

CS 4700: Foundations of Artificial Intelligence

Spring 2020
Prof. Haym Hirsh

Lecture 7
February 5, 2020

Reminder: Technology Policy

No technology except for first four rows of left and right sides

Jupyter Notebooks

“Primer”: Friday 5pm in Gates G01

Covers:

Jupyter Notebooks and .ipynb files

Google Colab (online Jupyter Notebook environment)

Useful Python features

Bring your laptop

Recap

- Uninformed search
 - Depth-first search
 - Breadth-first search
 - Iterative deepening search
- Informed search
 - Best-first search
 - A*
 - Hill climbing
 - Simulated annealing
- Genetic algorithms/evolutionary computation

Other Search Topics (not covered - yet)

- Continuous spaces / gradient-based methods
- Non-deterministic actions
 - Oops, I dropped the cup
- Partially observable states
 - What's behind me?

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Will come up in
machine learning

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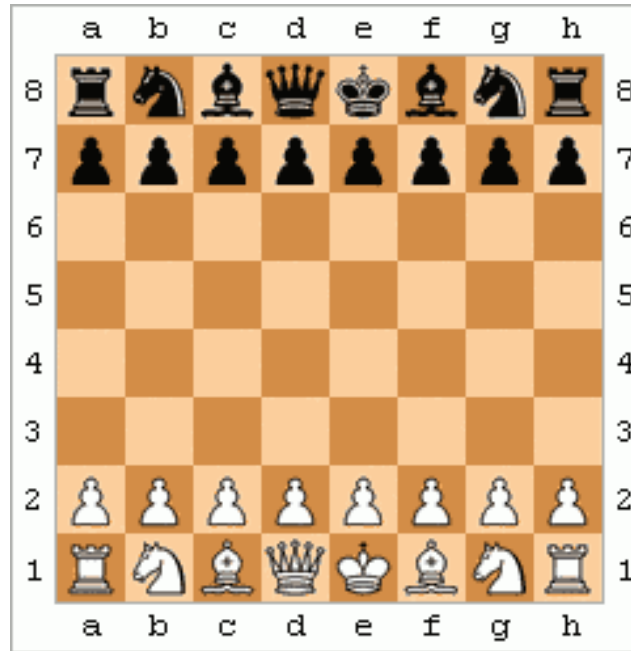
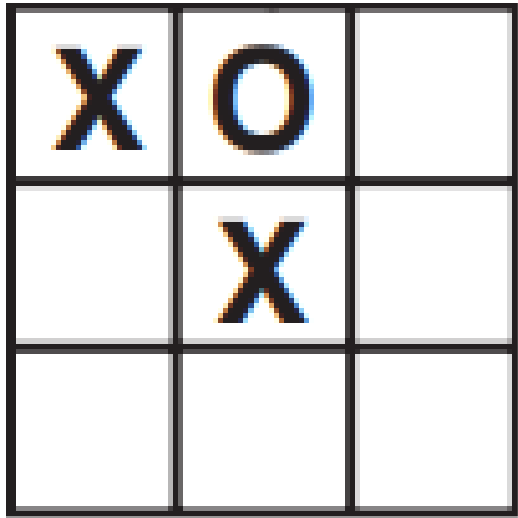
Will come up in
machine learning

Will come up in
reinforcement learning

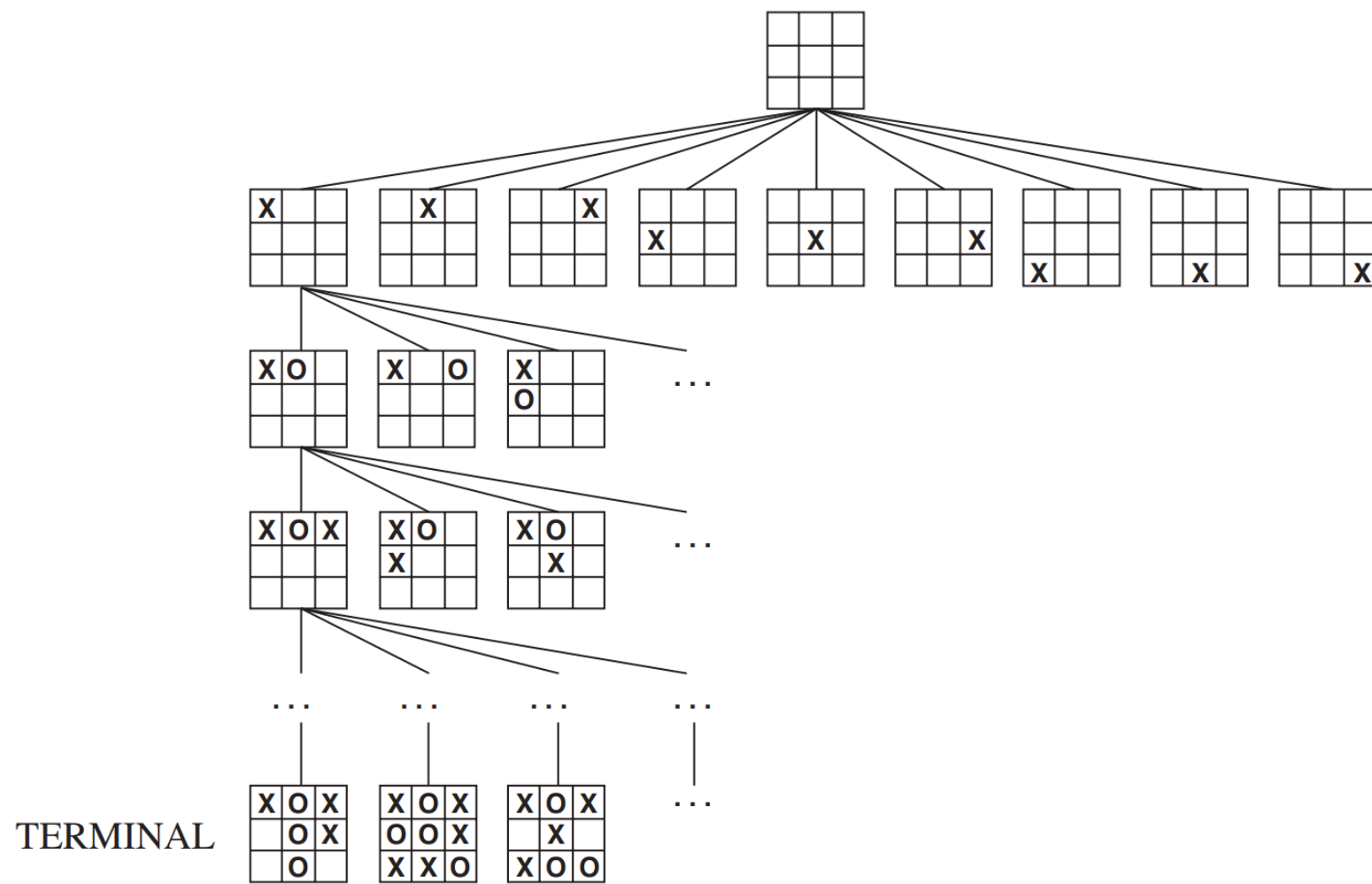
Adversarial Search (Playing Games)

Chapter 5

Two Player Zero-Sum Games

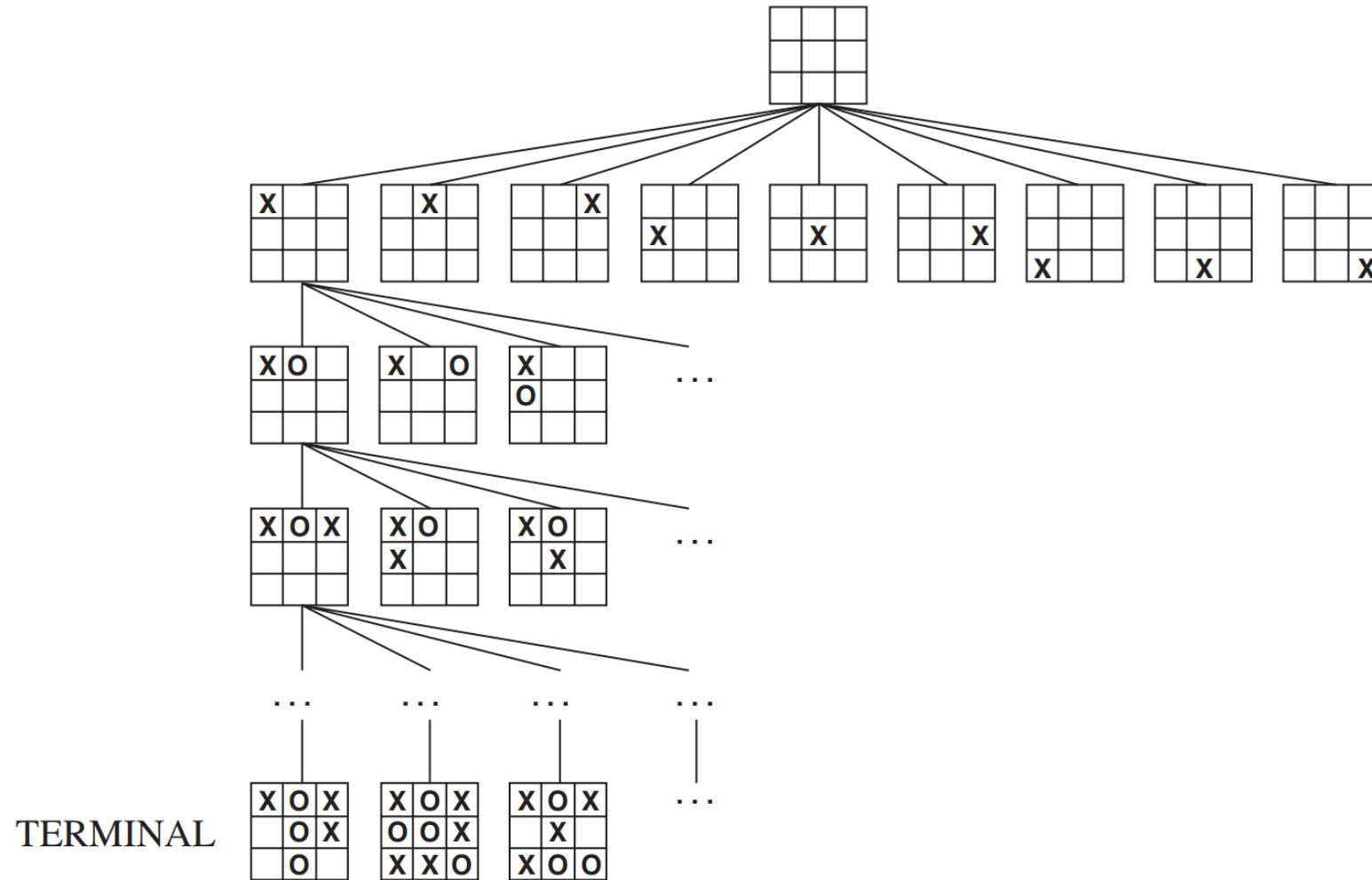


Tic Tac Toe

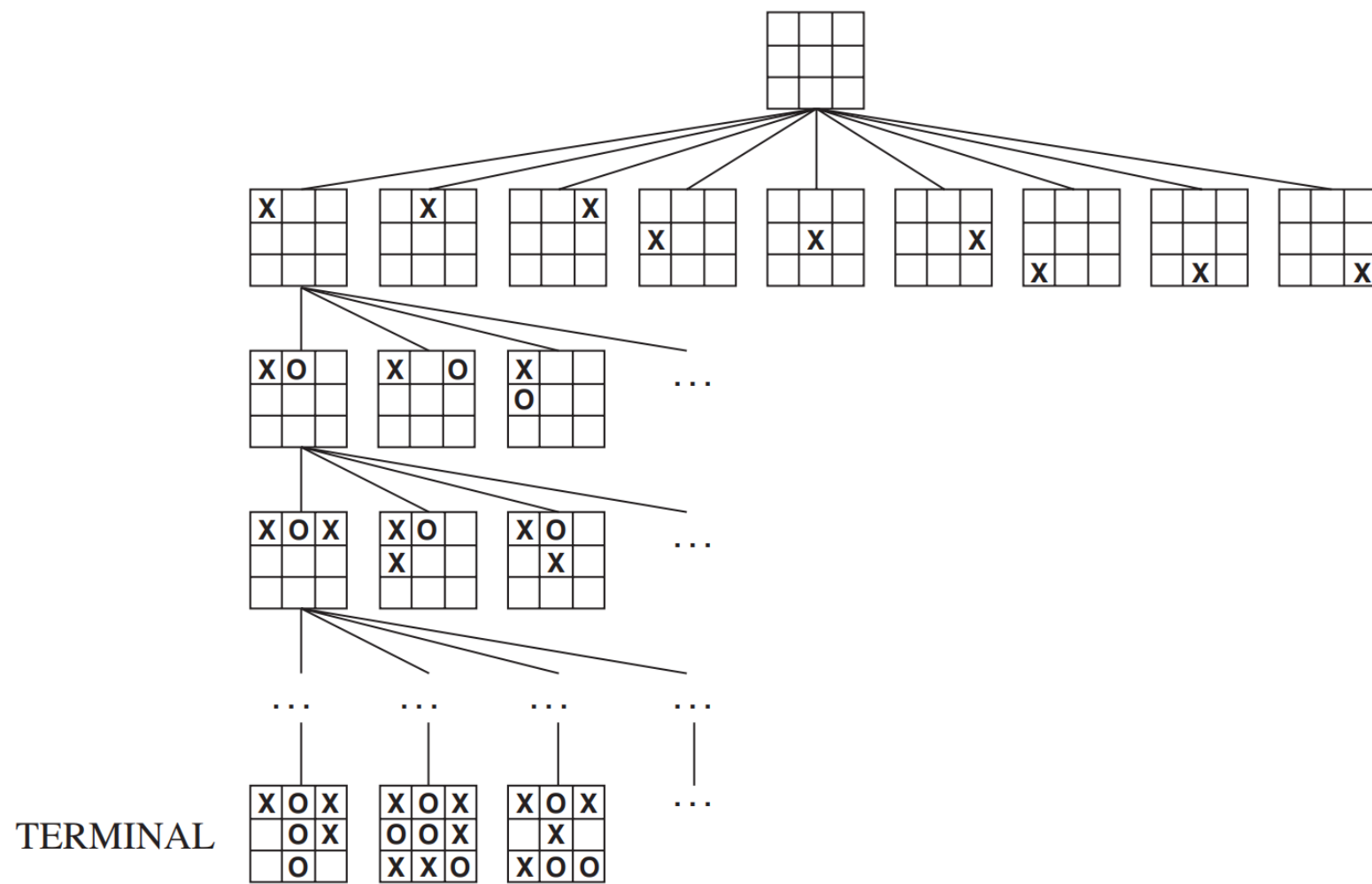


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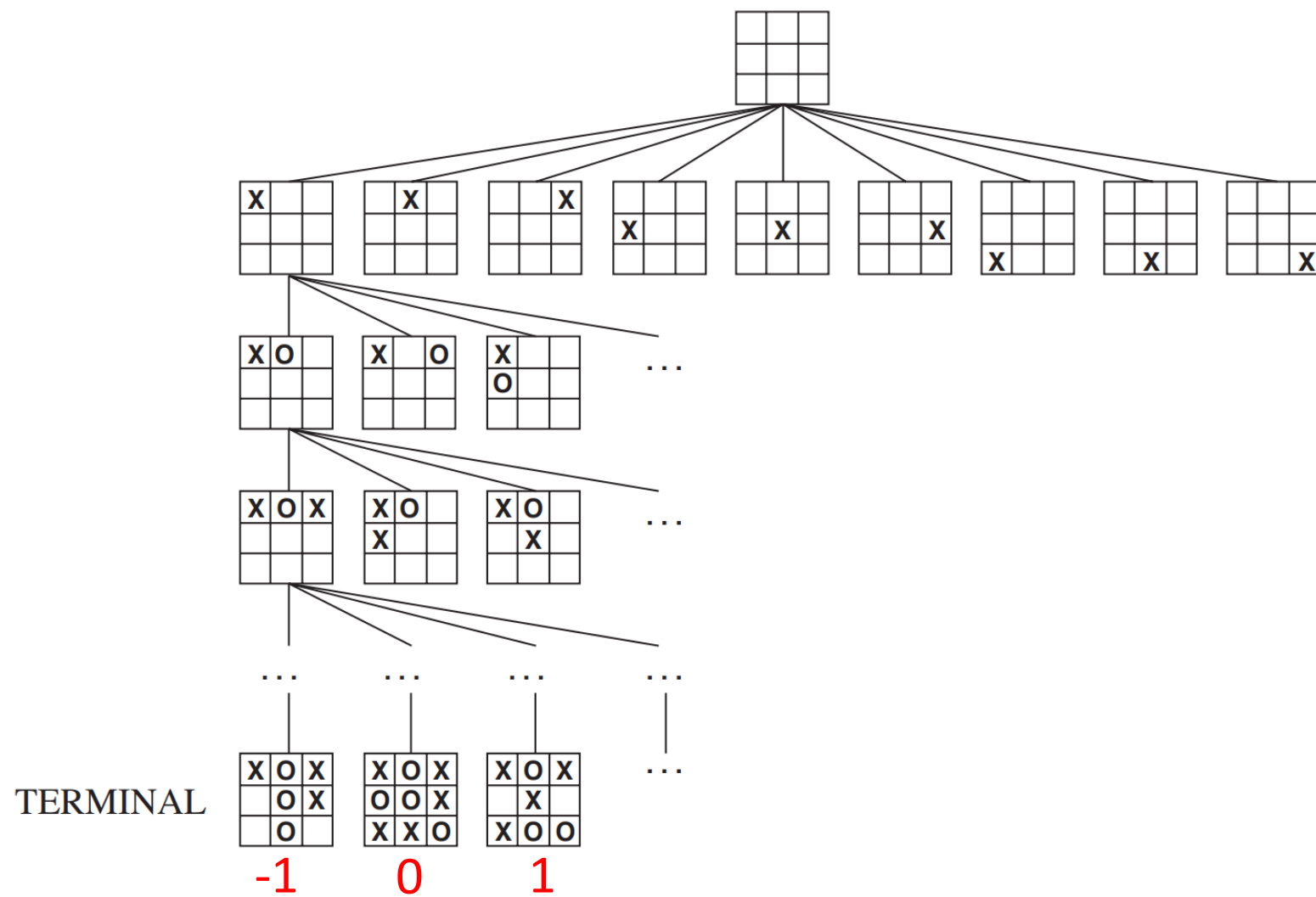
(Still using states and actions formulation)



