

Lecture 6

PLAN

1. Discuss Mini-projects
2. Note readings
3. Notation
4. Tableau proofs
5. Historical perspective on Computational Logic
6. Examples of computational content of propositions

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4. Tableau Proofs - examples using Smullyan rules p.17 (also see p.22)

The idea is to search for a counter example

1. $F(X \supset Y)$
 2. TX
 3. FY
 counter example
 open tableau

vs

1. $F(X \supset X)$ α rule
 2. TX x_1
 3. FX x_2
 * closed tableau

1. $F(X \supset (Y \supset X))$
 2. TX
 3. $F(Y \supset X)$
 4. TY
 5. FX
 * 2,5

1. $F((X \wedge (X \supset Y)) \supset Y)$
 2. $T(X \wedge (X \supset Y))$
 3. FY
 4. TX from 2.
 5. $T(X \supset Y)$ from 2 β -rule

| | |
|---------|---------|
| 6. FX | 7. TY |
| * 4,6 | * 3,7 |

Lecture 6 continued

5. Historical perspective

Euclid's Elements of Geometry is mostly constructive.

for an object to exist, it must be constructed
by ruler and compass, e.g. regular triangle, regular pentagon.

Note, on March 30, 1796 Gauss (Carl ~~Fried~~ Friedrich)
reported constructing a regular 17gon, solving a
2,000 year old geometry problem and starting a diary
ended only on July 9, 1814, after 146 mathematical
results.

When we say $\exists x. P(x)$ the question is always, can we
construct it. Consider this example.

$\exists x, y: \text{Irrational}. (x^y \text{ is rational})$

Here is a "non-constructive" proof.

Consider $\sqrt{2}^{\sqrt{2}}$, either this is rational or not, that is

$\text{Rational}(\sqrt{2}^{\sqrt{2}}) \vee \sim \text{Rational}(\sqrt{2}^{\sqrt{2}})$

If it is rational, then take $x = \sqrt{2}$ and $y = \sqrt{2}$.

If it is not rational, then it is irrational, so

take $x = \sqrt{2}^{\sqrt{2}}$ and $y = \sqrt{2}$.

Then $(\sqrt{2}^{\sqrt{2}})^{\sqrt{2}} = \sqrt{2}^2 = 2$ as required. QED.

But what are the first seven digits of x ?

are they 1.414213 or 1.189207?

We don't know because we haven't constructed a
~~any~~ unique x .

We see here that it is our understanding of or
which is the problem.

6. Examples of computational content of propositions.

1. The meaning of (evidence for) $X \Rightarrow X$ is a computable function taking evidence for X , say x , into evidence for X . Clearly the identity function works.

ML type: $'X \rightarrow 'X$

ML notation: $\backslash x. x$

Nuprl notation: $\lambda(x.x)$

2. Here is evidence for $X \Rightarrow (Y \Rightarrow X)$

ML type: $'X \rightarrow ('Y \rightarrow 'X)$

ML notation: $\backslash x. (\backslash y. x)$ for evidence, a program

Nuprl notation: $\lambda(x. \lambda(y.x))$

3. Here is evidence for $(A \wedge (B \vee C)) \Rightarrow ((A \wedge B) \vee (A \wedge C))$

ML notation for evidence:

$\backslash h$ let $h = \langle a, bc \rangle$ in

if $isl(bc)$ then $inl(a, outl(bc))$

else $inr(a, outr(bc))$

Nuprl notation:

$\lambda(h. \text{spread}(h; a, bc.$

$\text{decide}(bc; b. \text{inl}\langle a, b \rangle; c. \text{inr}\langle a, c \rangle))$)

Exercise. Let \perp be a constant (any in ML), find evidence

for $((A \vee (A \Rightarrow \perp)) \Rightarrow \perp) \Rightarrow \perp$. The ML type is

$((A + 'A \rightarrow \text{any}) \rightarrow \text{any}) \rightarrow \text{any}$.