CS4450

Computer Networks: Architecture and Protocols

Lecture 15
BGP

Spring 2018

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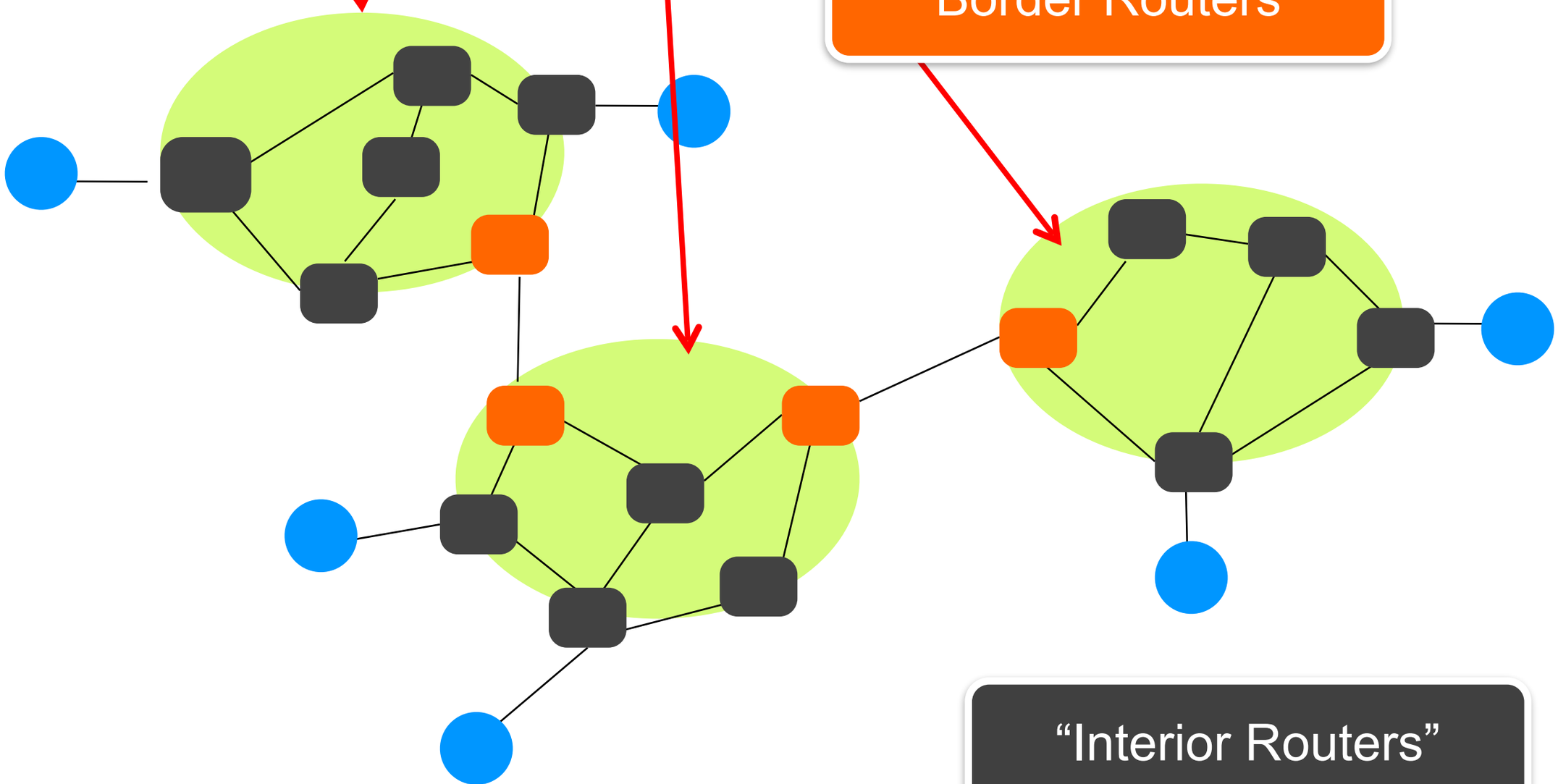
“Autonomous System (AS)” or “Domain”
Region of a network under a single administrative entity

