

## Game Playing

### An AI Favorite

- structured task
- not initially thought to require large amounts of knowledge
- focus on games of perfect information

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## Game Playing

**Initial State** is the initial board/position

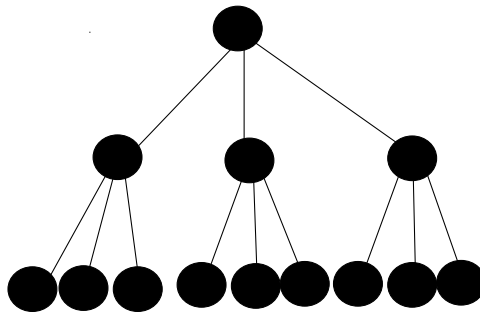
**Operators** define the set of legal moves from any position

**Terminal Test** determines when the game is over

**Utility Function** gives a numeric outcome for the game

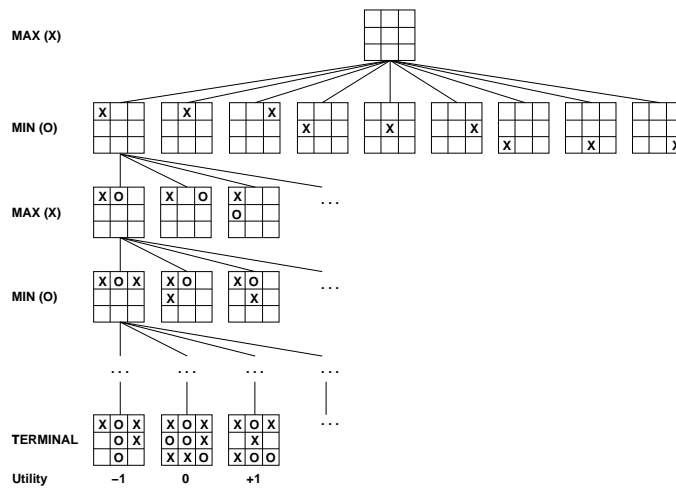
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### Game Playing as Search

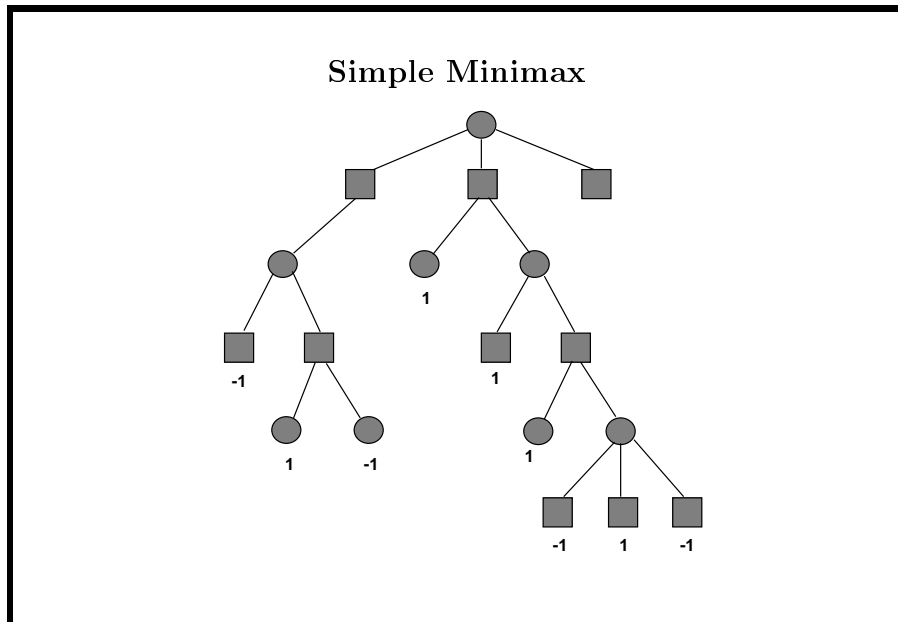


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### Partial Search Tree for Tic-Tac-Toe



