

Adjunctions

Ross Tate

October 5, 2014

Exercise 1. Prove that the inclusion functor $\mathbf{Set} \xrightarrow{I} \mathbf{Rel}$ has a right adjoint. You may use any of the equivalent definitions of adjunction. For clarification, I is the functor mapping each set X (an object of \mathbf{Set}) to the set X (also an object of \mathbf{Rel}) and each function $X \rightarrow Y$ (a morphism of \mathbf{Set}) to the relation $\lambda\langle x, y \rangle. f(x) = y$ (a morphism of \mathbf{Rel}).

Proof. There is a functor $\mathbb{P} : \mathbf{Rel} \rightarrow \mathbf{Set}$ mapping each set X to the set $\mathbb{P}X$ and each relation $R : X \times Y \rightarrow \mathbf{Prop}$ to the function $\lambda\vec{x}. \{y : Y \mid \exists x \in \vec{x}. x R y\}$. The identity relation $\lambda\langle x_1, x_2 \rangle. x_1 = x_2$ gets mapped to the function $\lambda\vec{x}. \{x_2 : X \mid \exists x_1 \in \vec{x}. x_1 = x_2\}$ which is simply the identity function. The composition relation $\lambda\langle x, z \rangle. \exists y : Y. x R_1 y \wedge y R_2 z$ gets mapped to the function $\lambda\vec{x}. \{z : Z \mid \exists x \in \vec{x}. \exists y : Y. x R_1 y \wedge y R_2 z\}$ which equals $\lambda\vec{x}. \{z : Z \mid \exists y \in \{y : Y \mid \exists x \in \vec{x}. x R_1 y\}. y R_2 z\}$, proving distributivity.

Given a \mathbf{Rel} -object Y , define the \mathbf{Rel} -morphism $\varepsilon_Y : I(\mathbb{P}(Y)) \rightarrow Y$ to be the binary relation $\lambda\langle \vec{y}, y \rangle. y \in \vec{y}$. Given another \mathbf{Rel} -morphism from some $I(X)$ to Y , i.e. a binary relation $R : X \times Y \rightarrow \mathbf{Prop}$, the unique corresponding \mathbf{Set} -morphism from X to $\mathbb{P}(Y)$ is the function $\lambda x. \{y : Y \mid x R y\}$. The \mathbf{Rel} -composition $I(\lambda x. \{y : Y \mid x R y\}); (\lambda\langle \vec{y}, y \rangle. y \in \vec{y})$ is by definition the binary relation $\lambda\langle x, y \rangle. \exists \vec{y} : \mathbb{P}(Y). \{y : Y \mid x R y\} = \vec{y} \wedge y \in \vec{y}$, which is equivalent to simply R . Furthermore, for any function $f : X \rightarrow \mathbb{P}Y$, the composition $\lambda\langle x, y \rangle. \exists \vec{y} : \mathbb{P}Y. f(x) = \vec{y} \wedge y \in \vec{y}$ is equivalent to $\lambda\langle x, y \rangle. y \in f(x)$, which is equivalent to R if and only if $f(x) = \{y : Y \mid \exists x : X. x R y\}$, making the function corresponding to R unique. \square

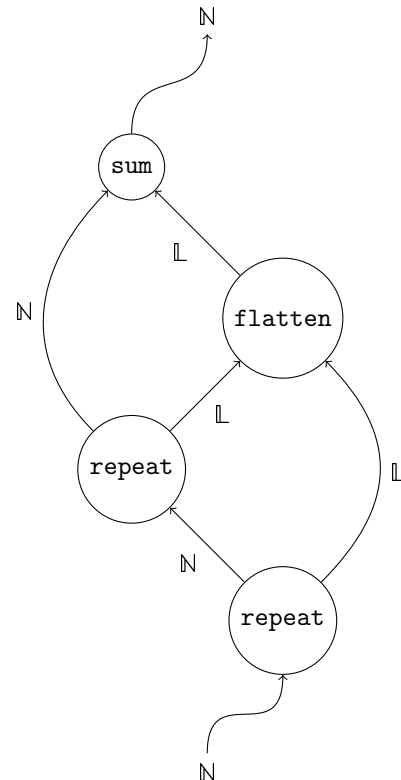
Exercise 2. There is a functor from $\mathbf{1}$ to \mathbf{Set} picking out the empty set, and another functor from $\mathbf{1}$ to \mathbf{Set} picking out the singleton set. One is the left adjoint to the unique functor from \mathbf{Set} to $\mathbf{1}$, and the other is the right adjoint to the unique functor from \mathbf{Set} to $\mathbf{1}$. Determine and prove which is which.

Proof. The functor $F : \mathbf{1} \rightarrow \mathbf{Set}$ picking out the empty set is the left adjoint to the unique functor $\langle \rangle$ from \mathbf{Set} to $\mathbf{1}$ (whose only object we call \star). For any $X : \mathbf{Set}$ and $\star : \mathbf{1}$, both $M_{\mathbf{Set}}(F(\star), X)$ and $M_{\mathbf{1}}(\star, \langle \rangle(X))$ have only one element, making them isomorphic. Furthermore, since $\mathbf{1}$ has only one morphism, this isomorphism is guaranteed to be natural, making this an adjunction.

The functor $G : \mathbf{1} \rightarrow \mathbf{Set}$ picking out the singleton set is the right adjoint to the unique functor $\langle \rangle$ from \mathbf{Set} to $\mathbf{1}$ (whose only object we call \star). For any $X : \mathbf{Set}$ and $\star : \mathbf{1}$, both $M_{\mathbf{1}}(\langle \rangle(X), \star)$ and $M_{\mathbf{Set}}(X, G(\star))$ have only one element, making them isomorphic. Furthermore, since $\mathbf{1}$ has only one morphism, this isomorphism is guaranteed to be natural, making this an adjunction. \square

Exercise 3. $\mathbb{N} : \mathbf{1} \rightarrow \mathbf{Set}$ maps the only object of $\mathbf{1}$ to the set \mathbb{N} . `repeat` is the natural transformation from \mathbb{N} to $\mathbb{N}; \mathbb{L}$ (i.e. $\mathbb{L}(\mathbb{N})$) mapping the sole object of $\mathbf{1}$ to the function mapping n to the length- n list $[n, \dots, n]$. `sum` is the natural transformation from $\mathbb{N}; \mathbb{L}$ to \mathbb{N} mapping the sole object of $\mathbf{1}$ to the function mapping a list of numbers and returns its sum.

The string diagram to the right denotes a natural transformation from the functor $\mathbb{N} : \mathbf{1} \rightarrow \mathbf{Set}$ to itself (\mathbb{N} maps the only object of $\mathbf{1}$ to the set \mathbb{N}). In particular, this means it describes a function from \mathbb{N} to \mathbb{N} . Determine what that function is in terms of basic arithmetic. (No proof necessary; the purpose of this is to learn the notation.)



Proof. The function is $\lambda n. n^3$. The program described by the diagram is $\lambda n. \text{sum}(\text{flatten}(\text{map}_{\text{repeat}}(\text{repeat}(n))))$. \square