

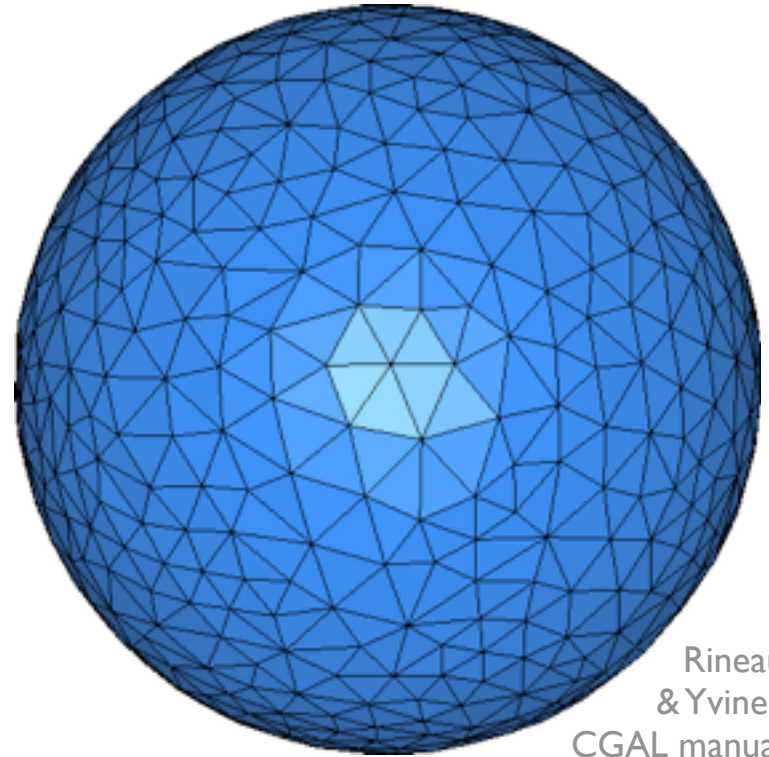
Triangle meshes

CS 4620 Lecture 7



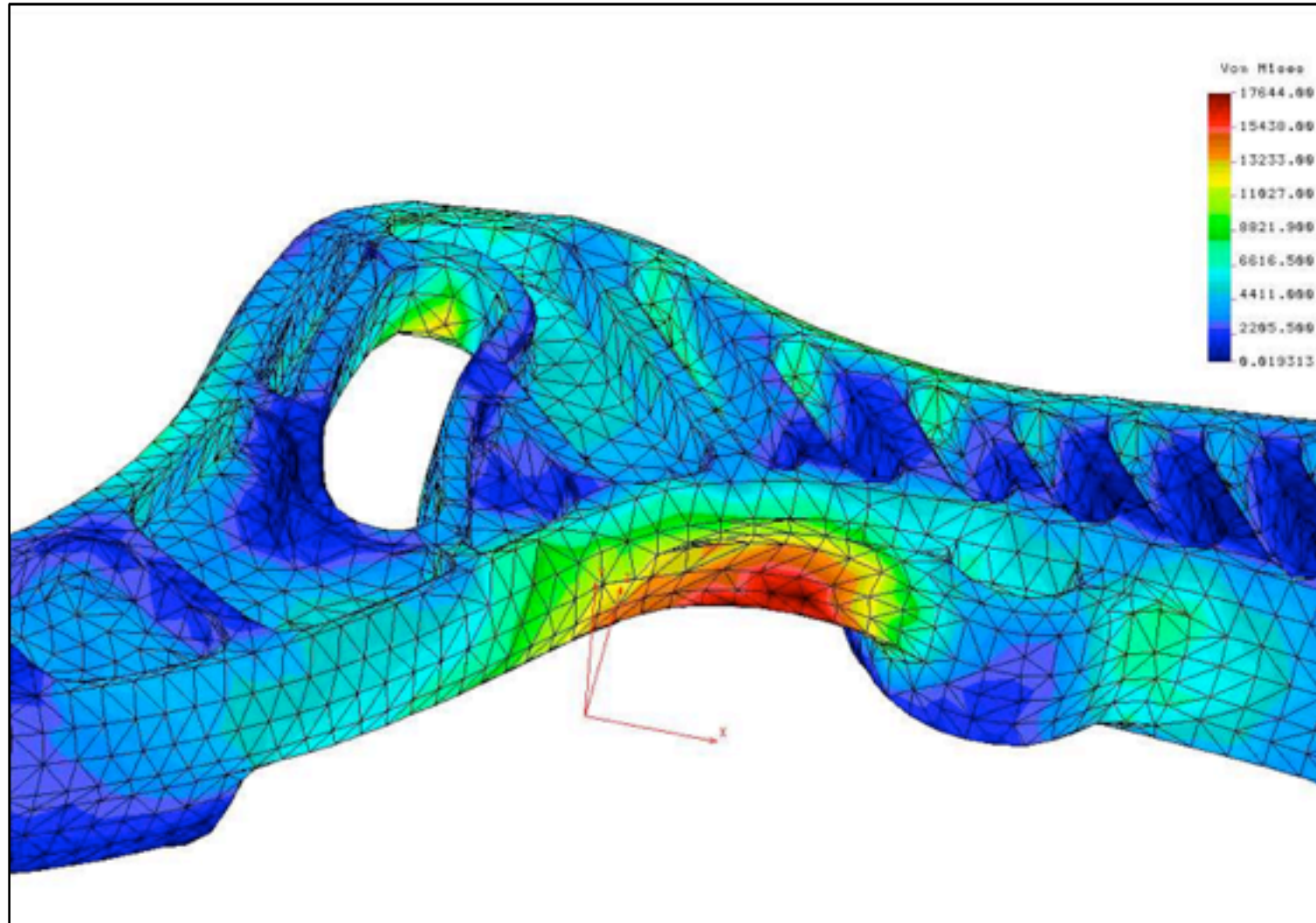
Andrzej Barabasz

spheres



Rineau
& Yvinec
CGAL manual

**approximate
sphere**



PATRIOT Engineering

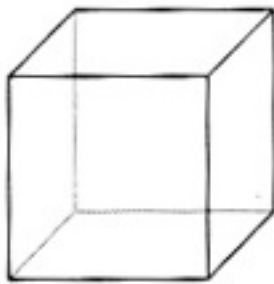
finite element analysis



Ottawa Convention Center

Notation

- $n_T = \#tris$; $n_V = \#verts$; $n_E = \#edges$
- Euler: $n_V - n_E + n_T = 2$ for a simple closed surface
 - and in general sums to small integer
 - argument for implication that $n_T:n_E:n_V$ is about 2:3:1



$$\begin{aligned} V &= 8 \\ E &= 12 \\ F &= 6 \end{aligned}$$



$$\begin{aligned} V &= 5 \\ E &= 8 \\ F &= 5 \end{aligned}$$



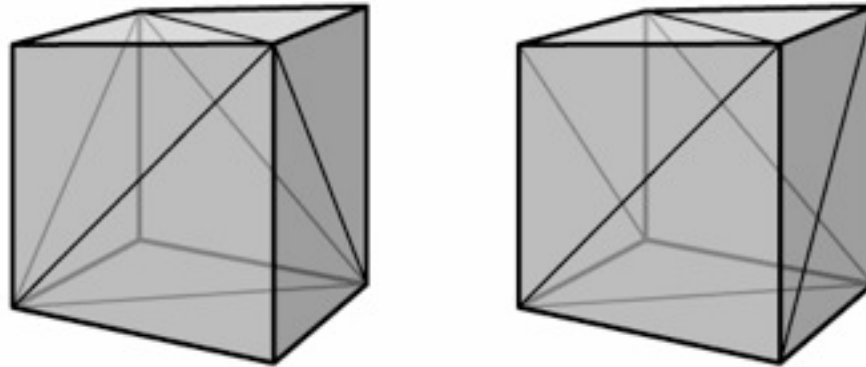
$$\begin{aligned} V &= 6 \\ E &= 12 \\ F &= 8 \end{aligned}$$

Validity of triangle meshes

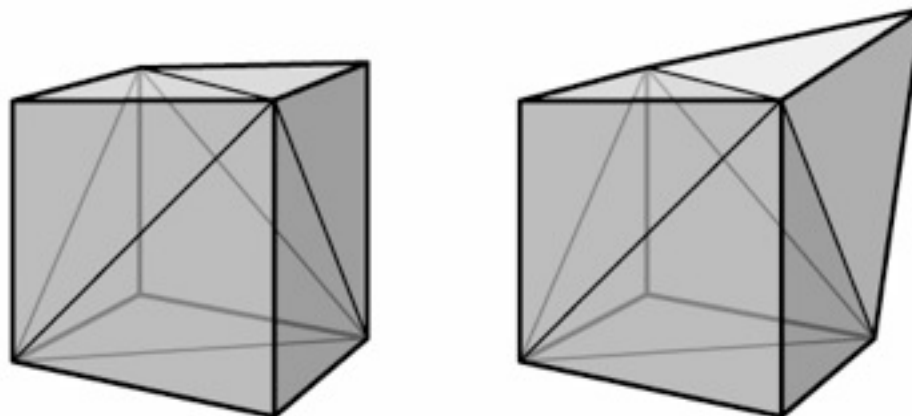
- in many cases we care about the mesh being able to bound a region of space nicely
- in other cases we want triangle meshes to fulfill assumptions of algorithms that will operate on them (and may fail on malformed input)
- two completely separate issues:
 - topology: how the triangles are connected (ignoring the positions entirely)
 - geometry: where the triangles are in 3D space

Topology/geometry examples

- same geometry, different mesh topology:



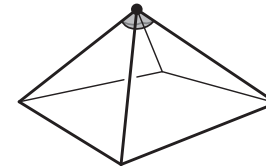
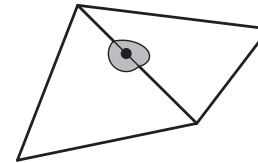
- same mesh topology, different geometry:



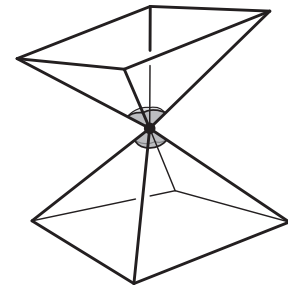
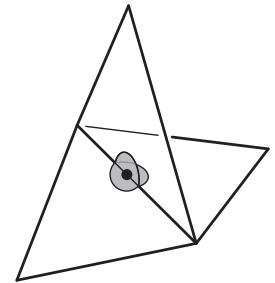
Topological validity

- strongest property: be a manifold
 - this means that no points should be "special"
 - interior points are fine
 - edge points: each edge must have exactly 2 triangles
 - vertex points: each vertex must have one loop of triangles
- slightly looser: manifold with boundary
 - weaken rules to allow boundaries