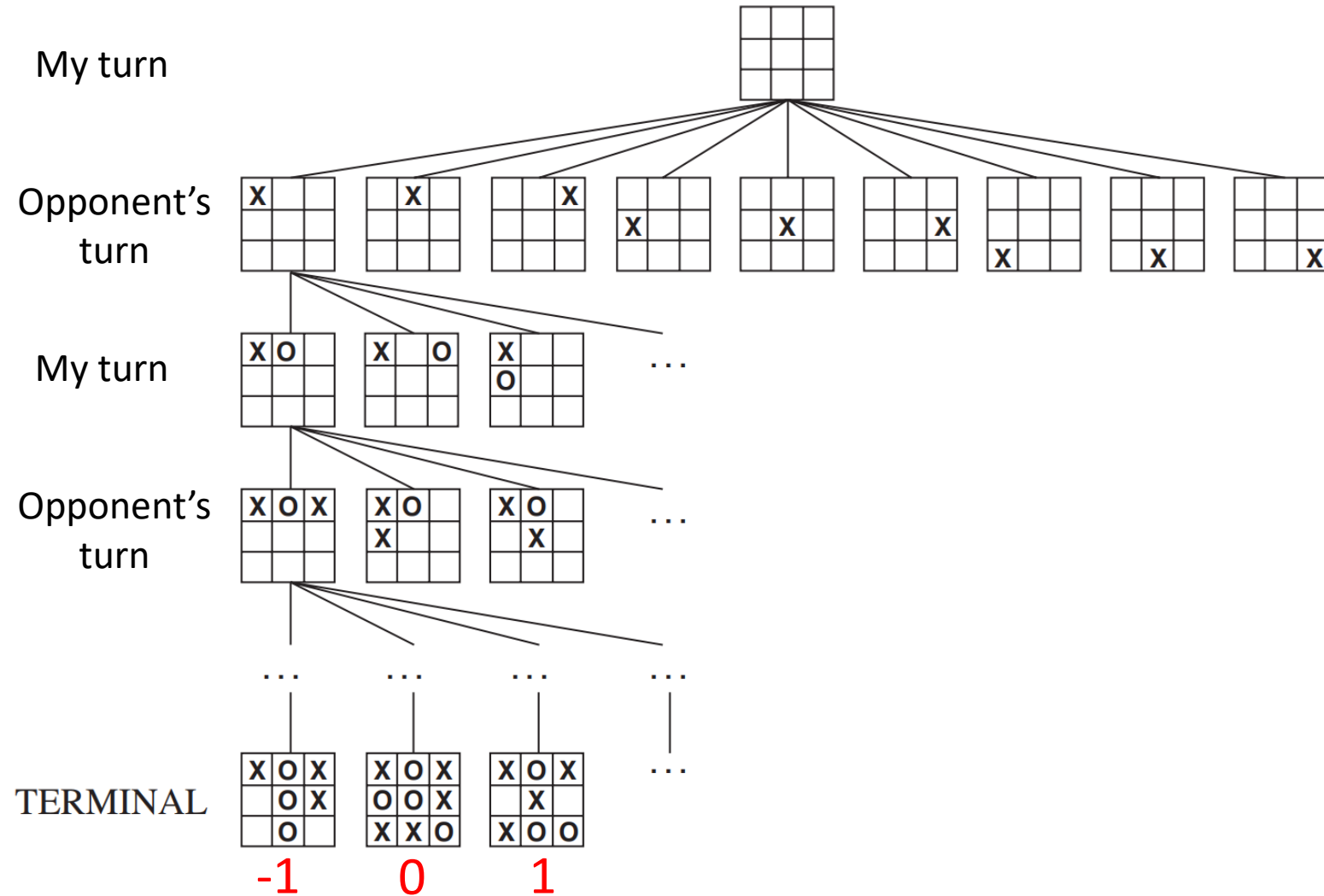
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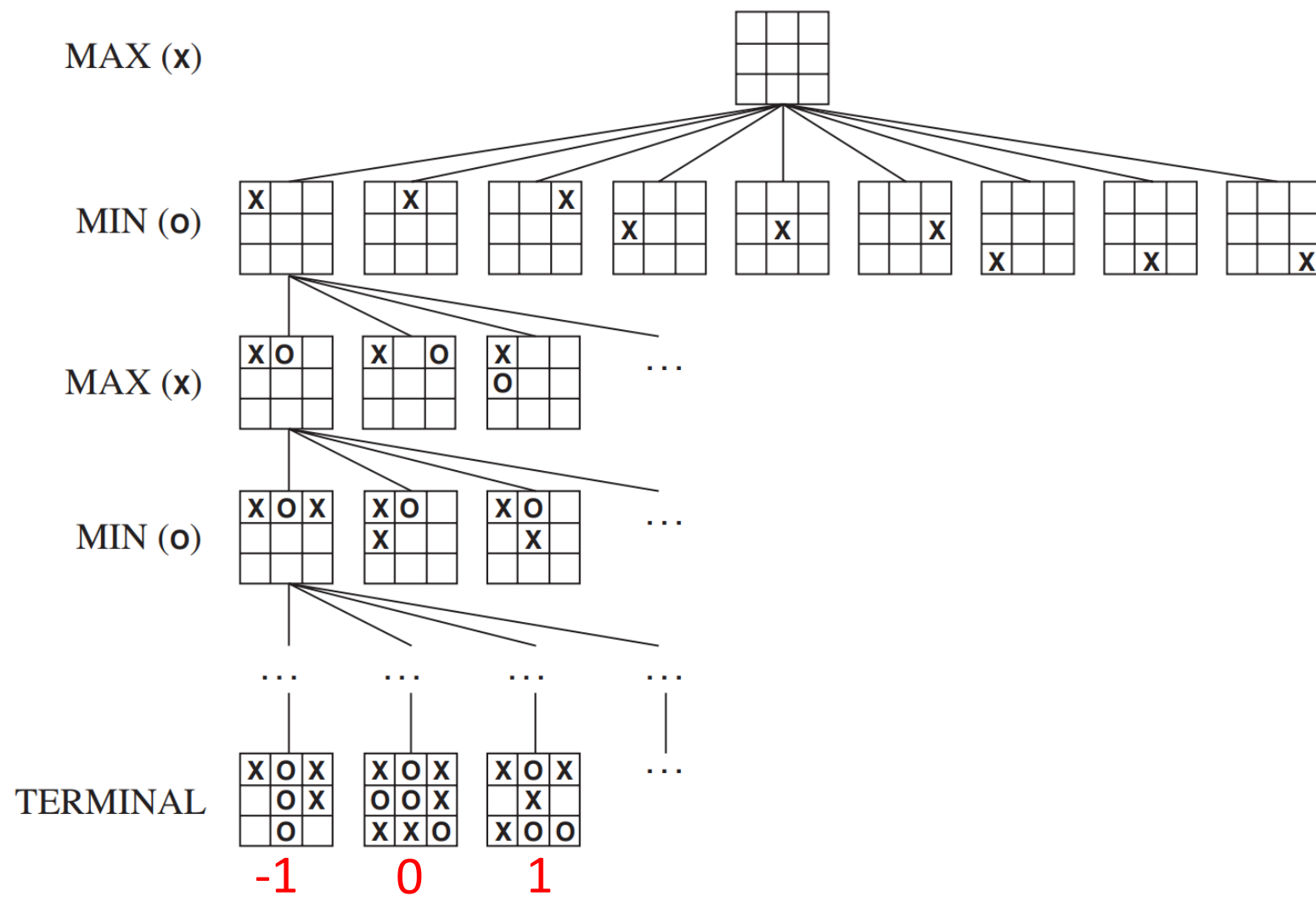
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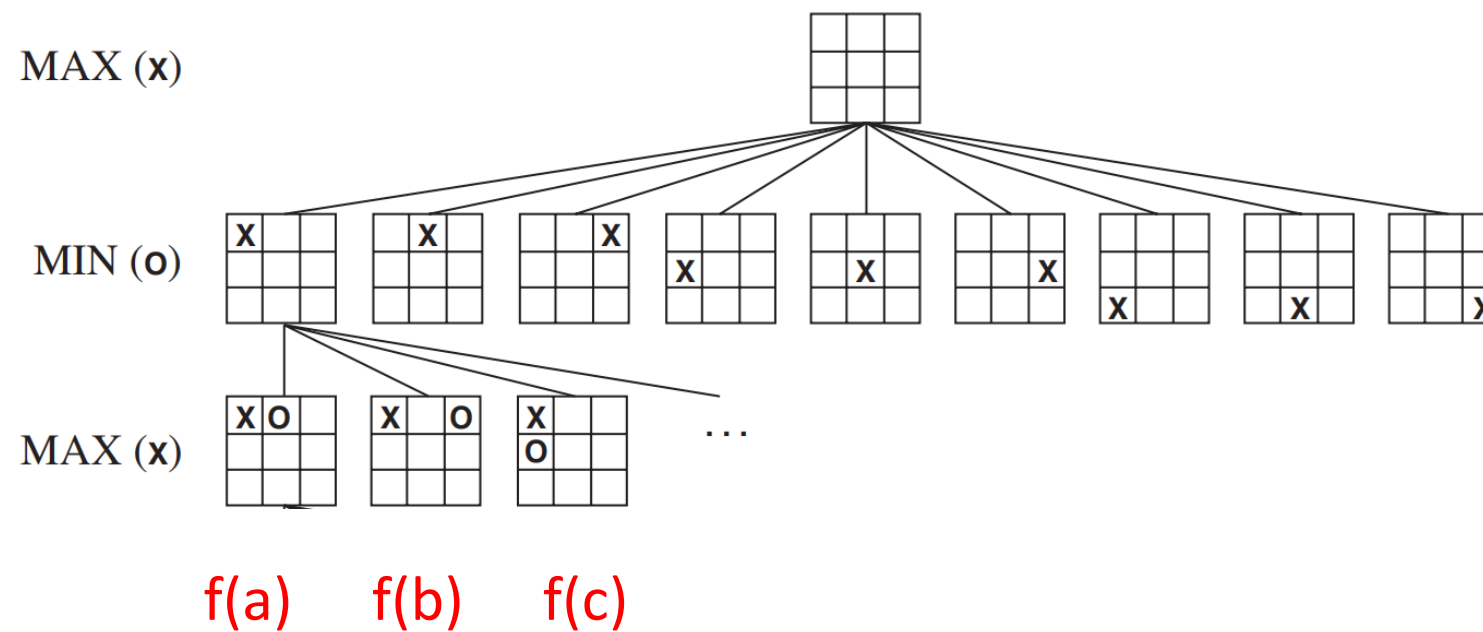
Tic Tac Toe



