

1. Give evidence semantics for each of the following propositional formulas in the Intuitionistic Propositional Calculus (IPC). For the stated examples, explain the formula as the statement of a programming problem or task.

*(a) $((A \& B) \Rightarrow C) \Rightarrow A \Rightarrow (B \Rightarrow C)$ (this has a proof)

*(b) $(A \Rightarrow (B \Rightarrow C)) \Rightarrow (A \& B) \Rightarrow C$ (so does this)

(c) $\neg(A \vee B) \Rightarrow \neg A \& \neg B$

(d) Recall that $(\text{False} \Rightarrow A)$ has evidence $\lambda(x.\text{any}(x))$,
 $(\neg A \vee B) \Rightarrow (A \Rightarrow B)$

*(e) $(A \Rightarrow C) \Rightarrow ((B \Rightarrow C) \Rightarrow (A \vee B) \Rightarrow C)$

(f) $(A \Rightarrow B) \Rightarrow ((A \Rightarrow (B \Rightarrow C)) \Rightarrow (A \Rightarrow C))$

*(g) $(A \Rightarrow B) \Rightarrow \neg B \Rightarrow \neg A$

(h) $\neg \neg(A \vee \neg A)$

2. What is an intuitive explanation for the invalidity of any of these formulas which is not valid? For valid ones (intuitionistically) provide the evidence.

(a) $\neg(A \& B) \Rightarrow \neg A \vee \neg B$ (b) $(\neg B \Rightarrow \neg A) \Rightarrow (A \Rightarrow B)$

(c) $(A \Rightarrow B) \Rightarrow \neg A \vee B$ (d) $\neg \neg A \Rightarrow A$

(e) $\neg(A \Rightarrow B) \Rightarrow \neg B$

(Handwritten notes from the back of the page, related to the last few questions)
 $(P \Rightarrow Q \wedge R) \Rightarrow L \Leftarrow (P \Rightarrow Q) \wedge (P \Rightarrow R) \Rightarrow L$
 $(P \Rightarrow E \wedge Q \Rightarrow A \wedge B) \Leftarrow (P \Rightarrow E \wedge (Q \Rightarrow A \wedge B)) \Rightarrow L$

(2)

CS6110/6116 Problem Set 5 cont.

3. Prove $(A \Rightarrow A)$ using a Hilbert style proof from these two axiom schemas and the rule modus ponens

1. $A \Rightarrow (B \Rightarrow A)$ as an axiom schema.

2. formula (i) from Problem 1 treated as an axiom schema.

$$\text{modus Ponens} \quad \frac{A \quad A \Rightarrow B}{B}$$

Instances of axiom schemas arise by substituting any propositional formula for the propositional variables A, B, C . For example, $A \Rightarrow (A \Rightarrow A)$ is an instance of 1.

4. Prove the following formulas using Refinement Logic rules from Lecture 30, Wed April 11.

$$(a) P \vee (Q \& R) \Rightarrow (P \vee Q) \& (P \vee R) \quad \text{"distributive"}$$

$$(b) (P \Rightarrow \neg P) \Rightarrow \neg P$$

$$(c) P \Rightarrow \neg \neg P$$

$$(d) \neg P \Rightarrow (P \Rightarrow Q)$$

$$(e) \neg \neg (P \vee \neg P)$$

5. There will be one more problem with four or five parts after we examine First-Order Logic on Monday April 16. One part will be to give the evidence for

$$(a) \forall x. (A(x) \Rightarrow B(x)) \Rightarrow (\forall x. A(x) \Rightarrow \forall x. B(x))$$

$$(b) \forall x. (A(x) \Rightarrow B(x)) \Rightarrow (\exists x. A(x) \Rightarrow \exists x. B(x)).$$

Also, for (c) and (d) below, give evidence if true, if not true explain why.

$$(c) (\exists x. A(x) \Rightarrow \exists x. B(x)) \Rightarrow \exists x. (A(x) \Rightarrow B(x))$$

$$(d) \exists y. (A(y) \Rightarrow B(y)) \Rightarrow \forall x. A(x) \Rightarrow \exists y. B(y)$$