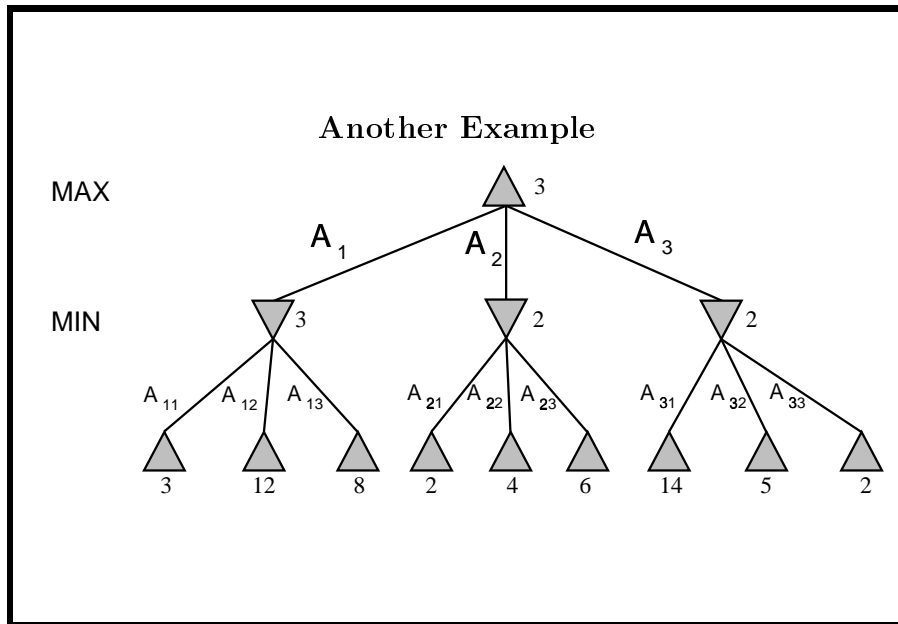
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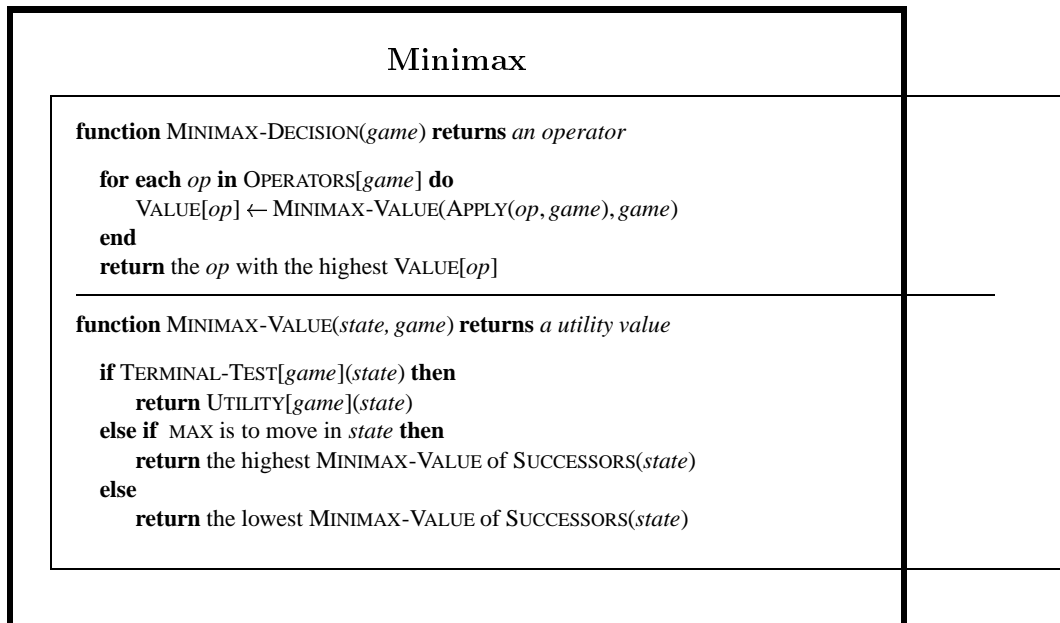
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- Simplified Minimax Algorithm**
1. Expand the entire tree below the root.
  2. Evaluate the terminal nodes as wins for the minimizer or maximizer.
  3. Select an unlabeled node,  $n$ , all of whose children have been assigned values. If there is no such node, we're done — return the value assigned to the root.
  4. If  $n$  is a minimizer move, assign it a value that is the minimum of the values of its children. If  $n$  is a maximizer move, assign it a value that is the maximum of the values of its children. Return to Step 3.