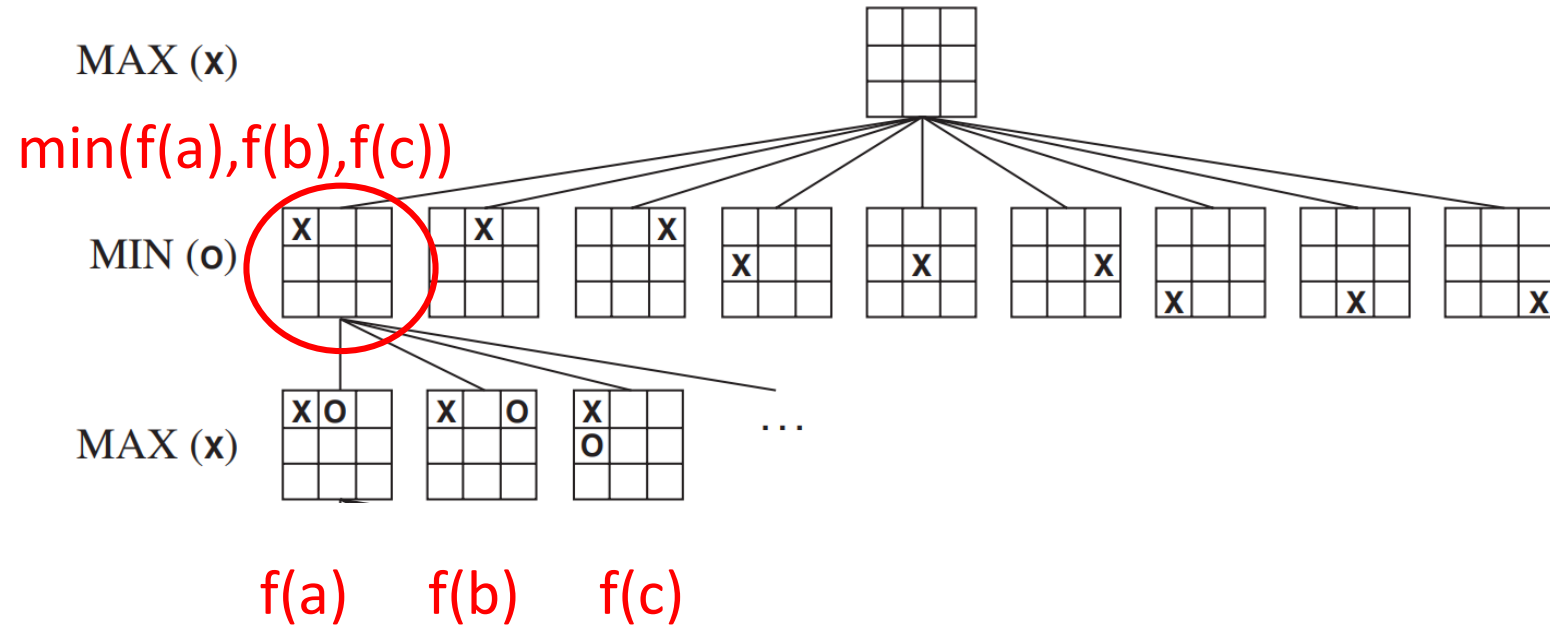
Tic Tac Toe



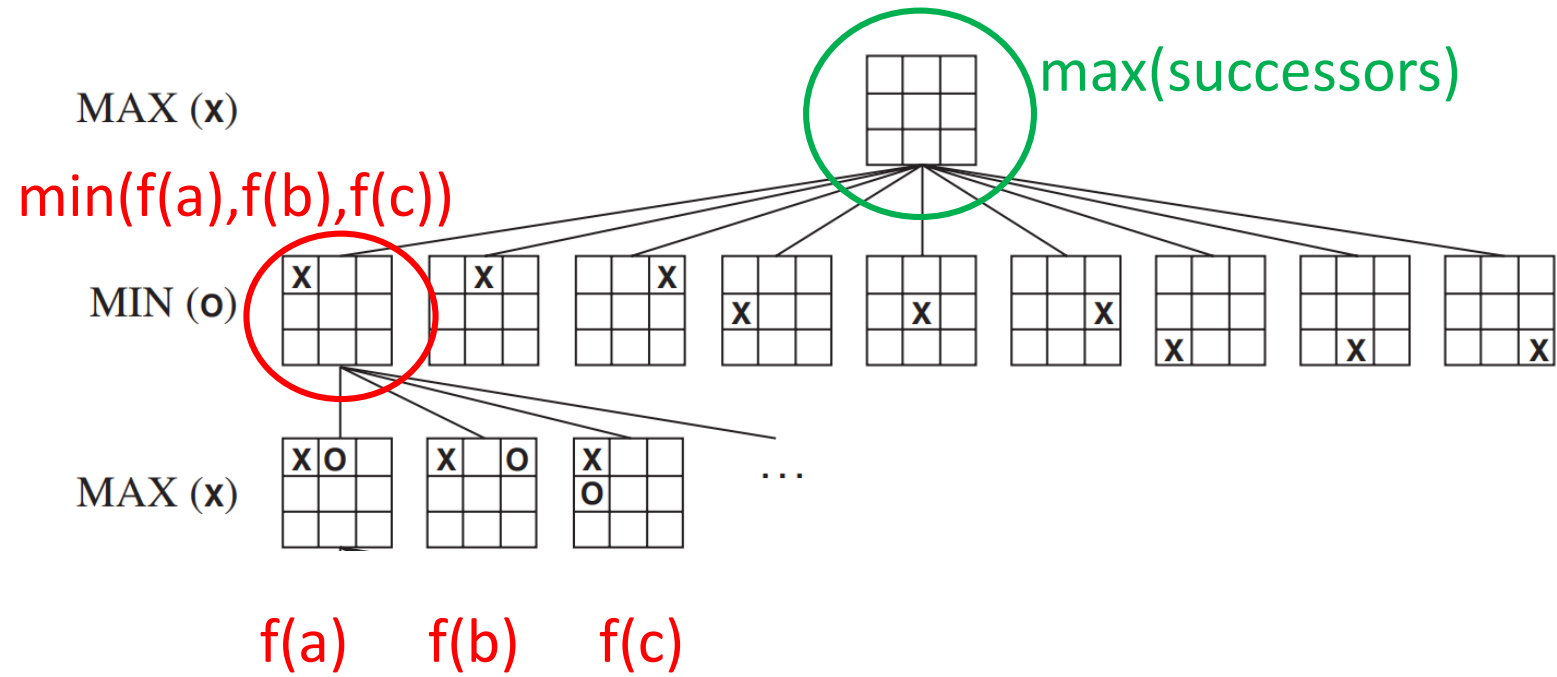
Tic Tac Toe



Tic Tac Toe



Tic Tac Toe



Adversarial Search

(Initial) assumptions:

- On my turn I consider all moves and pick the one that's best for me
- My opponent will consider all moves and the pick the one that's worst for me

$$f_{\text{opponent}}(s) = -f_{\text{me}}(s)$$

Adversarial Search

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We'll use $f(s)$ whenever we're talking about $f_{\text{me}}(s)$, and
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Adversarial Search

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- On my turn I consider all moves and pick the one that's best for me
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We'll use $f(s)$ whenever we're talking about $f_{\text{me}}(s)$, and

$-f(s)$ whenever we're talking about $f_{\text{opponent}}(s)$

My turn: $\max f(s)$ for all successors s (pick my best move)

Opponent's turn: $\min f(s)$ for all successors s (pick my worst move)

Minimax Value of a Game

“Ground truth”

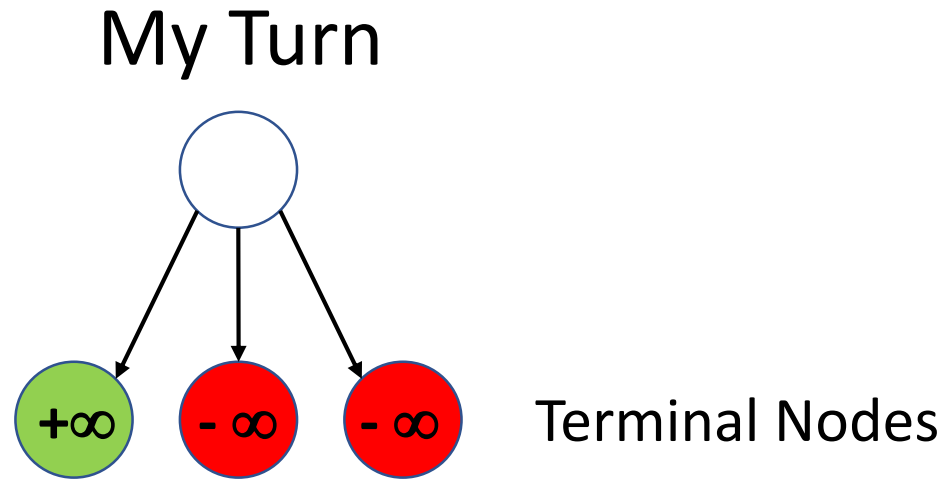
$V(s)$: What the true value of a game state is

Given for terminal nodes

Computed from non-terminal nodes

Minimax Value of a Game

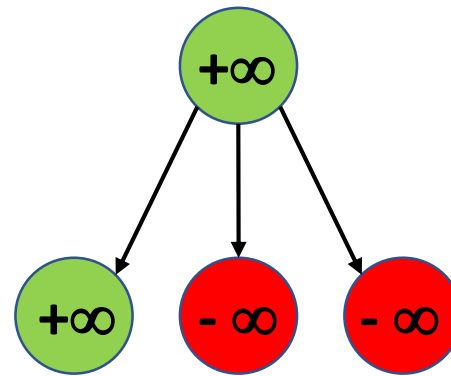
- I win = $+\infty$
- I lose = $-\infty$



Minimax Value of a Game

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My Turn

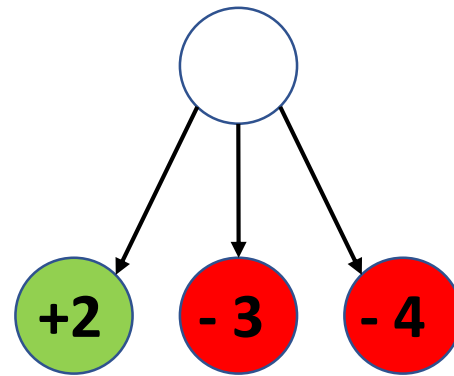


Terminal Nodes

Minimax Value of a Game

- $V(s)$ = value of win (to me)

My Turn

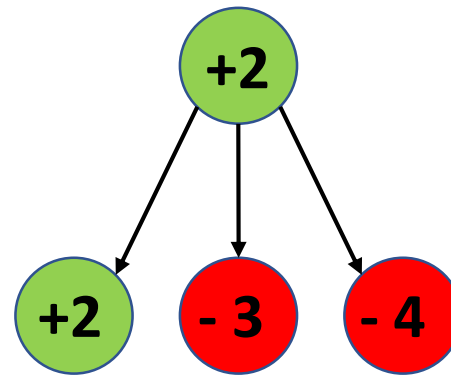


Terminal Nodes

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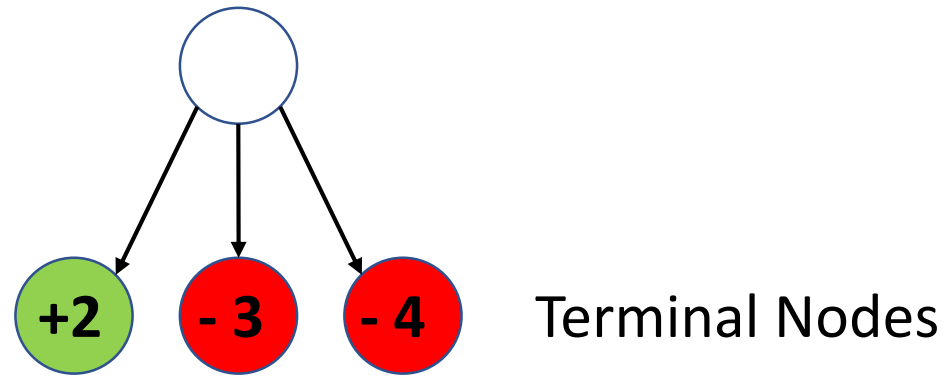


Terminal Nodes

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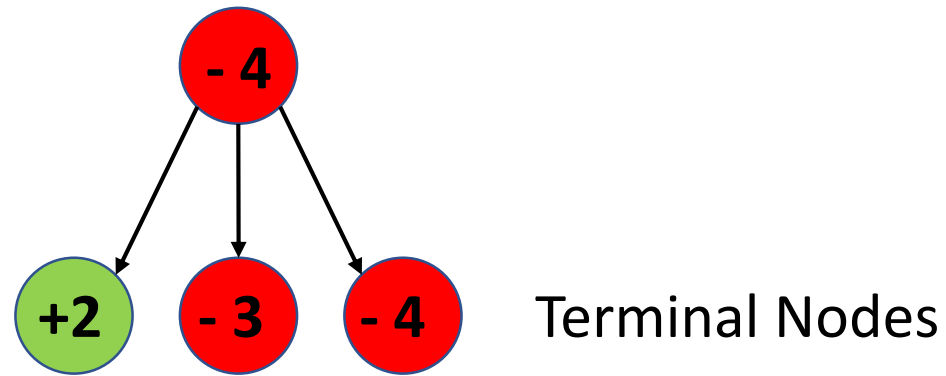
My Opponent's Turn



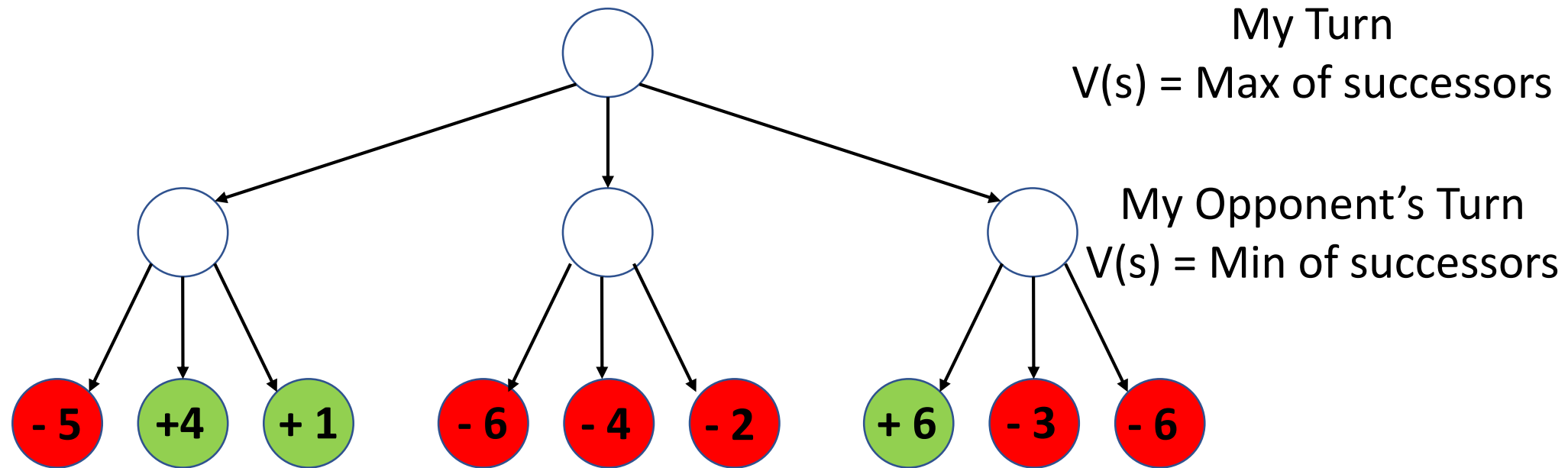
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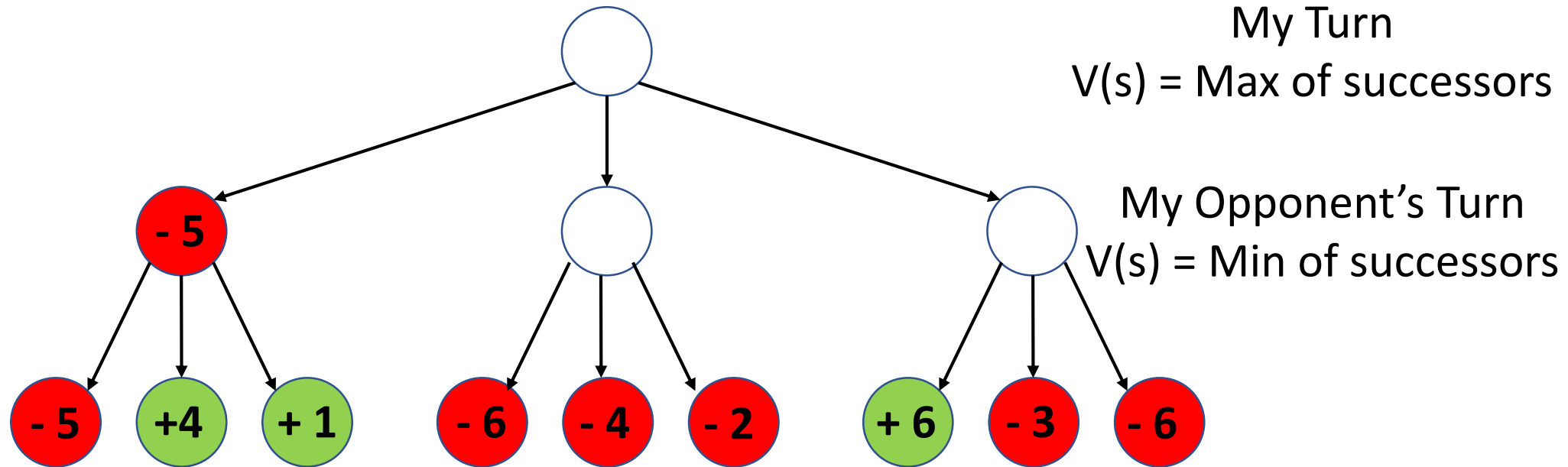
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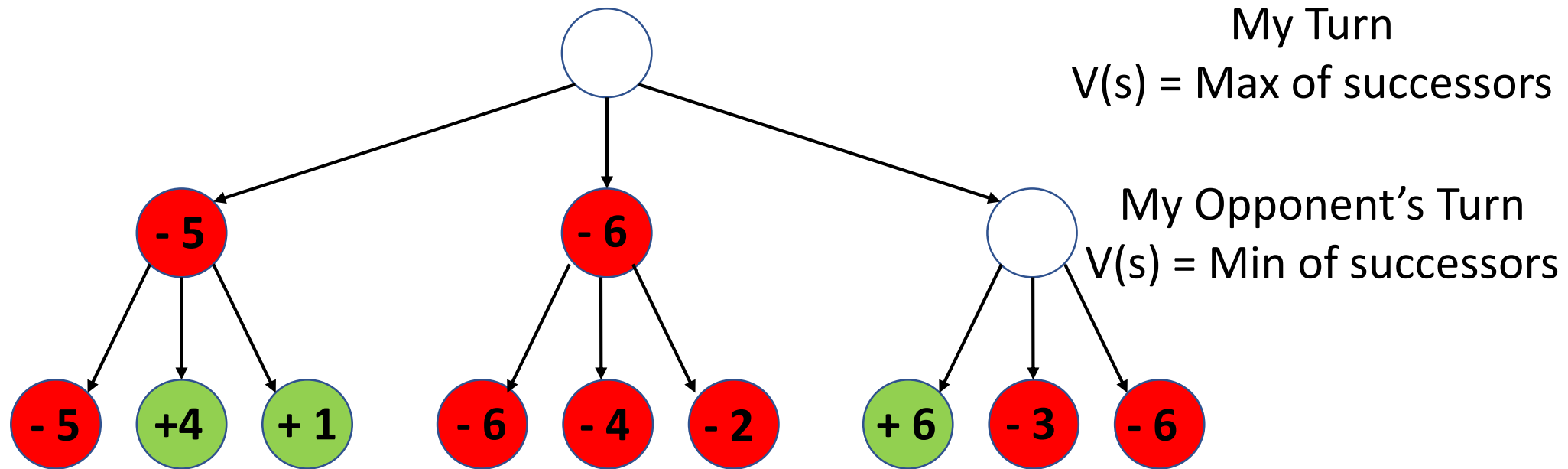
Minimax Value of a Game



