

CS 4700: Foundations of Artificial Intelligence

Spring 2020
Prof. Haym Hirsh

Lecture 9
February 12, 2020

Reminder: Technology Policy

No technology except for first four rows of left and right sides

Jupyter Notebooks **Postponed**

~~“Primer”: Friday 5pm in Gates G01~~

Friday Feb 14 5pm Gates G01

Covers:

Jupyter Notebooks and .ipynb files

Google Colab (online Jupyter Notebook environment)

Useful Python features

Bring your laptop

This Friday's Lecture **Preponed**

To be given

Today 3:30-4:20

Gates G01

Will be recorded

and available on course website

(must watch between Wed and following Mon)

(Please come to the live version!)

POSITION-PLAY VALUE

Each white piece has a certain position-play contribution and so has the black king. These must all be added up to give the position-play value.

For a Q, R, B, or Kt, count—

(a) The square root of the number of moves the piece can make from the position, counting a capture as two moves, and not getting that the king must not be left in check.

(b) (If not a Q) 1.0 if it is defended, and an additional 0.5 if twice defended.

For a K, count—

(c) For moves other than castling as (a) above.

(d) It is then necessary to make some allowance for the vulnerability of the K. This can be done by assuming it to be replaced by a friendly Q on the same square, estimating as in (a), but subtracting instead of adding.

(e) Count 1.0 for the possibility of castling later not being lost by moves of K or rooks, a further 1.0 if castling could take place on the next move, and yet another 1.0 for the actual performance of castling.

For a P, count—

(f) 0.2 for each rank advanced.

(g) 0.3 for being defended by at least one piece (not P).

For the black K, count—

(h) 1.0 for the threat of checkmate.

(i) 0.5 for check.

We can now state the rule for play as follows. The move chosen must have the greatest possible value, and, consistent with this, the greatest possible position-play value. If this condition admits of

1948/1951/1953

Alan Turing
Turochamp

Two-move lookahead
Evaluation function
One level “minimax”

Philosophical Magazine, Ser.7, Vol. 41, No. 314 - March 1950.

XXII. Programming a Computer for Playing Chess¹

By CLAUDE E. SHANNON

Bell Telephone Laboratories, Inc., Murray Hill, N.J.²

[Received November 8, 1949]

1. INTRODUCTION

This paper is concerned with the problem of constructing a computing routine or "program" for a modern general purpose computer which will enable it to play chess. Although perhaps of no practical importance, the question is of theoretical interest, and it is hoped that a satisfactory solution of this problem will act as a wedge in attacking other problems of a similar nature and of greater significance. Some possibilities in this direction are: -

- (1) Machines for designing filters, equalizers, etc.
- (2) Machines for designing relay and switching circuits.

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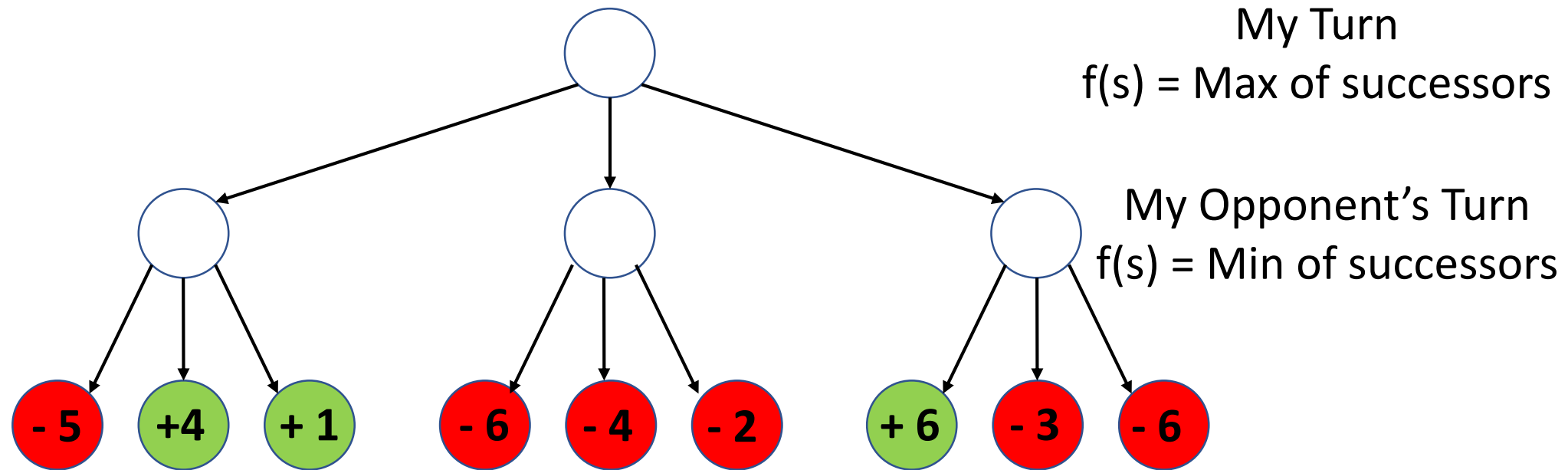
By CLAUDE E. SHANNON

$$f(P) = 200(K-K') + 9(Q-Q') + 5(R-R') + 3(B-B'+N-N') + (P-P') - \\ 0.5(D-D'+S-S'+I-I') + \\ 0.1(M-M') + \dots$$

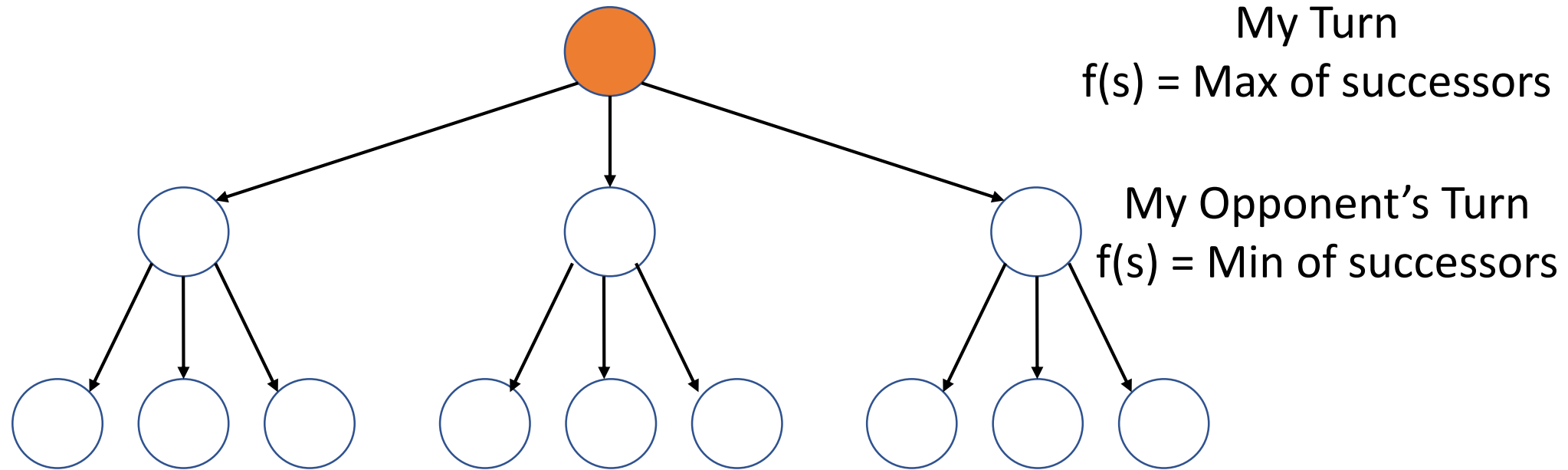
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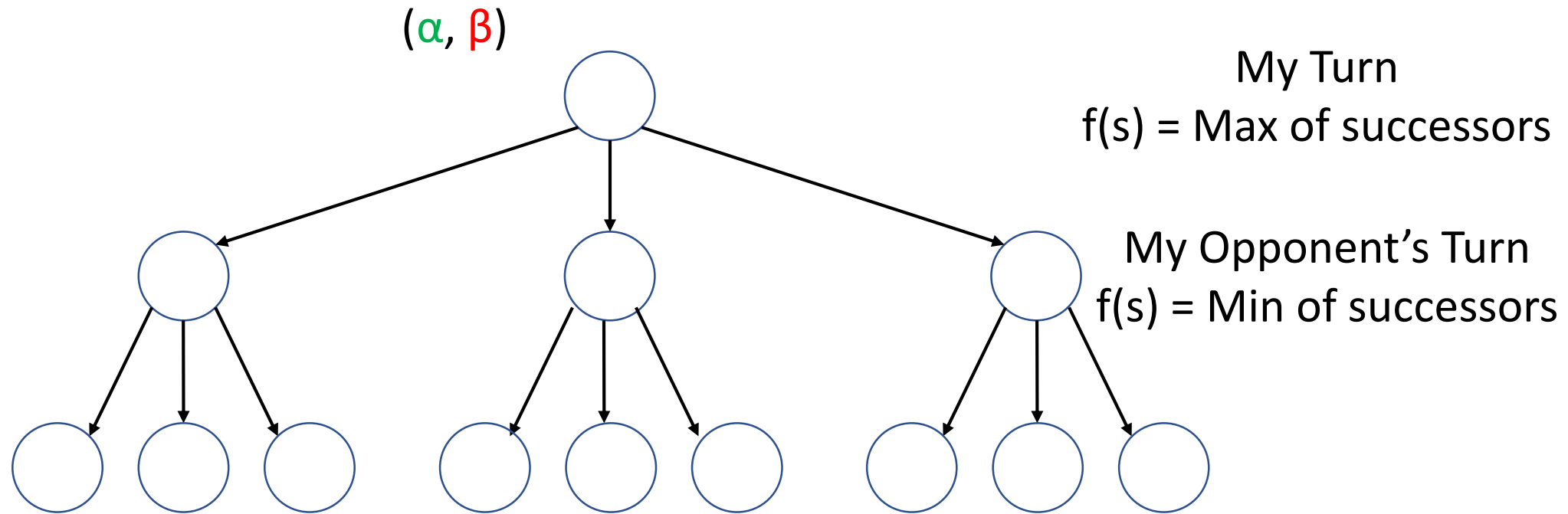
Alpha-Beta Pruning: Implementation



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Alpha-Beta Pruning: Implementation



Minimax Algorithm with Alpha-Beta Pruning

maxturn(s,ops,a,b): {my turn}

if cutoff(s,depth) then return f(s)

else

val $\leftarrow -\infty$;

foreach o $\in$ ops

val' $\leftarrow$ minturn(apply(s,o),ops,a,b);

if val' > val then

val $\leftarrow$ val';

bestop $\leftarrow$ o;

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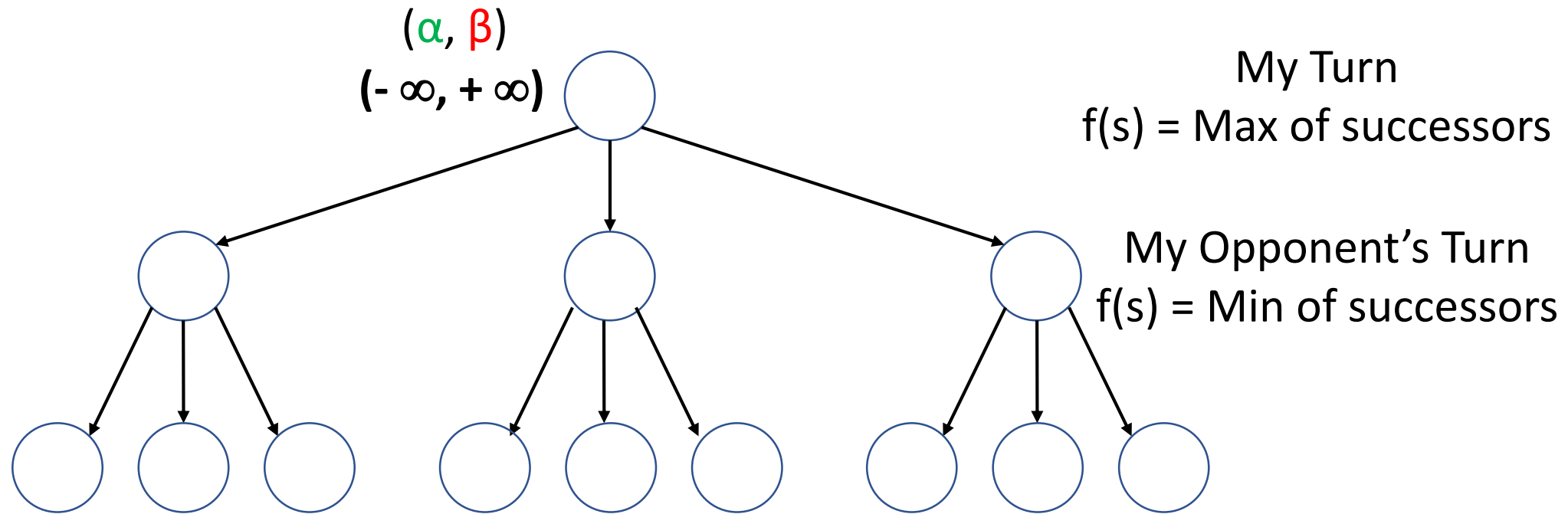
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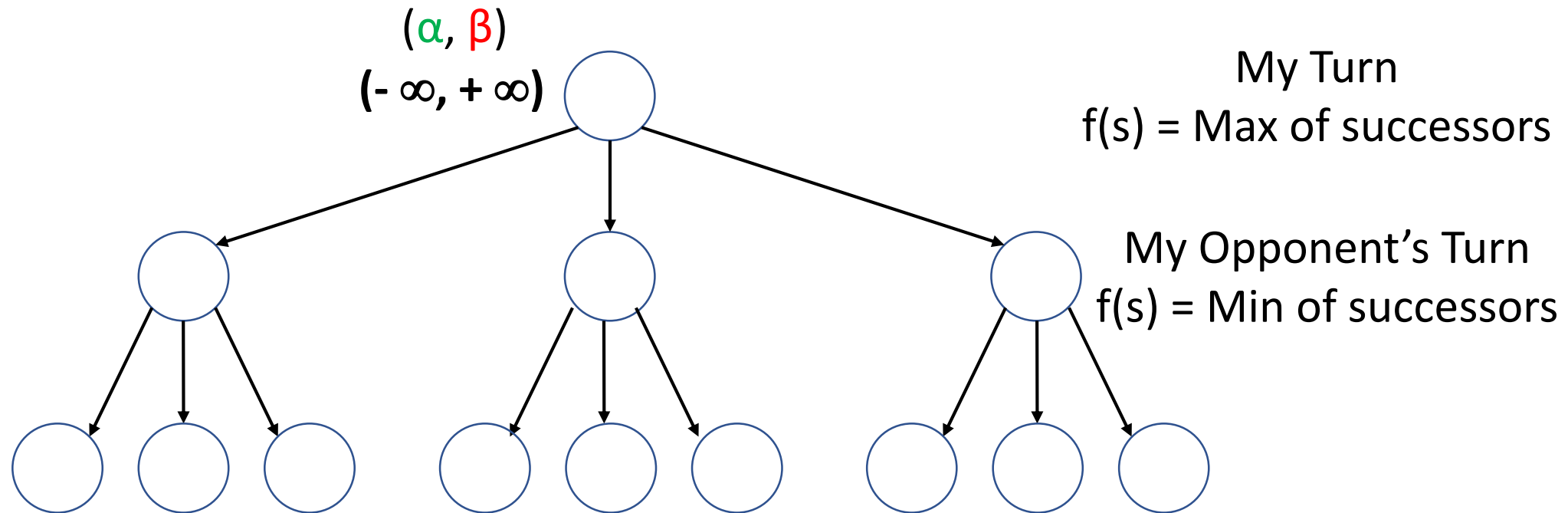
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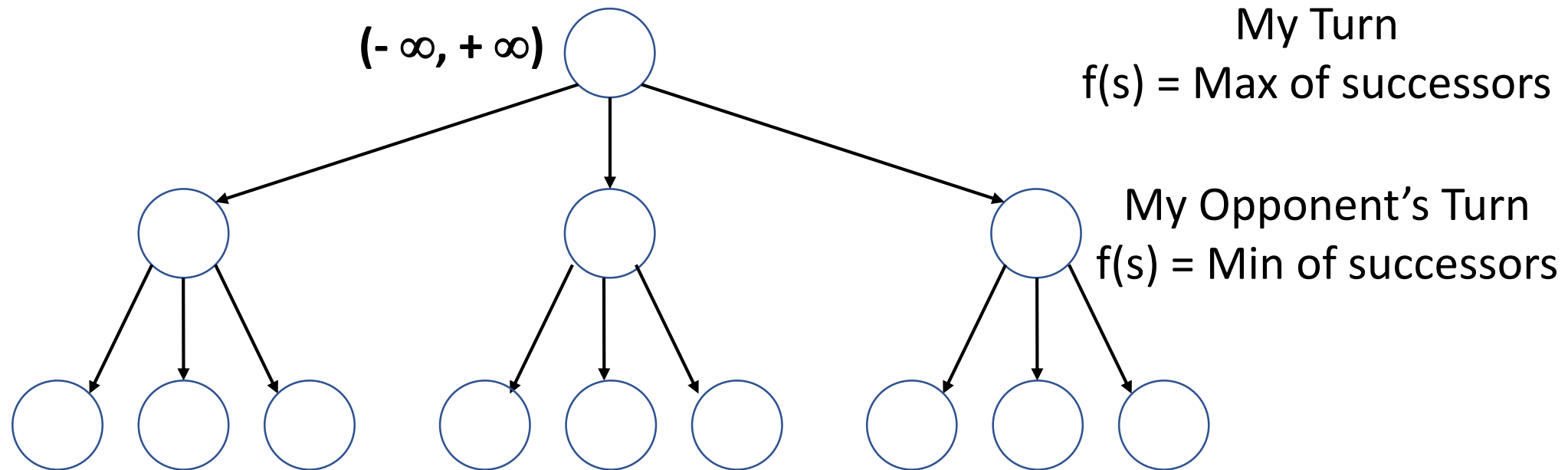


