## Limits

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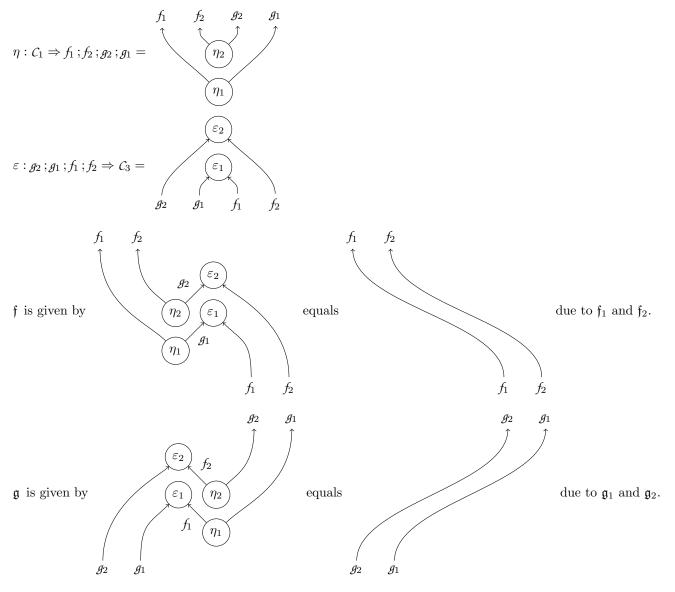
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**Exercise 1.** Prove that for any object  $\mathcal{A}$  of any category **C**, the object  $\mathcal{A} \& \top$  (if it exists) is isomorphic to  $\mathcal{A}$ .

*Proof.* Let  $f : \mathcal{A} \to \mathcal{A} \& \top$  be defined as  $\langle id_{\mathcal{A}}, \langle \rangle \rangle$ . Let  $g : \mathcal{A} \& \top \to \mathcal{A}$  be defined as  $\pi_{\mathcal{A}}$ .  $f ; g = \langle id_{\mathcal{A}}, \langle \rangle \rangle; \pi_{\mathcal{A}} = id_{\mathcal{A}}$  by the nature of projections, proving one direction.  $g ; f ; \pi_{\top} : \mathcal{A} \& \top \to \top$  equals  $id_{\mathcal{A}\&\top}; \pi_{\top} : \mathcal{A} \& \top \to \top$  since all morphisms from the same domain to  $\top$  are equal by the nature of terminal objects. Also,  $g ; f ; \pi_{\mathcal{A}} = \pi_{\mathcal{A}}; \langle id_{\mathcal{A}}, \langle \rangle \rangle; \pi_{\mathcal{A}} = \pi_{\mathcal{A}}; \langle id_{\mathcal{A}}, \langle \rangle \rangle; \pi_{\mathcal{A}} = \pi_{\mathcal{A}}; id_{\mathcal{A}} = \pi_{\mathcal{A}} = id_{\mathcal{A}\&\top}; \pi_{\mathcal{A}}$ . Thus, since all morphisms to a product that behave the same after being followed by both projections must be equal, we have  $g ; f = id_{\mathcal{A}\&\top}$ . So, f and g are inverses of each other, making  $\mathcal{A}$  isomorphic to  $\mathcal{A}\&\top$ .

**Exercise 2.** Prove that, in any 2-category, if morphisms  $C_1 \xrightarrow{f_1} C_2$  and  $C_2 \xrightarrow{f_2} C_3$  are both left adjoints, then their composition  $f_1$ ;  $f_2$  is also a left adjoint.

*Proof.* Let  $\langle C_1, C_2, f_1, g_1, \eta_1, \varepsilon_1, \mathfrak{f}_1, \mathfrak{g}_1 \rangle$  and  $\langle C_2, C_3, f_2, g_2, \eta_2, \varepsilon_2, \mathfrak{f}_2, \mathfrak{g}_2 \rangle$  be some adjunctions that  $f_1$  and  $f_2$  are the left adjoints of. Then  $\langle C_1, C_3, f_1; f_2, g_2; g_1, \eta, \varepsilon, \mathfrak{f}, \mathfrak{g} \rangle$  is an adjuction that  $f_1; f_2$  is the left adjoint of, where  $\eta, \varepsilon, \mathfrak{f}$ , and  $\mathfrak{g}$  are defined as follows:



**Exercise 3.** The monoid  $\mathcal{A} \& \mathcal{B}$  is commutative if both  $\mathcal{A}$  and  $\mathcal{B}$  are commutative, and in that case is (with the appropriate projection homomorphisms) also the product of  $\mathcal{A}$  and  $\mathcal{B}$  in **CommMon**. Prove that there are morphisms  $\kappa_{\mathcal{A}}$  and  $\kappa_{\mathcal{B}}$  demonstrating that  $\mathcal{A} \& \mathcal{B}$  is also the coproduct of  $\mathcal{A}$  and  $\mathcal{B}$  in **CommMon**. That is, prove that **CommMon** has *biproducts*, meaning it has products and coproducts and they coincide on objects.