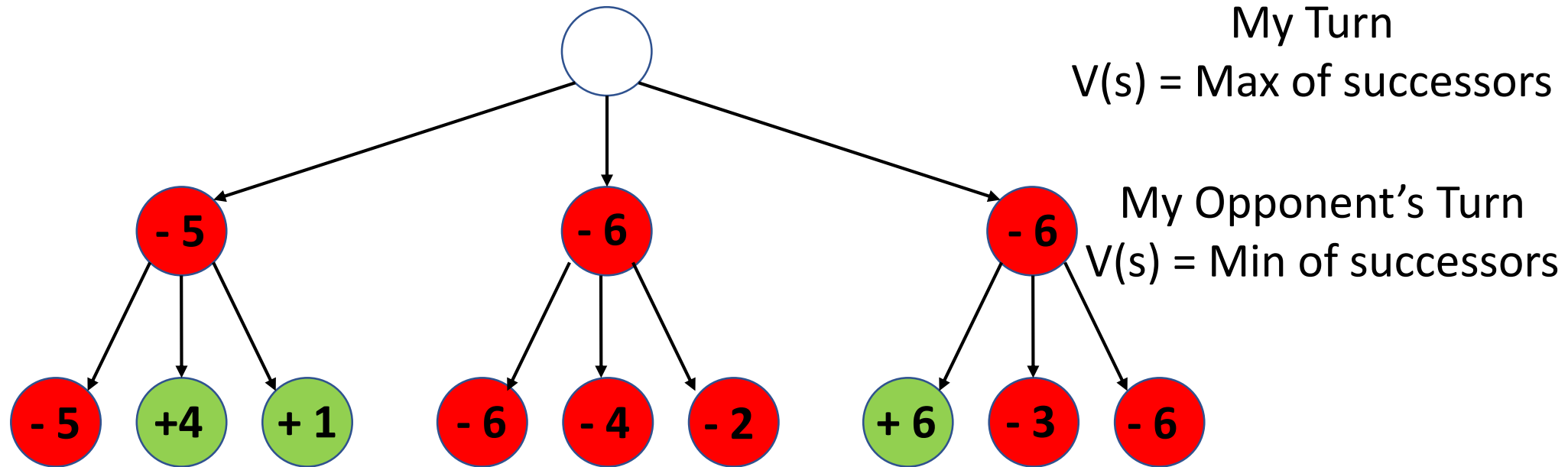
Minimax Value of a Game



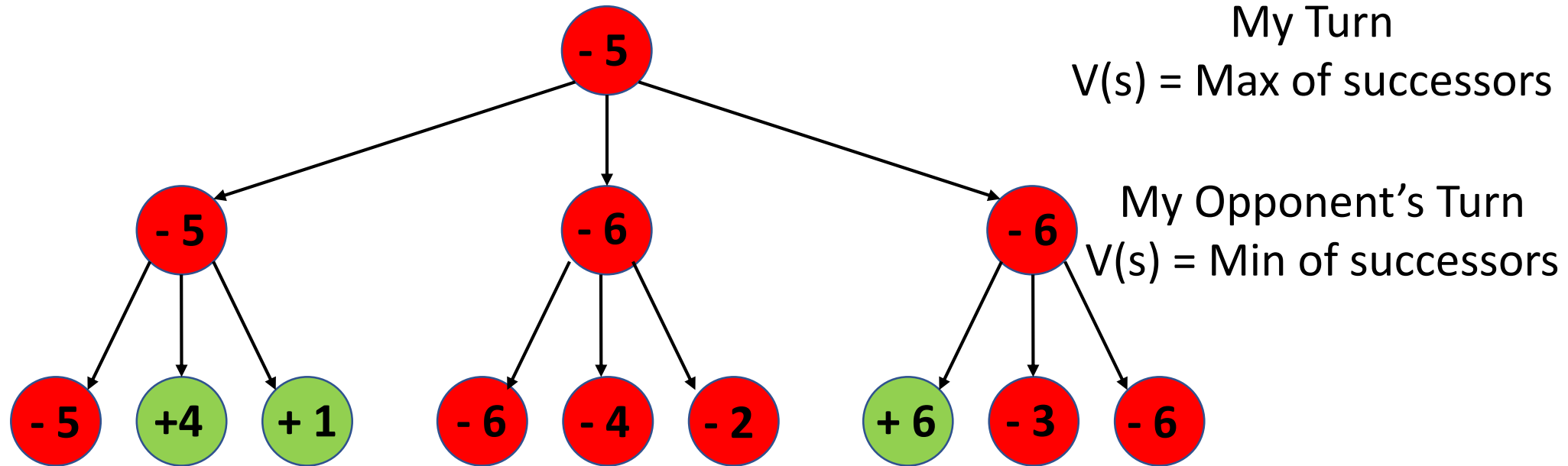
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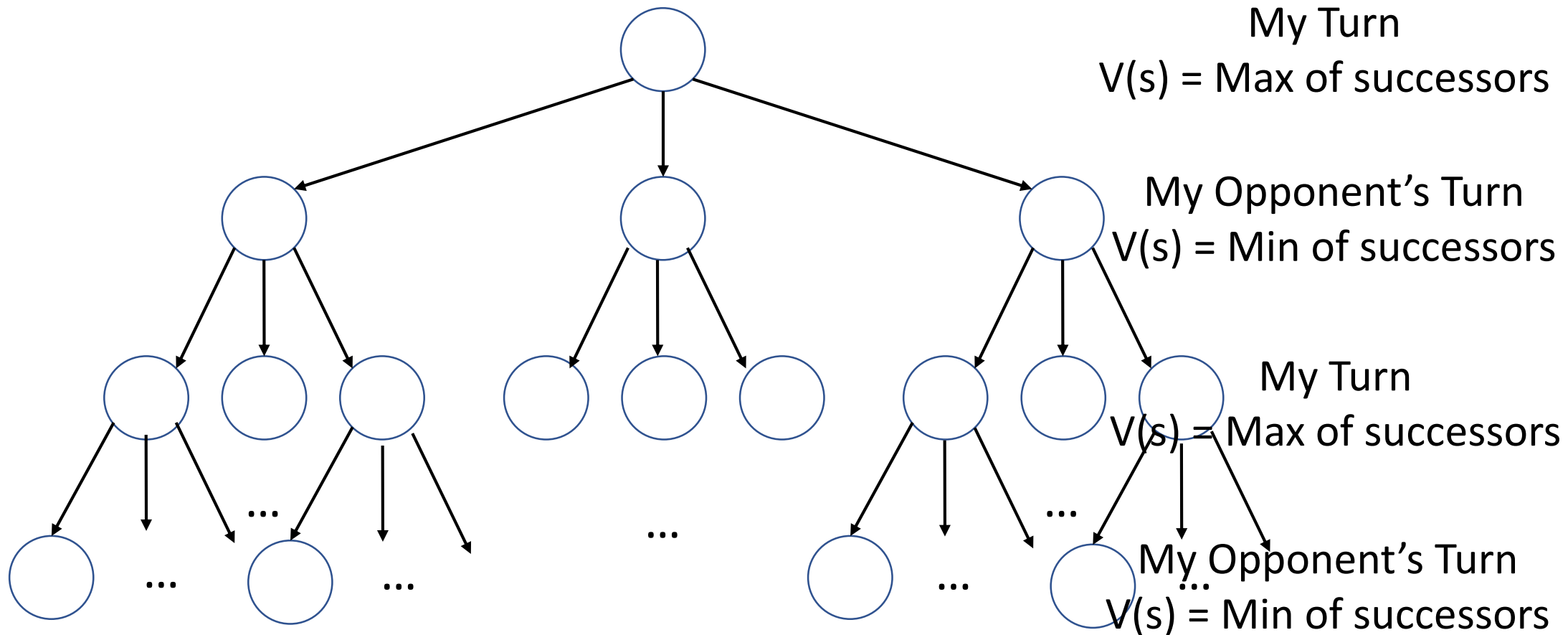
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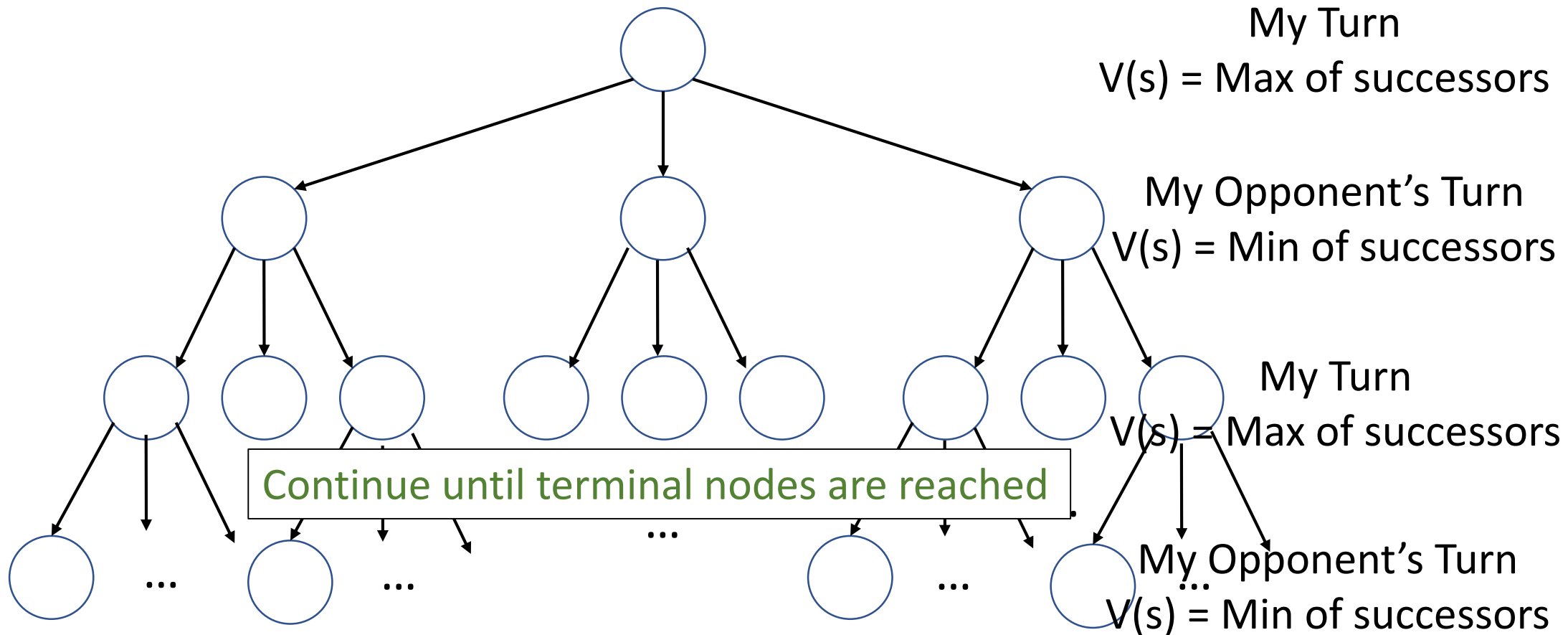
Minimax Value of a Game

- Current state: s
- Available operators: ops
- Value of a terminal state: $V(s)$
- My turn:
 - Value of state $s = V(s) = \max_{o \in ops} \{V(\text{apply}(s, o))\}$
 - Best move = $\operatorname{argmax}_{o \in ops} \{V(\text{apply}(s, o))\}$
- Opponent's turn:
 - Value of state $s = V(s) = \min_{o \in ops} \{V(\text{apply}(s, o))\}$
 - Best move = $\operatorname{argmin}_{o \in ops} \{V(\text{apply}(s, o))\}$

