Perspectives on Threaded and Asynchronous Programming

Oliver Kennedy Advanced Systems Lecture 2

Threaded vs Asynchronous

- CPUs are faster than most hardware
- Programs need to wait for hardware
- How does a program continue processing while waiting on hardware?

Poll and Process

- Continually poll the hardware for readiness
- Extremely inefficient

Threads

- Multiple threads processing data at once
- While one thread waits, the rest continue
- o Benefits
 - Intuitive, Multiprocessor Support,
 Modular code can't break the system
- o Drawbacks
 - o Inefficient, Synchronization issues

Asynchronous Programming

- o One thread processes events
- Hardware readiness treated as an event
- o Benefits
 - o Efficient, Linear Execution, No Starvation
- o Drawbacks
 - Unintuitive, Problems with Modular Code,
 Limited Concurrency, Tasks need to be short

Threads in Interactive Systems

- o Programmers use threads a lot
- A lot of modern problems have to do with threads
- What can we learn from this?
 - What do programmers use threads for?
 - What mistakes are made with threads?
 - How can threads be made more efficient?

Cedar and GVX

- A case study on two OSes
- Cedar and GVX use the Mesa language's thread system
 - Mesa supports standard primitives
 - ... and a strict priority scheduler

And they found...

- o A bird's eye view: The Profiler
- Three classes of threads
 - Eternal Threads
 - Worker Threads
 - Transient Threads
- Cedar vs GVX
 - Free and loose vs Small and Efficient

Thread Paradigms

- o Defer Work
- Pumps/Slack Processes
- Sleepers/One Shots
- Deadlock Avoiders
- Task Rejuvenation
- o Serializers
- Concurrency Exploiters
- Encapsulated Forks

Thread Problems-What works

- Sleepers, One shots, Pumps, Work-Deferrers all implemented properly
 - Yet these require little inter-thread interaction
- Concurrency Exploiters were new at the time
 - Work has been done since

Problems-Time Constraints

- High priority slack-processes can be hard to write for use with low priority threads
 - Yield and a strict priority scheduler don't play nice
- Solution: Add a YieldButNotToMe primitive

Problems-Priorities

- Synchronization primitives cause priority inversions in strict-priority schedulers
- Solution 1: High priority threads donate cycles to threads holding locks they need
- Solution 2: A high priority thread that periodically grants a time slice to a thread chosen at random

Problems-Misunderstandings

- Mesa implements locks in an unusual way
 - Programmers write code that might be correct in some circumstances
- Bugs introduced this way are hard to track (the code looks right)

Problems-Treating the Symptom

- A common problem
- Ex: Fixing a wait without a corresponding notify by adding a timeout to the wait
- Introduces delays and possibly bugs

Problems-Changing Hardware

- Magic Numbers
 - Timeouts and pause lengths based on one processor become invalid when a faster processor comes out
- Memory Ordering
 - Much code assumes strict memory ordering

Problems-Library Implementation

- o Notify, Yield, and Scheduling
- Strict priority scheduling sucks
- Time quantum
 - o (not a problem, but a consideration)

Future Work

- Analyze more systems!
- Come up with new scheduling techniques
- Keep analyzing known code

SEDA

- o The internet is big... really big
- Loads are getting bigger
- Dynamic content becoming prevalent
- o Services need to adapt to these loads

But how can we adapt?

- o Can't we use threads?
- What if we only used so many threads?
- Weren't you talking about some asynchronous nonsense earlier?
- How about a mix?

