# Problem Set 3 Due Friday Mar 8 CS6116 students should solve Problems 2 and 5

#### Problem 1

Revisiting Euclid's Theorem in Prim Rec Arith.

Euclid's Theorem symbolically:

 $\forall a, b: \mathbb{N}. \exists d: \mathbb{N}. div(d, a) \& div(d, b) \& \forall z: \mathbb{N}. (div(z, a) \& div(z, b) \Rightarrow z \leq d)$ 

If we let  $div(x,y) = E_k^y(k * x) = y$ , then using Goodstein's notation, Euclid's Theorem is

 $A^{n}_{a}(A^{a}_{b}(E^{a}_{d}(div(d,a) \& div(d,b)))) \& A^{n}_{z}(A^{n}_{a}(A^{a}_{b}(div(z,a) \& div(z,b)) \Rightarrow (z \doteq d = 0)))$ 

- (a) The key step in proving this theorem is witnessing  $E_d^a$ . A function gcd(a, b) computing the greatest common divisor is a witness for the existential quantifier. Write a primitive recursive function to compute gcd.
- (b) Write a while program to compute *gcd* and add assertions that document the code and provide a basis for proving that the program is correct.
- (c) Prove that the while program is correct either using the structural operational semantics or the Binary Relation Semantics both in CS6110 2010 Lecture 17 posted with Monday's lecture.
- (d) Sketch a Hoare Logic asserted program *total correctness* proof.

#### Problem 2

Define the integer square root of a non-negative number n to be the least non-negative number r such that  $r^2 \leq n < (r+1)^2$ . Call this relation Root(r, n).

- (a) Write an IMP program to compute r.
- (b) Prove that your program is *totally correct*, either using asserted programs with a termination predicate (see notes on PL/CV) or using structured operational semantics or Relational Semantics (on-line notes for Lect 16) or positional semantics.

### Problem 3

Write a careful proof using Relational Semantics, that the Hoare while rule is valid for partial correctness.

### Problem 4

Pick an imperative programming language that you know well (e.g. Java, Python, C,  $C^{++}$ , etc.) and explain at least four caveats that need to be stated in order to make the Hoare assignment rule valid.

## Problem 5

- (a) Define a function  $\lambda(x.big(x))$ , related to the functions  $f_n$  in Def 3.2 page 3 of Meyer and Ritchie's paper, which is not Loop computable, thus not primitive recursive.
- (b) Sketch a proof that  $\lambda(x.big(x))$  is not Loop computable.
- (c) For extra credit, write the continuation passing form of this function.

#### Problem 6 - More Extra Credit

Write your intuitive explanation of why continuation passing computations are tail recursive.