

Nulls

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Definition ((Biased) Semigroup). A tuple $\langle S, *, \mathbf{a} \rangle$ where the components have the following types:

Underlying Set S : Type

Operator $*$: $S \times S \rightarrow S$ (infix)

Associativity \mathbf{a} : $\forall s_1, s_2, s_3 : S. (s_1 * s_2) * s_3 = s_1 * (s_2 * s_3)$

Example. $\mathbb{N}_{\min} = \langle \mathbb{N}, \min, * \rangle$

$(\mathbb{L}_+T)_{++} = \langle \mathbb{L}_+T, ++, \cdot \rangle$ where \mathbb{L}_+T denotes nonempty (finite) lists of T

$(\mathbb{M}_+T)_+ = \langle \mathbb{M}_+T, +, \cdot \rangle$ where \mathbb{M}_+T denotes nonempty multisets of T

$(\mathbb{S}_+T)_\cup = \langle \mathbb{S}_+T, \cup, \cdot \rangle$ where \mathbb{S}_+T denotes nonempty finite subsets of T

$(\mathbb{S}_+T)_\cap = \langle \mathbb{S}_+T, \cap, \cdot \rangle$

$(\mathbb{P}_+T)_\cup = \langle \mathbb{P}_+T, \cup, \cdot \rangle$ where \mathbb{P}_+T denotes nonempty subsets of T

$(\mathbb{P}_+T)_\cap = \langle \mathbb{P}_+T, \cap, \cdot \rangle$

Definition ((Biased) Semigroup Homomorphism from $\langle S, *, \cdot \rangle$ to $\langle T, +, \cdot \rangle$). A tuple $\langle f, \mathfrak{d} \rangle$ where the components have the following types:

Underlying Function f : $S \rightarrow T$

Distributivity \mathfrak{d} : $\forall s_1, s_2 : S. f(s_1) + f(s_2) = f(s_1 * s_2)$

Example. $\langle \lambda n. n + 1, \cdot \rangle$ from \mathbb{N}_{\min} to \mathbb{N}_{\min} (and from \mathbb{N}_{\max} to \mathbb{N}_{\max}).

Exercise 1. Prove that **Mon** is a non-full subcategory of **Sgr**, the category of semigroups and semigroup homomorphisms (with the obvious composition and identity), via the obvious inclusion functor.

Definition ((Biased) Group). A tuple $\langle G, *, \mathbf{a}, e, \mathbf{i}, {}^{-1}, \mathbf{inv} \rangle$ where the components have the following types:

Underlying Set G : Type

Operator $*$: $G \times G \rightarrow G$ (infix)

Associativity \mathbf{a} : $\forall g_1, g_2, g_3 : G. (g_1 * g_2) * g_3 = g_1 * (g_2 * g_3)$

Identity Element e : G

Identity \mathbf{i} : $\forall g : G. e * g = g = g * e$

Inverse Operator ${}^{-1}$: $G \rightarrow G$ (postfix)

Inverse \mathbf{inv} : $\forall g : G. g * g^{-1} = e = g^{-1} * g$

Definition ((Biased) Group Homomorphism from $\langle G, *, \cdot, e, \cdot, {}^{-1}, \cdot \rangle$ to $\langle H, +, \cdot, i, \cdot, -, \cdot \rangle$). A tuple $\langle f, \mathfrak{d}, \mathbf{i}, \mathbf{inv} \rangle$ where the components have the following types:

Underlying Function $f: G \rightarrow H$

Distributivity $\delta: \forall g_1, g_2 : G. f(g_1) + f(g_2) = f(g_1 * g_2)$

Identity $i: i = f(e)$

Inverse inv: $\forall g : G. -f(g) = f(g^{-1})$

Exercise 2. Prove that **Grp**, the category of groups and group homomorphism (with the obvious composition and identity), is a full subcategory of **Mon**.

Definition (Reflection Arrow for $\mathbf{S} \xrightarrow{I} \mathbf{C}$). An object C of \mathbf{C} and object \mathcal{R} of \mathbf{S} with a morphism $C \xrightarrow{r} I(\mathcal{R})$ of \mathbf{C} such that for every object S of \mathbf{S} with a morphism $C \xrightarrow{m} I(S)$ of \mathbf{C} there exists a unique morphism $\mathcal{R} \xrightarrow{m'} S$ of \mathbf{S} with $r; I(m') = m$.

Exercise 3. Prove that, for $\mathbf{Grp} \hookrightarrow \mathbf{Mon}$, the identity on a monoid is a reflection arrow if and only if that monoid is a group.

Exercise 4. Prove that, for $\mathbf{Mon} \hookrightarrow \mathbf{Sgr}$, the identity on a semigroup is never a reflection arrow even if that semigroup is a monoid.

Definition (Reflective Subcategory). A subcategory $\mathbf{S} \xrightarrow{I} \mathbf{C}$ with a reflection arrow $C \xrightarrow{r_C} I(\mathcal{R}_C)$ for every object C of \mathbf{C} .

Exercise 5. Prove that **Grp** is a reflective subcategory of **Mon**, and that **Mon** is a reflective subcategory of **Sgr**.

Exercise 6. Prove that every reflective subcategory $\mathbf{S} \xrightarrow{I} \mathbf{C}$ has a unique way to extend a function $R(C) = \mathcal{R}_C$ to a functor so that the following diagram commutes for every morphism $C_1 \xrightarrow{m} C_2$ of \mathbf{C} (meaning all paths are equal):

$$\begin{array}{ccc} C_1 & \xrightarrow{r_{C_1}} & I(\mathcal{R}_{C_1}) \\ m \downarrow & & \downarrow I(R(m)) \\ C_2 & \xrightarrow{r_{C_2}} & I(\mathcal{R}_{C_2}) \end{array}$$

Exercise 7. Prove that, for a reflective subcategory $\mathbf{S} \xrightarrow{I} \mathbf{C}$, the subcategory $\langle \mathbf{S}, I \rangle$ is full if and only if for every object S of \mathbf{S} the identity on $I(S)$ is a reflection arrow.