“Border Routers”



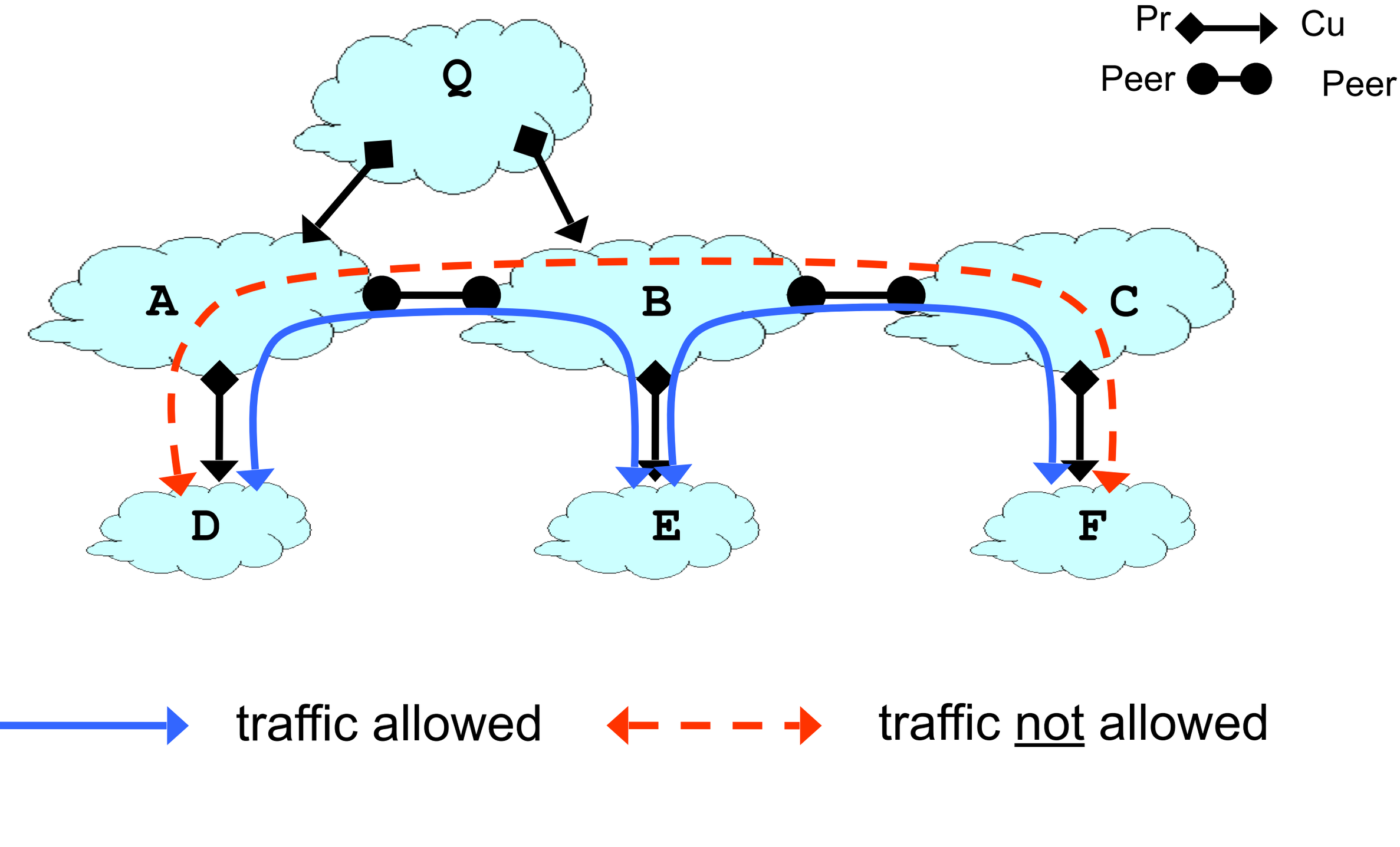
“Interior Routers”