*Proof.* Define the underlying function of  $\kappa_{\mathcal{A}} : \mathcal{A} \to \mathcal{A} \& \mathcal{B}$  to be  $\lambda a. \langle a, e_{\mathcal{B}} \rangle$ . This clearly preserves identity and multiplication, making  $\kappa_{\mathcal{A}}$  a monoid homomorphism. Similarly,  $\kappa_{\mathcal{B}} = \langle \lambda b. \langle e_{\mathcal{A}}, b \rangle, ., . \rangle$ .

Given a monoid C and monoid homomorphisms  $f_{\mathcal{A}} : \mathcal{A} \to C$  and  $f_{\mathcal{B}} : \mathcal{B} \to C$ , define the underlying function of  $[f] : \mathcal{A} \& \mathcal{B} \to C$  to be  $\lambda \langle a, b \rangle$ .  $f_{\mathcal{A}}(a) +_{\mathcal{C}} f_{\mathcal{B}}(b)$ . We have to show this is a monoid morphism. It distributes since  $[f](\langle a_1, b_1 \rangle +_{\mathcal{A}\&\mathcal{B}} \langle a_2, b_2 \rangle) = [f](\langle a_1 +_{\mathcal{A}} a_2, b_1 +_{\mathcal{B}} b_2 \rangle) = f_{\mathcal{A}}(a_1 +_{\mathcal{A}} a_2) +_{\mathcal{C}} f_{\mathcal{B}}(b_1 +_{\mathcal{B}} b_2) = f_{\mathcal{A}}(a_1) +_{\mathcal{C}} f_{\mathcal{A}}(a_2) +_{\mathcal{C}} f_{\mathcal{B}}(b_1) +_{\mathcal{C}} f_{\mathcal{B}}(b_1) +_{\mathcal{C}} f_{\mathcal{B}}(b_1) +_{\mathcal{C}} f_{\mathcal{B}}(b_1) +_{\mathcal{C}} f_{\mathcal{B}}(b_1) +_{\mathcal{C}} f_{\mathcal{B}}(b_2) = [f](\langle a_1, b_1 \rangle) +_{\mathcal{C}} [f](\langle a_2, b_2 \rangle)$ . It preserves identity since  $[f](e_{\mathcal{A}\&\mathcal{B}}) = [f](\langle e_{\mathcal{A}}, e_{\mathcal{B}} \rangle) = f_{\mathcal{A}}(e_{\mathcal{A}}) +_{\mathcal{C}} f_{\mathcal{B}}(e_{\mathcal{B}}) = e_{\mathcal{C}} +_{\mathcal{C}} e_{\mathcal{C}} = e_{\mathcal{C}}$ .

Given an a: A,  $[f](\kappa_{\mathcal{A}}(a)) = [f](\langle a, e_{\mathcal{B}} \rangle) = f_{\mathcal{A}}(a) + f_{\mathcal{B}}(e_{\mathcal{B}}) = f_{\mathcal{A}}(a) + e_{\mathcal{C}} = f_{\mathcal{A}}(a)$ , so  $\kappa_{\mathcal{A}}$ ; [f] equals  $f_{\mathcal{A}}$ . Similarly,  $\kappa_{\mathcal{B}}$ ; [f] equals  $f_{\mathcal{B}}$ .

Lastly, suppose  $g : \mathcal{A} \& \mathcal{B} \to \mathcal{C}$  also has the property that  $\kappa_{\mathcal{A}}; g$  equals  $f_{\mathcal{A}}$  and  $\kappa_{\mathcal{B}}; g$  equals  $f_{\mathcal{B}}$ . To be a monoid homomorphism, since for any  $a : \mathcal{A}$  and  $b : \mathcal{B}$  the sum  $\langle a, e_{\mathcal{B}} \rangle +_{\mathcal{A}\&\mathcal{B}} \langle e_{\mathcal{A}}, b \rangle$  equals  $\langle a, b \rangle$ , we know that  $g(\langle a, b \rangle) = g(\langle a, e_{\mathcal{B}} \rangle +_{\mathcal{A}\&\mathcal{B}} \langle e_{\mathcal{A}}, b \rangle) = g(\kappa_{\mathcal{A}}(a) +_{\mathcal{A}\&\mathcal{B}} \kappa_{\mathcal{B}}(b)) = g(\kappa_{\mathcal{A}}(a)) +_{\mathcal{C}} g(\kappa_{\mathcal{B}}(b)) = f_{\mathcal{A}}(a) +_{\mathcal{C}} f_{\mathcal{B}}(b) = [f](\langle a, b \rangle).$  Thus, any such g must equal [f], proving uniqueness.