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(best I can force thus far, worst my opponent can force thus far)



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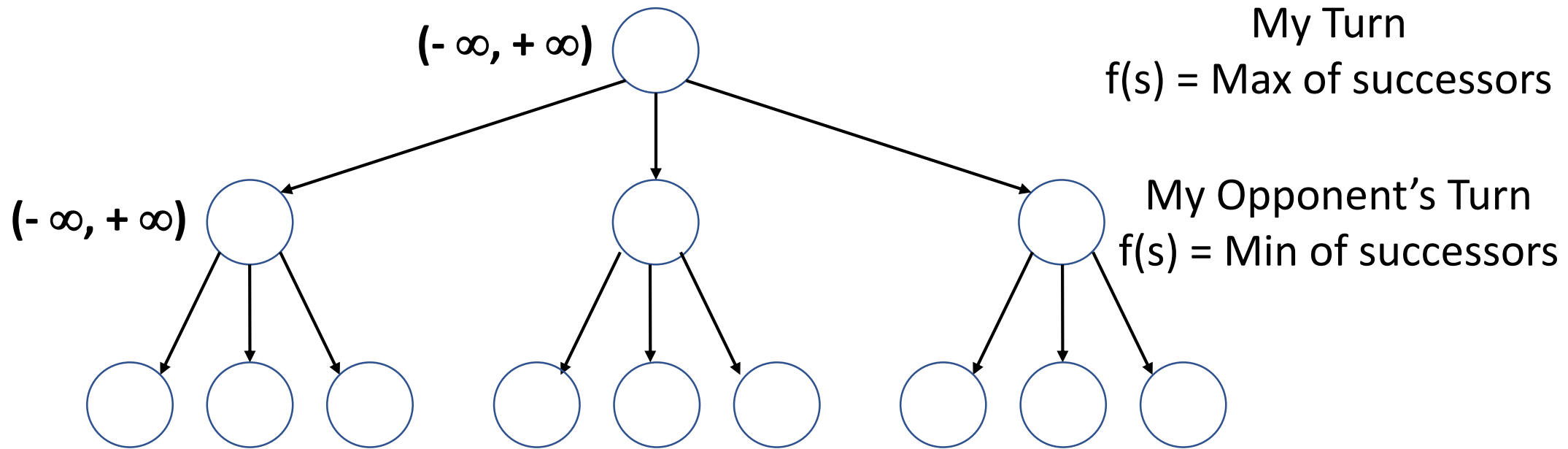
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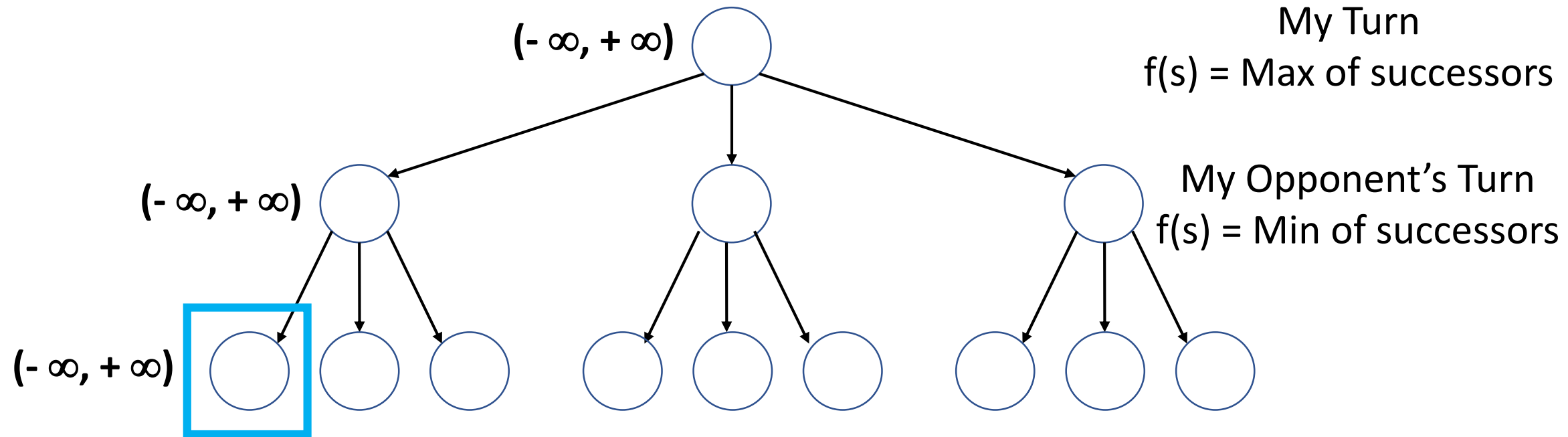
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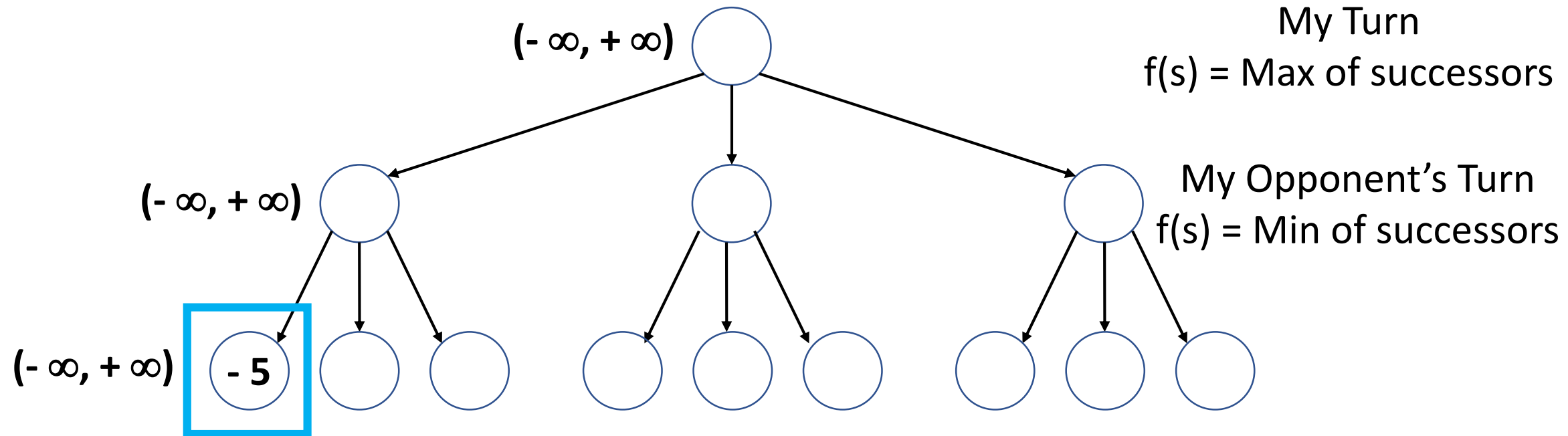
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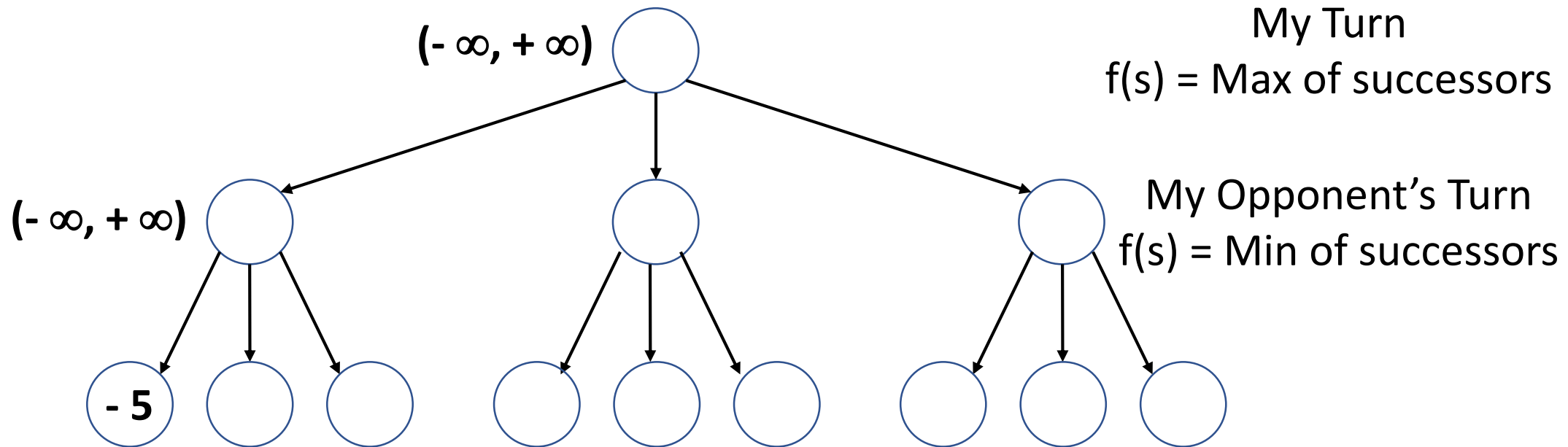
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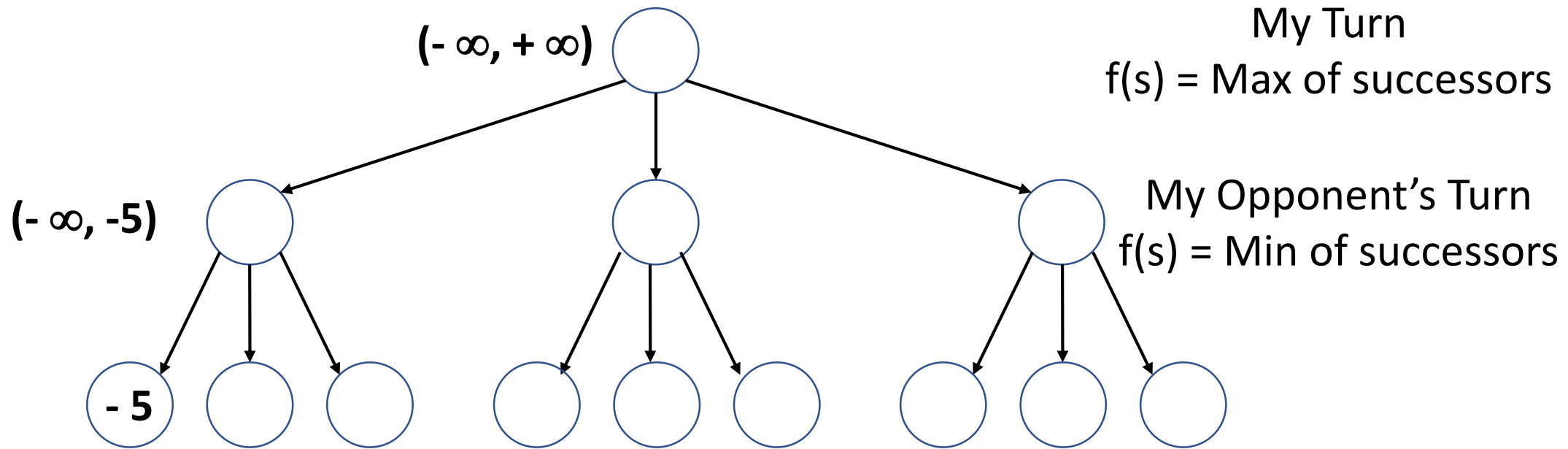
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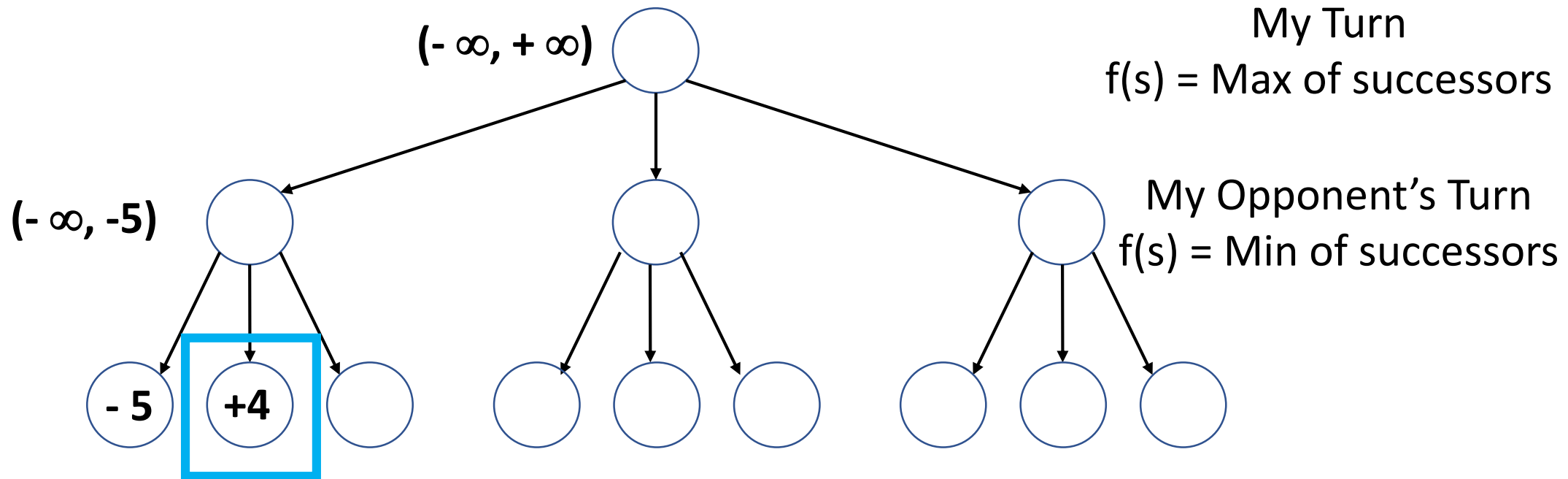
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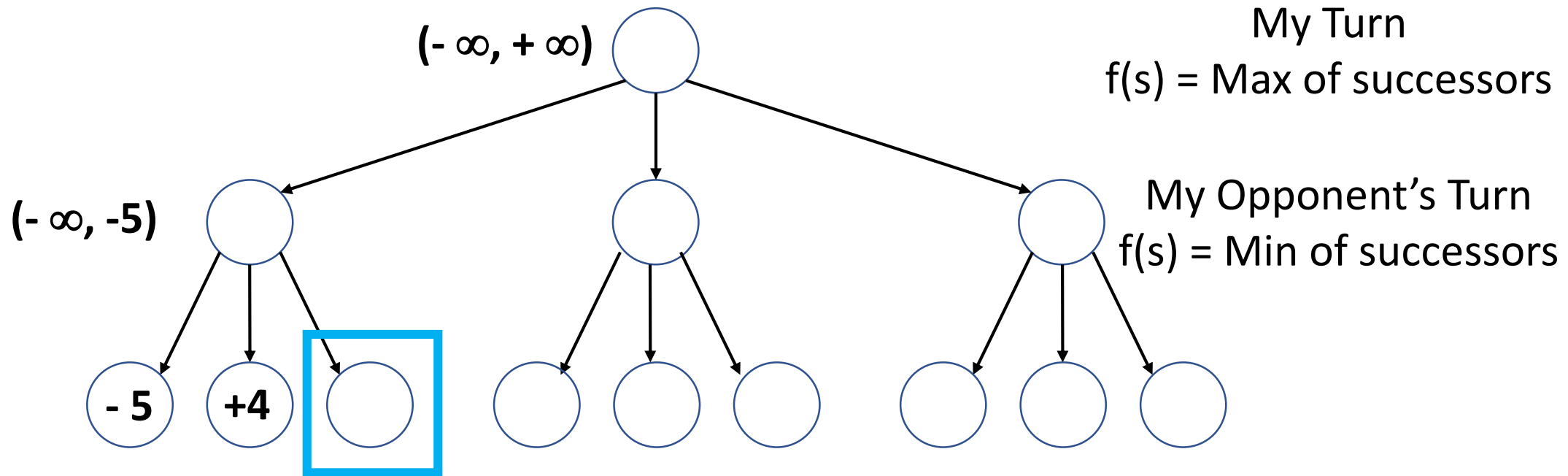
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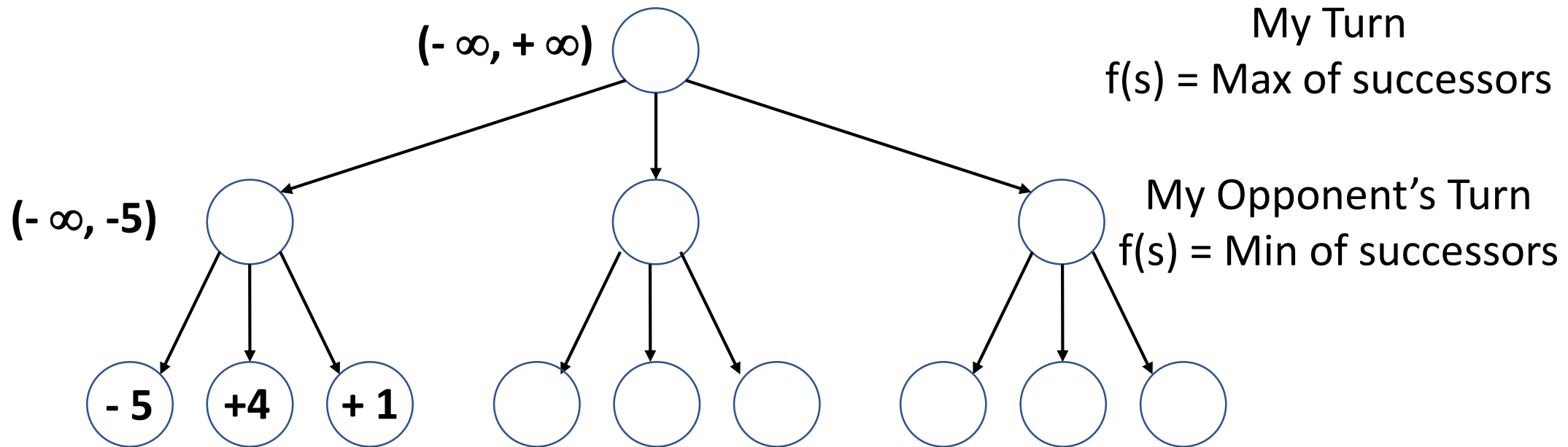