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### **The Need for Imperfect Decisions**

**Problem:** Minimax assumes the program has time to search to the terminal nodes.

**Solution:** Cut off search earlier and apply a heuristic evaluation function to the leaves.

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### **Static Evaluation Functions**

Minimax depends on the translation of board quality into a single, summarizing number. Difficult. Expensive.

- Add up values of pieces each player has (weighted by importance of piece).
- Isolated pawns are bad.
- How well protected is your king?
- How much maneuverability to you have?
- Do you control the center of the board?

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- Strategies change as the game proceeds.

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### Design Issues of Heuristic Minimax

**Evaluation Function:** What features should we evaluate and how should we use them? An evaluation function should:

- 1.
- 2.
- 3.

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### Linear Evaluation Functions

- $w_1f_1 + w_2f_2 + \dots + w_nf_n$
- This is what most game playing programs use
- Steps in designing an evaluation function:
  1. Pick informative features
  2. Find the weights that make the program play well

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### Design Issues of Heuristic Minimax

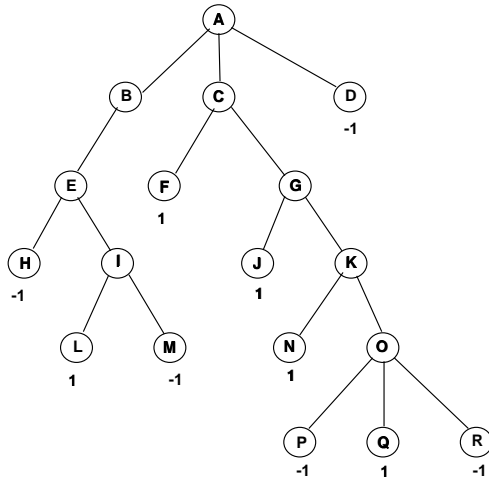
**Search:** search to a constant depth

Problems:

- Some portions of the game tree may be “hotter” than others. Should search to *quiescence*. Continue along a path as long as one move’s static value stands out (indicating a likely capture).
- *Horizon effect*
- Secondary search. (*singular extension heuristic*)

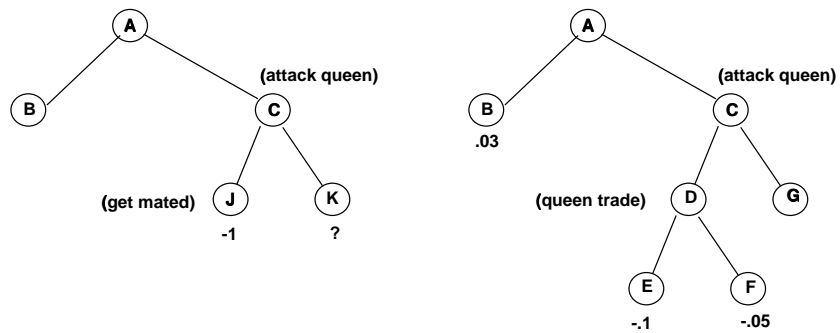
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### Improving Minimax — $\alpha - \beta$ pruning



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### Two More Examples



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### Algebraic Solution

Let  $g' = e(g)$ . Then  $c' = \min(-.05, g')$ .

