

CS6110/6116 Friday February 10, 2012

CS6116 Lect 3
CS6110 Lect 9

Evaluation function using closures

We use the lecture notation for closures $\{\lambda(x.b), e\}$
where e is an environment, say a list of $\langle v, d \rangle$
pairs where v is a variable, d is a value. Values will
be closures. We assume $FV(\lambda(x.b)) \subseteq \text{dom}(e)$, i.e. a v in $\langle v, d \rangle$.

We use the auxiliary functions:

$$\text{lookup}(v, \text{env}) = \begin{cases} \text{env}(v) & \text{if } v \text{ is in the domain of env} \\ \text{abort} & \text{otherwise} \end{cases}$$

$$\text{ext}(\text{env}, \langle v, d \rangle) = \begin{cases} \text{add } \langle v, d \rangle \text{ to env if } v \text{ is new} \\ \text{replace } \langle v, d \rangle \text{ by } \langle d, d \rangle \text{ otherwise} \end{cases}$$

$$\text{eval}(t, \text{env}) =$$

case $\text{isvar}(t)$ then $\text{lookup}(t, \text{env})$

$\text{isabs}(t)$ then let $\lambda(x.b) = \lambda(\text{bvar}(t), \text{body}(t))$
in $\{\lambda(x.b), \text{env}\}$

$\text{isap}(t)$ then let $c1 = \text{eval}(\text{fun}(t), \text{env})$
 $a = \text{eval}(\text{arg}(t), \text{env})$
in

$\text{eval}(c1 \text{ body}(c1), \text{ext}(e, \langle v, a \rangle))$

where $c1 \text{ body}(c1)$ is the body of the closure
 v is $\text{bvar}(c1)$
 e is $\text{env}(c1)$
(this is the environment of the $c1$)

Show how to extend this if we take closures $\{\lambda(x.b), e\}$
as inputs