SEDA

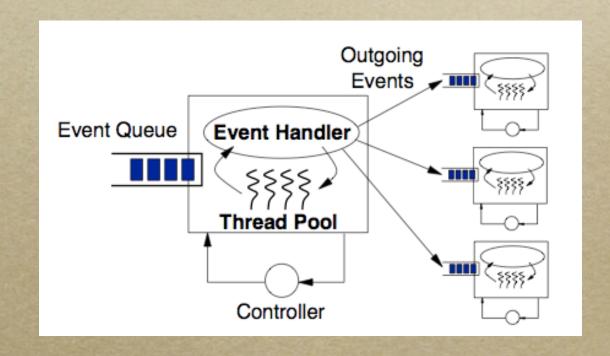
- A means of building scalable web services
 - Has to support concurrency
 - Has to be easy to program
 - Has to let the application manage load
 - Has to tune itself

Stages

- Any task can be broken down
- SEDA breaks tasks down into stages
- A stage has an input queue
- A full web-service is multiple stages networked together

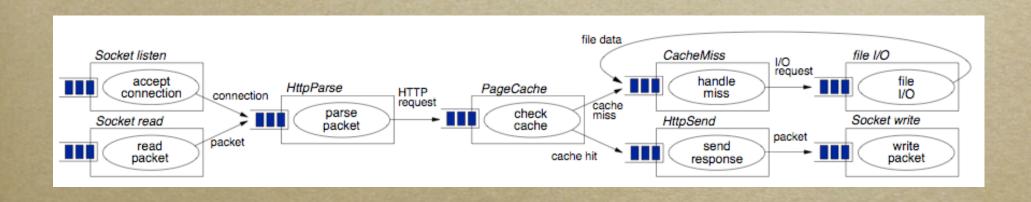
Stages

- o Event Queue
- o Thread Pool
 - Automatic tuning
- Controller
 - Feedback
- Event Handler



Stages

Sales of the sales



Why stages?

- Allows for isolation
- Fine grained tuning
- Easier debugging

Self-Tuning

- Each stage has an associated controller
- o Thread Pool Size
 - More threads = More concurrency (up to a point)
- Request Batching
 - o Cache Locality, Task Aggregation

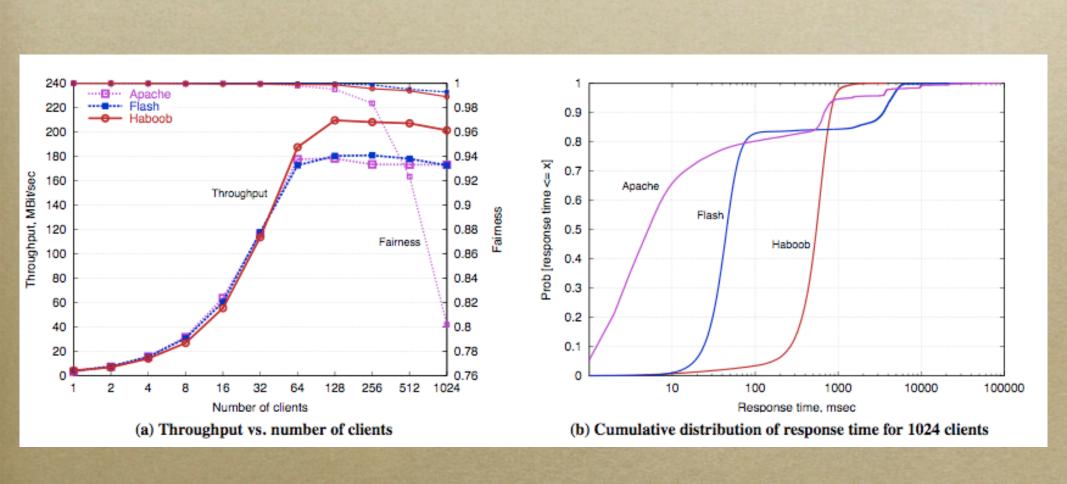
Sandstorm

- Java based implementation of SEDA
 - Simple memory management
- o Provides APIs
 - o Queue and stage management
 - o Profiling/debugging

Haboob/GnutellaServer

- An implementation of common web services on top of Sandstorm
- Both performed admirably
- Haboob (despite being written in Java) had better performance characteristics than Apache under high loads.

Haboob vs Apache vs Flash



Threads vs Hybrids

- The age old conflict
 - Monolithic vs Microprogramming
- Threads expose more
 - o It's 12 years later, the tech is here
- SEDA is more elegant
 - o But it's in Java...

Any Questions?