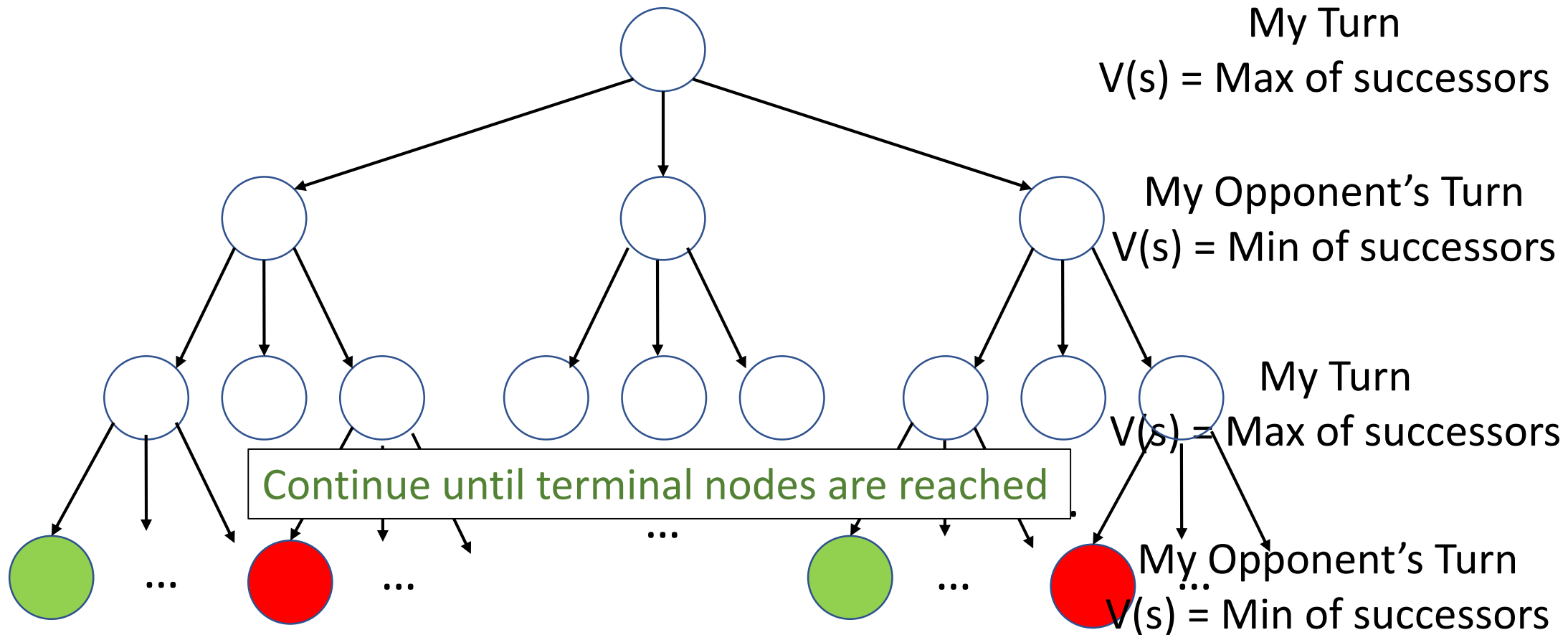
Minimax Value of a Game



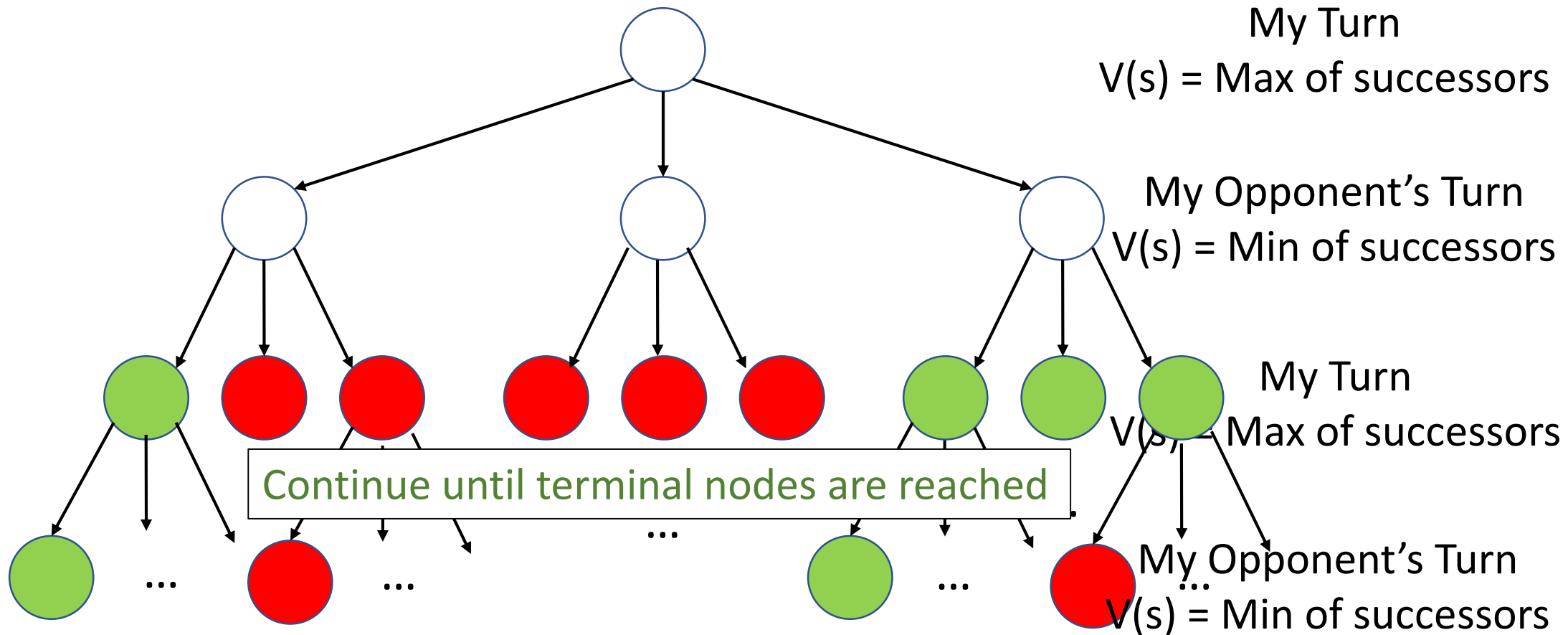
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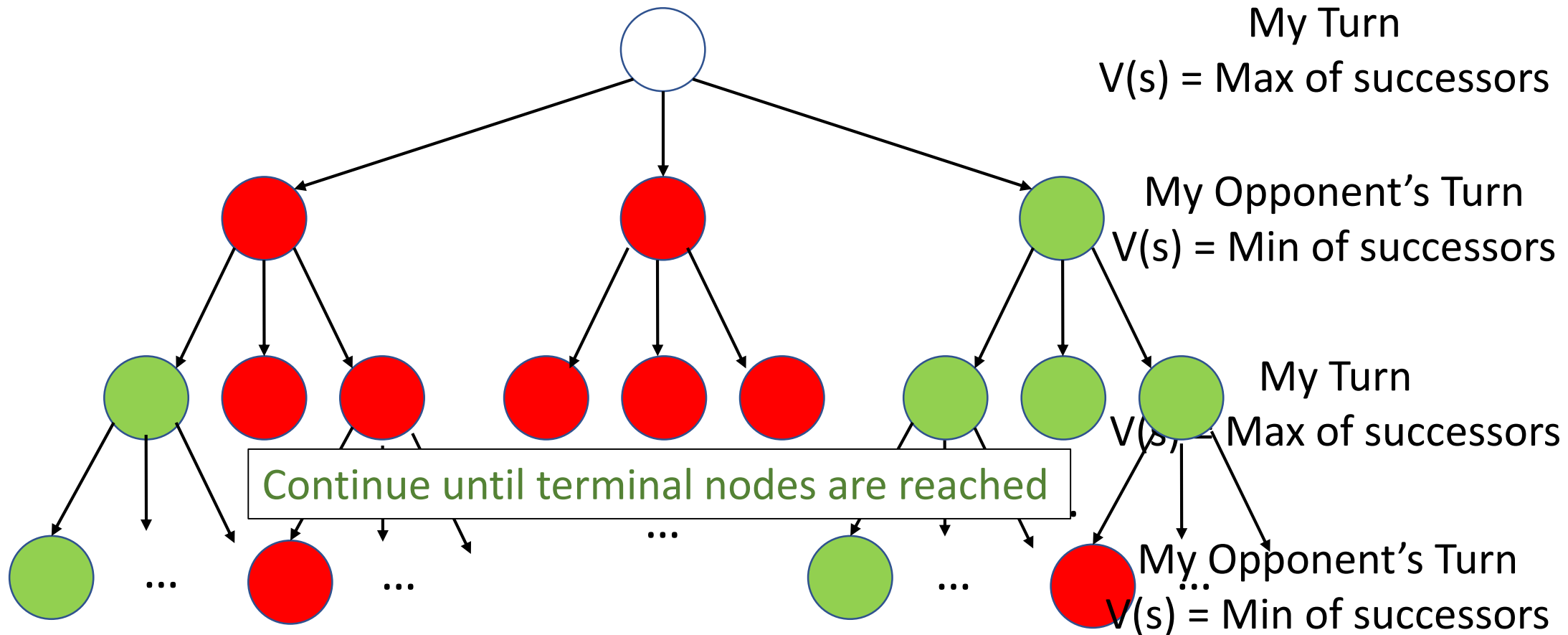
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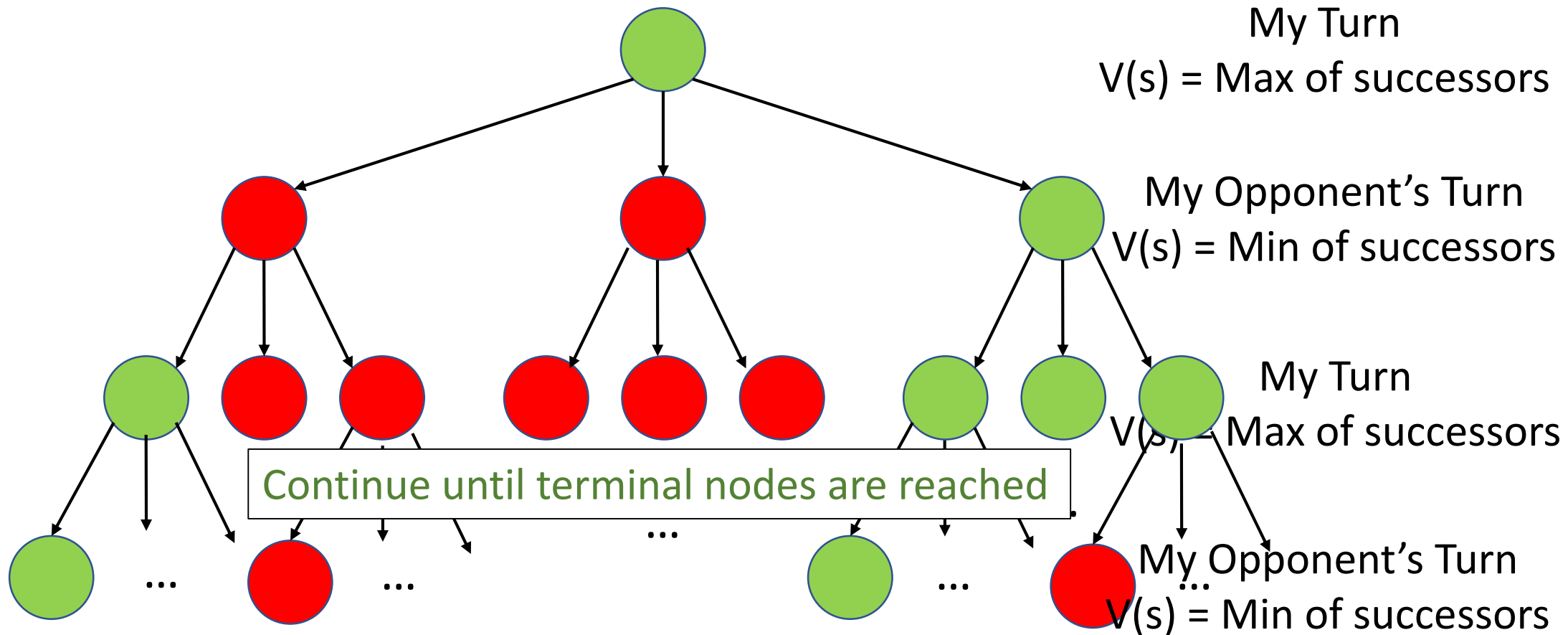
Minimax Value of a Game



Minimax Value of a Game



Minimax Value of a Game



Minimax Algorithm

Initial call:

- My turn: `maxturn(initial-state,ops)`
- Opponent's turn: `minturn(initial-state,ops)`

Minimax Algorithm

maxturn(s,ops):

 if terminal(s) then return $V(s)$

 else

$val \leftarrow -\infty$;

 foreach $o \in ops$

$val' \leftarrow \text{minturn}(\text{apply}(s,o),ops)$;

 if $val' > val$ then $val \leftarrow val'$; bestop $\leftarrow o$;

 return val

Minimax Algorithm

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maxturn(s,ops):  
  if terminal(s) then return V(s)  
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  return val
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Shorthand for return a data structure with the value,
the operator that took you here, possibly other info

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minturn(s,ops):

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Minimax Algorithm (Complete Search)

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Minimax Search

- Complete search:
Generally intractable to go all the way to terminal nodes
- Key idea:
Use a heuristic evaluation function $f(s)$ that applies to intermediate states and returns a number that estimates the value of s

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Initial call:

- If I go first: maxturn(initial-state,ops)
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Minimax Algorithm (Heuristic Search)

maxturn(s,ops):

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if cutoff(s) then return f(s)
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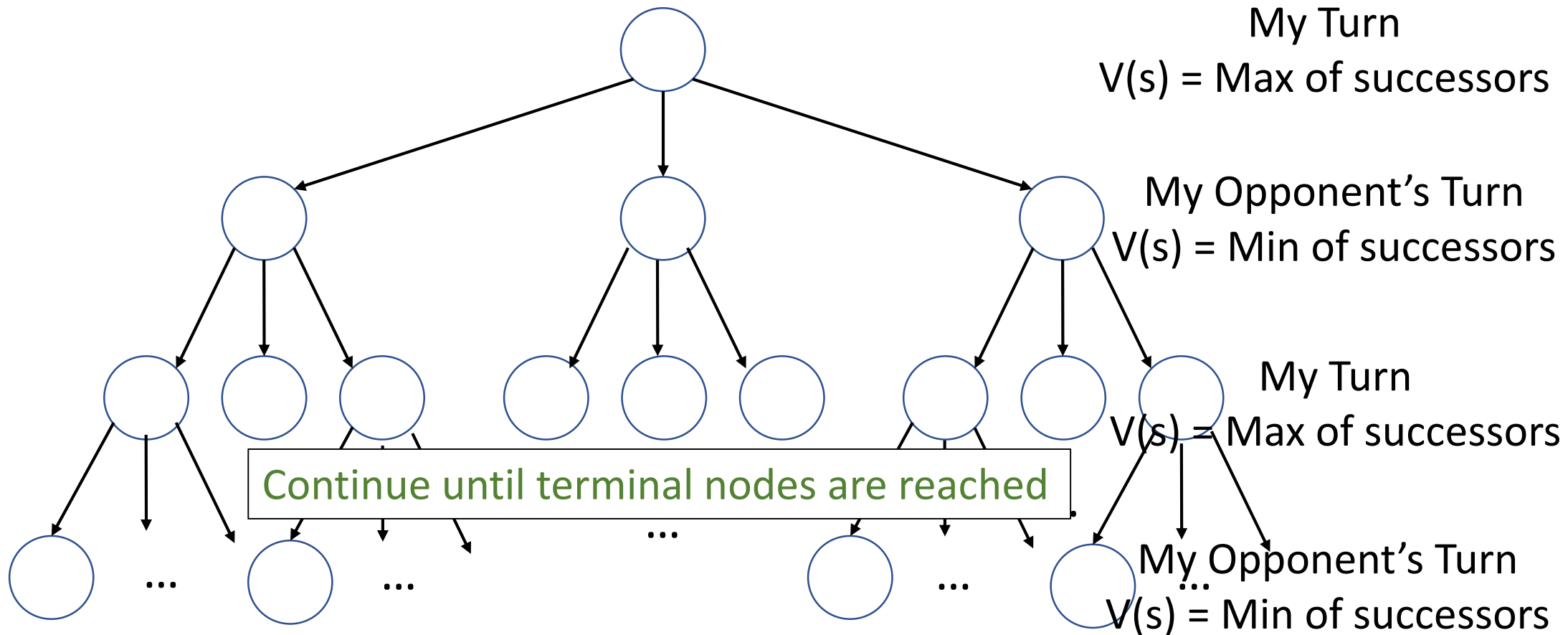
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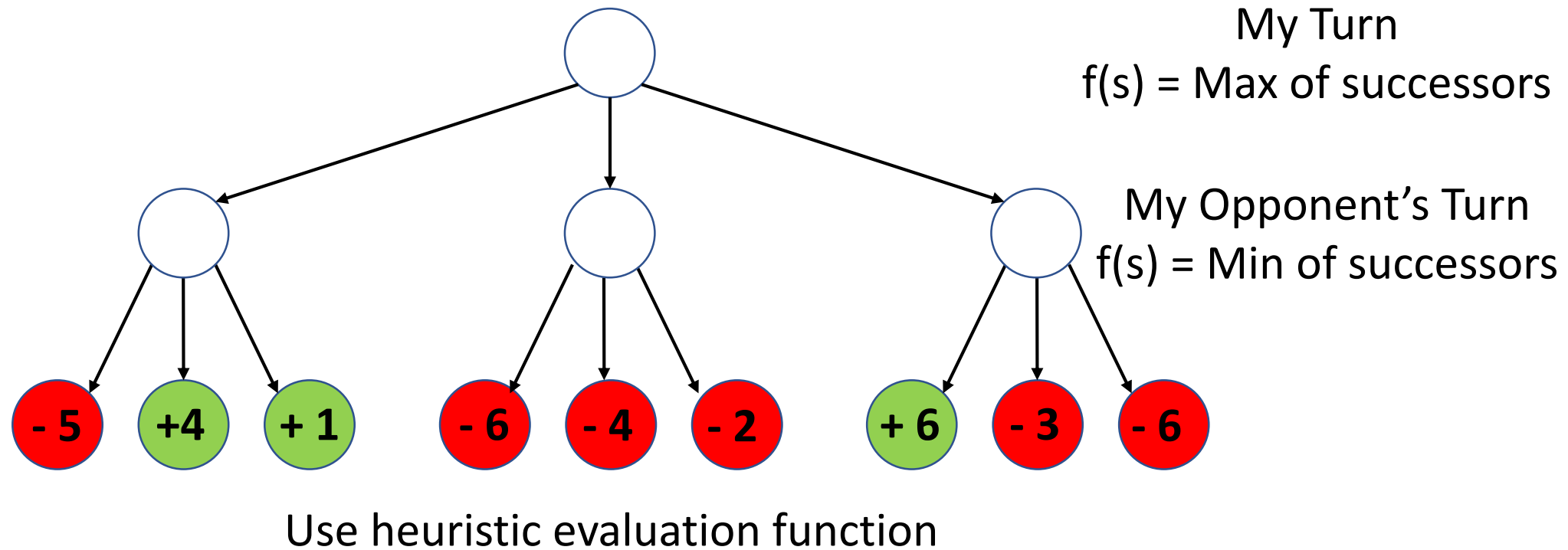
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cutoff(s) could be a depth-bound
or something more sophisticated