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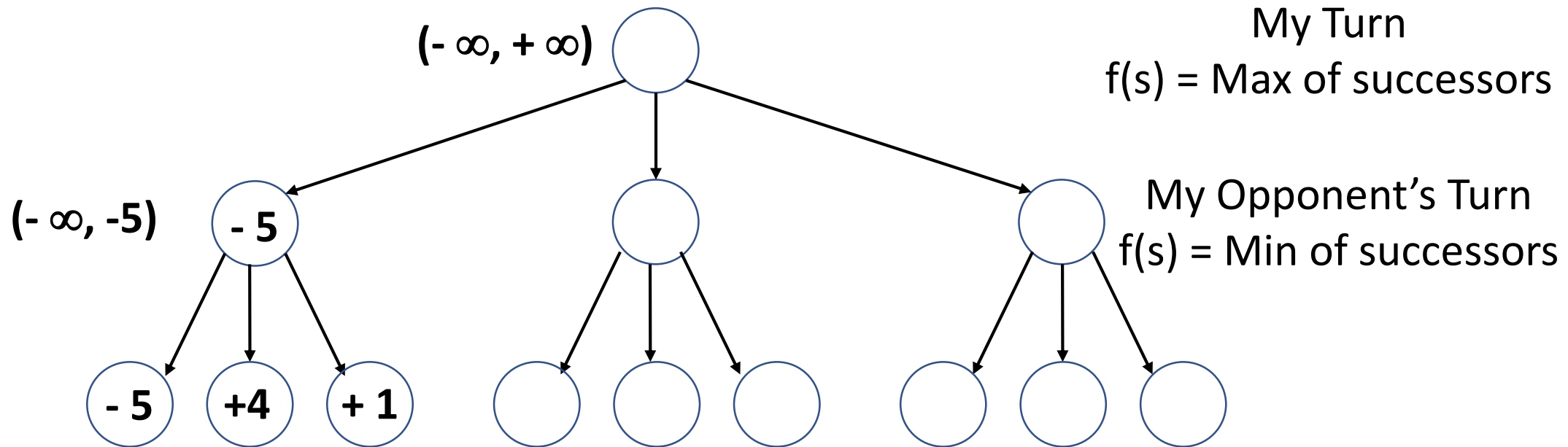
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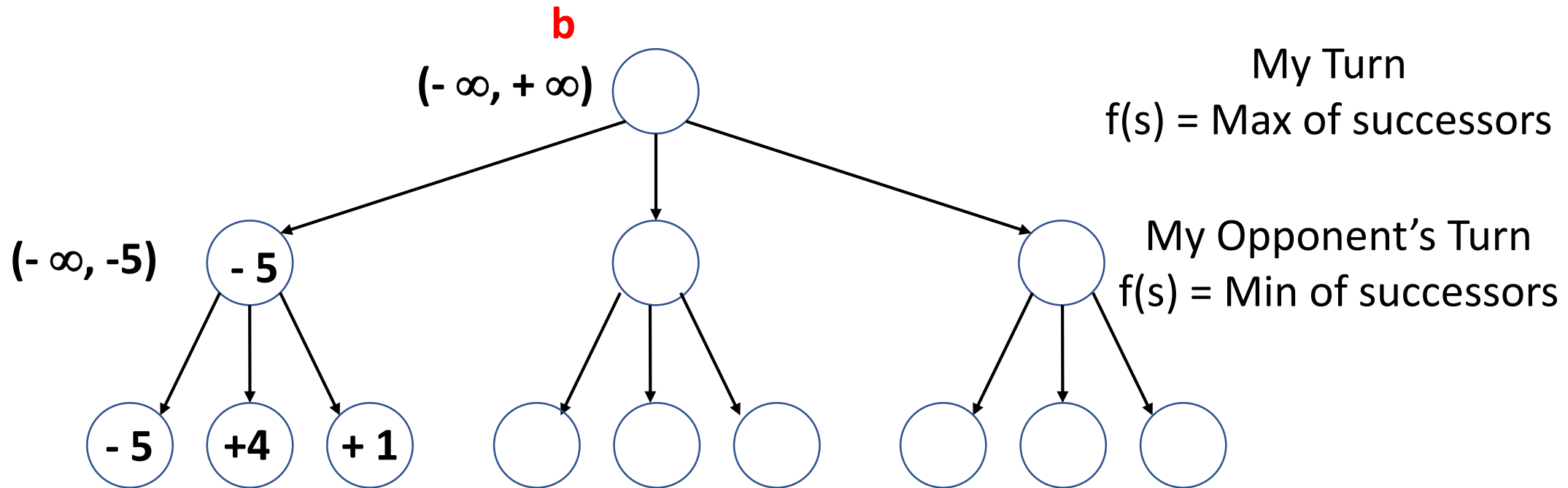
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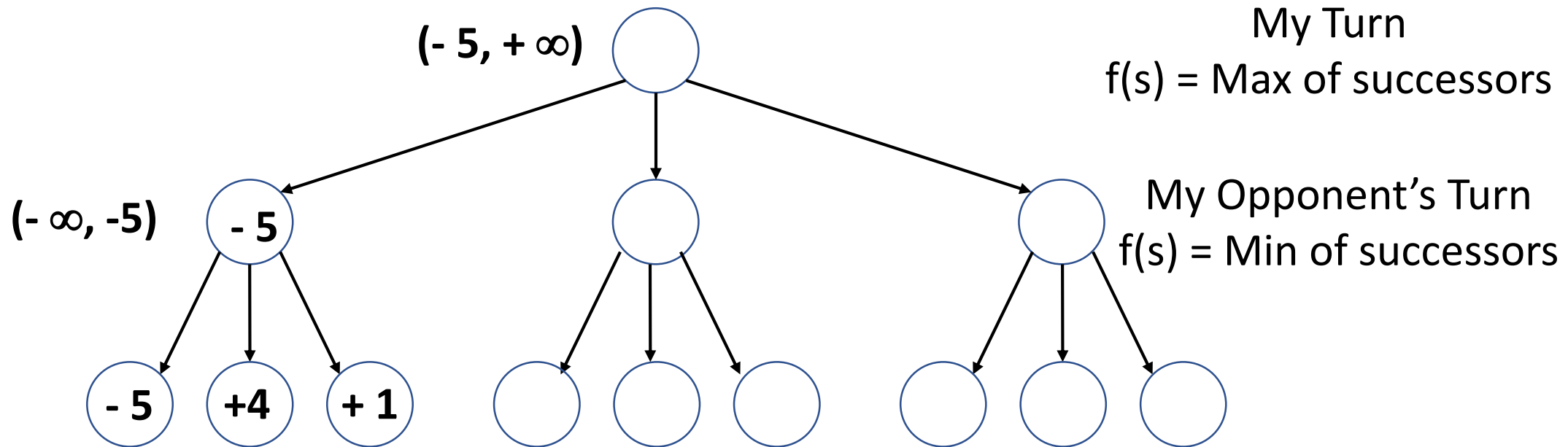
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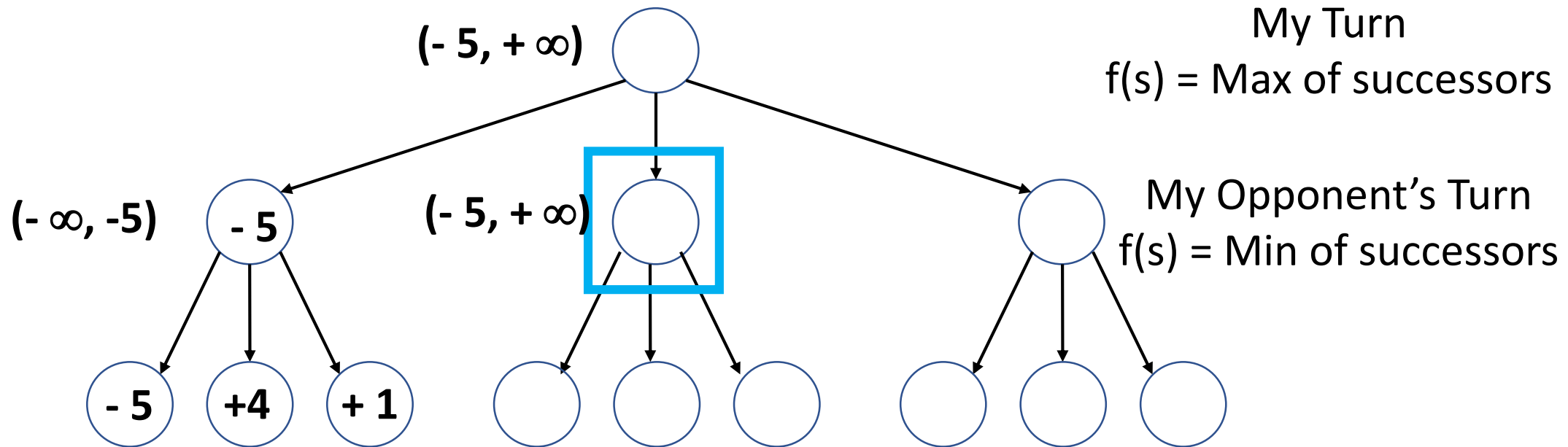
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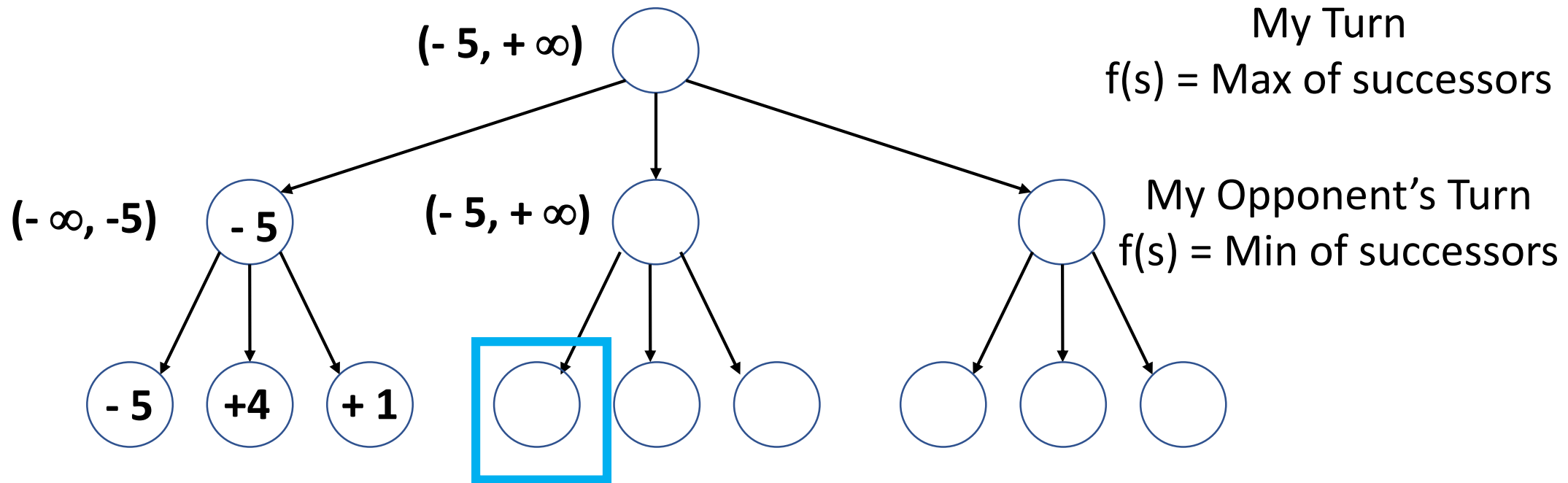
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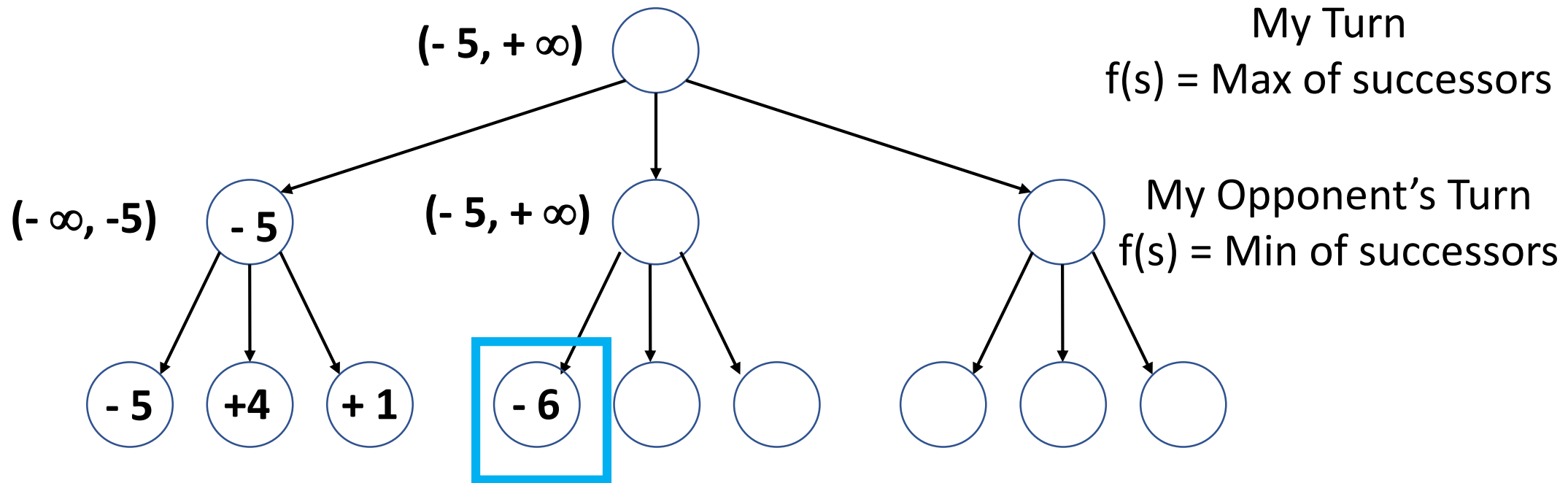
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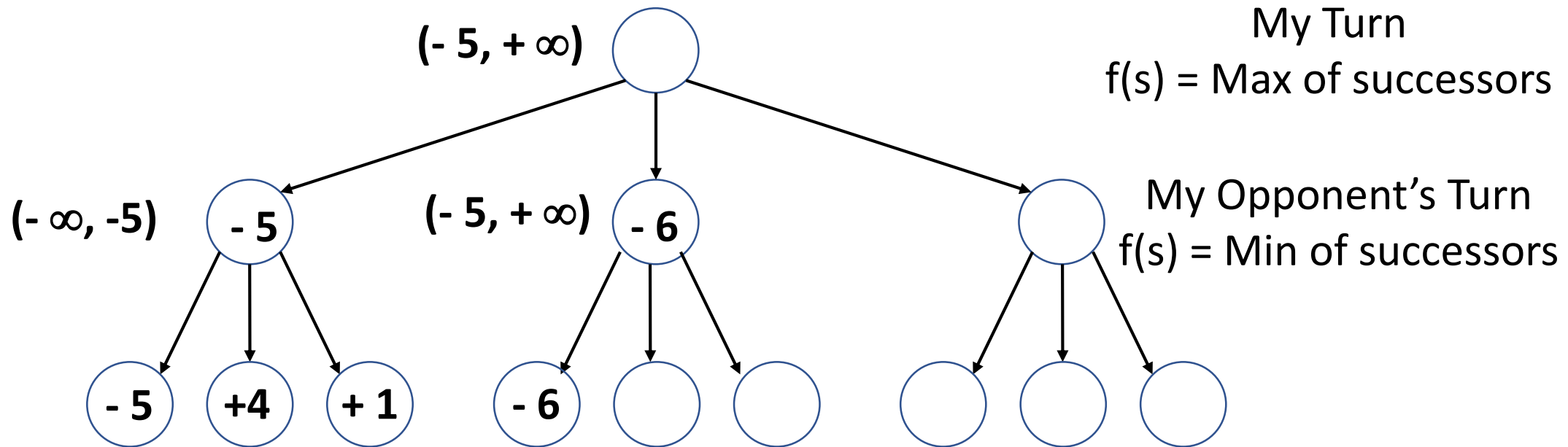
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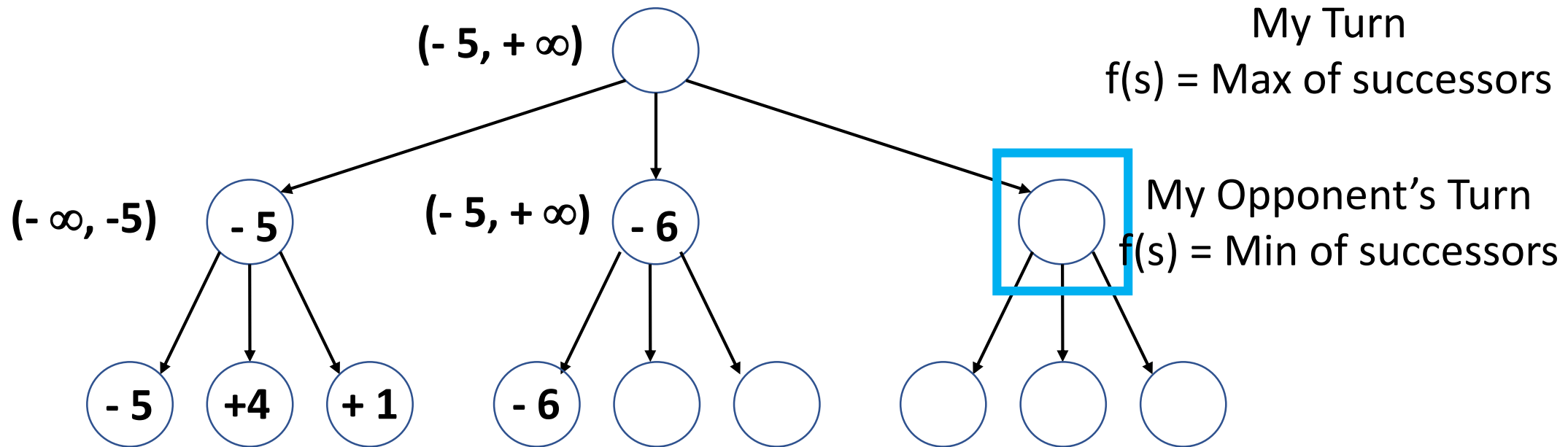
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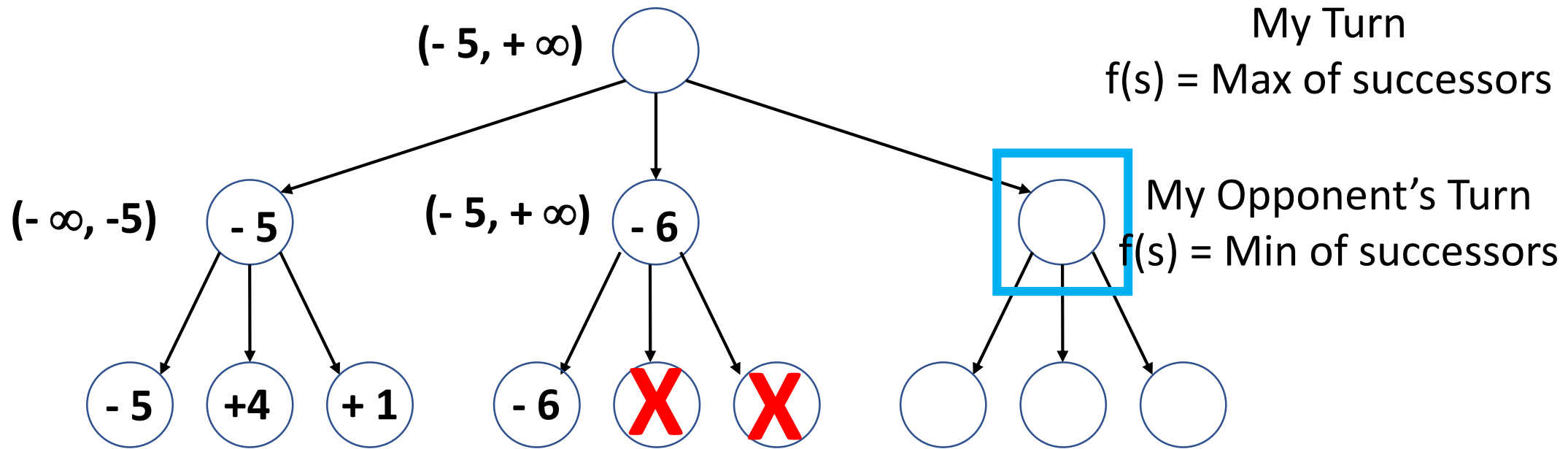
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And so on