The value assigned to the root node  $a$  is

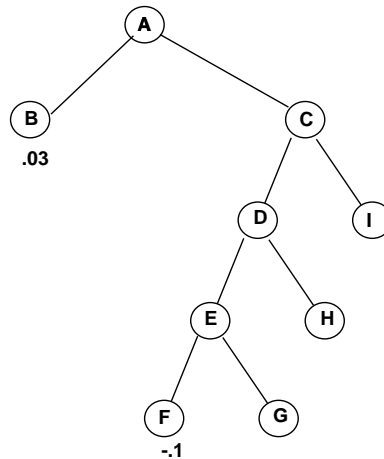
$$a' = \max(.03, \min(-.05, g')) = .03$$

because  $\min(-.05, g') \leq -.05 < .03$ .

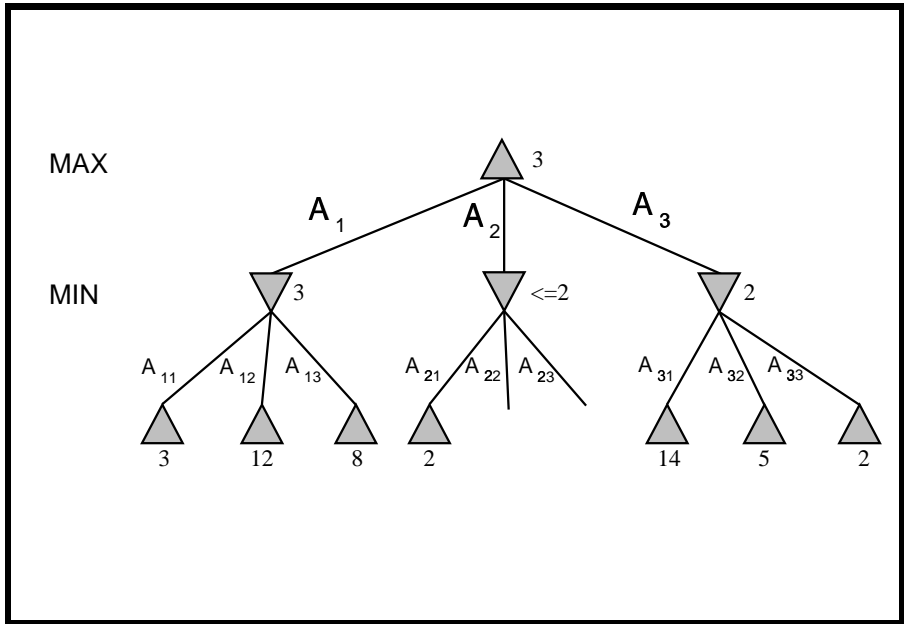
The value assigned to  $a$  is independent of the value assigned to  $g$ .

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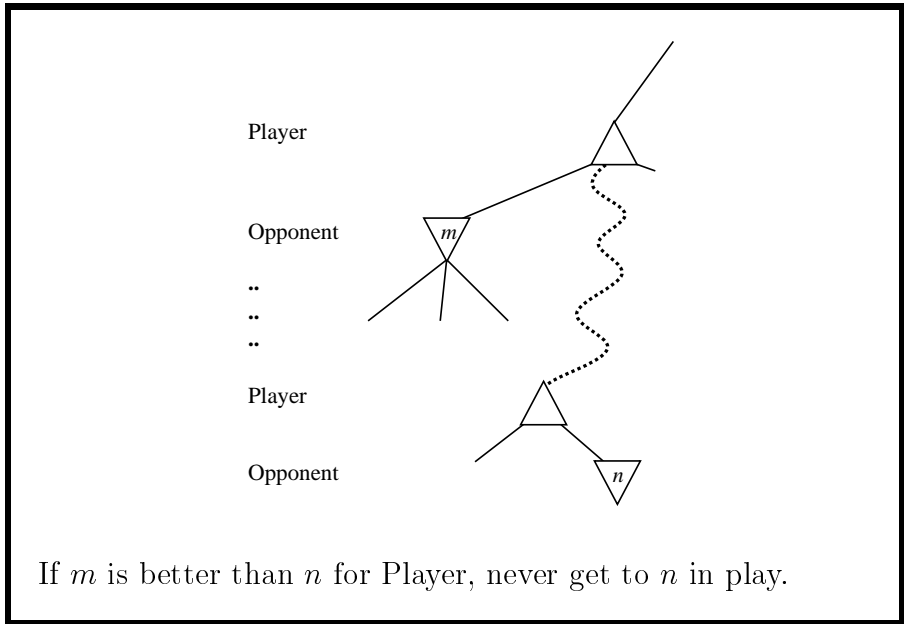
### A deep $\alpha - \beta$ cutoff