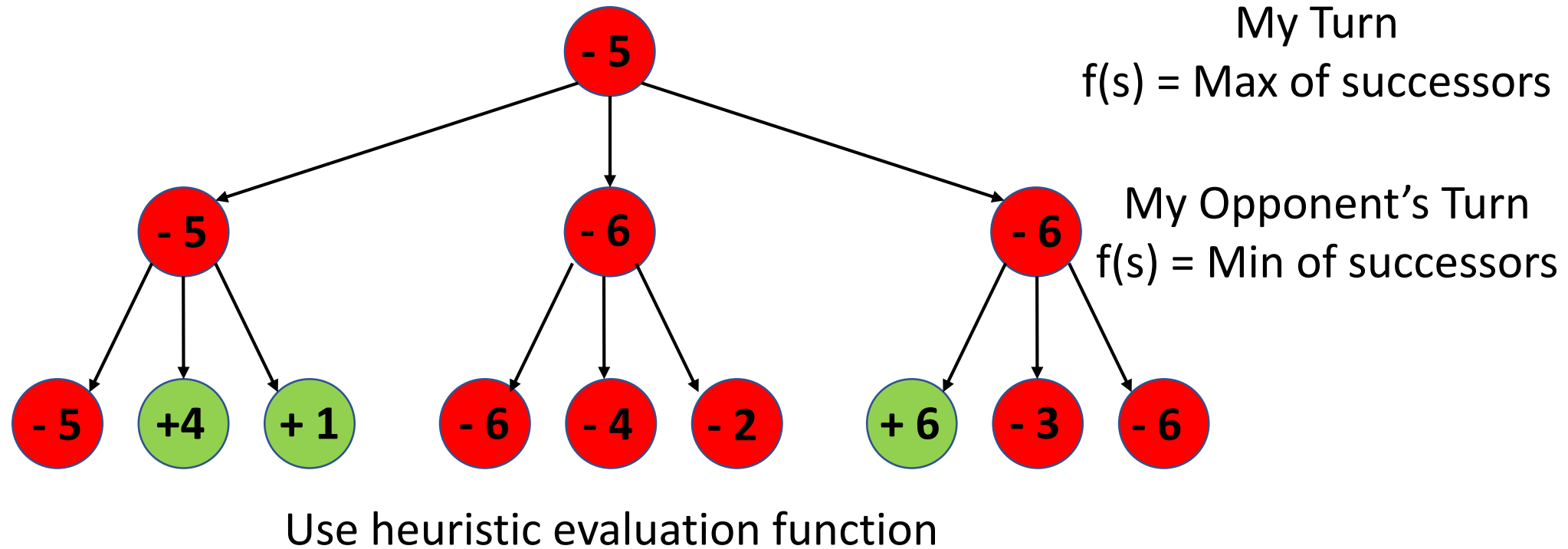
Minimax Value of a Game



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Minimax Value of a Game



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Alpha-Beta Pruning

New idea to improve efficiency:

Can prune branches that are guaranteed never to be used

Analogy

Given $\text{AND}(p, q, r)$

If p is FALSE don't need to evaluate q or r

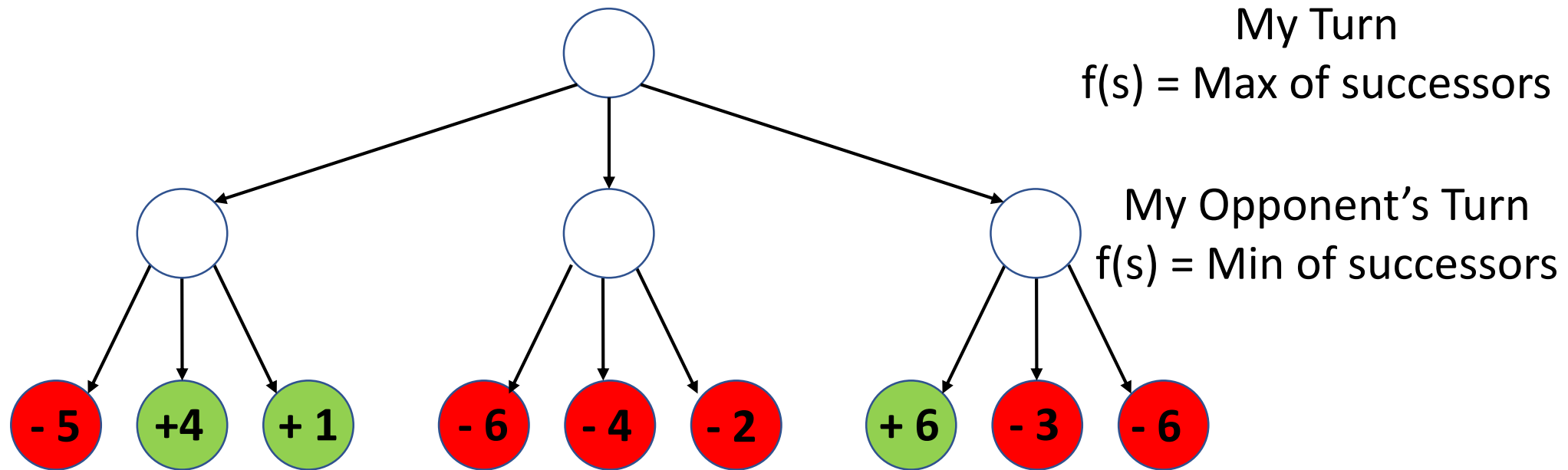
Given $\text{OR}(p, q, r)$

If q is TRUE don't need to evaluate q or r

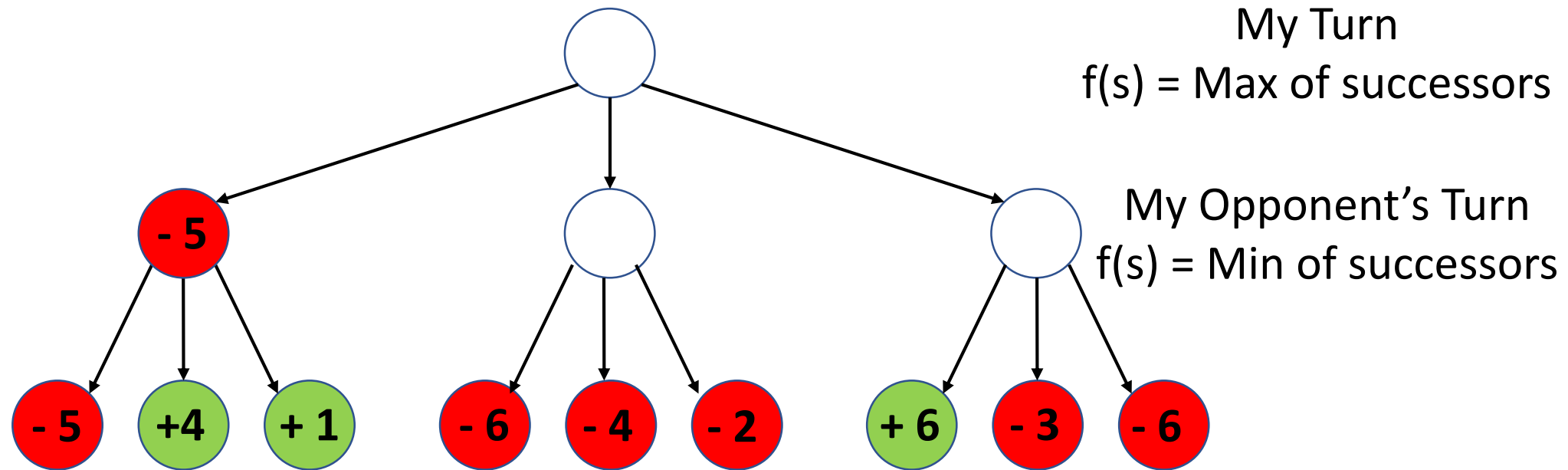
Alpha-Beta Pruning

Don't explore branches that will have no effect on the outcome

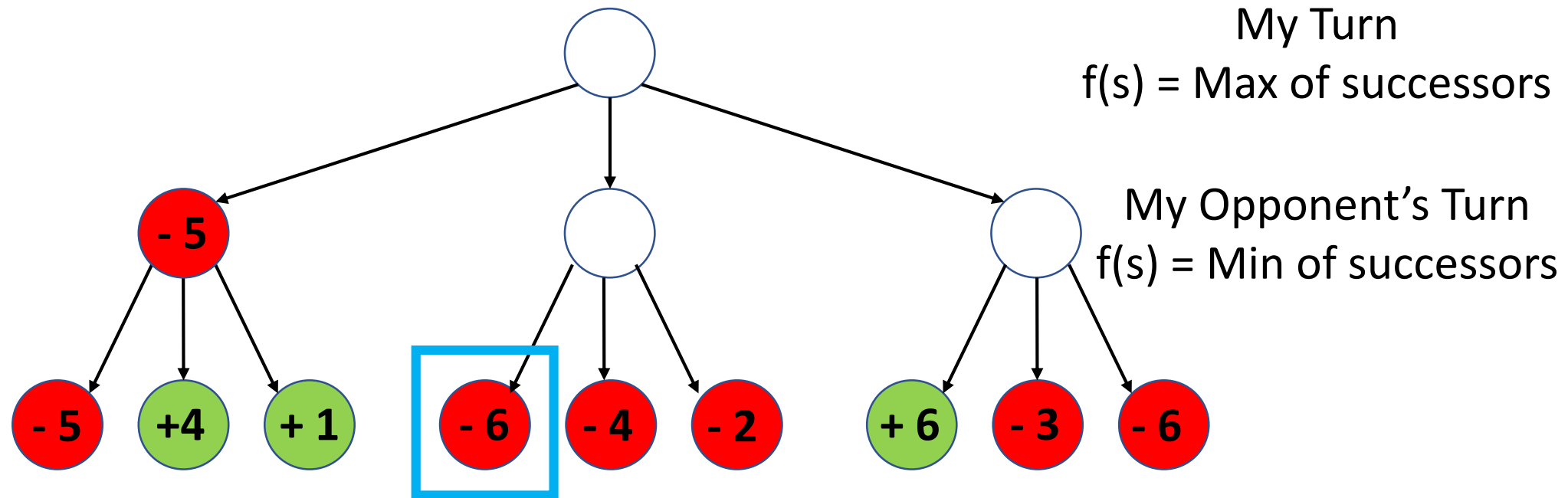
Alpha-Beta Pruning



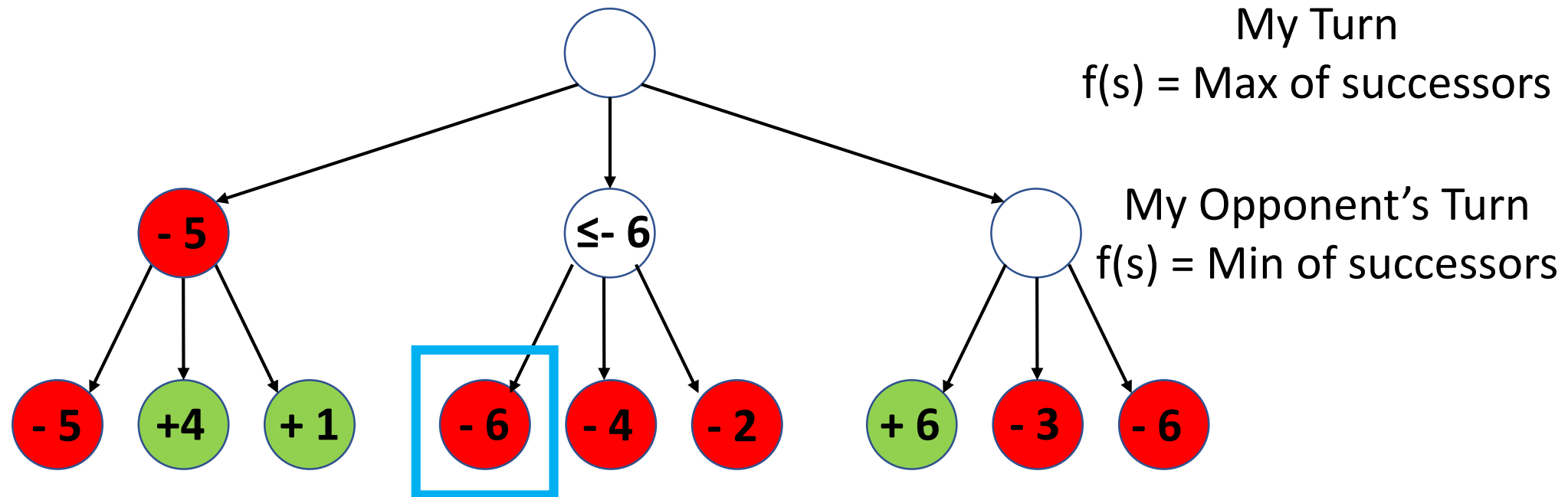
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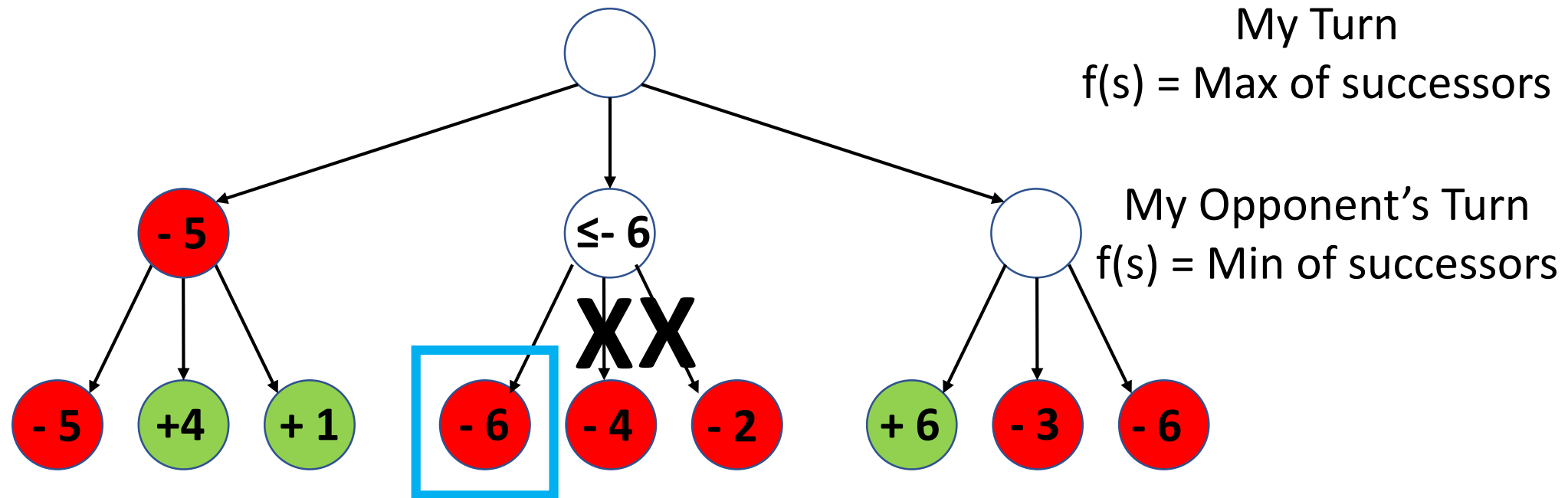
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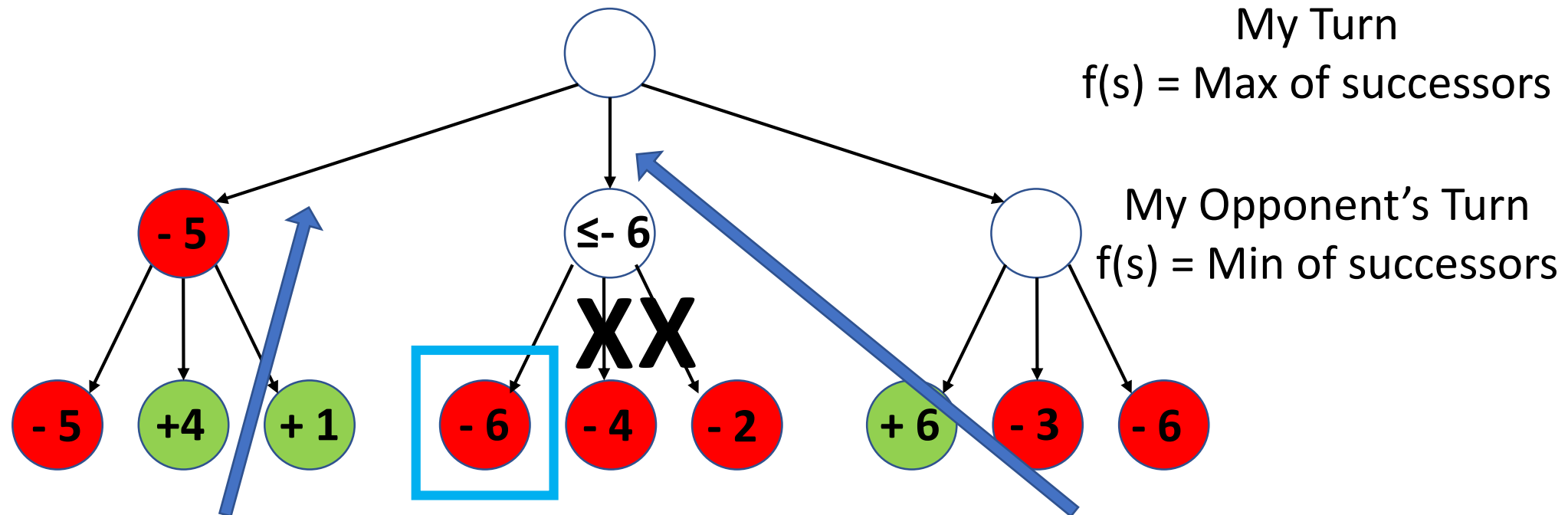
Alpha-Beta Pruning



