CS 4700:
Foundations of Artificial Intelligence

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Module:
Adversarial Search
R&N: Chapter 5

Part I
Outline

Adversarial Search
Optimal decisions
Minimax
$\alpha$-$\beta$ pruning
Case study: Deep Blue
UCT and Go
Adversarial Reasoning: Games

Mathematical Game Theory

Branch of economics that views any multi-agent environment as a game, provided that the impact of each agent on the others is “significant”, regardless of whether the agents are cooperative or competitive.

First step:
– Deterministic
– Turn taking
– 2-player
– Zero-sum game of perfect information (fully observable) “my win is your loss” and vice versa; utility of final states opposite for each player. My +10 is your -10.
Multi-agent game vs. single-agent search problem

"Unpredictable" opponent need a strategy: specifies a move for each possible opponent reply. E.g with “huge” lookup table.

Time limits → unlikely to find optimal response, must approximate

Rich history of game playing in AI, in particular in the area of chess.

Both Turing and Shannon viewed chess as an important challenge for machine intelligence because playing chess appears to require some level of intelligence.
A Brief History of Computer Chess

1912

1950s

1970s

1997

Today
Human-computer hybrid most exciting new level of play. Computers as smart assistants are becoming accepted. Area referred to as “Assisted Cognition.”

Another example: mind-reading binoculars for 10 km vision.
Why is Game-Playing a Challenge for AI?

Competent game playing is a mark of some aspects of “intelligence”
- Requires planning, reasoning and learning

Proxy for real-world decision making problems
- Easy to represent states & define rules
- Obtaining good performance is hard

“Adversary” can be nature

PSPACE-complete (or worse)
- Computationally equivalent to hardware debugging, formal verification, logistics planning
- PSPACE believed to be harder than NP.
Traditional Board Games

Finite
Two-player
Zero-sum
Deterministic
Perfect Information
Sequential
Tic-tac-toe (or Noughts and crosses, Xs and Os)

Key Idea: Look Ahead

We start 3 moves per player in:

3x3 Tic-Tac-Toe optimal play

X’s turn

O’s turn

loss

loss
Look-ahead based Tic-Tac-Toe

X’s turn

O’s turn

X

Tie

Tie

Tie

Tie
Look-ahead based Tic-Tac-Toe

X’s turn

O’s turn

Win for O

Win for O

Tie

Tie

Tie

Tie
Look-ahead based Tic-Tac-Toe

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Tie

Win for O
Each board in game tree gets unique game tree value (utility; -1/0/+1) under optimal rational play. (Convince yourself.)

Approach: Look first at bottom tree. Label bottom-most boards. Then label boards one level up, according result of best possible move. … and so on. Moving up layer by layer.

Termed the **Minimax Algorithm**

– Implemented as a depth-first search

E.g. 0 for top board.

What if our opponent does not play optimally?
Aside: Game tree learning

Can (in principle) store all board values in large table. $3^{19} = 19.683$ for tic-tac-toe.

Can use table to try to train classifier to predict “win”, “loss”, or “draw.”

Issue: For real games, one can only look at tiny, tiny fragment of table.

Reinforcement learning builds on this idea.

See eg Irvine Machine Learning archive.
archive.ics.uci.edu/ml/datasets/Tic-Tac-Toe+Endgame
Look-ahead based Chess

But there’s a catch…
How big is this tree?

Approx. $10^{120} >$ Number of atoms in the observable universe ($10^{80}$)

We can really only search a tiny, miniscule faction of this tree!

Around 60 x $10^9$ nodes for 5 minute move. Approx. $1 / 10^{70}$ fraction.
Don’t search to the very end

- Go down 10-12 levels (still deeper than most humans)
- But now what?
- Compute an estimate of the position’s value
  - This heuristic function is typically designed by a domain expert

Consider a game tree with leaf utilities (final boards) $+1 / 0 / -1$ (or $+\infty / 0 / -\infty$). What are the utilities of intermediate boards in the game tree?

$+1 / 0 / -1$
(or $+\infty / 0 / -\infty$)

The board heuristics is trying to estimate these values from a quick calculation on the board. Eg, considering material won/loss on chess board or regions captures in GO. Heuristic value of e.g. $+0.9$, suggests true value may be $+1$. 

What’s the work-around?
What is a problem for the board heuristics (or evaluation functions) at the beginning of the game?

(Consider a heuristics that looks at lost and captured pieces.)

What will the heuristic values be near the top?

Close to 0! Not much has happened yet…..

Other issue: children of any node are mostly quite similar. Gives almost identical heuristic board values. Little or no information about the right move.