Improving Minimax Search

For vanilla two-player zero-sum games:

Improving Minimax Search

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- Ordering (or pruning) the operators to try at each node:

Improving Minimax Search

For vanilla two-player zero-sum games:

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 - Use $f(s)$
 - Learn from game play (can be simple statistics on move outcomes)

Improving Minimax Search

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- Dealing with the “horizon effect” – the effect of a move is past depth-bound “horizon”:

Improving Minimax Search

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 - Use $f(s)$
 - Learn from game play (can be simple statistics on move outcomes)
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Improving Minimax Search

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- Random game elements (stochastic games) – example: dice

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1959

Some Studies in Machine Learning Using the Game of Checkers

Arthur L. Samuel

Improving play by learning
Book games
Self-play

Abstract: Two machine-learning procedures have been investigated in some detail using the game of checkers. Enough work has been done to verify the fact that a computer can be programmed so that it will learn to play a better game of checkers than can be played by the person who wrote the program. Furthermore, it can learn to do this in a remarkably short period of time (8 or 10 hours of machine-playing time) when given only the rules of the game, a sense of direction, and a redundant and incomplete list of parameters which are thought to have something to do with the game, but whose correct signs and relative weights are unknown and unspecified. The principles of machine learning verified by these experiments are, of course, applicable to many other situations.