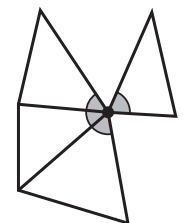
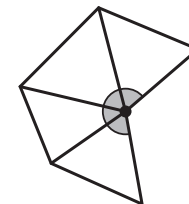
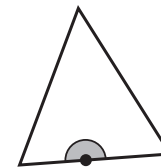
manifold



not manifold

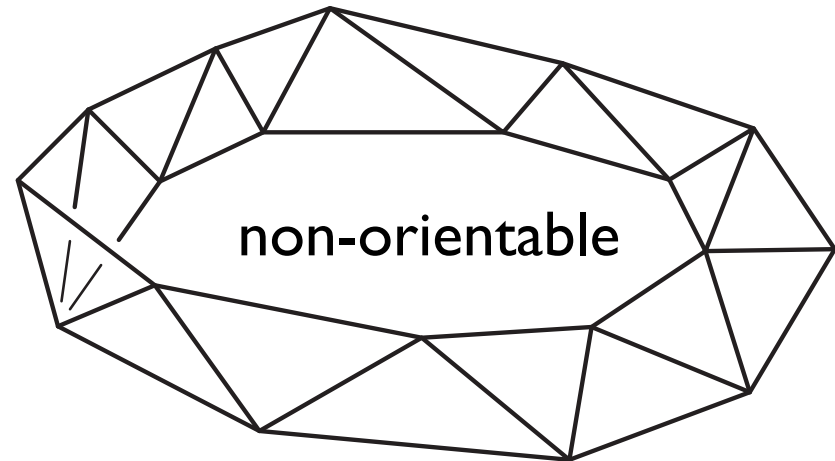
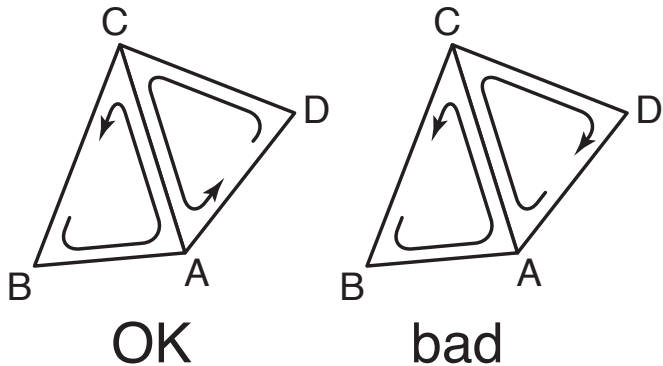


with boundary



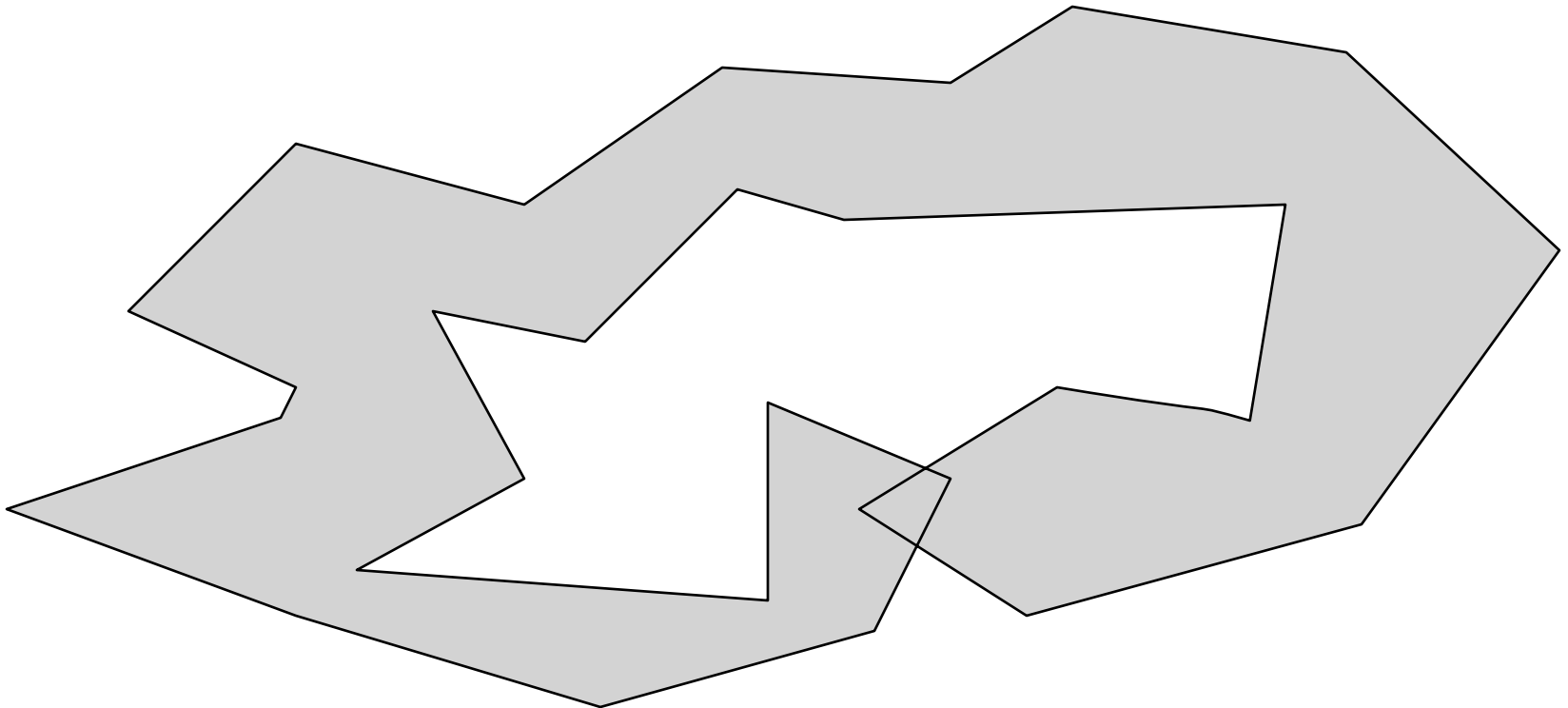
Topological validity

- Consistent orientation
 - Which side is the “front” or “outside” of the surface and which is the “back” or “inside?”
 - rule: you are on the outside when you see the vertices in counter-clockwise order
 - in mesh, neighboring triangles should agree about which side is the front!
 - caution: not always possible



Geometric validity

- generally want non-self-intersecting surface
- hard to guarantee in general
 - because far-apart parts of mesh might intersect



Representation of triangle meshes

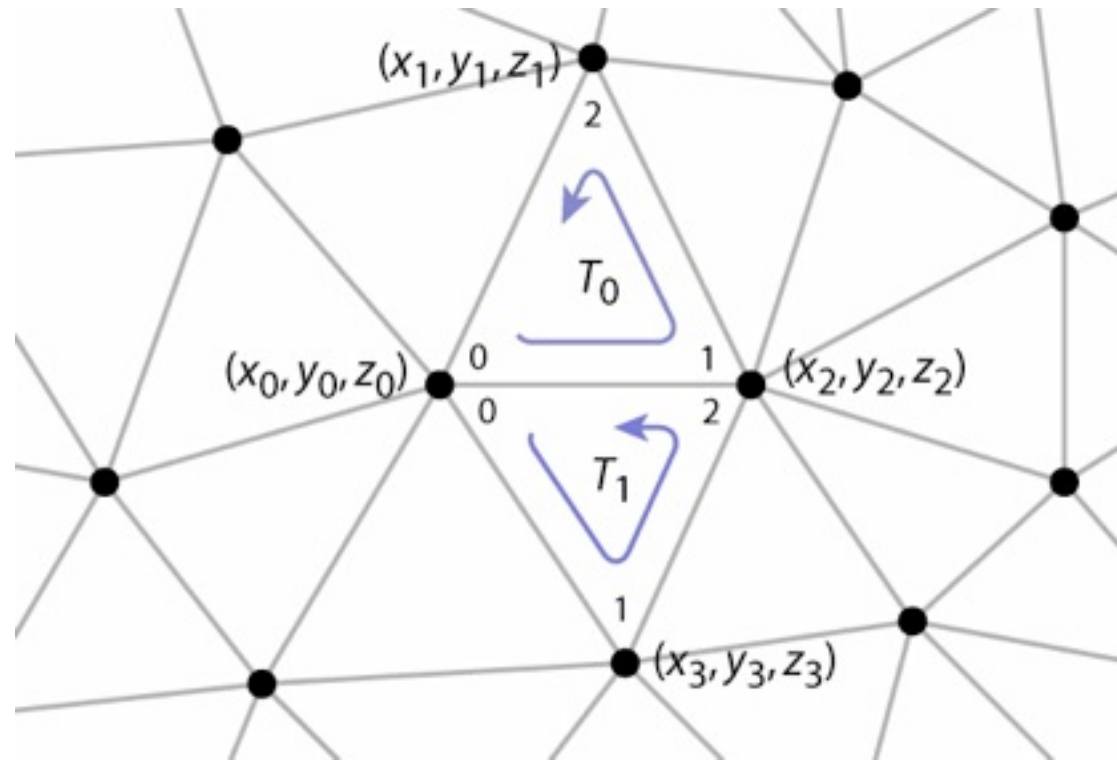
- Compactness
- Efficiency for rendering
 - enumerate all triangles as triples of 3D points
- Efficiency of queries
 - all vertices of a triangle
 - all triangles around a vertex
 - neighboring triangles of a triangle
 - (need depends on application)
 - finding triangle strips
 - computing subdivision surfaces
 - mesh editing

Representations for triangle meshes

- Separate triangles
- Indexed triangle set
 - shared vertices
- Triangle strips and triangle fans
 - compression schemes for transmission to hardware
- Triangle-neighbor data structure
 - supports adjacency queries
- Winged-edge data structure
 - supports general polygon meshes

Separate triangles

	[0]	[1]	[2]
tris[0]	x_0, y_0, z_0	x_2, y_2, z_2	x_1, y_1, z_1
tris[1]	x_0, y_0, z_0	x_3, y_3, z_3	x_2, y_2, z_2
	⋮	⋮	⋮



Separate triangles

- array of triples of points
 - $\text{float}[n_T][3][3]$: about 72 bytes per vertex
 - 2 triangles per vertex (on average)
 - 3 vertices per triangle
 - 3 coordinates per vertex
 - 4 bytes per coordinate (float)
- various problems
 - wastes space (each vertex stored 6 times)
 - cracks due to roundoff
 - difficulty of finding neighbors at all

Indexed triangle set

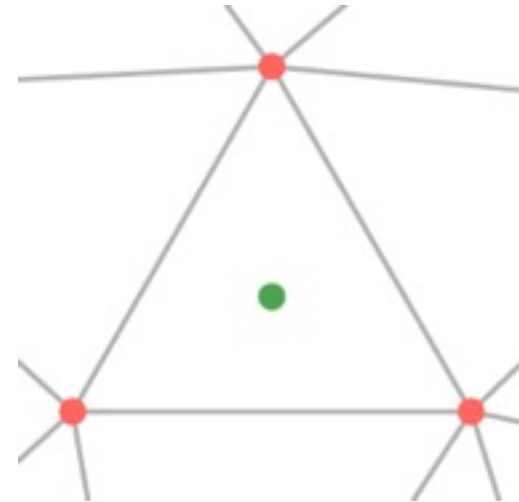
- Store each vertex once
- Each triangle points to its three vertices

```
Triangle {  
    Vertex vertex[3];  
}
```

```
Vertex {  
    float position[3]; // or other data  
}
```

```
// ... or ...
```

```
Mesh {  
    float verts[nv][3]; // vertex positions (or other data)  
    int tInd[nt][3]; // vertex indices  
}
```



Indexed triangle set

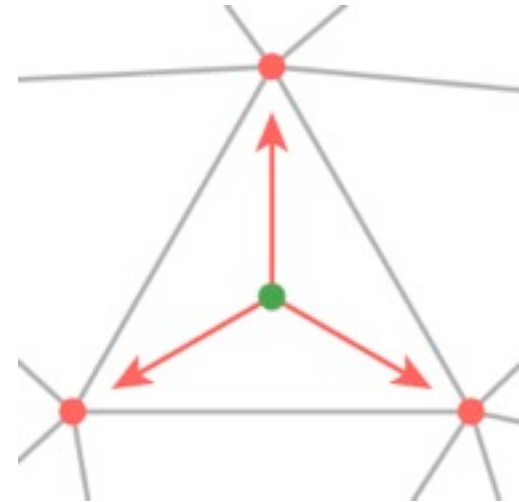
- Store each vertex once
- Each triangle points to its three vertices

```
Triangle {  
    Vertex vertex[3];  
}
```

```
Vertex {  
    float position[3]; // or other data  
}
```

```
// ... or ...
```