Alpha-Beta Pruning

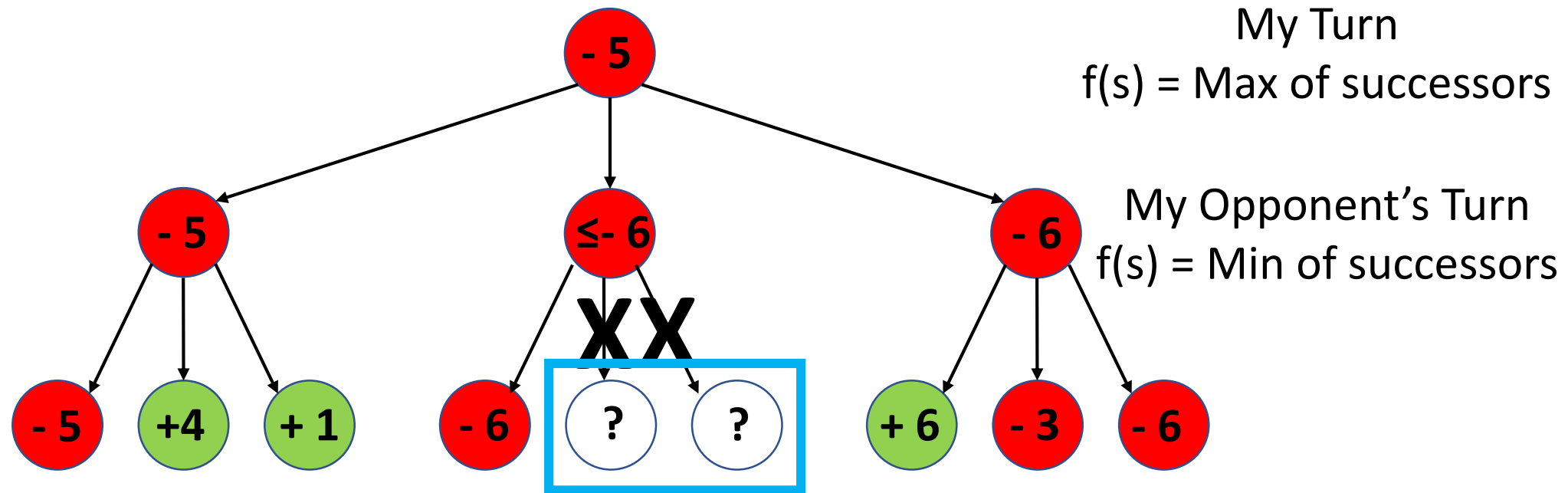


Alpha-Beta Pruning



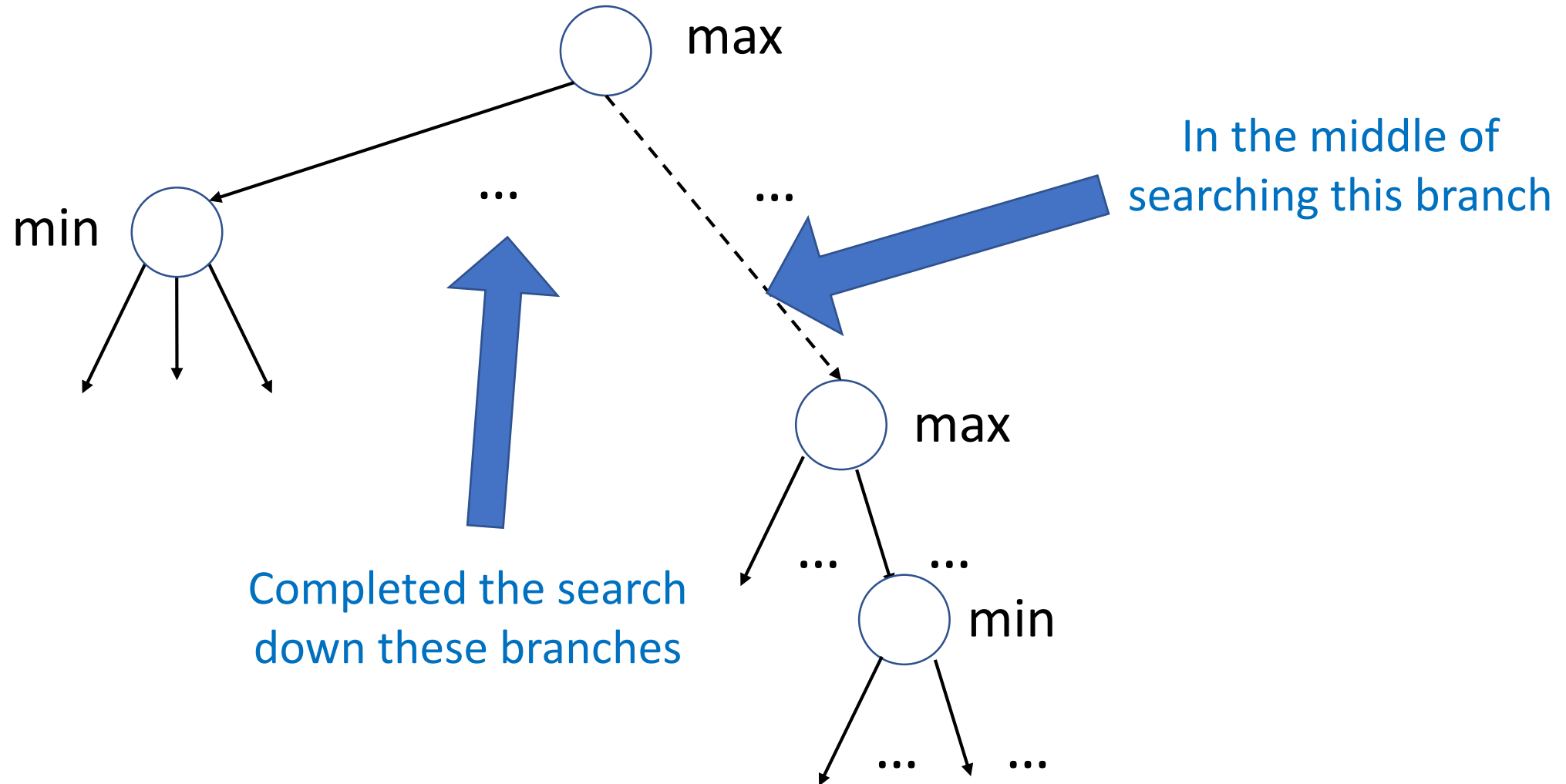
This is guaranteed to be a better move than this
regardless of what results from the other actions

Alpha-Beta Pruning

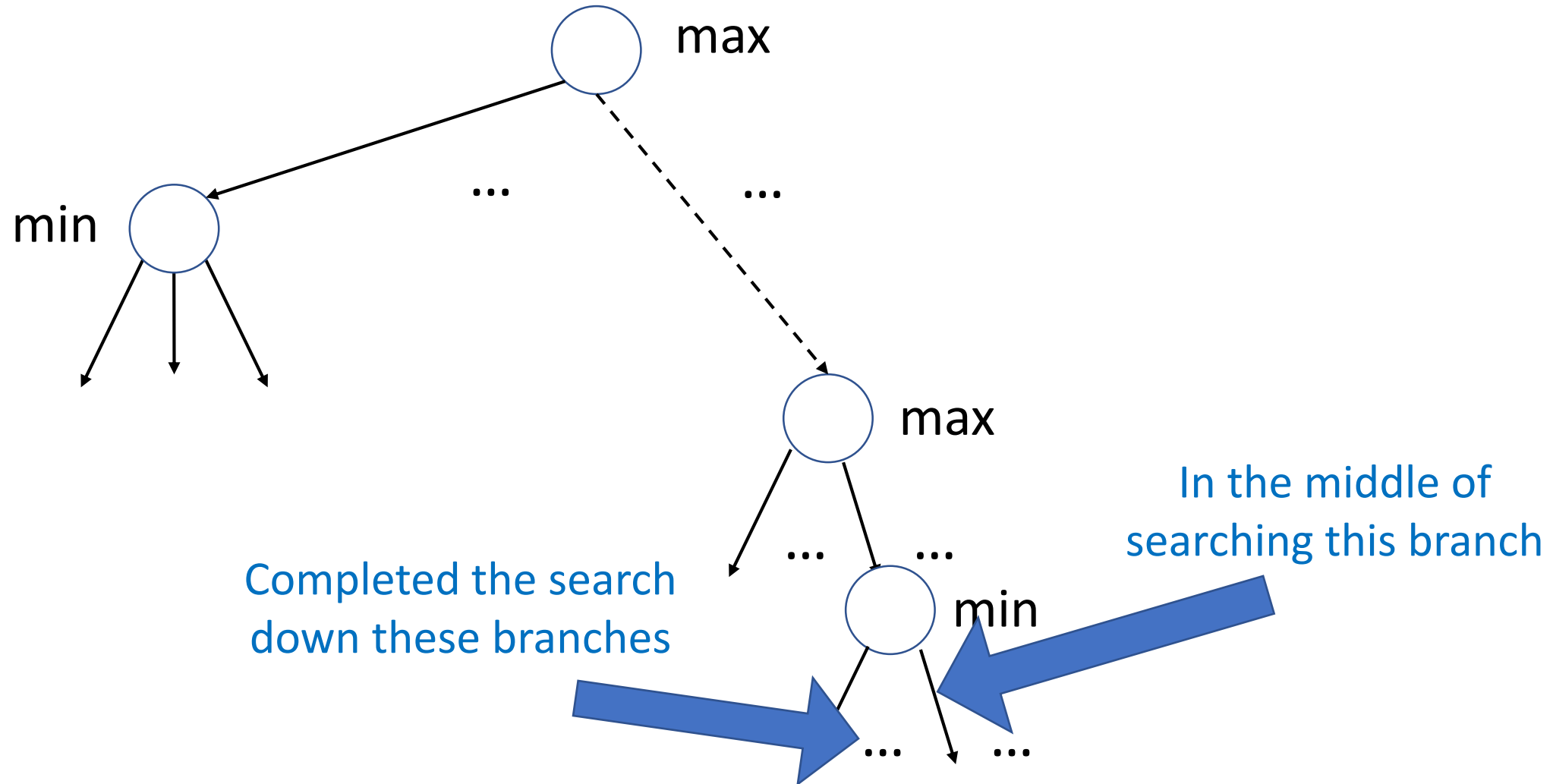


There is no way to set a value to either of these two states that would change the outcome