Introduction

The studies reported here have been concerned with the programming of a digital computer to behave in a way which, if done by human beings or animals, would be described as involving the process of learning. While this is not the place to dwell on the importance of machine-learning procedures, or to discourse on the philosophical aspects, there is obviously a very large amount

method should lead to the development of general-purpose learning machines. A comparison between the size of the switching nets that can be reasonably constructed or simulated at the present time and the size of the neural nets used by animals, suggests that we have a long way to go before we obtain practical devices.² The second procedure requires reprogramming for each new applica-

1989

A Parallel Network that Learns to Play Backgammon

G. Tesauro*

*Center for Complex Systems Research, University of Illinois
at Urbana-Champaign, 508 S. Sixth St., Champaign,
IL 61820, U.S.A.*

T.J. Sejnowski**

*Biophysics Department, The Johns Hopkins University,
Baltimore, MD 21218, U.S.A.*

ABSTRACT

A class of connectionist networks is described that has learned to play backgammon at an intermediate-to-advanced level. The networks were trained by back-propagation learning on a large set of sample positions evaluated by a human expert. In actual match play against humans and conventional computer programs, the networks have demonstrated substantial ability to generalize on the basis of expert knowledge of the game. This is possibly the most complex domain yet studied with connectionist learning. New techniques were needed to overcome problems due to the scale and complexity of the task. These include techniques for intelligent design of training set examples and efficient coding schemes, and procedures for escaping from local minima. We suggest how these techniques might be used in applications of network learning to general large-scale, difficult "real-world" problem domains.

Function approximation
for evaluation function
Repeatedly simulate games

Improving Minimax Search

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- Games of hidden information – example: other players' hands

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Much of this is captured in Monte Carlo Tree Search (coming later)

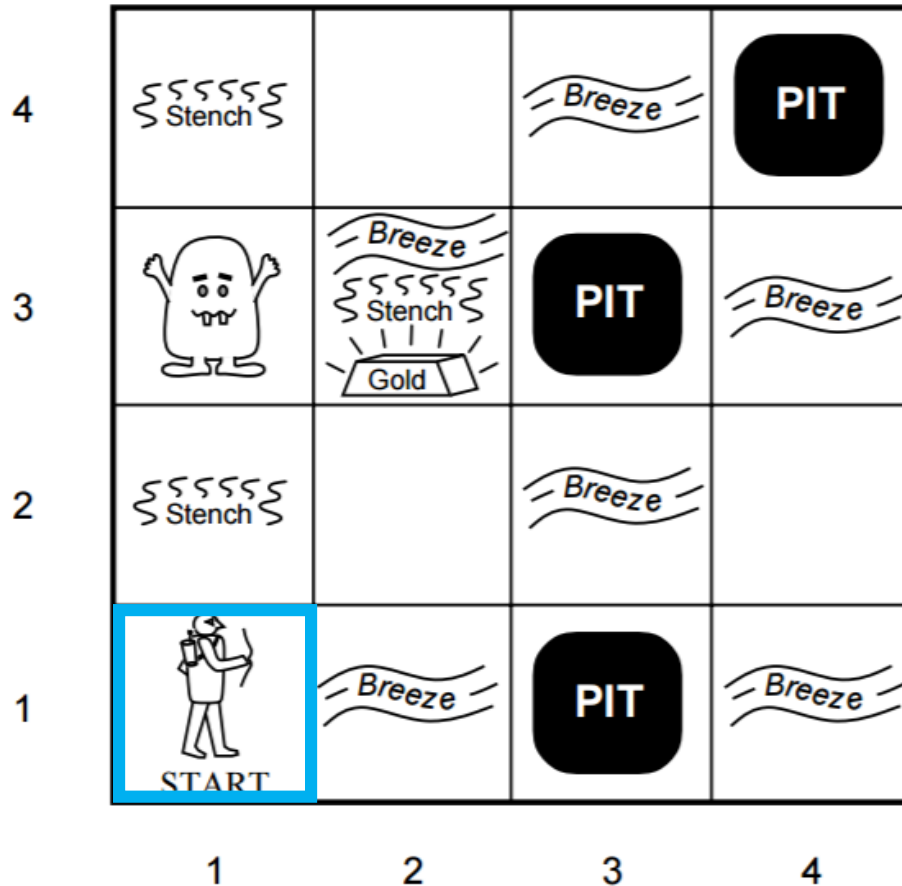
Next Topic:
Knowledge Representation and Reasoning
Textbook Chapter 7

Knowledge Representation and Reasoning

- Agents:
 - Represent what they know about the world
 - Use inference to derive new information

Sometimes called the “Logicist Approach” to AI

Wumpus World



Get points for taking gold

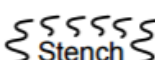
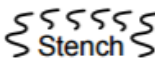
Die if in the same square as pit or wumpus

Can move Up, Down, Left, Right

Sensors:

- Stench: Wumpus is 1 away
- Breeze: Pit is 1 away
- Glitter: Gold is in current room
- States: $[x, y, \text{Stench?}, \text{Breeze?}, \text{Glitter?}]$
 - Initial state: $[1, 1, \text{False}, \text{False}, \text{False}]$

Wumpus World

4	 Stench		 Breeze	
3		 Breeze  Stench  Gold		 Breeze
2	 Stench		 Breeze	
1	 START	 Breeze		 Breeze
	1	2	3	4

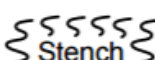
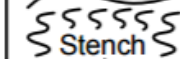
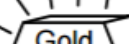
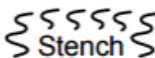
Definition of world
(not known to agent)

1,4	2,4	3,4	4,4
1,3	2,3	3,3	4,3
1,2	2,2	3,2	4,2
1,1  OK	2,1	3,1	4,1

State = [1,1,False,False,False]

What the agent knows

Wumpus World

4	 Stench		 Breeze	
3		 Breeze  Stench  Gold		 Breeze
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1	 START	 Breeze		 Breeze
	1	2	3	4

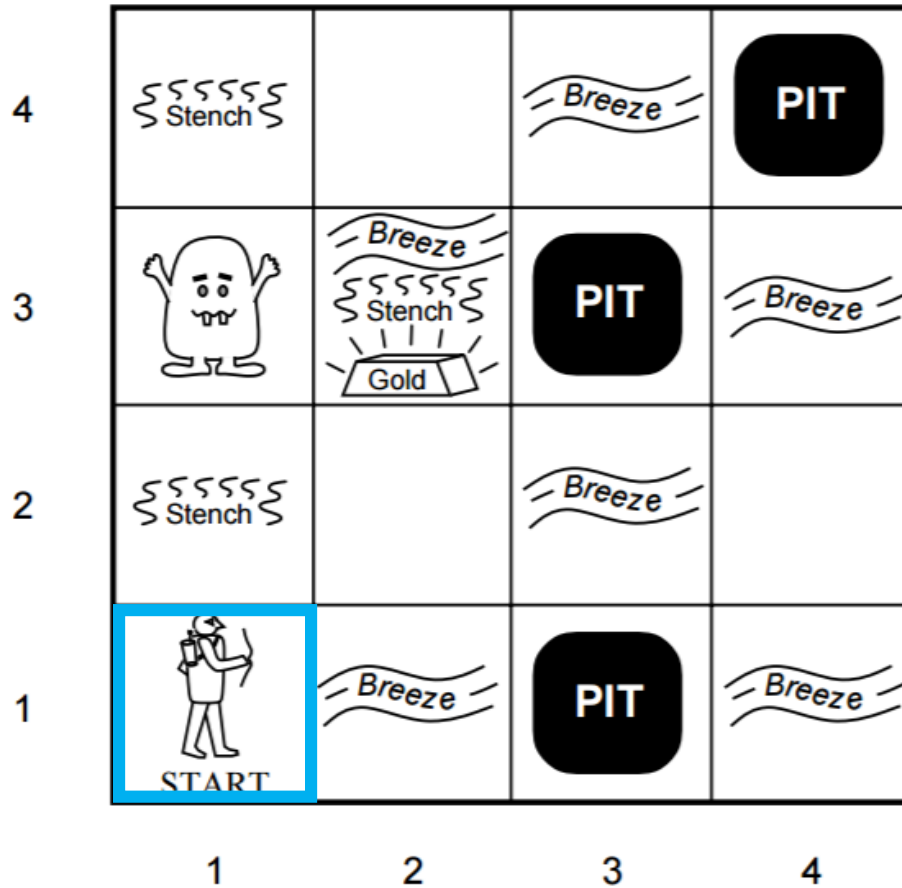
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1,4	2,4	3,4	4,4
1,3	2,3	3,3	4,3
1,2	2,2	3,2	4,2
1,1  OK	2,1	3,1	4,1

OK means "safe" (won't die)

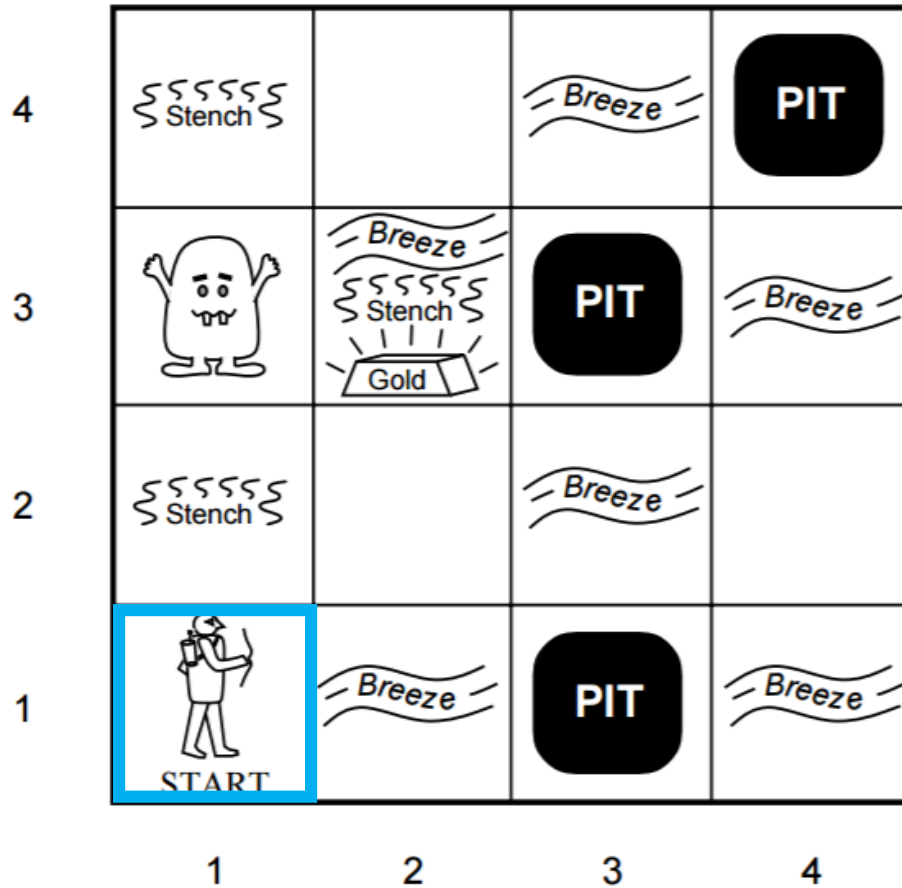
What the agent knows

Wumpus World



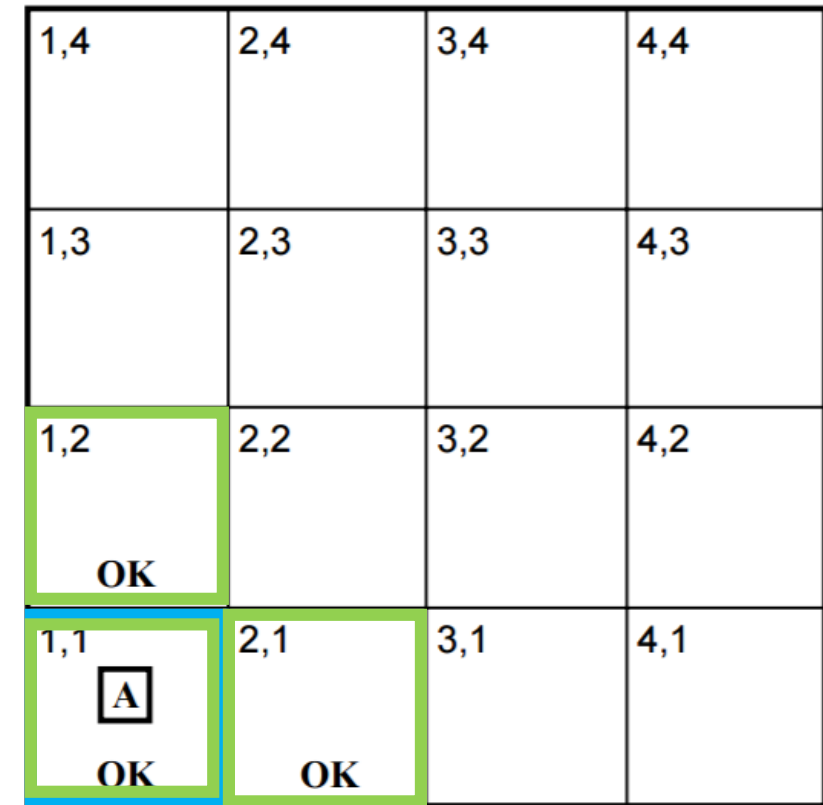
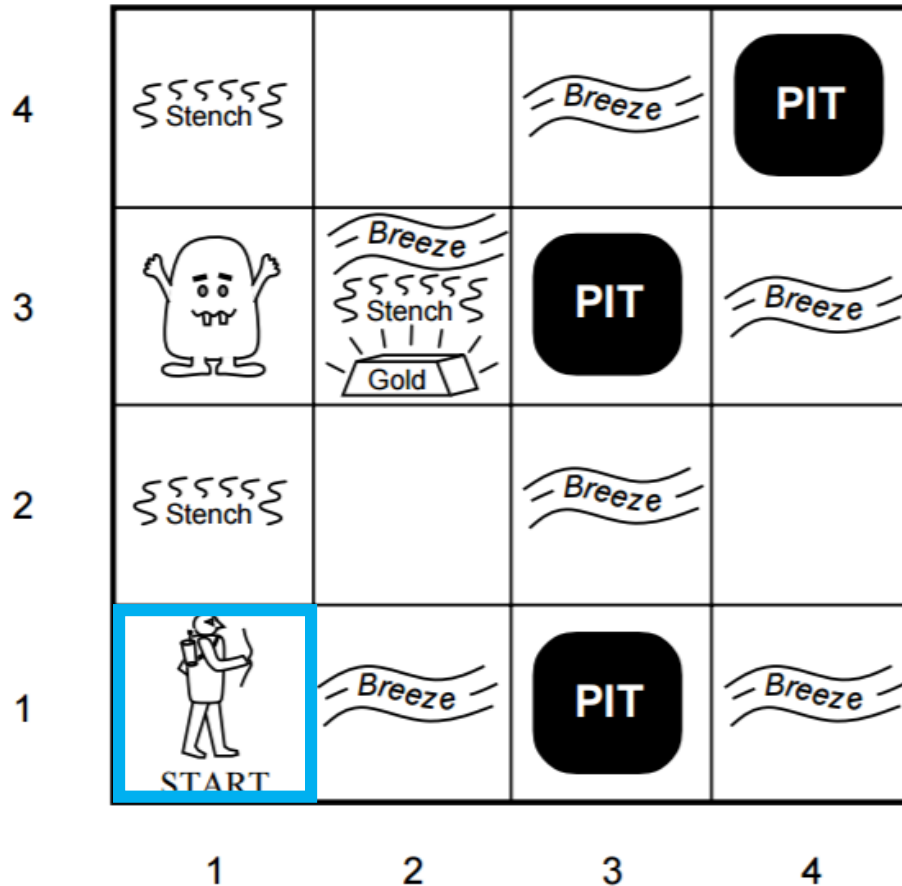
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Wumpus World