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If  $m$  is better than  $n$  for Player, never get to  $n$  in play.

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### $\alpha - \beta$ Search

$c$  = search cutoff

$\alpha$  = lower bound on Max's outcome; initially set to  $-\infty$

$\beta$  = upper bound on Min's outcome ; initially set to  $+\infty$

We'll call  $\alpha - \beta$  procedure recursively with a narrowing range between  $\alpha$  and  $\beta$ .

Maximizing levels may reset  $\alpha$  to a higher value; Minimizing levels may reset  $\beta$  to a lower value.

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### $\alpha - \beta$ Search Algorithm

1. If the limit of search has been reached, compute  $e(n)$  and report the result.
2. Otherwise, if the level is a **minimizing** level,
  - Until no more children or  $\alpha \geq \beta$ ,
    - Use  $\alpha - \beta$  search on child with current values of  $\alpha$  and  $\beta$ ; note the value,  $v$ , returned.
    - If  $v < \beta$ , reset  $\beta$  to  $v$ .
  - Report  $\beta$ .

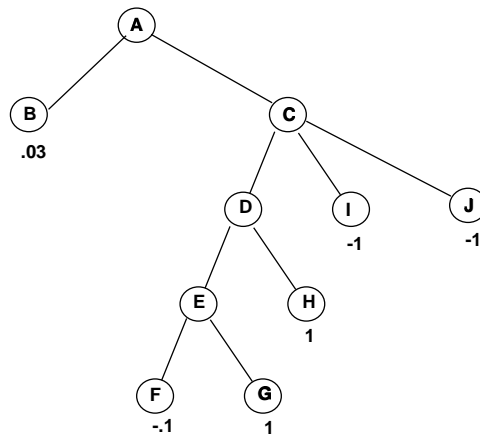
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3. Otherwise, the level is a **maximizing** level:

- Until no more children or  $\alpha \geq \beta$ ,
  - Use  $\alpha - \beta$  search on child with current values of  $\alpha$  and  $\beta$ ; note the value,  $v$ , returned.
  - If  $v > \alpha$ , reset  $\alpha$  to  $v$ .
- Report  $\alpha$ .

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Example

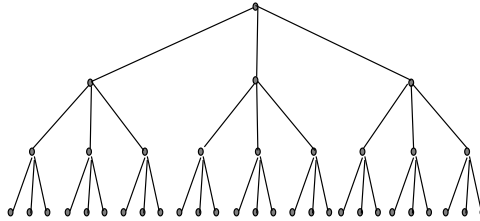


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### Search Space Size Reductions

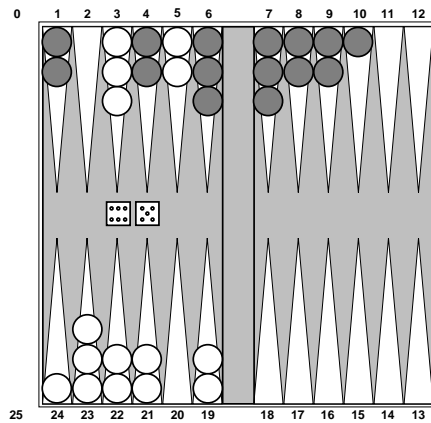
**Worst Case:** In an ordering where worst options evaluated first, all nodes must be examined.