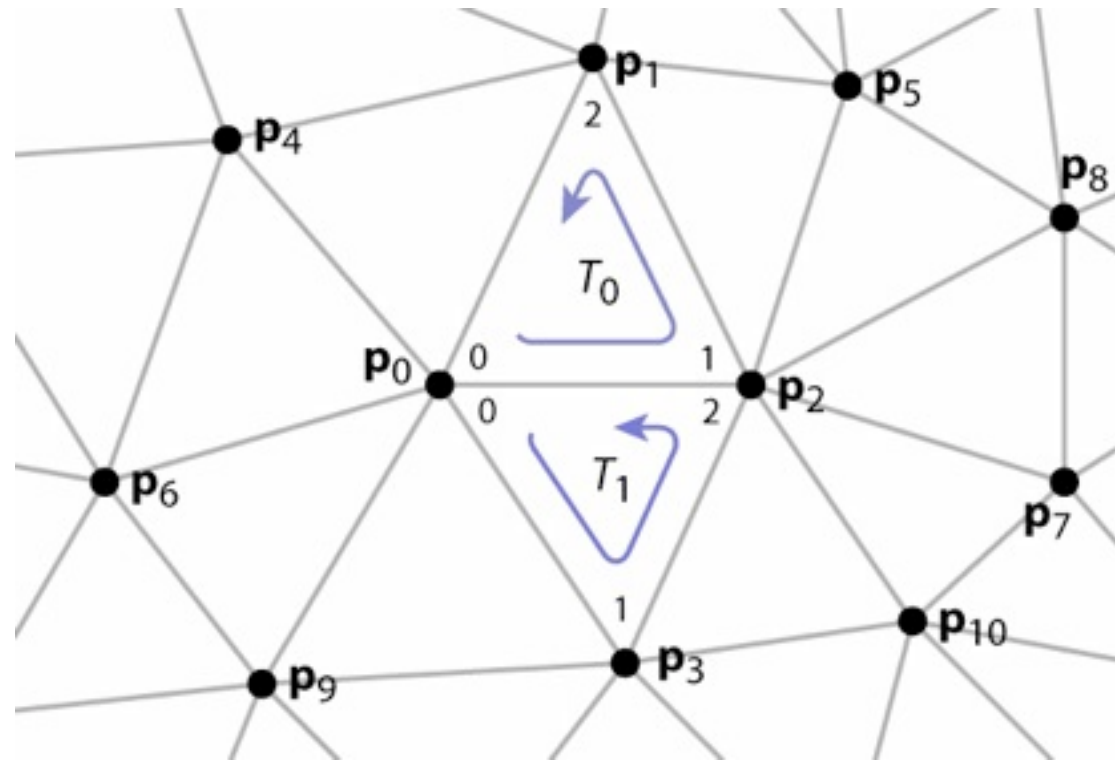
```
Mesh {  
    float verts[nv][3]; // vertex positions (or other data)  
    int tInd[nt][3]; // vertex indices  
}
```



Indexed triangle set

verts[0]	x_0, y_0, z_0
verts[1]	x_1, y_1, z_1
	x_2, y_2, z_2
	x_3, y_3, z_3
	\vdots

tInd[0]	0, 2, 1
tInd[1]	0, 3, 2
	\vdots

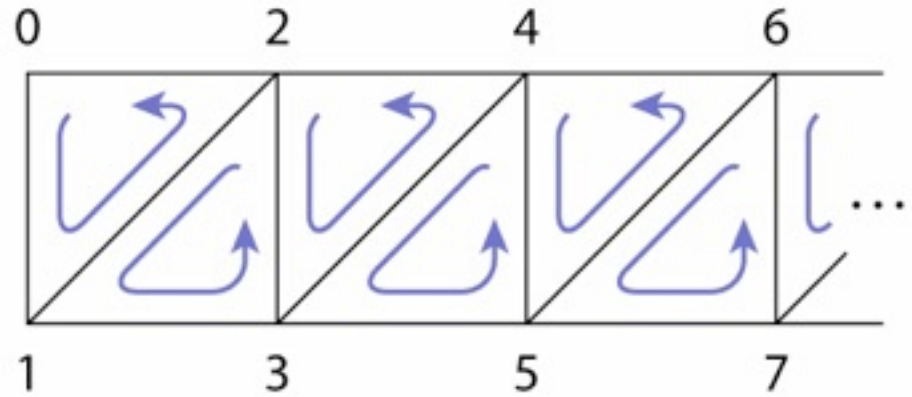


Indexed triangle set

- array of vertex positions
 - $\text{float}[n_V][3]$: 12 bytes per vertex
 - (3 coordinates x 4 bytes) per vertex
- array of triples of indices (per triangle)
 - $\text{int}[n_T][3]$: about 24 bytes per vertex
 - 2 triangles per vertex (on average)
 - (3 indices x 4 bytes) per triangle
- total storage: 36 bytes per vertex (factor of 2 savings)
- represents topology and geometry separately
- finding neighbors is at least well defined

Triangle strips

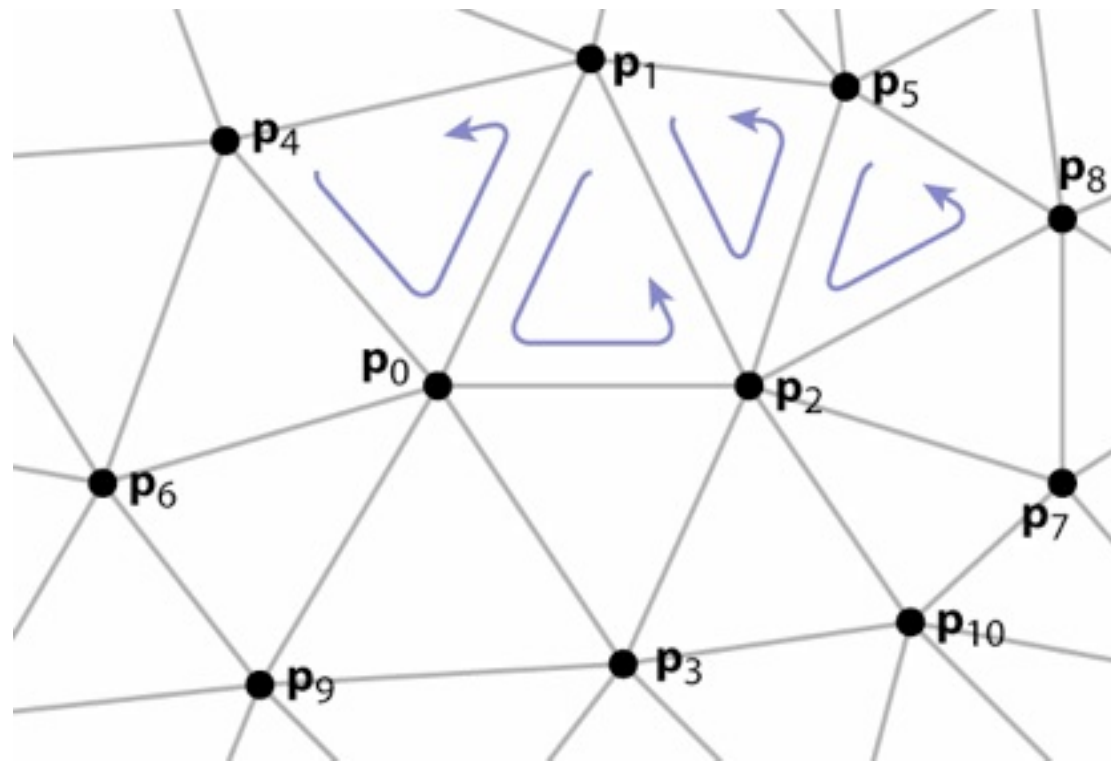
- Take advantage of the mesh property
 - each triangle is usually adjacent to the previous
 - let every vertex create a triangle by reusing the second and third vertices of the previous triangle
 - every sequence of three vertices produces a triangle (but not in the same order)
 - e. g., 0, 1, 2, 3, 4, 5, 6, 7, ... leads to
(0 1 2), (2 1 3), (2 3 4), (4 3 5), (4 5 6), (6 5 7), ...
 - for long strips, this requires about one index per triangle



Triangle strips

verts[0]	x_0, y_0, z_0
verts[1]	x_1, y_1, z_1
	x_2, y_2, z_2
	x_3, y_3, z_3
	\vdots

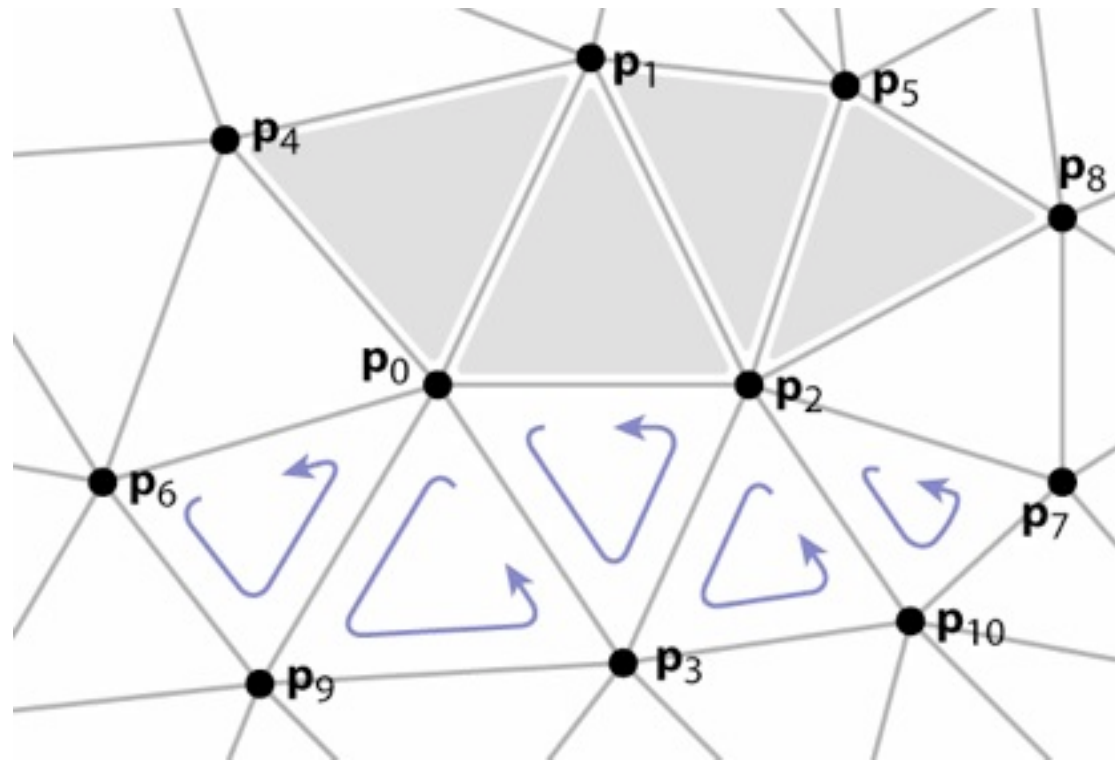
tStrip[0]	4, 0, 1, 2, 5, 8
tStrip[1]	6, 9, 0, 3, 2, 10, 7
	\vdots



Triangle strips

verts[0]	x_0, y_0, z_0
verts[1]	x_1, y_1, z_1
	x_2, y_2, z_2
	x_3, y_3, z_3
	\vdots

tStrip[0]	4, 0, 1, 2, 5, 8
tStrip[1]	6, 9, 0, 3, 2, 10, 7
	\vdots

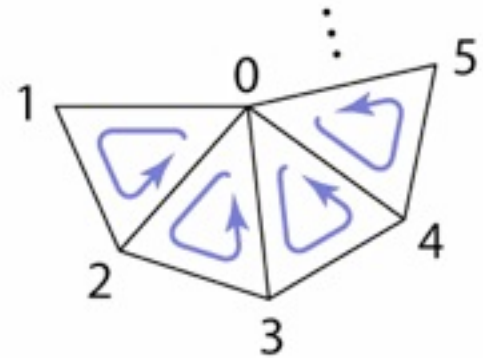


Triangle strips

- array of vertex positions
 - $\text{float}[n_V][3]$: 12 bytes per vertex
 - (3 coordinates x 4 bytes) per vertex
- array of index lists
 - $\text{int}[n_S][\text{variable}]$: 2 + n indices per strip
 - on average, $(1 + \varepsilon)$ indices per triangle (assuming long strips)
 - 2 triangles per vertex (on average)
 - about 4 bytes per triangle (on average)
- total is 20 bytes per vertex (limiting best case)
 - factor of 3.6 over separate triangles; 1.8 over indexed mesh

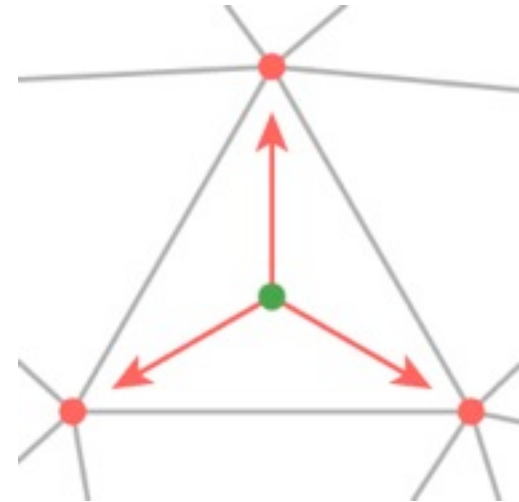
Triangle fans

- Same idea as triangle strips, but keep oldest rather than newest
 - every sequence of three vertices produces a triangle
 - e. g., 0, 1, 2, 3, 4, 5, ... leads to
(0 1 2), (0 2 3), (0 3 4), (0 3 5),
...
 - for long fans, this requires
about one index per triangle
- Memory considerations exactly the same as triangle strip



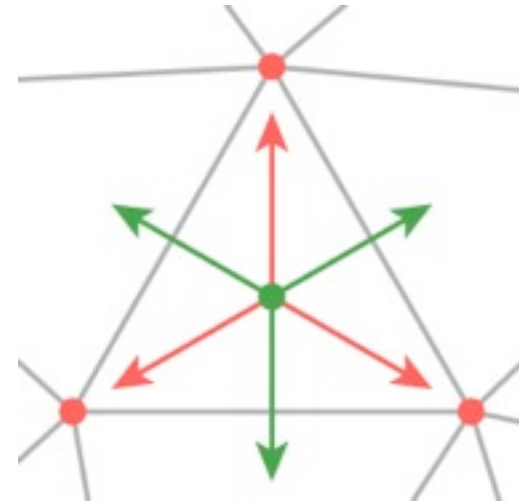
Triangle neighbor structure

- Extension to indexed triangle set
- Triangle points to its three neighboring triangles
- Vertex points to a single neighboring triangle
- Can now enumerate triangles around a vertex



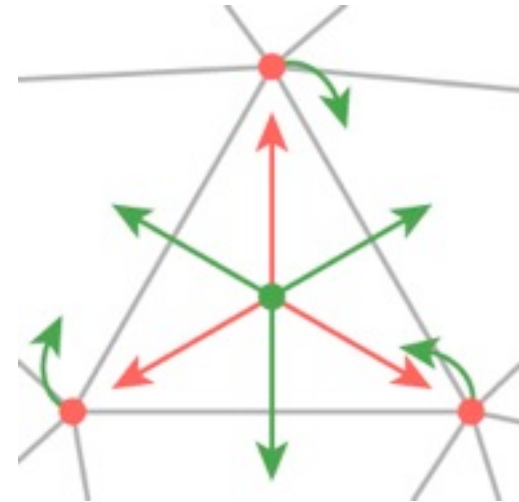
Triangle neighbor structure

- Extension to indexed triangle set
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Triangle neighbor structure

- Extension to indexed triangle set
- Triangle points to its three neighboring triangles
- Vertex points to a single neighboring triangle
- Can now enumerate triangles around a vertex



Triangle neighbor structure

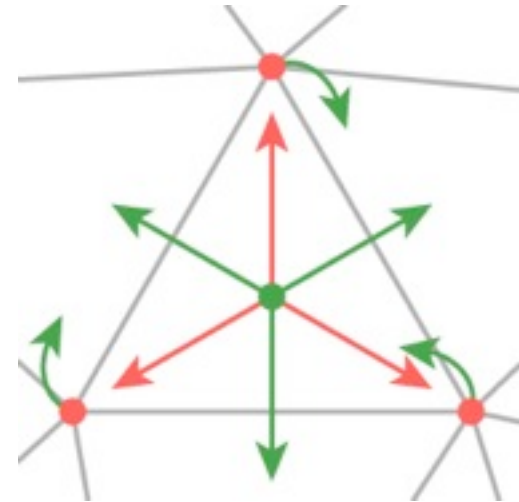
```
Triangle {  
    Triangle nbr[3];  
    Vertex vertex[3];  
}
```

```
// t.neighbor[i] is adjacent  
// across the edge from i to i+1
```

```
Vertex {  
    // ... per-vertex data ...  
    Triangle t; // any adjacent tri  
}
```

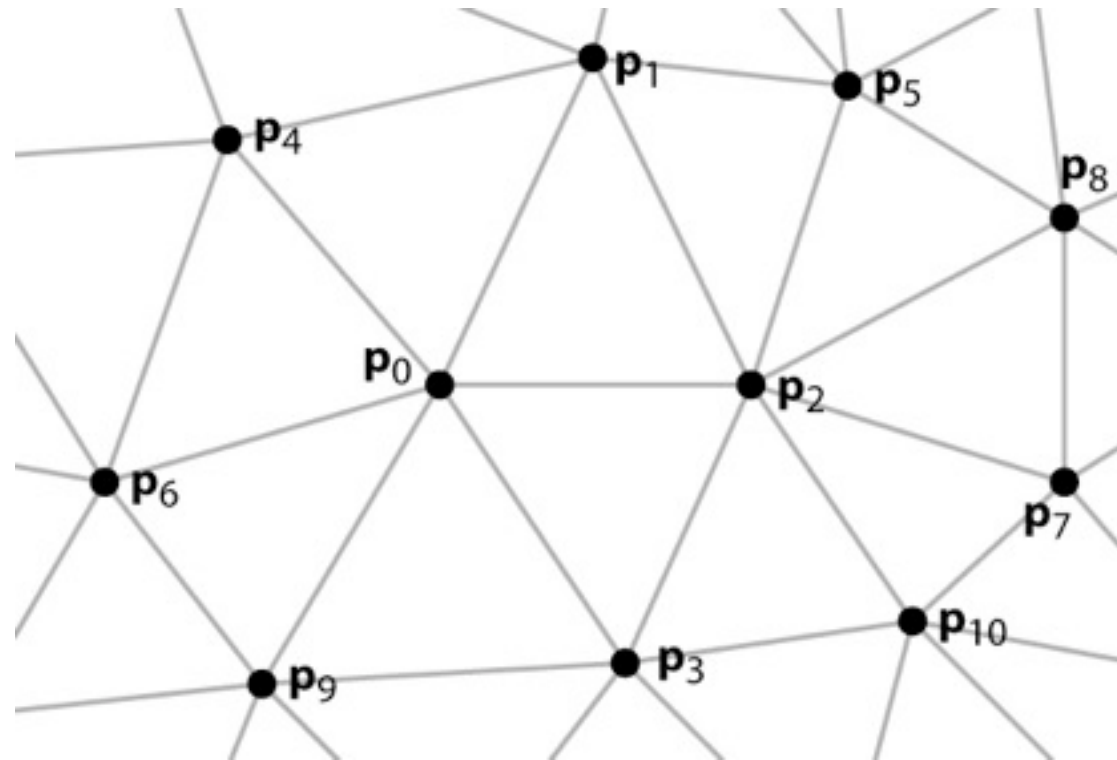
```
// ... or ...
```

```
Mesh {  
    // ... per-vertex data ...  
    int tInd[nt][3]; // vertex indices  
    int tNbr[nt][3]; // indices of neighbor triangles  
    int vTri[nv]; // index of any adjacent triangle  
}
```



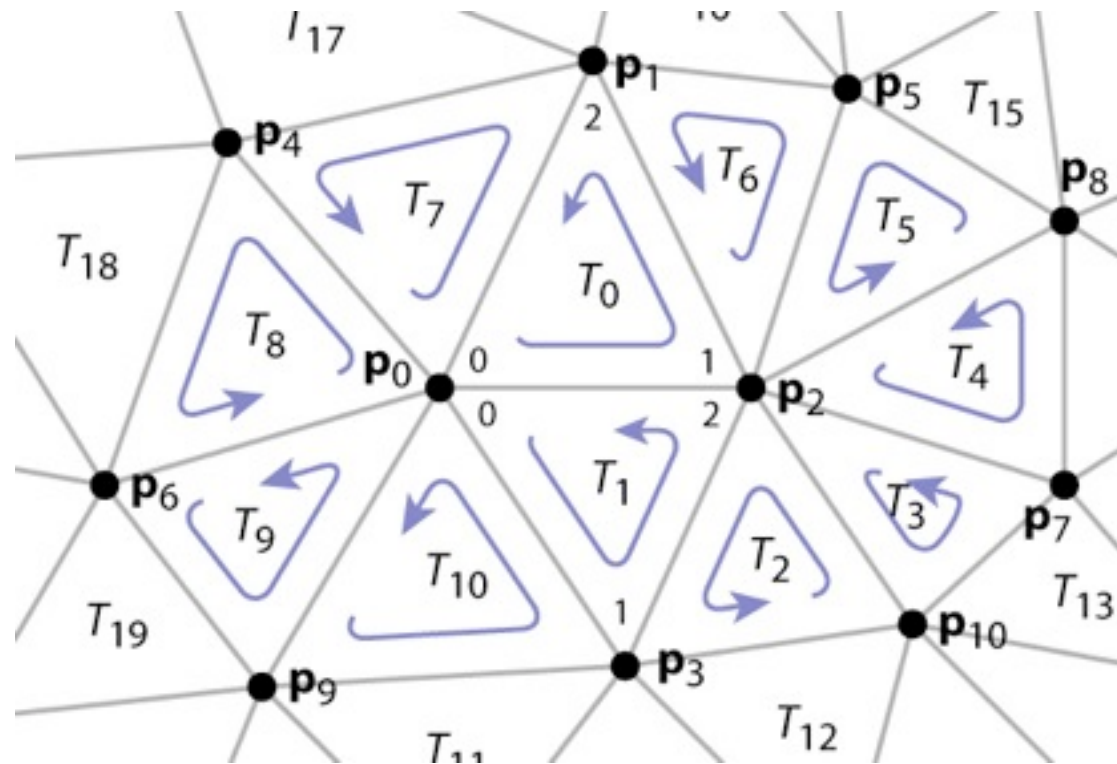
Triangle neighbor structure

$vTri[0]$	0	$tNbr[0]$	1, 6, 7
$vTri[1]$	6	$tNbr[1]$	10, 2, 0
$vTri[2]$	1	$tNbr[2]$	3, 1, 12
$vTri[3]$	1	$tNbr[3]$	2, 13, 4
	\vdots		\vdots
		$tInd[0]$	0, 2, 1
		$tInd[1]$	0, 3, 2
		$tInd[2]$	10, 2, 3
		$tInd[3]$	2, 10, 7
			\vdots



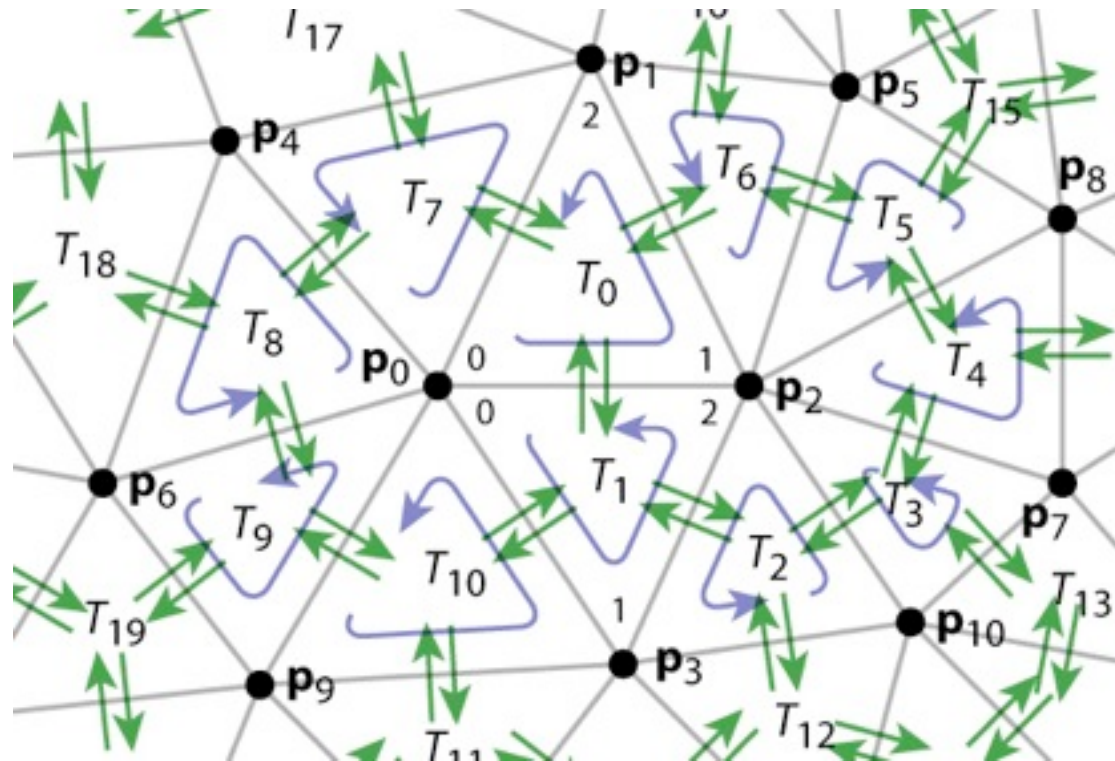
Triangle neighbor structure

vTri[0]	0	tNbr[0]	1, 6, 7
vTri[1]	6	tNbr[1]	10, 2, 0
vTri[2]	1	tNbr[2]	3, 1, 12
vTri[3]	1	tNbr[3]	2, 13, 4
	⋮		⋮
		tInd[0]	0, 2, 1
		tInd[1]	0, 3, 2
		tInd[2]	10, 2, 3
		tInd[3]	2, 10, 7
			⋮



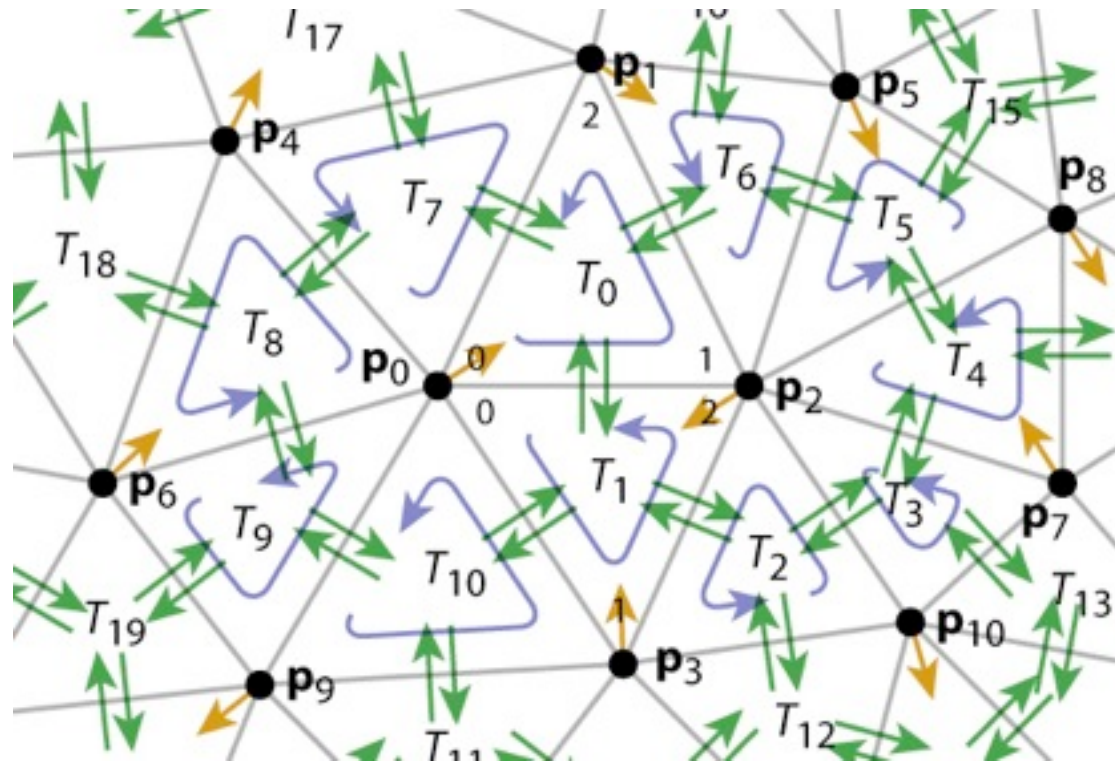
Triangle neighbor structure

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vTri[3]	1	tNbr[3]	2, 13, 4
	⋮		⋮
		tInd[0]	0, 2, 1
		tInd[1]	0, 3, 2
		tInd[2]	10, 2, 3
		tInd[3]	2, 10, 7
			⋮



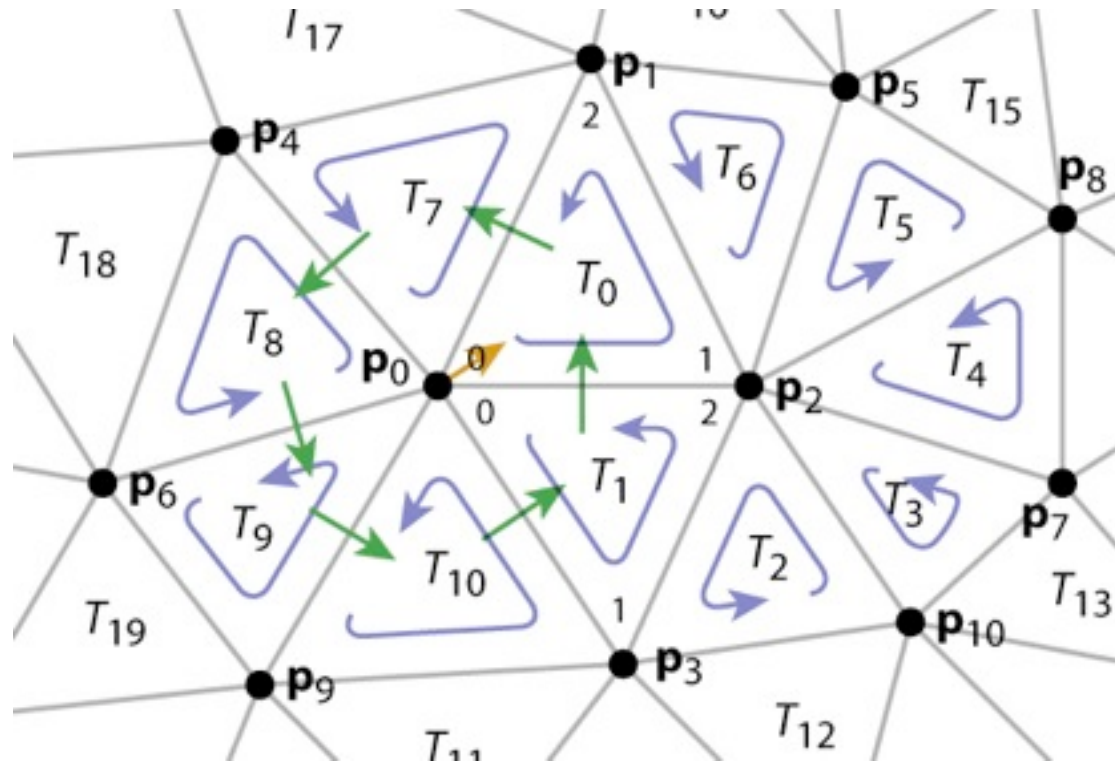
Triangle neighbor structure

$vTri[0]$	0	$tNbr[0]$	1, 6, 7
$vTri[1]$	6	$tNbr[1]$	10, 2, 0
$vTri[2]$	1	$tNbr[2]$	3, 1, 12
$vTri[3]$	1	$tNbr[3]$	2, 13, 4
	⋮		⋮
		$tInd[0]$	0, 2, 1
		$tInd[1]$	0, 3, 2
		$tInd[2]$	10, 2, 3
		$tInd[3]$	2, 10, 7
			⋮



Triangle neighbor structure

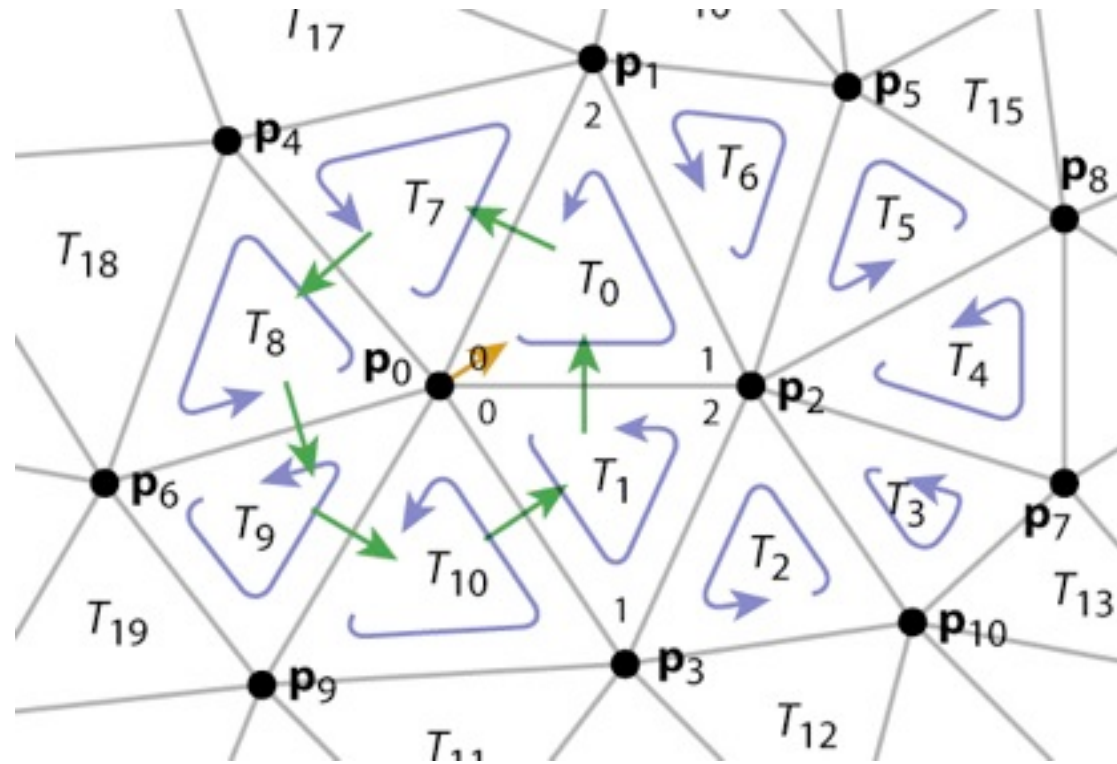
vTri[0]	0	tNbr[0]	1, 6, 7
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vTri[2]	1	tNbr[2]	3, 1, 12
vTri[3]	1	tNbr[3]	2, 13, 4
	⋮		⋮
		tInd[0]	0, 2, 1
		tInd[1]	0, 3, 2
		tInd[2]	10, 2, 3
		tInd[3]	2, 10, 7
			⋮



Triangle neighbor structure

```
TrianglesOfVertex(v) {  
  t = v.t;  
  do {  
    find t.vertex[i] == v;  
    t = t.nbr[pred(i)];  
  } while (t != v.t);  
}
```

```
pred(i) = (i+2) % 3;  
succ(i) = (i+1) % 3;
```



Triangle neighbor structure

- indexed mesh was 36 bytes per vertex
- add an array of triples of indices (per triangle)
 - $\text{int}[n_T][3]$: about 24 bytes per vertex
 - 2 triangles per vertex (on average)
 - (3 indices x 4 bytes) per triangle
- add an array of representative triangle per vertex
 - $\text{int}[n_V]$: 4 bytes per vertex
- total storage: 64 bytes per vertex
 - still not as much as separate triangles

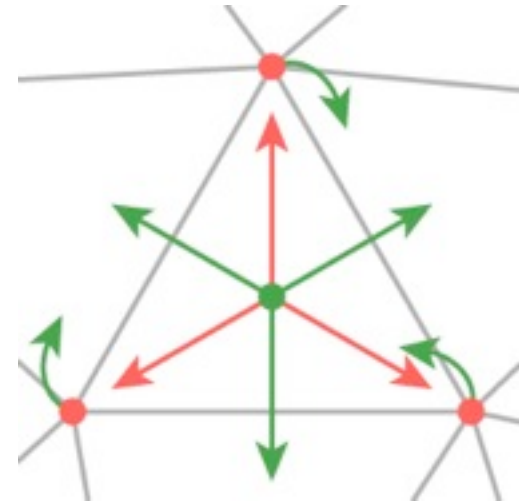
Triangle neighbor structure—refined

```
Triangle {  
    Edge nbr[3];  
    Vertex vertex[3];  
}
```

```
// if t.nbr[i].i == j  
// then t.nbr[i].t.nbr[j] == t
```

```
Edge {  
    // the i-th edge of triangle t  
    Triangle t;  
    int i; // in {0,1,2}  
    // in practice t and i share 32 bits  
}
```

```
Vertex {  
    // ... per-vertex data ...  
    Edge e; // any edge leaving vertex  
}
```



$T_0.nbr[0] = \{ T_1, 2 \}$

$T_1.nbr[2] = \{ T_0, 0 \}$

$V_0.e = \{ T_1, 0 \}$

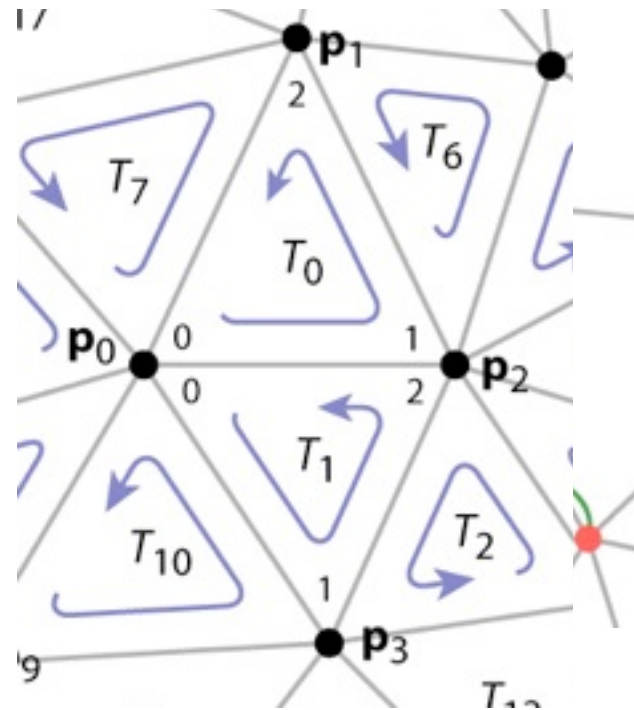
Triangle neighbor structure—refined

```
Triangle {  
    Edge nbr[3];  
    Vertex vertex[3];  
}
```

```
// if t.nbr[i].i == j  
// then t.nbr[i].t.nbr[j] == t
```

```
Edge {  
    // the i-th edge of triangle t  
    Triangle t;  
    int i; // in {0,1,2}  
    // in practice t and i share 32 bits  
}
```

```
Vertex {  
    // ... per-vertex data ...  
    Edge e; // any edge leaving vertex  
}
```



$T_0.nbr[0] = \{ T_1, 2 \}$

$T_1.nbr[2] = \{ T_0, 0 \}$

$V_0.e = \{ T_1, 0 \}$

Triangle neighbor structure

```

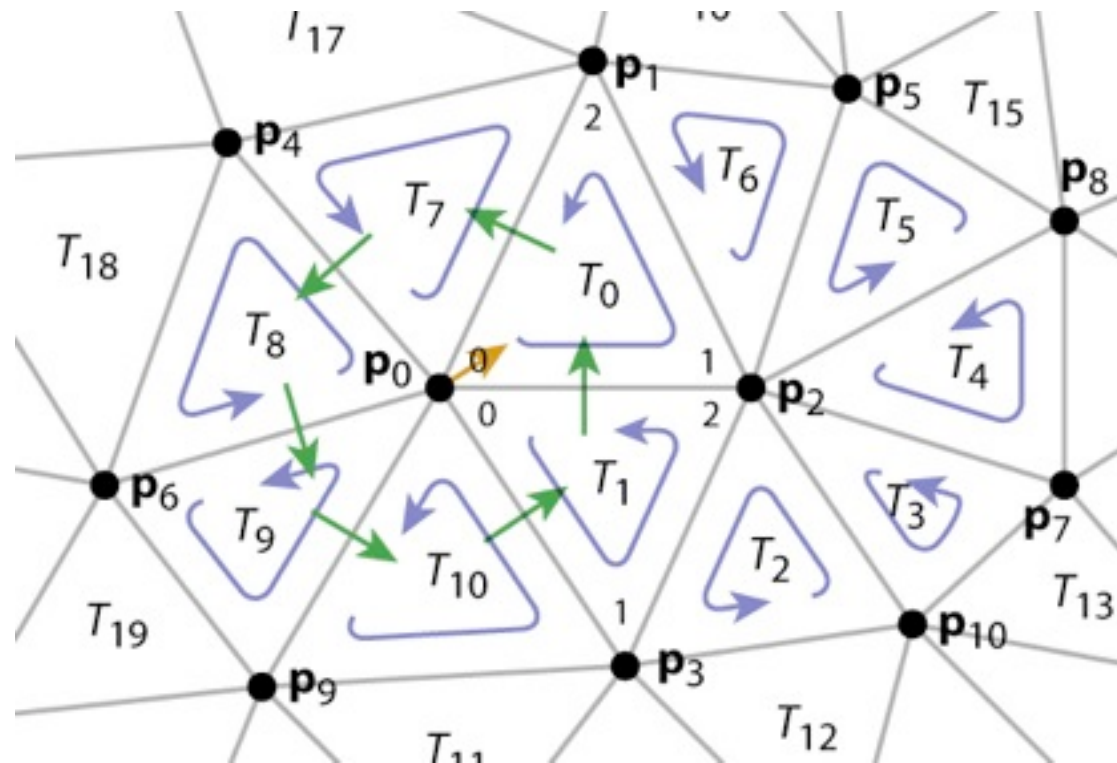
TrianglesOfVertex(v) {
  {t, i} = v.e;
  do {
    {t, i} = t.nbr[pred(i)];
  } while (t != v.t);
}

```

```

pred(i) = (i+2) % 3;
succ(i) = (i+1) % 3;

```



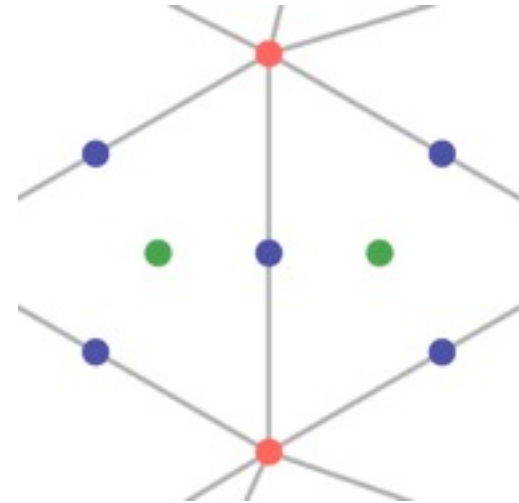
$T_0.nbr[0] = \{ T_1, 2 \}$

$T_1.nbr[2] = \{ T_0, 0 \}$

$V_0.e = \{ T_1, 0 \}$

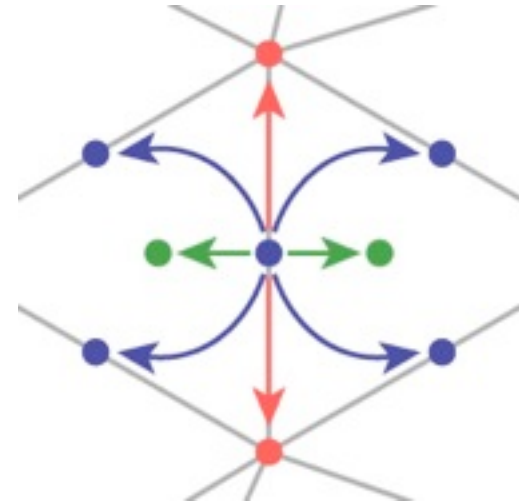
Winged-edge mesh

- Edge-centric rather than face-centric
 - therefore also works for polygon meshes
- Each (oriented) edge points to:
 - left and right forward edges
 - left and right backward edges
 - front and back vertices
 - left and right faces
- Each face or vertex points to one edge



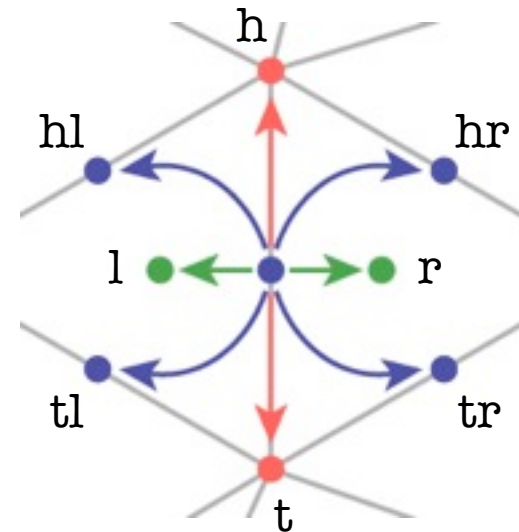
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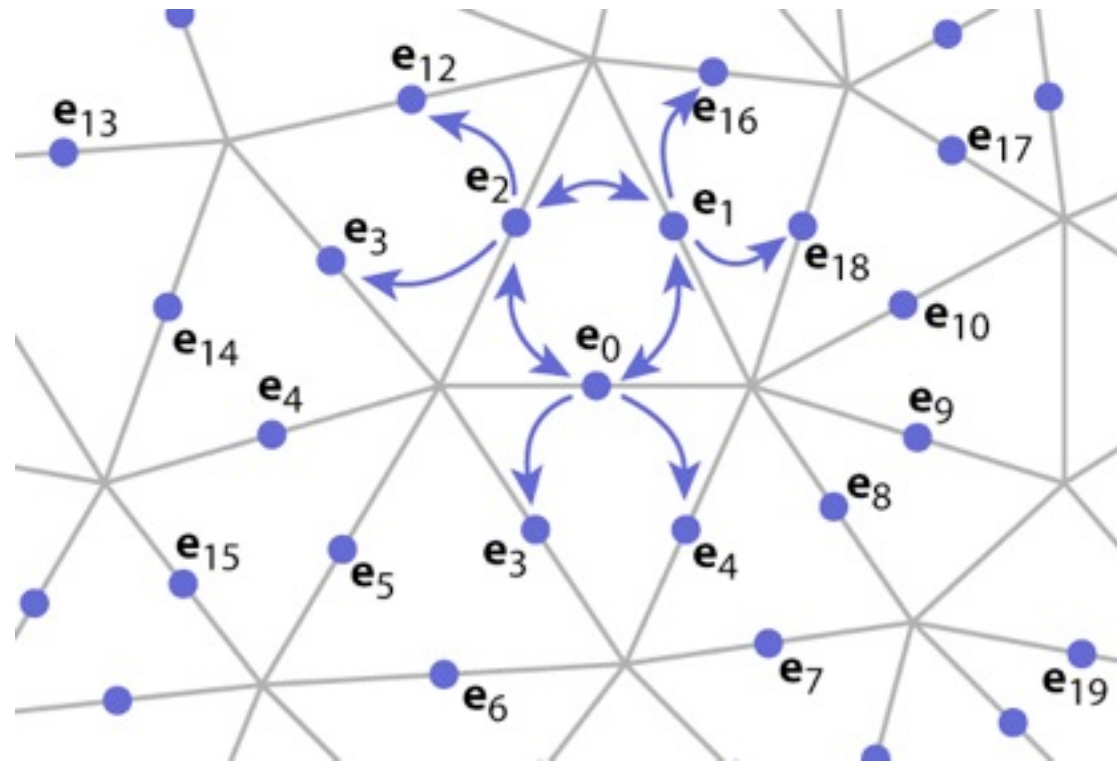
Winged-edge mesh

```
Edge {  
    Edge hl, hr, tl, tr;  
    Vertex h, t;  
    Face l, r;  
}  
  
Face {  
    // per-face data  
    Edge e; // any adjacent edge  
}  
  
Vertex {  
    // per-vertex data  
    Edge e; // any incident edge  
}
```



Winged-edge structure

	hl	hr	tl	tr
edge[0]	1	4	2	3
edge[1]	18	0	16	2
edge[2]	12	1	3	0
	⋮			



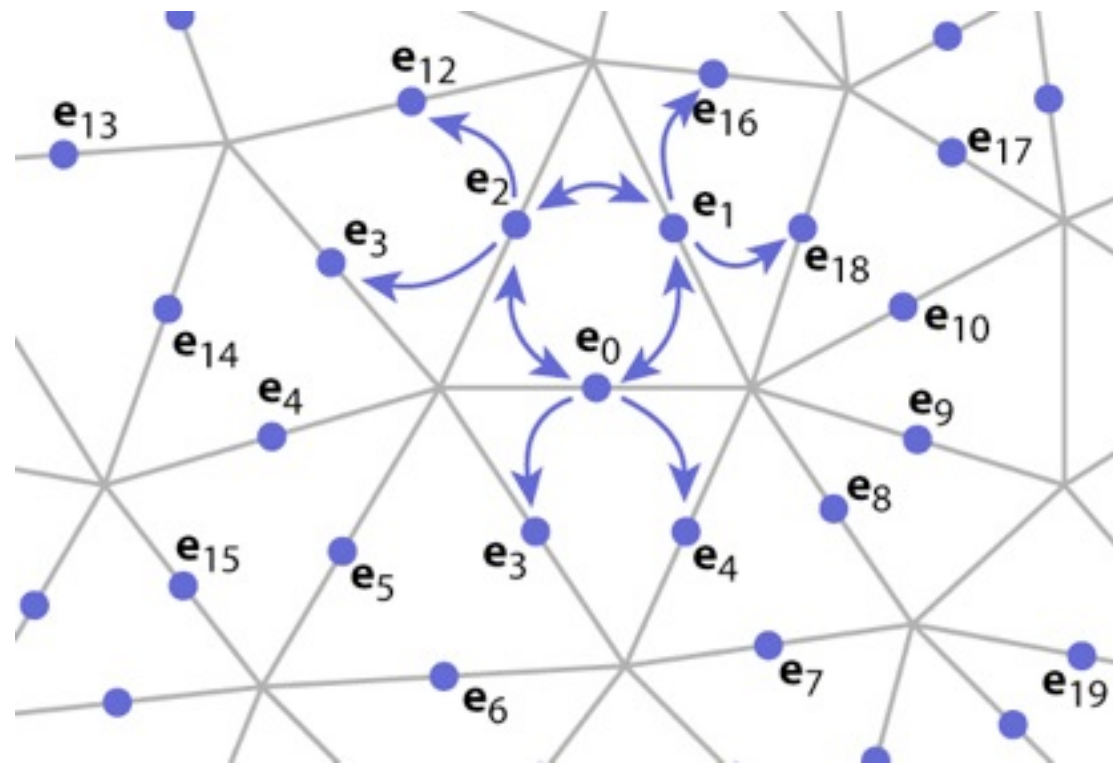
Winged-edge structure

```

EdgesOfFace(f) {
  e = f.e;
  do {
    if (e.l == f)
      e = e.hl;
    else
      e = e.tr;
  } while (e != f.e);
}

```

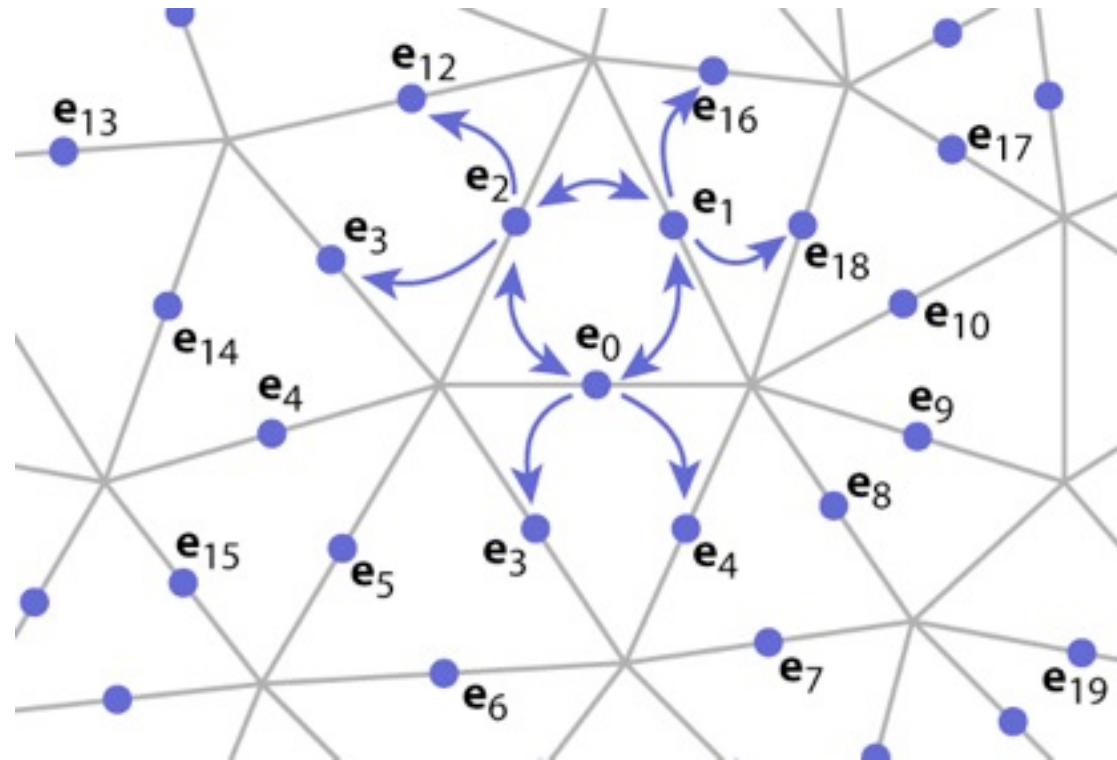
	hl	hr	tl	tr
edge[0]	1	4	2	3
edge[1]	18	0	16	2
edge[2]	12	1	3	0
	⋮			



Winged-edge structure

```
EdgesOfVertex(v) {
  e = v.e;
  do {
    if (e.t == v)
      e = e.tl;
    else
      e = e.hr;
  } while (e != v.e);
}
```

	hl	hr	tl	tr
edge[0]	1	4	2	3
edge[1]	18	0	16	2
edge[2]	12	1	3	0
	⋮			

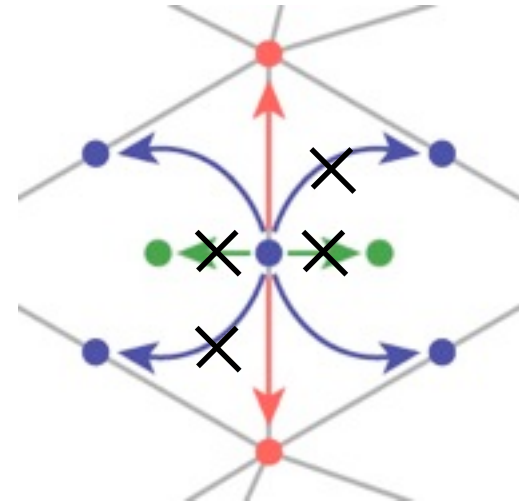


Winged-edge structure

- array of vertex positions: 12 bytes/vert
- array of 8-tuples of indices (per edge)
 - head/tail left/right edges + head/tail verts + left/right tris
 - $\text{int}[n_E][8]$: about 96 bytes per vertex
 - 3 edges per vertex (on average)
 - (8 indices x 4 bytes) per edge
- add a representative edge per vertex
 - $\text{int}[n_V]$: 4 bytes per vertex
- total storage: 112 bytes per vertex
 - but it is cleaner and generalizes to polygon meshes

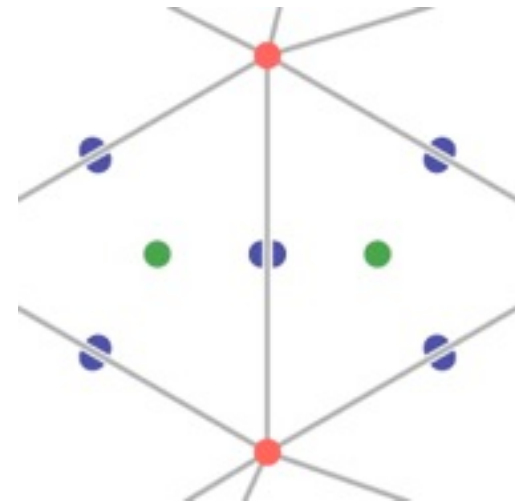
Winged-edge optimizations

- Omit faces if not needed
- Omit one edge pointer on each side
 - results in one-way traversal



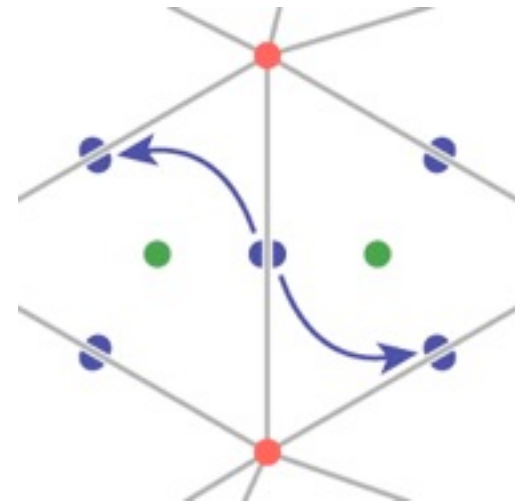
Half-edge structure

- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge



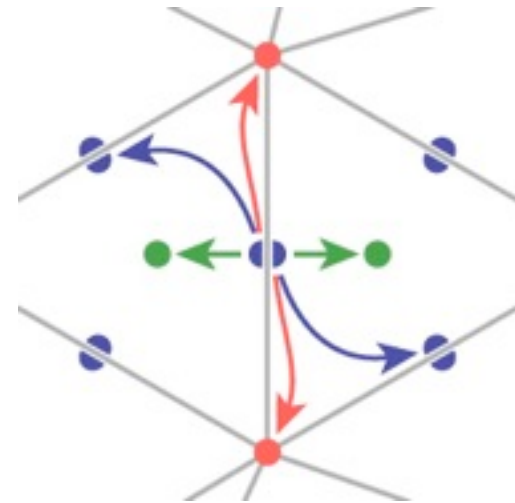
Half-edge structure

- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge



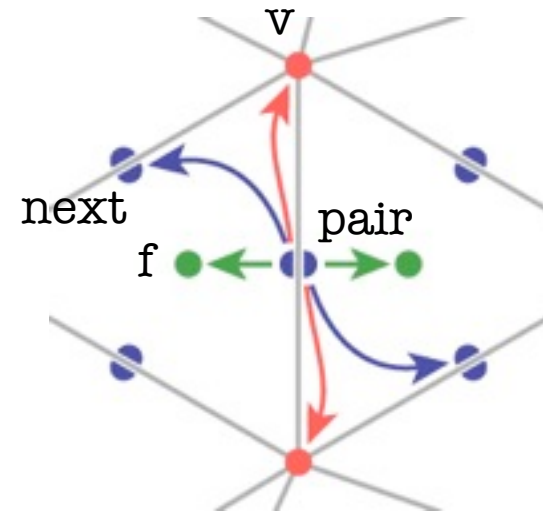
Half-edge structure

- Simplifies, cleans up winged edge
 - still works for polygon meshes
- Each half-edge points to:
 - next edge (left forward)
 - next vertex (front)
 - the face (left)
 - the opposite half-edge
- Each face or vertex points to one half-edge



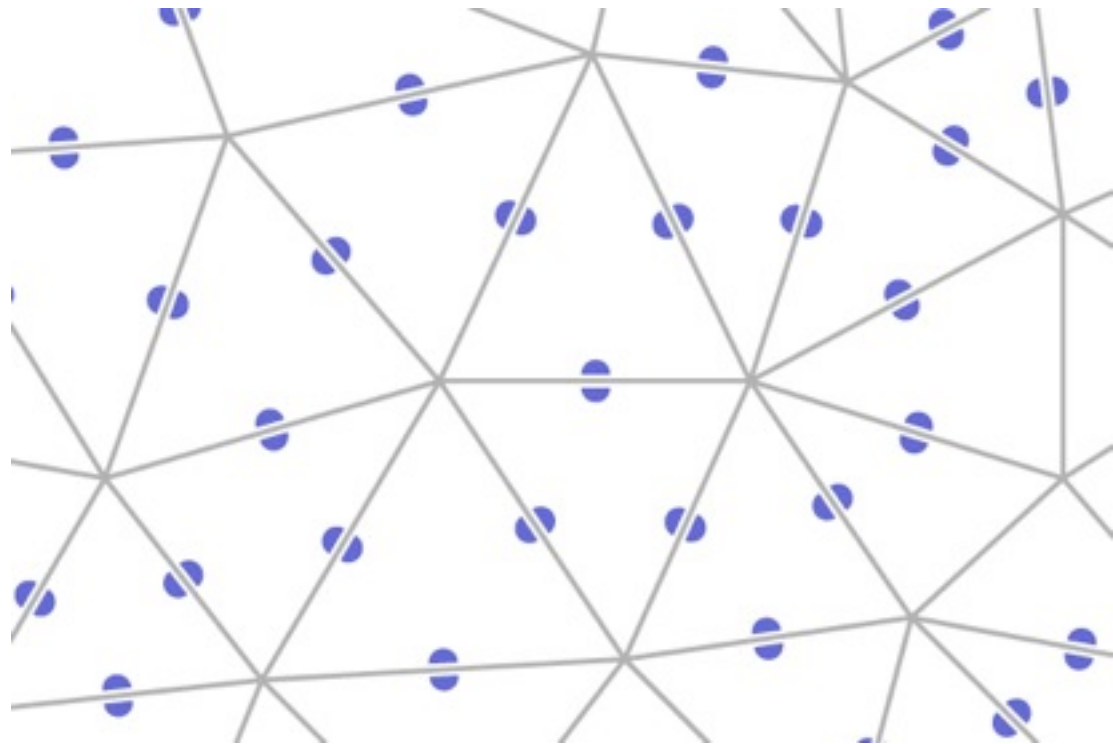
Half-edge structure

```
HEdge {  
    HEdge pair, next;  
    Vertex v;  
    Face f;  
}  
  
Face {  
    // per-face data  
    HEdge h; // any adjacent h-edge  
}  
  
Vertex {  
    // per-vertex data  
    HEdge h; // any incident h-edge  
}
```



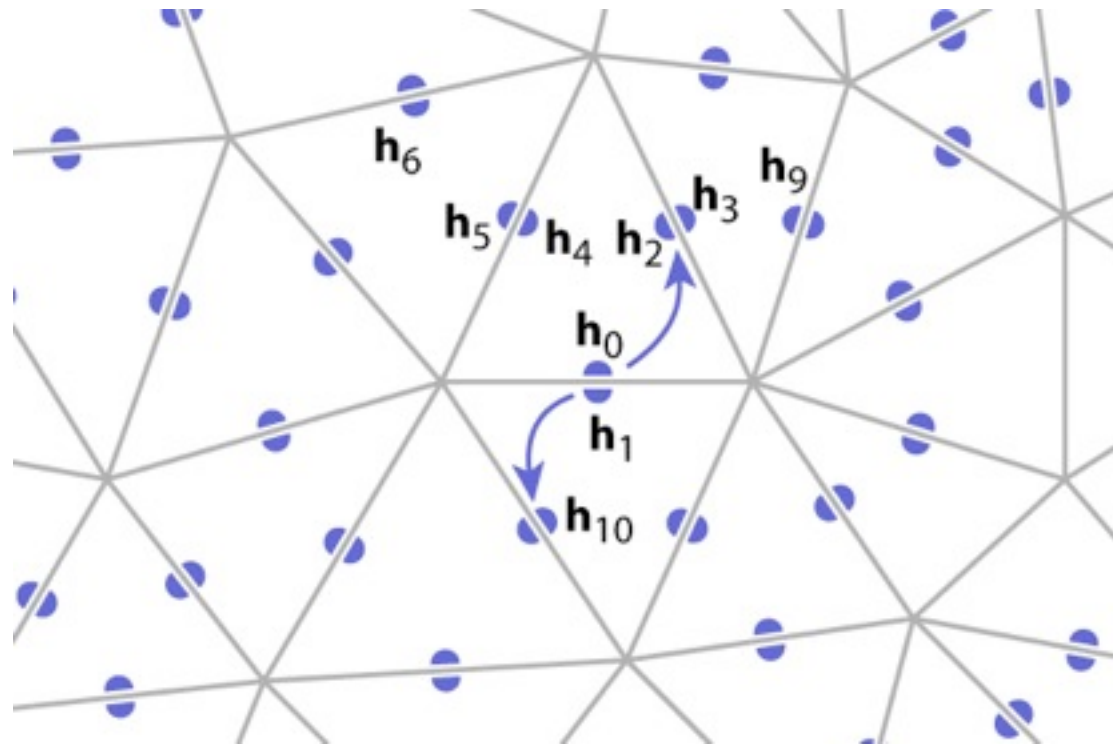
Half-edge structure

	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	⋮	



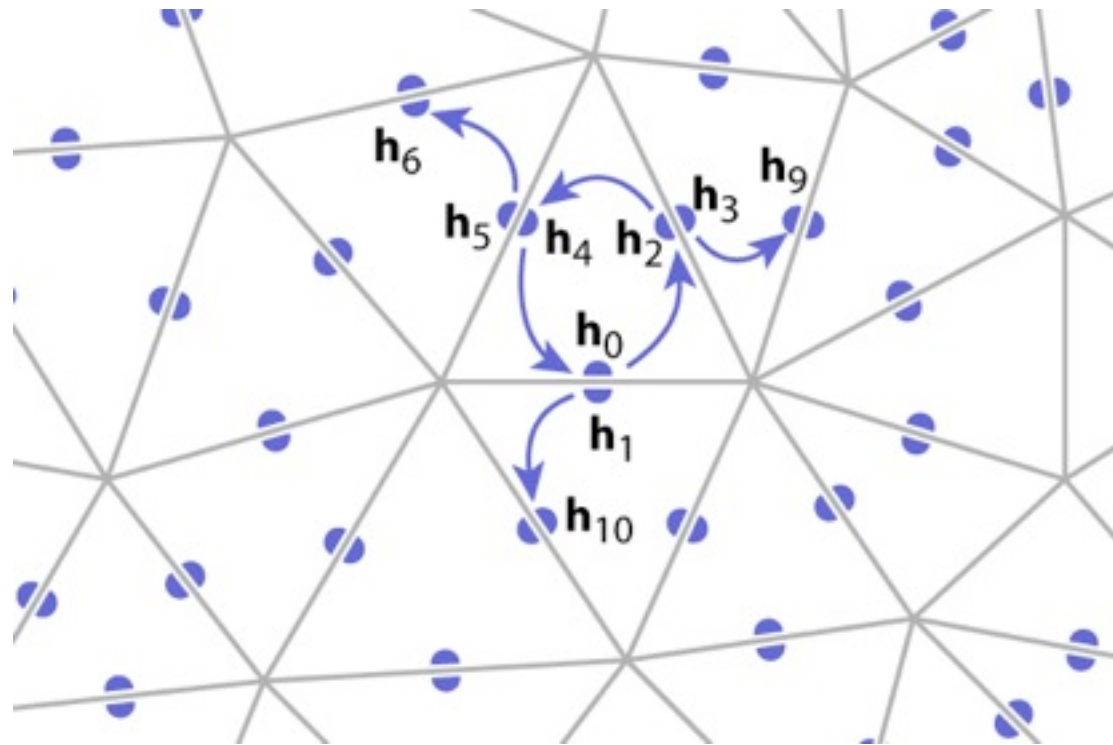
Half-edge structure

	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	⋮	



Half-edge structure

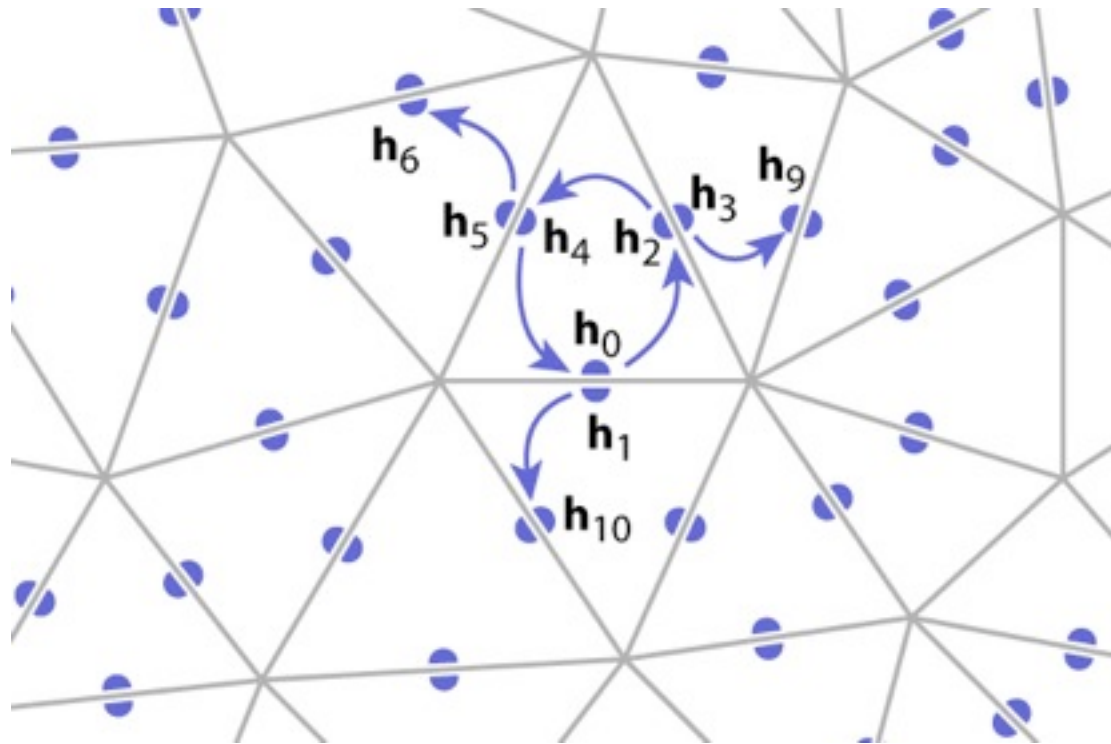
	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	⋮	



Half-edge structure

```
EdgesOfFace(f) {  
  h = f.h;  
  do {  
    h = h.next;  
  } while (h != f.h);  
}
```

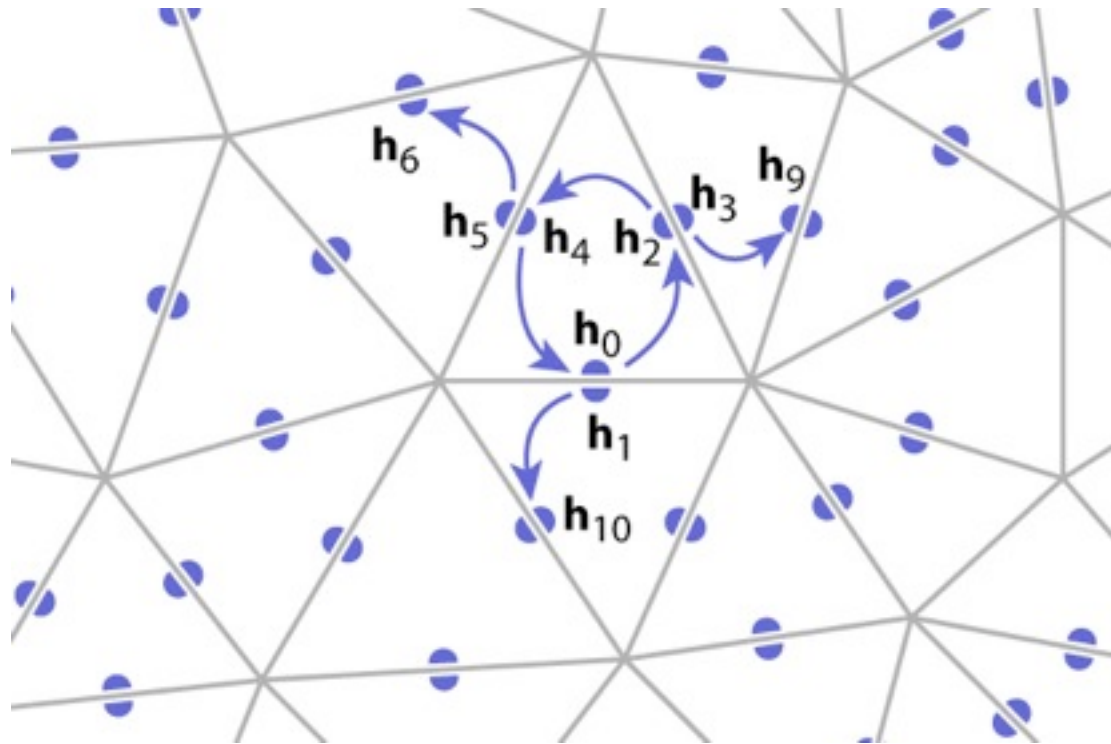
	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	⋮	



Half-edge structure

```
EdgesOfVertex(v) {  
  h = v.h;  
  do {  
    h = h.next.pair;  
  } while (h != v.h);  
}
```

	pair	next
hedge[0]	1	2
hedge[1]	0	10
hedge[2]	3	4
hedge[3]	2	9
hedge[4]	5	0
hedge[5]	4	6
	⋮	



Half-edge structure

- array of vertex positions: 12 bytes/vert
- array of 4-tuples of indices (per h-edge)
 - next, pair h-edges + head vert + left tri
 - $\text{int}[2n_E][4]$: about 96 bytes per vertex
 - 6 h-edges per vertex (on average)
 - (4 indices x 4 bytes) per h-edge
- add a representative h-edge per vertex
 - $\text{int}[n_V]$: 4 bytes per vertex
- total storage: 112 bytes per vertex

Half-edge optimizations

- Omit faces if not needed
- Use implicit pair pointers
 - they are allocated in pairs
 - they are even and odd in an array