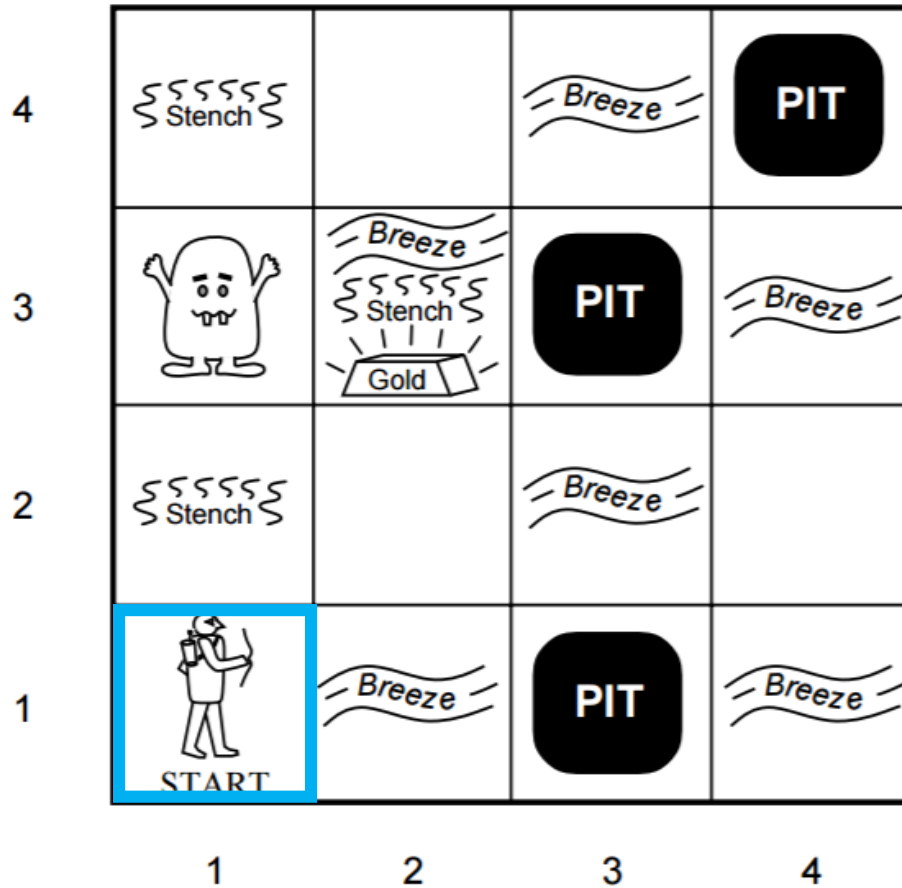
1,4	2,4	3,4	4,4
1,3	2,3	3,3	4,3
1,2 OK	2,2	3,2	4,2
1,1 A OK	2,1 OK	3,1	4,1

Wumpus World



No stench => no Wumpus in adjacent squares
 No breeze => no pit in adjacent squares

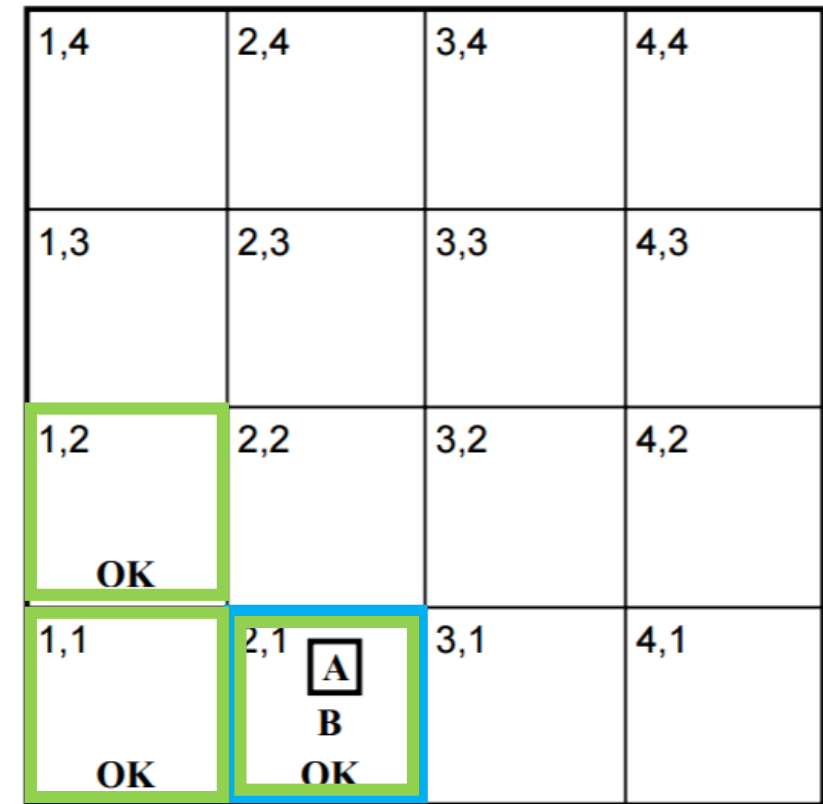
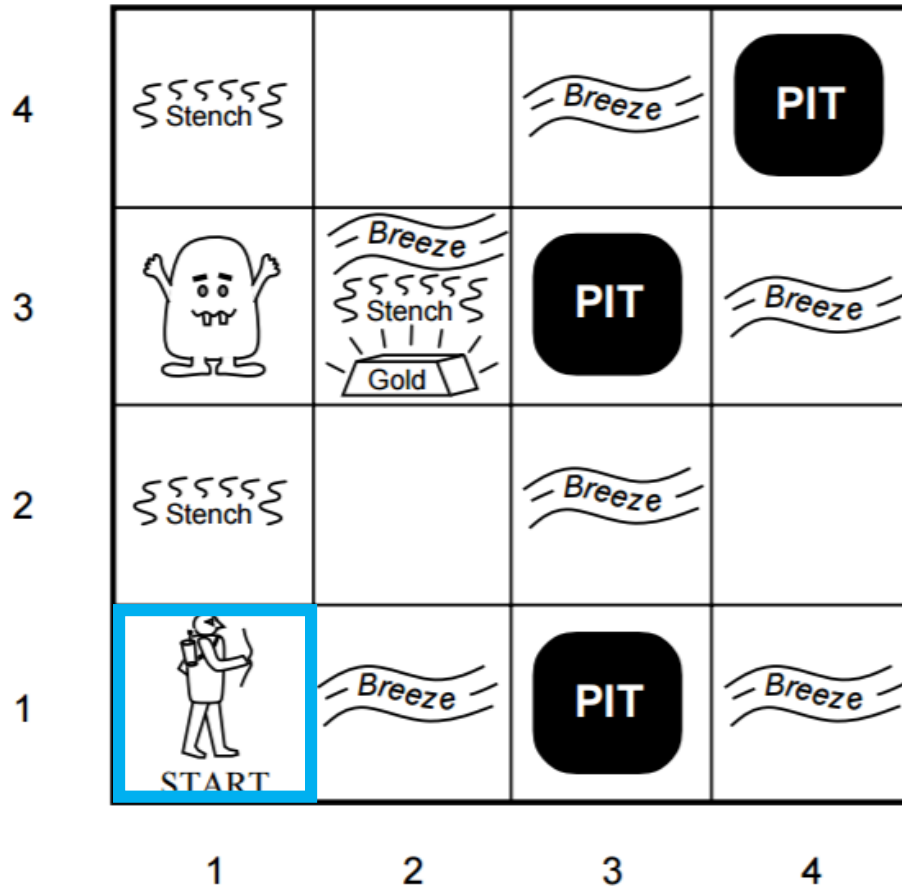
Wumpus World



1,4	2,4	3,4	4,4
1,3	2,3	3,3	4,3
1,2 OK	2,2	3,2	4,2
1,1 A OK	2,1 OK	3,1	4,1

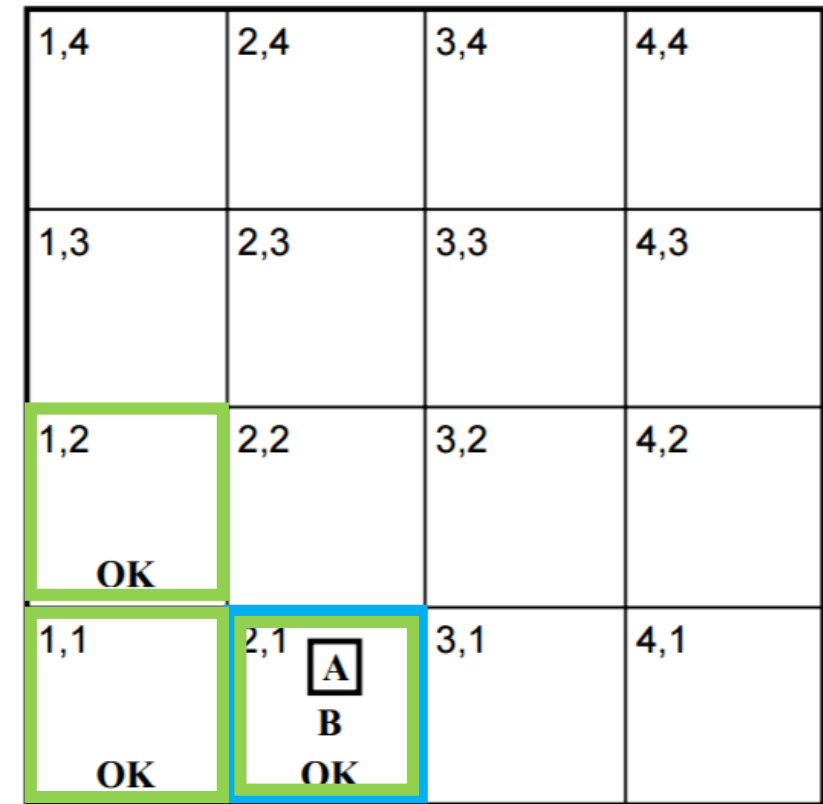
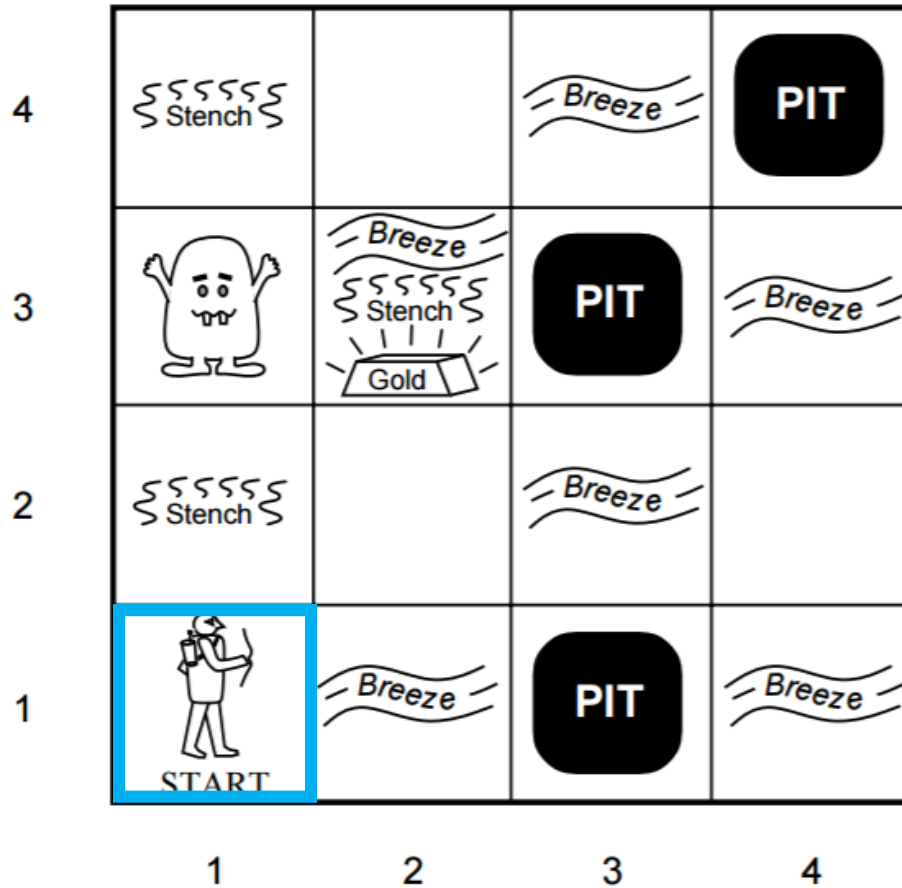
No glitter => no gold here

