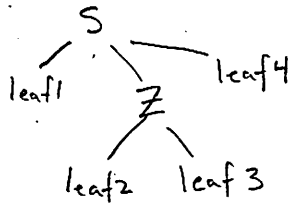


1. Let  $w = w_1 w_2 \dots w_n$  be a length- $n$  sentence. For now, assume no adjunction.

When is an initial tree  $\alpha$  be the root of a derivation tree for  $w$ ?



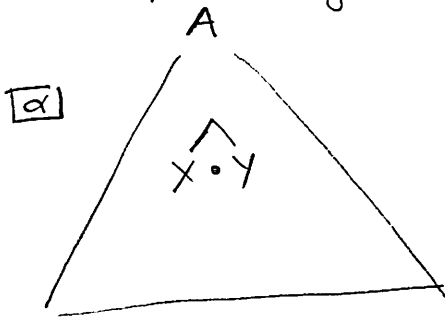
... if  $w = w_1 w_2 w_3 w_4$  for some substrings  $w_i$ , and each  $\text{leaf}_i \Rightarrow^* w_i$ .  
 (Either  $\text{leaf}_i$  is a sequence of terminals or substitution of some tree  $\alpha$  into  $\text{leaf}_i$  "eventually" yields a fringe that is  $w_i$ .)

This suggests a "left-to-right" search of the leaves of a possible start tree:

for  $i = 1 \dots \# \text{ of leaves}$ :

given what  $\text{leaf}_1 \dots \text{leaf}_{i-1}$  cover, figure out what tree could substitute into  $\text{leaf}_i$  and what  $w_i$  is a section of  $w$  is.

2. Dot notation: for tracking search in an elementary tree.



"everything to the left has been checked"  
 "everything to the right must be verified"

state:  $[\alpha, \text{address of } X \text{ in } \alpha, \text{right}, \quad ]$   
 $\equiv [\alpha, \text{address of } Y \text{ in } \alpha, \text{left}, \quad ]$

3. Dynamic programming idea:

