \cdot q \cdot q = p q^2$$

$$P_0(\text{"police police police"}) = [p \cdot p \cdot q^3] + [p \cdot q \cdot p \cdot q^2] = 2 p^2 q^3$$

$$\sum_s P_0(s) = \sum_{i=1}^{\infty} \# \text{ of 'bracketings' of length-} i \text{ sequence} \cdot q^i \cdot p^{i-1}$$

Let $X_d = \text{prob of all trees w/ depth } \leq d$ (fringe contains only words)

$$X_{d=1} = q$$

$$X_2 = q + p q^2$$

$$X_i = q + p X_{i-1}^2$$

$$X_3 = q + p q^2 + p (q + p q^2)^2 (q + p q^2)^2$$

apparently converges to $\min(1, \frac{q}{p})$

$$= \sum_{i=1}^{\infty} (\# \text{ bracketings}(i)) \cdot q^i (q p)^{i-1}$$

$$\sum_{i=1}^{\infty} \frac{1}{i} \binom{2(i-1)}{i-1} q (q p)^{i-1}$$

(presumably the same geometric sum?)

$$\approx q \sum_{i=1}^{\infty} \frac{1}{i} \frac{4^{i-1}}{\sqrt{\pi(i-1)}} (q p)^{i-1}$$

generate i leaves
got i "parents" of those leaves.

abandon this

if $p > 1/2$, this will not be proper (too much probability on the "expansion" p as opposed to the "generation of words").

ex: content model: [Bergsby & Lee '04]

→ (learn) how topics relate to each other w/in docs.

→ potentially relevant to projects:
< probably have fixed set of topics >

(etc. stat)

talked about LM on just top-most frequent words

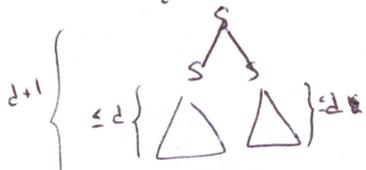
stylistic

what do you do w/ the 'other' words to make them not dominate?

entropy → fixed-max-length strings.

2nd attempt, using Chi & Gema's approach:

let $X_d = \text{prob of all trees w/ depth } \leq d$ (fringe contains only words)



→ $p \cdot X_d \cdot X_d$ and then $\frac{S}{\text{police}} = q$

$$X_{d+1} = q + p X_d^2$$

guess that $X_{d+1} \geq X_d \geq 0$ ~~...~~ $\leq q + p X_{d+1}^2$

If we assume convergence, then we solve for $X = q + p X^2$, or $0 = p X^2 - X + q$.

By quadratic formula, roots are:

$$\Rightarrow 0 = (p x - (1-p))(x-1)$$

$$= (p x - q)(x-1)$$

solutions are $x = \frac{q}{p}$ or $x = 1$.

(thanks to Chanhoo Tan for idea of assuming convergence)