Wumpus World



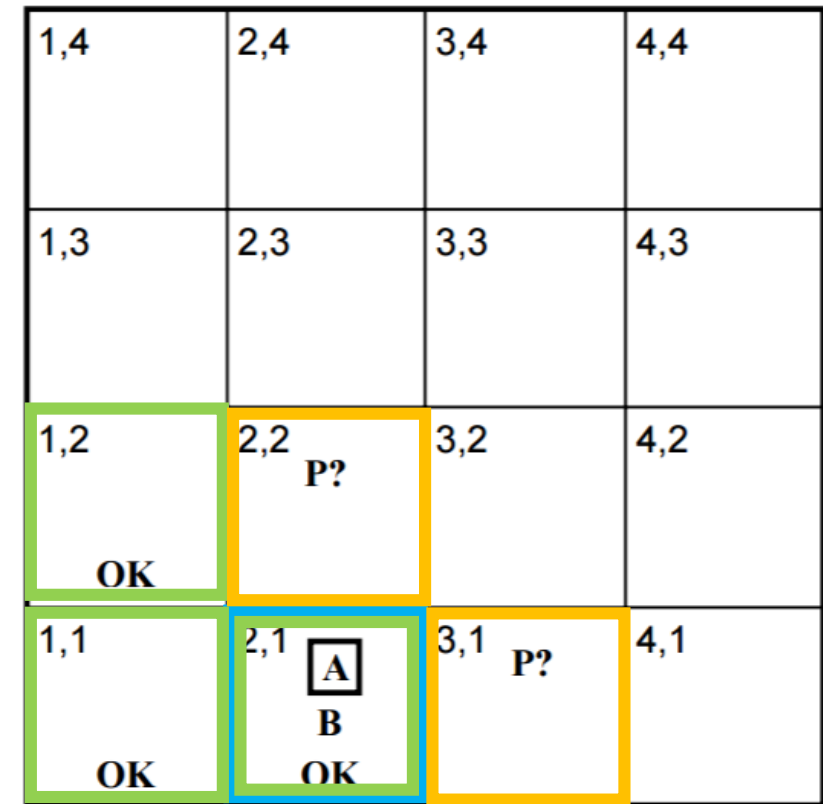
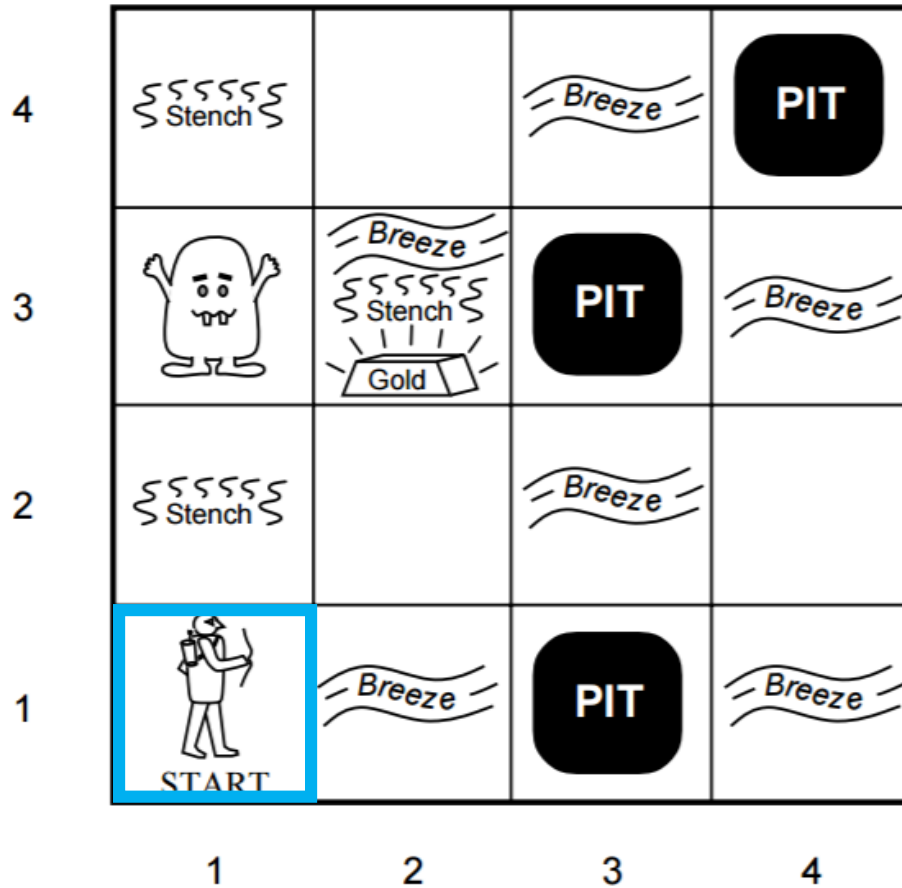
Wumpus moves Right

Wumpus World

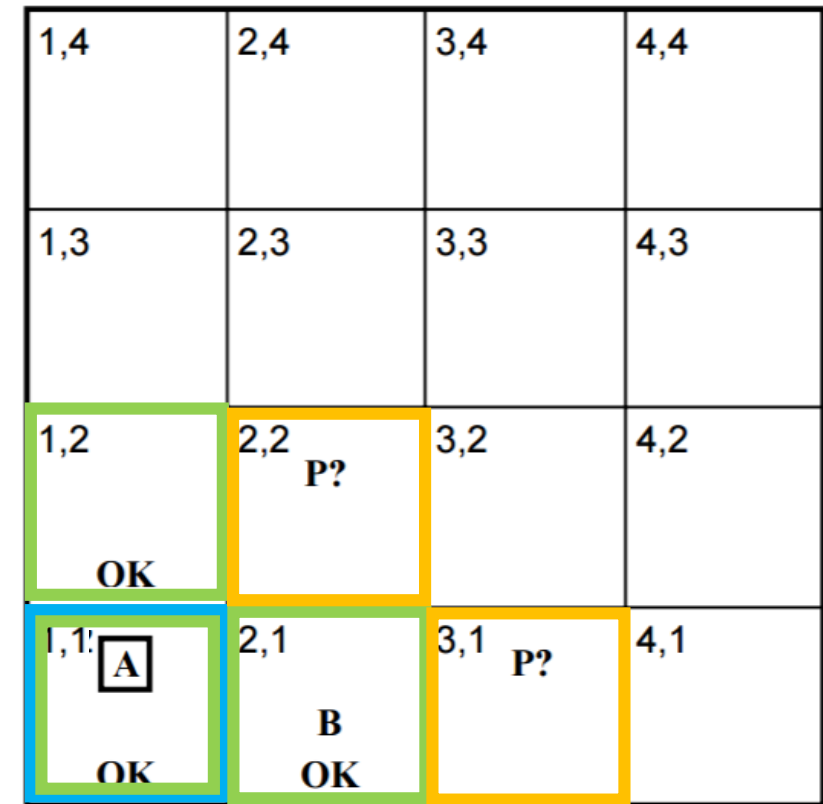
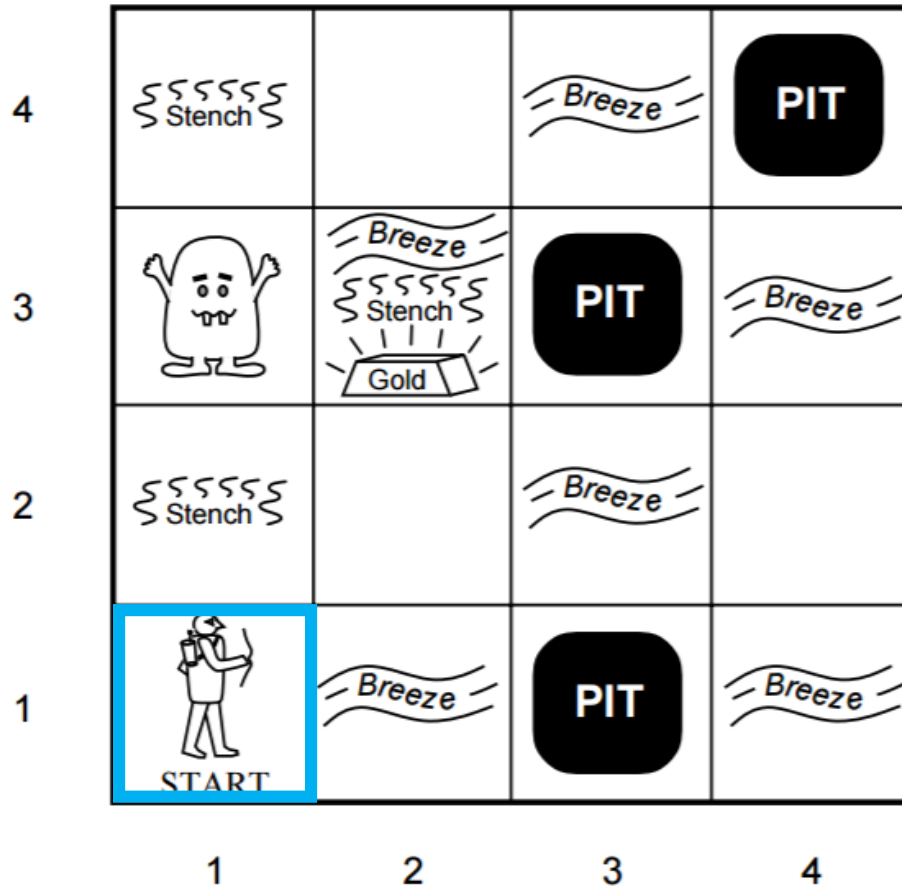


State = [2,1,False,True,False]