**Best Case:** If nodes ordered so that the best options are evaluated first, then what?



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### Backgammon – Board



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### Backgammon – Rules

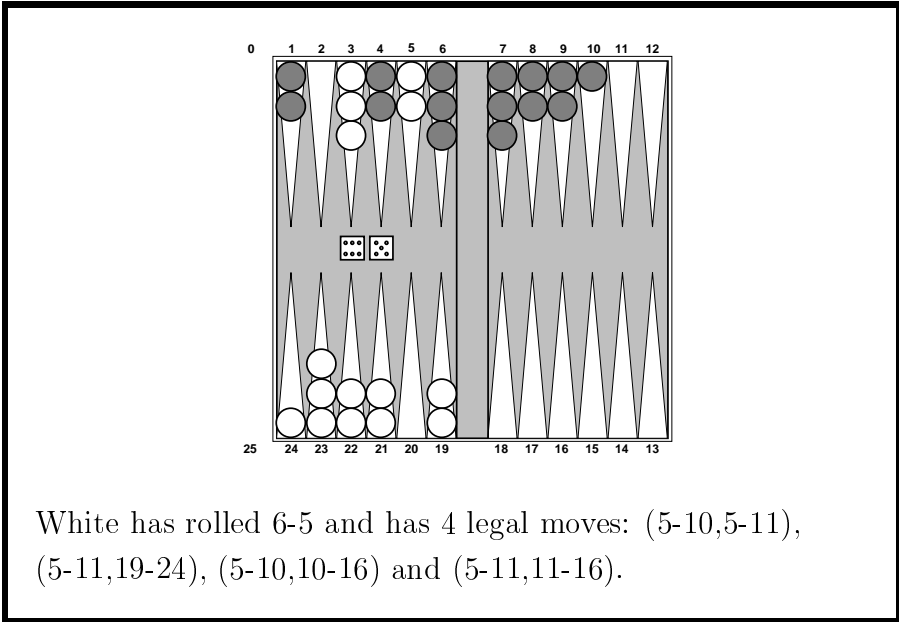
- Goal: move all of your pieces off the board before your opponent does.
- White moves counterclockwise toward 0.
- Black moves clockwise toward 25.
- A piece can move to any position except one where there are two or more of the opponent's pieces.
- If it moves to a position with one opponent piece, that piece is captured and has to start its journey from the beginning.

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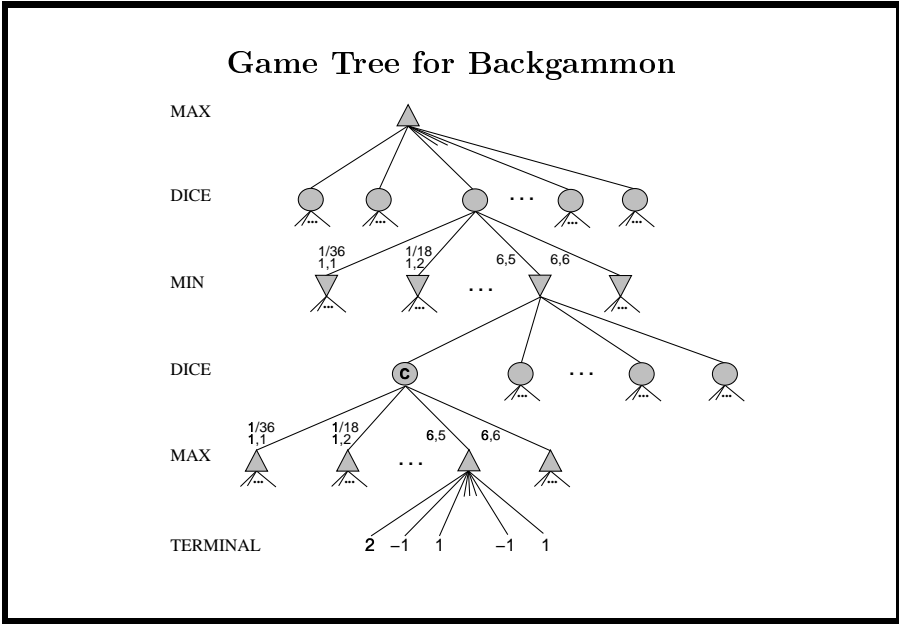
### Backgammon – Rules

- If you roll doubles you take 4 moves (example: roll 5,5, make moves 5,5,5,5).
- Moves can be made by one or two pieces (in the case of doubles by 1, 2, 3 or 4 pieces)
- And a few other rules that concern *bearing off* and *forced moves*.