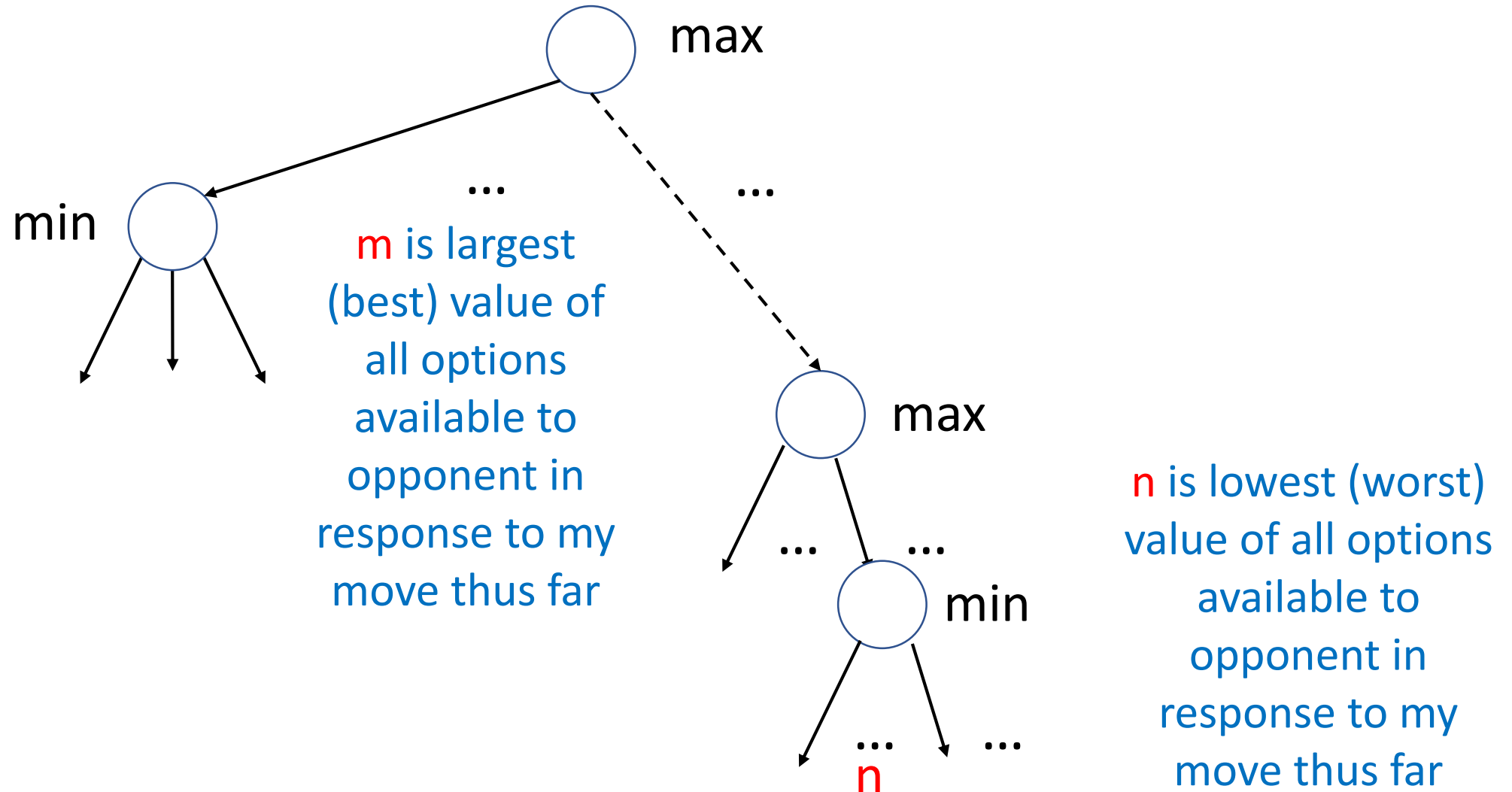
Alpha-Beta Pruning: General Case



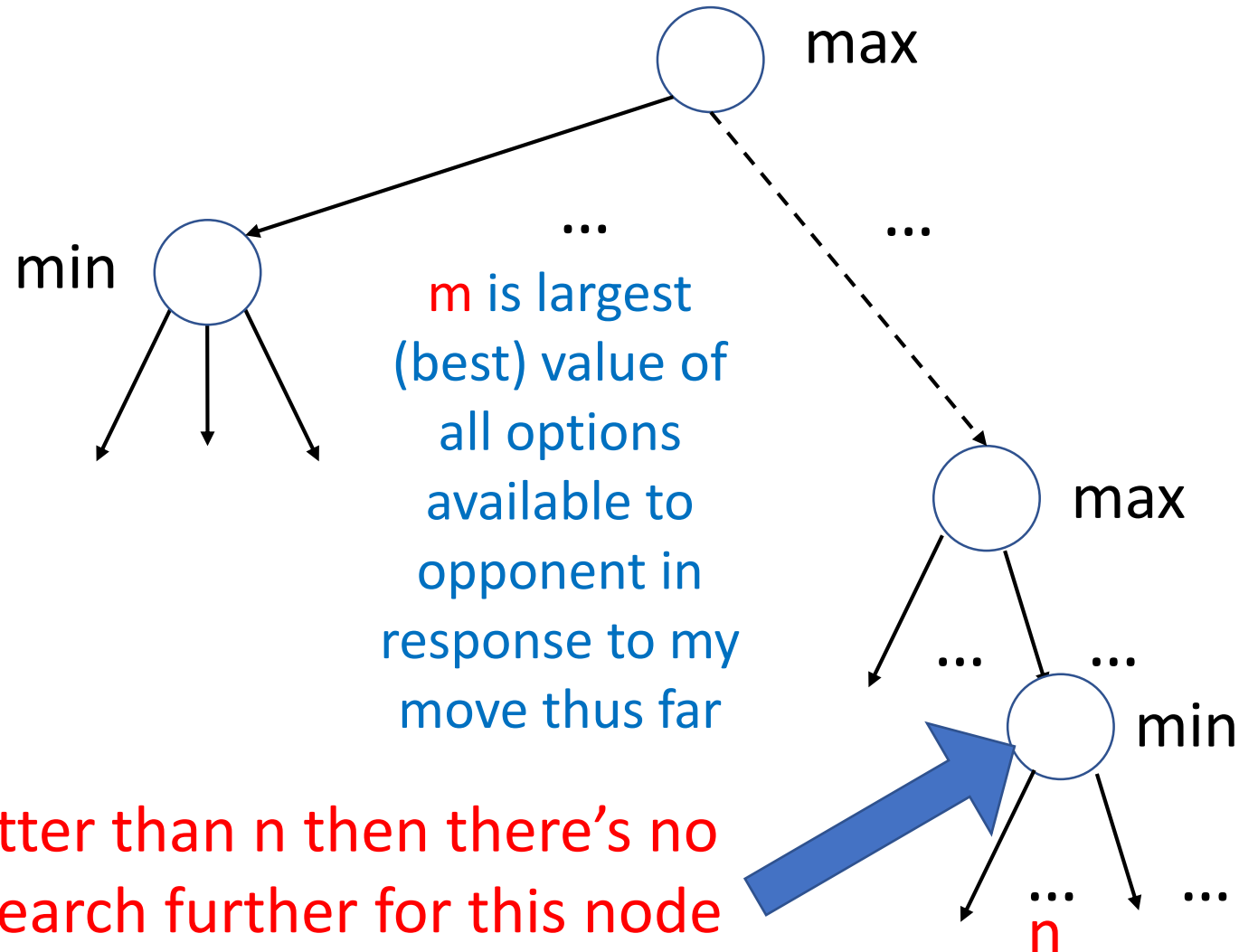
Alpha-Beta Pruning: General Case



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Alpha-Beta Pruning: General Case

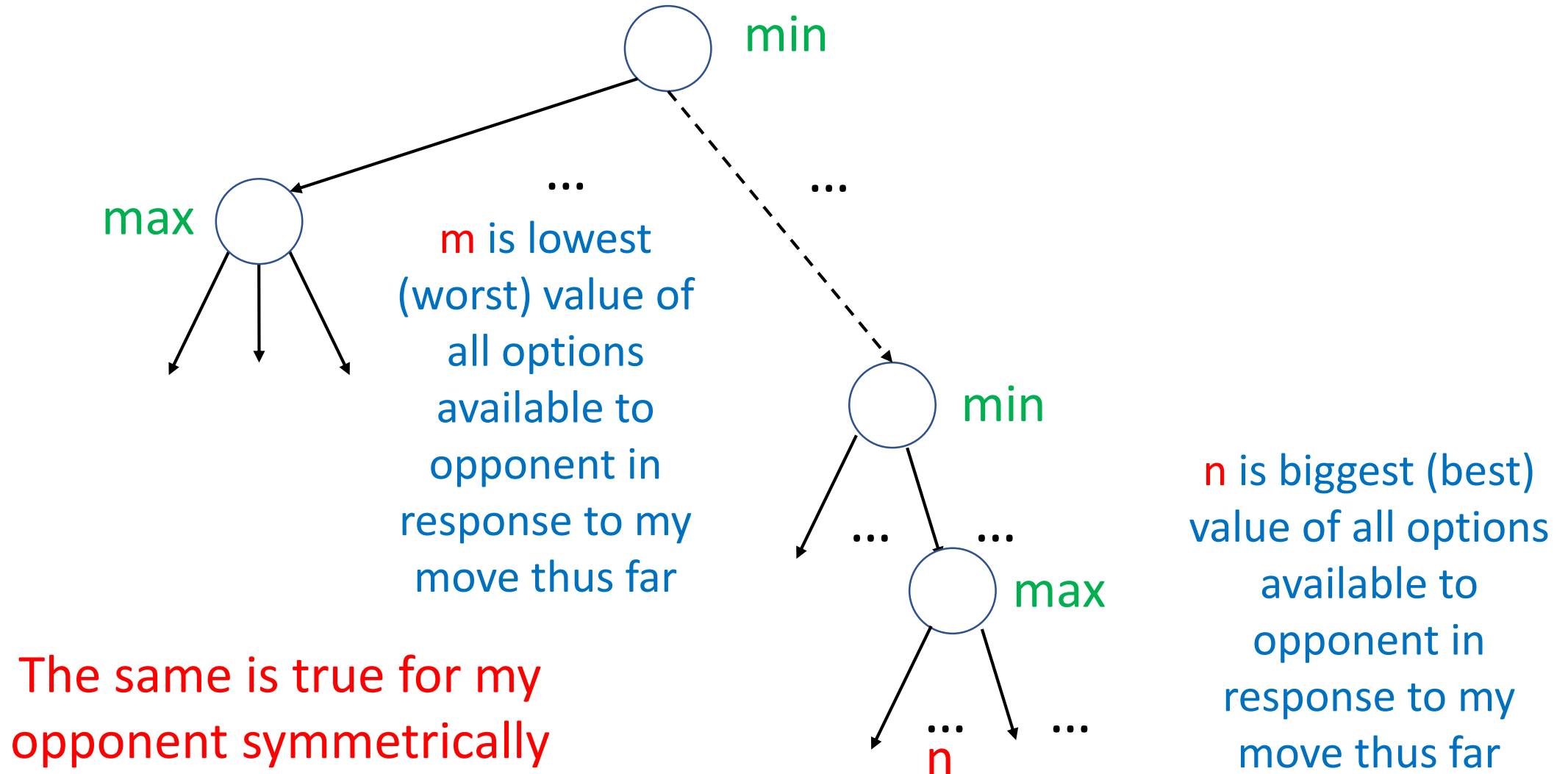


m is largest
(best) value of
all options
available to
opponent in
response to my
move thus far

n is lowest (worst)
value of all options
available to
opponent in
response to my
move thus far

If **m** is better than **n** then there's no
need to search further for this node

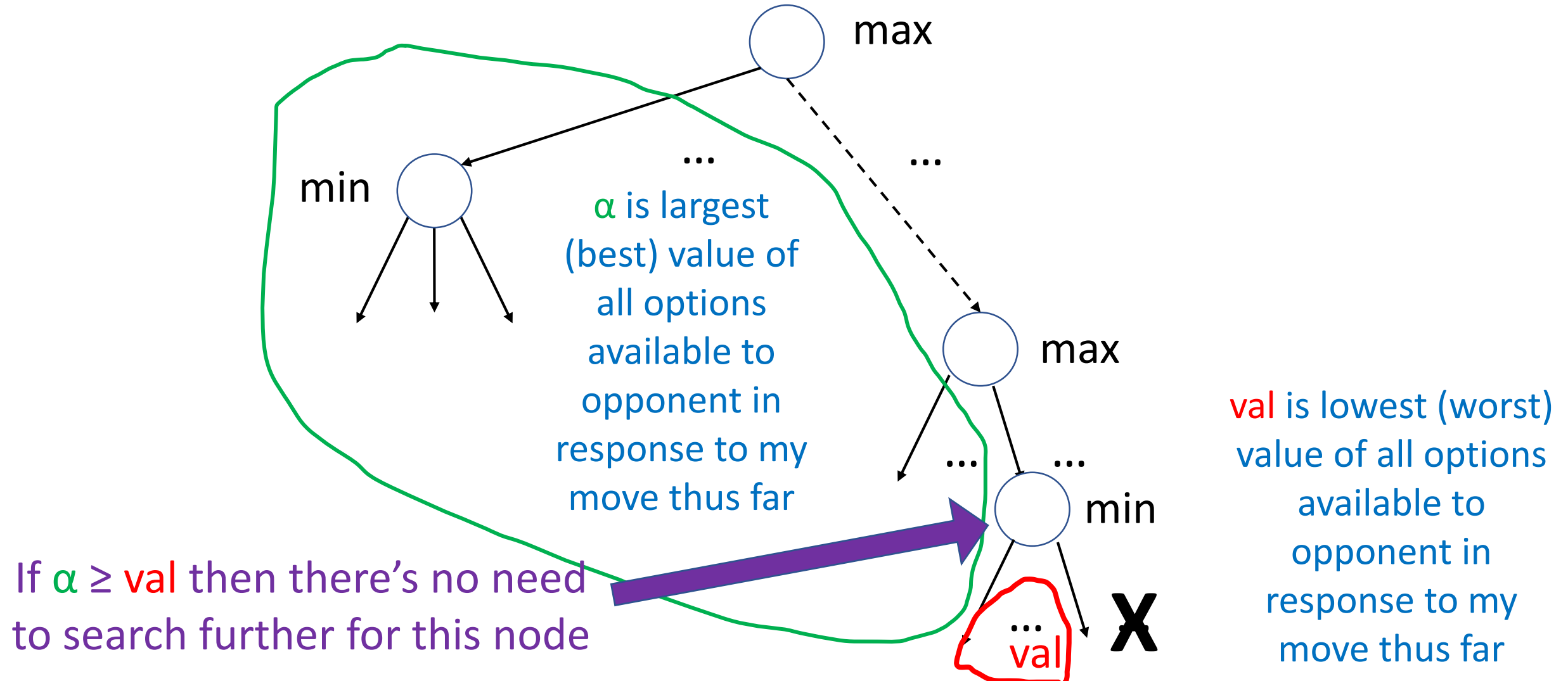
Alpha-Beta Pruning: General Case



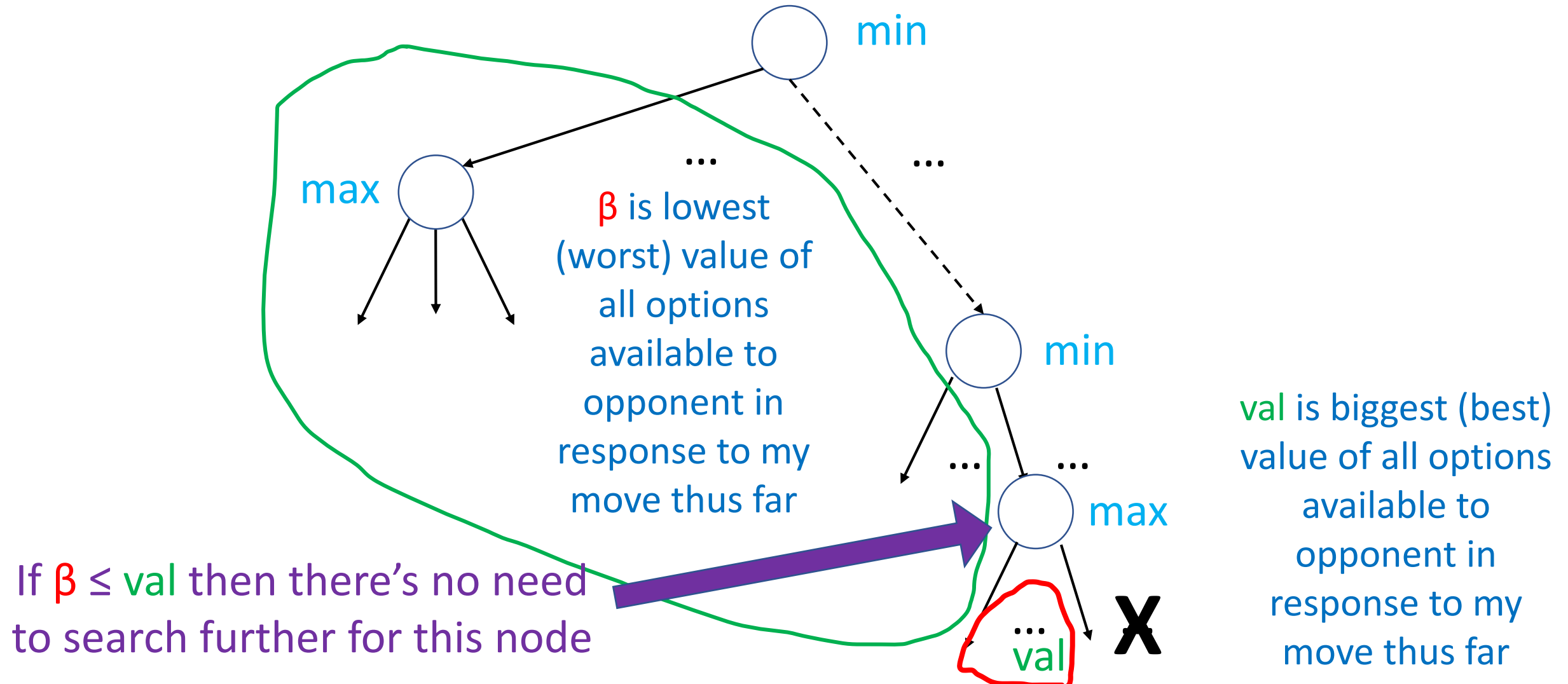
Alpha-Beta Pruning: Implementation

- Key idea:
 - At each node keep track of two parameters, α and β (I'll use a and b in the algorithm)
 - α is the best value I can guarantee achieving via some earlier move
 - β is the worst value my opponent can guarantee achieving via some earlier move
- Don't explore further nodes at this level if
 - You're at a min node and it's already guaranteed that this will be worse than some other move I can force (worst so far means $< \alpha$)
 - You're at a max node and it's already guaranteed that this will be better than some other move my opponent can force on me (best so far means $< \beta$)

Alpha-Beta Pruning: General Case



Alpha-Beta Pruning: General Case



Minimax Algorithm

maxturn(s,ops):

```
if cutoff(s) then return f(s)
else
  val  $\leftarrow -\infty$ ;
  foreach o  $\in$  ops
    val'  $\leftarrow$  minturn(apply(s,o),ops);
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maxturn(s,ops):

```
if cutoff(s) then return f(s)
else
  val  $\leftarrow +\infty$ ;
  foreach o  $\in$  ops
    val'  $\leftarrow$  minturn(apply(s,o),ops);
    if val' < val then
      val  $\leftarrow$  val';
      bestop  $\leftarrow$  o;
  return val
```

Initial call:

- If I go first: maxturn(initial-state,ops,0)
- If opponent goes first: minturn(initial-state,ops,0)

Minimax Algorithm with Alpha-Beta Pruning

maxturn(s,ops,**a**,**b**): {my turn}

```
if cutoff(s,depth) then return f(s)
else
  val  $\leftarrow -\infty$ ;
  foreach o  $\in$  ops
    val'  $\leftarrow$  minturn(apply(s,o),ops,a,b);
    if val' > val then
      val  $\leftarrow$  val';
      bestop  $\leftarrow$  o;
      if val  $\geq$  b then return val;
      a  $\leftarrow$  max(a,val)
  return val
```

maxturn(s,ops,**a**,**b**): {opponent's turn}

```
if cutoff(s,depth) then return f(s)
else
  val  $\leftarrow +\infty$ ;
  foreach o  $\in$  ops
    val'  $\leftarrow$  minturn(apply(s,o),ops,a,b);
    if val' < val then
      val  $\leftarrow$  val';
      bestop  $\leftarrow$  o;
      if val  $\leq$  a then return val;
      b  $\leftarrow$  min(b,val)
  return val
```

Initial call:

- If I go first: maxturn(initial-state,ops, **$-\infty$** , **$+\infty$**)
- If opponent goes first: minturn(initial-state,ops, **$-\infty$** , **$+\infty$**)