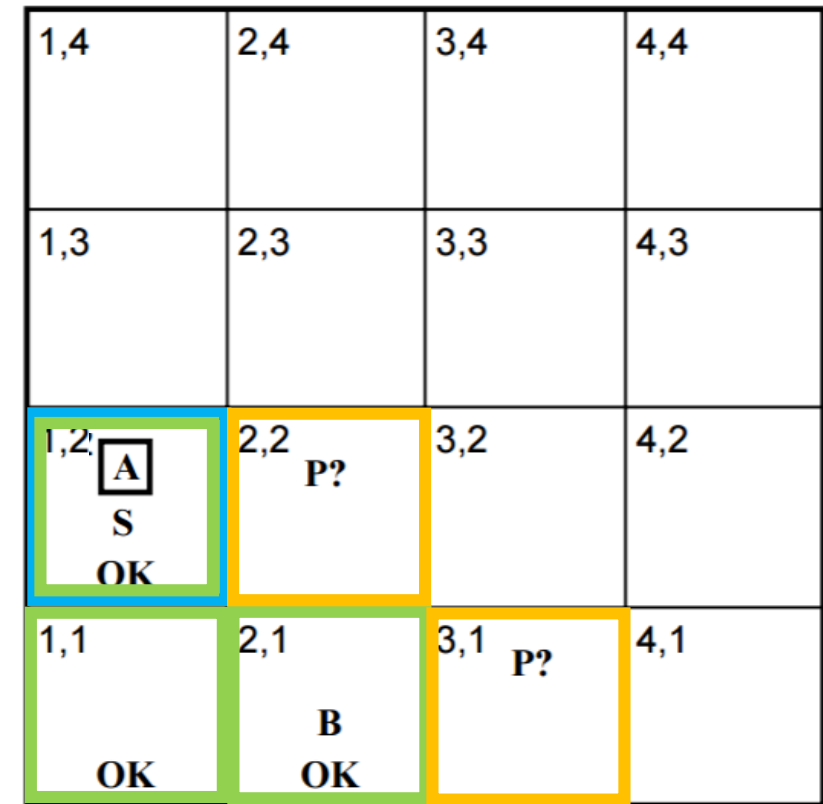
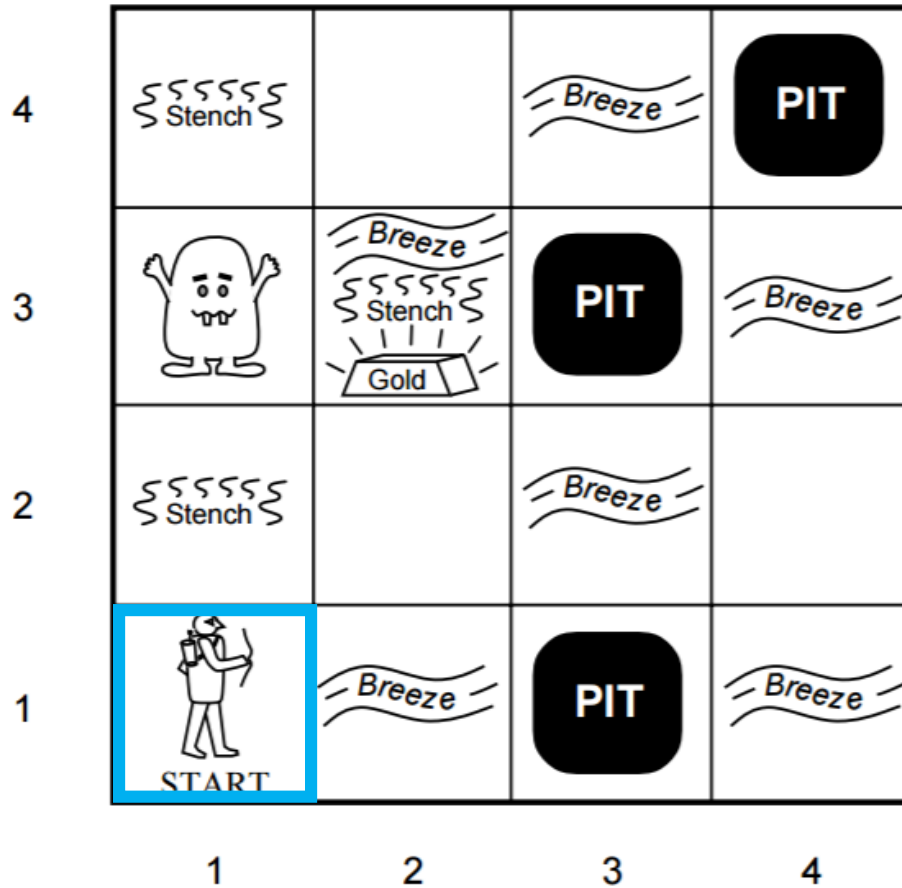
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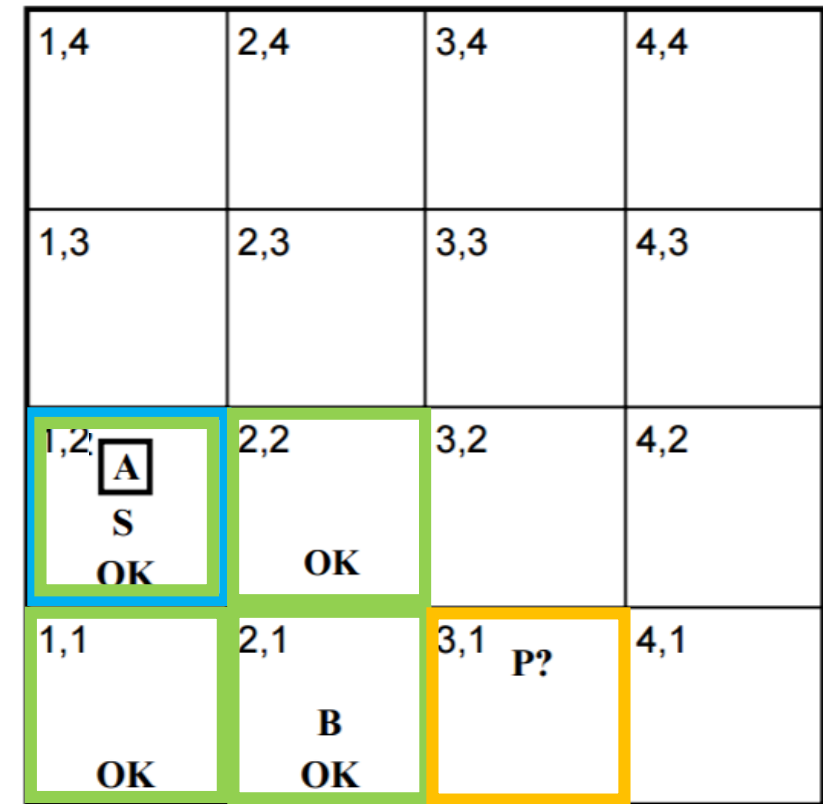
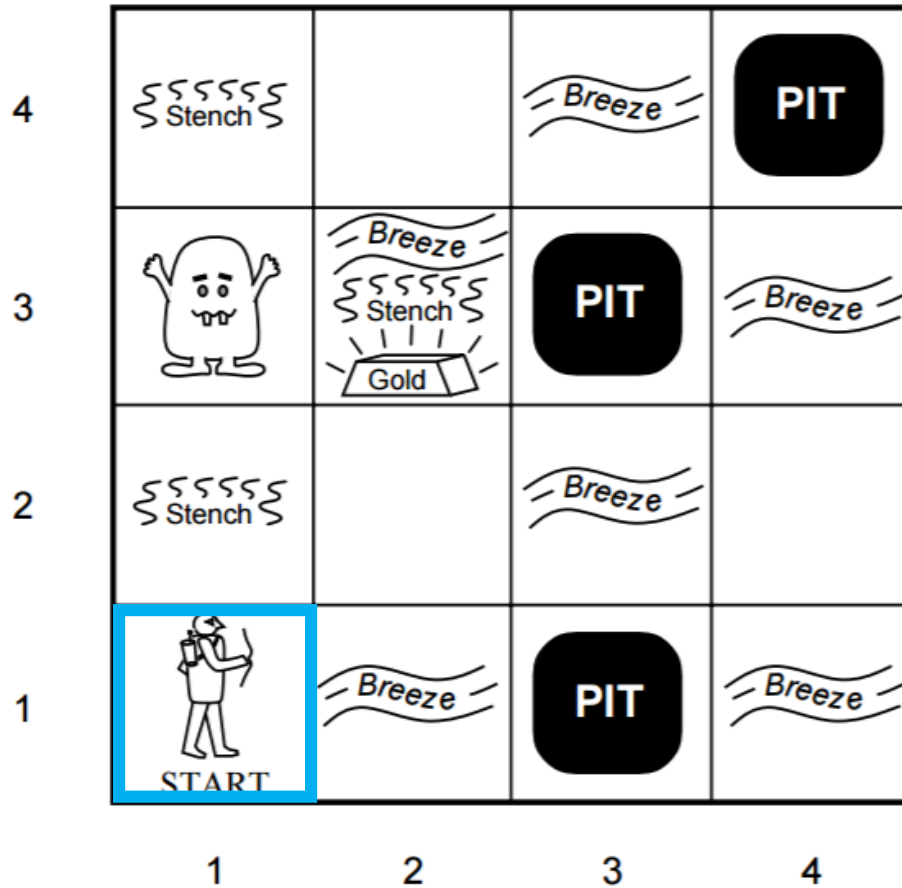
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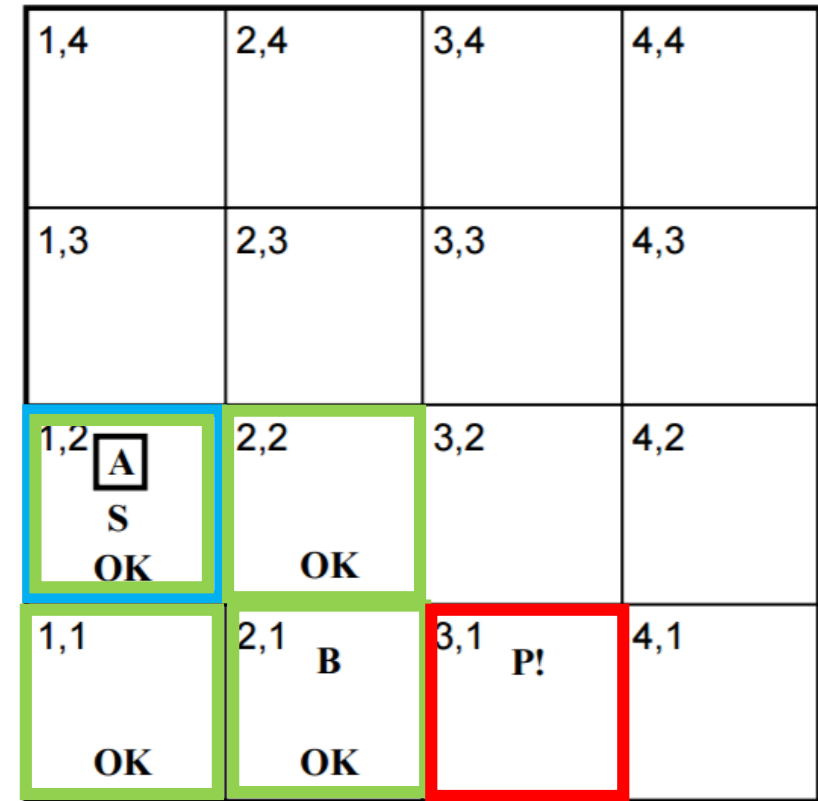
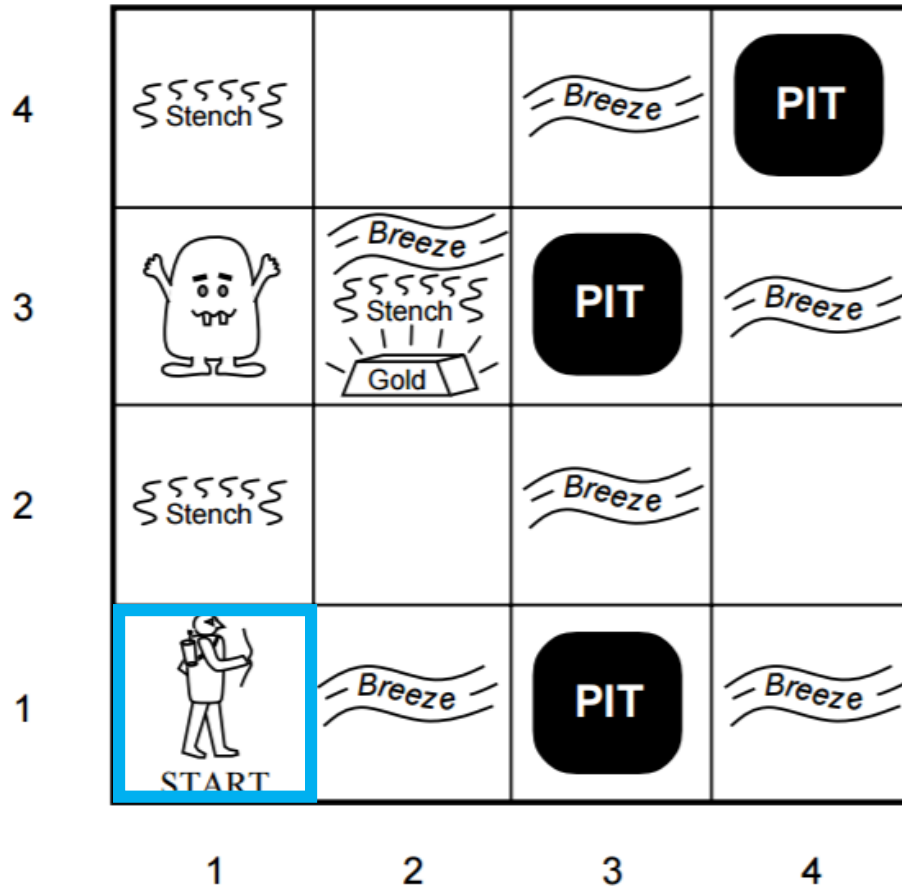
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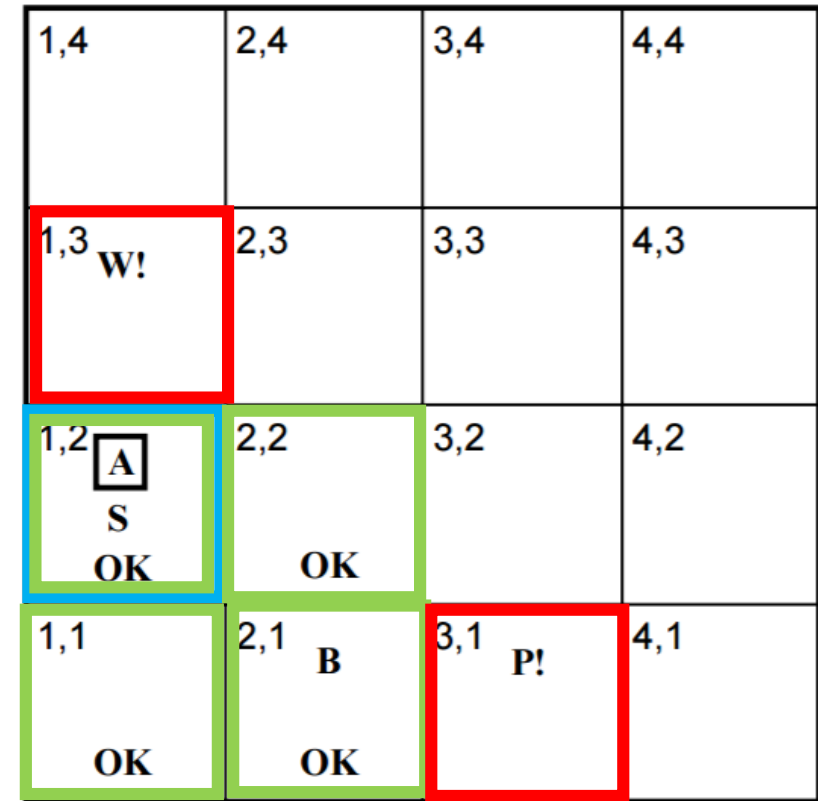
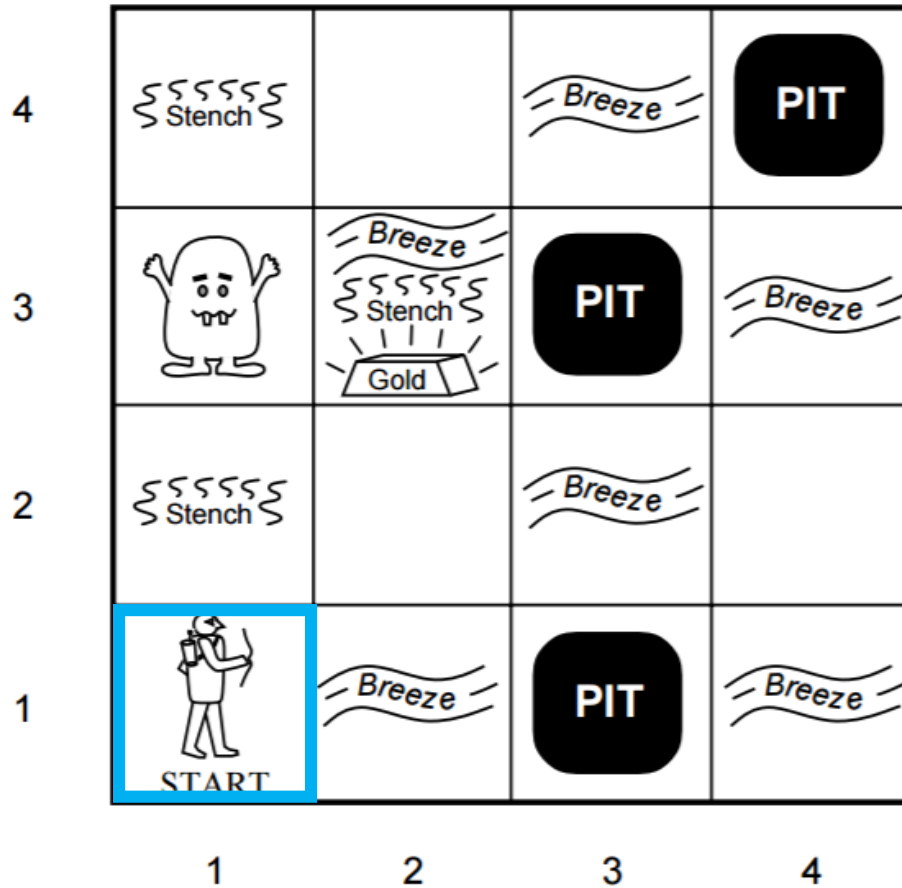
Wumpus World



Wumpus World



Wumpus World



Knowledge Representation and Reasoning

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Focus here will be on formal logic