## Expansion Moves, Approximation Bound, and Slanted Surfaces

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Swap moves not equal to an approximation algorithm 1 2 3 a,b,c  $d(a,b) = d(b,c) = \frac{k}{2}$  d(a,c) = kCost of having label a next to label b 1 2 3 а 0 k k b k 0 k 2 2 0 с a b c Total cost to have this setup is k Minium cost would be to have this arrangement c c c Cost = 4**α**, f find f'=avgerage minium of E( f ) f when  $\alpha$  - expansion of f а α а  $\alpha$  compete with everyone else α d-link costs not a problem q р n-link costs are a problem а

In event of assign p & q to a no n-link cost for same terminal

When  $f(p) \neq f(q)$  then there is a cost

No Cost or Free f'(p) = f'(q) =  $\alpha$ Should be expensive when f'(p) $\neq$ f'(q) Should be expensive when f'(p) = f'(q) =  $\overline{a}$  if and only if f(p)  $\neq$  f(q)



 $f^{a}$  expand  $\hat{f}$  to make A= $\alpha$  $\hat{f}$  local minimum  $E(f) \leq E(f^{a})$  $E_{A}(f) =$  Energ involving A data + smoothness c(p,l) discontinuities p,q  $p \in A$  p or  $q \in A$ 

$$E_{A}(\hat{f}) \leq E_{A}(f^{a}) \leq E_{A}(f^{*})$$

$$\sum_{A} E_{A}(\hat{f}) \leq E_{A}(f^{a})$$

$$\hat{f} \leq E(f^{*}) + \mathbf{l} \# disc(f^{*})$$

$$E(\hat{f}) \leq 2E(f^{*})$$

## POTTS MODEL

Potts model prefers piecewise constant solutions



Ideal for the potts model

L varies Set L =  $Z^+$  {0,1,...,16}

Sloped surface looks like connected components with segmentation



Consider each segment



Find a plane that best explains the intensities of the segment Hypothesize a plane J that will fit this segment



Eventually, find a plane equation for the segment with a slope These plane equations become L

Now run again using the new label set.

One of the plane equations using the potts model will win.