Solution: Look ahead. I.e., build search tree several levels deep (hopefully 10 or more levels). Boards at bottom of tree more diverse. Use minimax search to determine value of starting board, assuming optimal play for both players.
Figure 6.23. Relationship between the level of play by chess programs
Will deeper search give stronger play?  Always? And why?

Very counterintuitive: there are “artificial games” where searching deeper leads to worse play! (Nau and Pearl 1980) Not in natural games! 
Game tree anomaly.

Heuristic board eval value is sometimes informally referred to as the “chance of winning” from that position.

That’s a bit odd, because in a deterministic game with perfect information and optimal play, there is no “chance” at all! Each board has a fixed utility: -1, 0, or +1 (a loss, draw, or a win). (result from game theory)

Still, “chance of winning” is an informally useful notion. But no clear semantics to heuristic values.

What if board eval gives true board utility? How much search is needed to make a move? We’ll see that using machine learning and “self play,” we can get close for backgammon.
Two important factors for success:

- Deep look ahead
- Good heuristic function

Are there games where this is not feasible?
Limitations?

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Are there games where this is not feasible?

Looking 14 levels ahead in Chess ≈
Looking 4 levels ahead in Go
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Minimax players for GO were very weak until 2007…but now play at master level.
Limitations?

Two important factors for success:

- Deep look ahead
- Good heuristic function

Are there games where this is not feasible?

Looking 14 levels ahead in Chess \( \approx \) Looking 4 levels ahead in Go

Moves have extremely delayed effects

New sampling based search method: Upper Confidence bounds applied to Trees (UCT)
Well... Why not use a strategy / knowledge, as humans do?

Consider for Tic-Tac-Toe:

1. If there is a winning move, make it.
2. If the opponent can win at a square by his next move, play that move.
3. Taking the central square is more important than taking other squares.
4. Taking corner squares is more important than taking squares on the edges.

Oops!!

Consider

Black uses the strategy...
So, although one can capture strategic knowledge of many games in high-level rules (at least to some extent), in practice any interesting game will revolve precisely around the exceptions to those rules!

Issue has been studied for decades but research keeps coming back to game tree search (or most recently, game tree sampling).

Currently only one exception: reinforcement learning for backgammon. (discussed later)

A very strong board evaluation function was learned in self-play. Represented as a neural net. Almost no search remained.
Formal definition of a game:

- Initial state
- Successor function: returns list of \((move, state)\) pairs
- Terminal test: determines when game over
  Terminal states: states where game ends
- Utility function (objective function or payoff function): gives numeric value for terminal states

We will consider games with 2 players (Max and Min)

Max moves first.
Game Tree Example: Tic-Tac-Toe

Tree from Max’s perspective
Minimax Algorithm

**Minimax algorithm**
- Perfect play for deterministic, 2-player game
- Max tries to maximize its score
- Min tries to minimize Max’s score (Min)
- Goal: Max to move to position with highest *minimax value*
  - Identify best achievable payoff against best play
Minimax Algorithm

Payoff for Max
Minimax Algorithm (cont’d)

Payoff for Max
Minimax Algorithm (cont’d)

Payoff for Max
What if
payoff(Q) = 100
payoff(R) = 200

Starting DFS, left to right, do we need to know eval(H)?

Do DFS. Real games: use iterative deepening. (gives “anytime” approach.)

$$\geq 3 \text{ (DFS left to right)}$$

$$\leq 0$$

$$\leq 2$$

Prune!

Prune!

Payoff for Max
Properties of minimax algorithm:

**Complete?** Yes (if tree is finite)

**Optimal?** Yes (against an optimal opponent)

**Time complexity?** $O(b^m)$

**Space complexity?** $O(bm)$ (depth-first exploration, if it generates all successors at once)

For chess, $b \approx 35$, $m \approx 80$ for "reasonable" games

$\rightarrow$ exact solution completely infeasible

$m$ – maximum depth of the tree; $b$ – legal moves
Minimax Algorithm

Limitations
  – Generally not feasible to traverse entire tree
  – Time limitations

Key Improvements
  – Use evaluation function instead of utility (discussed earlier)
    • Evaluation function provides estimate of utility at given position
  – Alpha/beta pruning
Can we improve search by reducing the size of the game tree to be examined?

→ Yes! Using alpha-beta pruning

Principle
- If a move is determined worse than another move already examined, then there is no need for further examination of the node.

Analysis shows that will be able to search almost twice as deep. *Really is what makes game tree search practically feasible.*

E.g. Deep Blue 14 plies using alpha-beta pruning. Otherwise only 7 or 8 (weak chess player). (plie = half move / one player)
$\alpha$-$\beta$ Pruning Example
MAX

MIN

3

12

8

2

≥ 3

≤ 2

X

X
What gives best pruning?

