

- Announcements

- Send references to bhoward@cs.cornell.edu

- Random walks on expanders

- Real world examples: Finding random websites

Comparing search engines

- Regular degree d graphs

$P = (P_1, P_2, \dots, P_n)^T$ where P_i is the probability of being at vertex i

Initially, $P = (1, 0, \dots, 0)^T \leftarrow$ start at vertex 1

$A \cdot \frac{P}{d} \rightarrow P'$ (P' is the new prob)

Suggest normalizing A by dividing by d : $\frac{A}{d}$

Stationary probability: $\left(\frac{1}{n}, \dots, \frac{1}{n}\right) = \frac{u}{n}$ where $u = (1, \dots, 1)$

Theorem: Let P be an arbitrary probability distribution.

$$\text{Then } \left| \left(\frac{A}{d} \right)^s P - \frac{u}{n} \right|_1 \leq \sqrt{n} \left(\frac{\lambda}{d} \right)^s \quad \text{where } \lambda = \max\{|\lambda_2|, |\lambda_n|\}$$

and s is the number of steps.

Proof: Since $\frac{u}{n}$ is an eigenvector of $\frac{A}{d}$ with eigenvalue 1,

$$\frac{A}{d} \cdot \frac{u}{n} = \frac{u}{n} \quad \text{and} \quad \left(\frac{A}{d} \right)^s \cdot \frac{u}{n} = \frac{u}{n}$$

$$\text{So, } \left(\frac{A}{d} \right)^s P - \frac{u}{n} = \left(\frac{A}{d} \right)^s P - \left(\frac{A}{d} \right)^s \frac{u}{n} = \left(\frac{A}{d} \right)^s \left(P - \frac{u}{n} \right)$$

$$\left| \left(\frac{A}{d} \right)^s P - \frac{u}{n} \right|_2 = \left| \left(\frac{A}{d} \right)^s \left(P - \frac{u}{n} \right) \right|_2$$

The coordinates of P sum to 1 as do the coordinates of $\frac{u}{n}$. Thus the coordinates

of $\left(P - \frac{u}{n} \right)$ sum to zero. Therefore, $\left(P - \frac{u}{n} \right)$ is orthogonal to u and hence is in

the subspace spanned by the remaining eigenvectors v_2, v_3, \dots, v_n .

$$\left| \left(\frac{A}{d} \right)^s \left(P - \frac{u}{n} \right) \right|_2 \leq \left(\frac{\lambda_2(A)}{d} \right)^s \left| P - \frac{u}{n} \right|_2$$

$$\begin{aligned} p - \frac{u}{n} &= \left| (p_1, \dots, p_n) - \left(\frac{1}{n}, \dots, \frac{1}{n} \right) \right|_2 \leq \left| (1, 0, \dots, 0) - \left(\frac{1}{n}, \dots, \frac{1}{n} \right) \right|_2 \\ &= \left(1 - \frac{1}{n} \right)^2 + \frac{1}{n^2} + \dots + \frac{1}{n^2} \\ &= 1 - \frac{2}{n} + \frac{1}{n} + \frac{n-1}{n^2} = 1 - \frac{2}{n} + \frac{1}{n} \end{aligned}$$

Since $\left| p - \frac{u}{n} \right|_2 \leq 1$

$$\left| \left(\frac{A}{d} \right)^s \left(P - \frac{\mu}{n} \right) \right|_2 \leq \left(\lambda \left(\frac{A}{d} \right) \right)^s$$

To convert this to a one-norm, use $|v|_1 \leq \sqrt{n}|v|_2$

Thus $\left| \left(\frac{A}{d} \right)^s \left(P - \frac{\mu}{n} \right) \right|_1 \leq \sqrt{n} \left(\lambda \left(\frac{A}{d} \right) \right)^s = \sqrt{n} \left(\frac{\lambda(A)}{d} \right)^s$

Corollary: $\left| \left(\frac{A}{d} \right)^s \left(P - \frac{\mu}{n} \right) \right|_1 \leq \sqrt{n} \left(1 - \frac{h^2(G)}{2d^2} \right)^s$

Where $h(G) = \min_{|S| \leq \frac{n}{2}} \frac{|\partial S|}{|S|}$ * $|\partial S|$ is the number of edges out of S

$\Delta(G) = \lambda_1 - \lambda_2$ * Recall $h^2(G) \leq 2d\Delta(G)$

Proof: $\lambda \left(\frac{A}{d} \right) \geq \lambda_2 \left(\frac{A}{d} \right) = \frac{\lambda_2(A)}{d} = \frac{\lambda_1 - \Delta}{d} = 1 - \frac{\Delta}{d}$

$$\left| \left(\frac{A}{d} \right)^s \left(P - \frac{\mu}{n} \right) \right|_1 \leq \sqrt{n} \left(1 - \frac{\Delta}{d} \right)^s \quad \boxed{h^2(G) \leq 2d\Delta(G) \Rightarrow \frac{h^2(G)}{2d} \leq \Delta(G)}$$

$$\therefore \left| \left(\frac{A}{d} \right)^s \left(P - \frac{\mu}{n} \right) \right|_1 \leq \sqrt{n} \left(1 - \frac{h^2(G)}{2d^2} \right)^s$$