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### State of the Art in Checkers

- 1952: Samuel developed a checkers program that learned its own evaluation function through self play.
- 1992: *Chinook* (J. Schaeffer) wins the U.S. Open. At the world championship, Marion Tinsley beat *Chinook*.

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### State of the Art in Backgammon

- 1980: *BKG* using one-ply search and lots of luck defeated the human world champion.
- 1992: Tesauro combines Samuel's learning method with neural networks to develop a new evaluation function, resulting in a program ranked among the top 3 players in the world.

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### State of the Art in Go

\$2,000,000 prize available for first computer program to defeat a top level player.

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### History of Chess in AI

500	legal chess
1200	occasional player
2000	world-ranked
2900	Gary Kasparov

**Early 1950's** Shannon and Turing both had programs that (barely) played legal chess (500 rank).

**1950's** Alex Bernstein's system, (500+ $\epsilon$ ).

**1957** Herb Simon claims that a computer chess program would be world chess champion in 10 years...yeah, right.

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- 1966** McCarthy arranges chess match, Stanford vs. Russia. Long, drawn-out match. Russia wins.
- 1967** Richard Greenblatt, MIT. First of the modern chess programs, *MacHack* (1100 ranking).
- 1968** McCarthy, Michie, Papert bet Levy (rated 2325) that a computer program would beat him within 10 years.
- 1970** ACM started running chess tournaments. Chess 3.0-6 (rated 1400).
- 1973** By 1973...Slate: "It had become too painful even to look at Chess 3.6 any more, let alone work on it."

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- 1973** Chess 4.0: smart plausible-move generator rather than speeding up the search. Improved rapidly when put on faster machines.
- 1976** Chess 4.5: ranking of 2070.
- 1977** Chess 4.5 vs. Levy. Levy wins.
- 1980's** Programs depend on search speed rather than knowledge (2300 range).
- 1993** DEEP THOUGHT: Sophisticated special-purpose computer;  $\alpha - \beta$  search; searches 10 ply; singular extensions; rated about 2600.

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**1995** DEEP BLUE: searches 14-ply; considers 100–200 billion positions per move.

**1997** DEEP BLUE: first match won against world-champion (Kasparov).

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### Concludes “Search”

- **Problem Solving as Search**
- **Uninformed search:** DFS / BFS / Uniform cost search  
time / space complexity  
size search space: up to approx.  $10^{11}$  nodes  
special case: **Constraint Satisfaction / CSPs**  
generic framework: variables & constraints  
backtrack search (DFS); propagation (forward-checking / arc-consistency, variable / value ordering  
(but **incomplete**)

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- **Informed Search:** use heuristic function guide to goal

*Greedy search*

*A\* search / provably optimal*

Search space up to approximately  $10^{25}$

**Local search** (incomplete)

*Greedy / Hillclimbing*

*Simulated annealing*

*Tabu search*

*Genetic Algorithms / Genetic Programming*

search space  $10^{100}$  to  $10^{1000}$

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- **Adversary Search / Game Playing**

**minimax**

Up to around  $10^{10}$  nodes, 6 — 7 ply in chess.

**alpha-beta pruning**

Up to around  $10^{20}$  nodes, 14 ply in chess.

*provably optimal*

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## Search and AI

Why such a central role?

A. Basically, because lots of tasks in AI are **intractable**.  
Search is “only” way to handle them.

Many applications of search, in e.g.,  
Learning / Reasoning / Planning / NLU / Vision

Good thing: much recent progress ( $10^{30}$  quite feasible;  
sometimes up to  $10^{1000}$ ). **Qualitative difference**  
**from only a few years ago!**