Visit most promising (from min/max perspective) first.
Rules:

− \( \alpha \) is the best (highest) found so far along the path for \textbf{Max}
− \( \beta \) is the best (lowest) found so far along the path for \textbf{Min}
− Search below a \textbf{MIN} node may be \textit{alpha-pruned} if the its \( \beta \leq \alpha \) of some \textbf{MAX} ancestor
− Search below a \textbf{MAX} node may be \textit{beta-pruned} if the its \( \alpha \geq \beta \) of some \textbf{MIN} ancestor.
α is the value of the best (i.e., highest-value) choice found so far at any choice point along the path for \textit{max}.

If \( v \) is worse than \( \alpha \), \textit{max} will avoid it.

\[ \rightarrow \] prune that branch

Define \( \beta \) similarly for \textit{min}. 

More abstractly
Properties of $\alpha$-$\beta$ Prune

Pruning does not affect final result

Good move ordering improves effectiveness of pruning (e.g., chess, try captures first, then threats, forward moves, then backward moves…)

With "perfect ordering," time complexity = $O(b^{m/2})$

$\rightarrow$ doubles depth of search that alpha-beta pruning can explore

Example of the value of reasoning about which computations are relevant (a form of metareasoning)
A few quick approx. numbers for Chess:

\[ b = 35 \]
\[ 200 \text{M nodes / second} \implies 5 \text{ mins} = 60 \text{ B nodes in search tree} \]
\[ (2 \text{ M nodes / sec. software only, fast PC} \implies 600 \text{ M nodes in tree}) \]

\[ 35^7 = 64 \text{ B} \]
\[ 35^5 = 52 \text{ M} \]

So, basic minimax: around 7 plies deep. (5 plies)
With, alpha-beta \[ 35^{(14/2)} = 64 \text{ B} \]. Therefore, 14 plies deep. (10 plies)

Aside:
4-ply \approx \text{human novice}
8-ply / 10-ply \approx \text{typical PC, human master}
14-ply \approx \text{Deep Blue, Kasparov (+ depth 25 for “selective extensions”) / 7 moves by each player.}
Can’t go to all the way to the “bottom:”

**evaluation function**

= estimated desirability of position

cutoff test:

e.g., depth limit
(Use Iterative Deepening)

“What is the problem with that?”

**Horizon effect.**

“Unstable positions:”

Search deeper.

Selective extensions.

E.g. exchange of several pieces in a row.

⇒ add quiescence search:

⇒ **quiescent position**: position where next move unlikely to cause large change in players’ positions
Evaluation Function

- Performed at search cutoff point
- Must have same terminal/goal states as utility function
- Tradeoff between accuracy and time → reasonable complexity
- Accurate
  - Performance of game-playing system dependent on accuracy/goodness of evaluation
  - Evaluation of nonterminal states strongly correlated with actual chances of winning
Evaluation functions

For chess, typically **linear** weighted sum of **features**

\[ Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \ldots + w_n f_n(s) \]

e.g., \( w_1 = 9 \) with

\[ f_1(s) = (\text{number of white queens}) - (\text{number of black queens}), \text{ etc.} \]

**Key challenge – find a good evaluation features:**

- Not just material! (as used by novice)
- Isolated pawns are bad.
- How well protected is your king?
- How much maneuverability to you have?
- Do you control the center of the board?
- Strategies change as the game proceeds

Features are a form of chess knowledge. Hand-coded in eval function.

Knowledge tightly integrated in search.

Feature weights: can be automatically tuned (“learned”).

Standard issue in machine learning:

- Features, generally hand-coded; weights tuned automatically.
When Chance is involved:
Backgammon Board
Expectiminimax

Generalization of minimax for games with chance nodes

Examples: Backgammon, bridge

Calculates expected value where probability is taken over all possible dice rolls/chance events
- Max and Min nodes determined as before
- Chance nodes evaluated as weighted average
Expectiminimax

Expectiminimax(n) =

\[ \text{Utility}(n) \]

for n, a terminal state

\[ \max_{s \in \text{Succ}(n)} \text{expectiminimax}(s) \]

for n, a Max node

\[ \min_{s \in \text{Succ}(n)} \text{expectiminimax}(s) \]

for n, a Min node

\[ \sum_{s \in \text{Succ}(n)} P(s) \times \text{expectiminimax}(s) \]

for n, a chance node
.9 * 2 + .1 * 3 = 2.1

Small chance at high payoff wins. But, not necessarily the best thing to do!