Business Relationships Shape Topology and Policy

- Three basic kinds of relationships between ASes
  - AS A can be AS B's *customer*
  - AS A can be AS B's *provider*
  - AS A can be AS B's *peer*

- Business implications
  - Customer pays provider
  - Peers don't pay each other
    - Exchange roughly equal traffic

Routing Follows the Money



Interdomain Routing: Setup

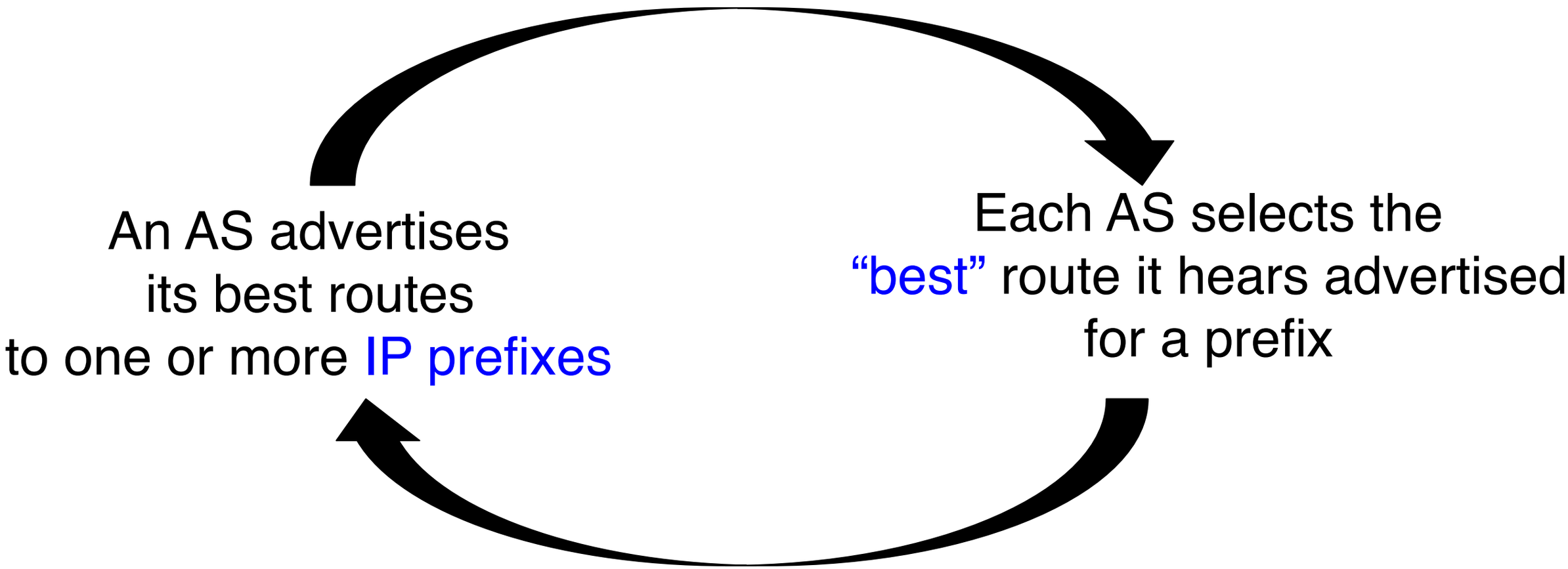
- Destinations are IP prefixes (12.0.0.0/8)

- Nodes are Autonomous Systems (ASes)
  - Internals of each AS are hidden

- Links represent both physical links and business relationships

- BGP (Border Gateway Protocol) is the Interdomain routing protocol
  - Implemented by AS border routers

BGP



Sound familiar?

BGP Inspired by Distance Vector

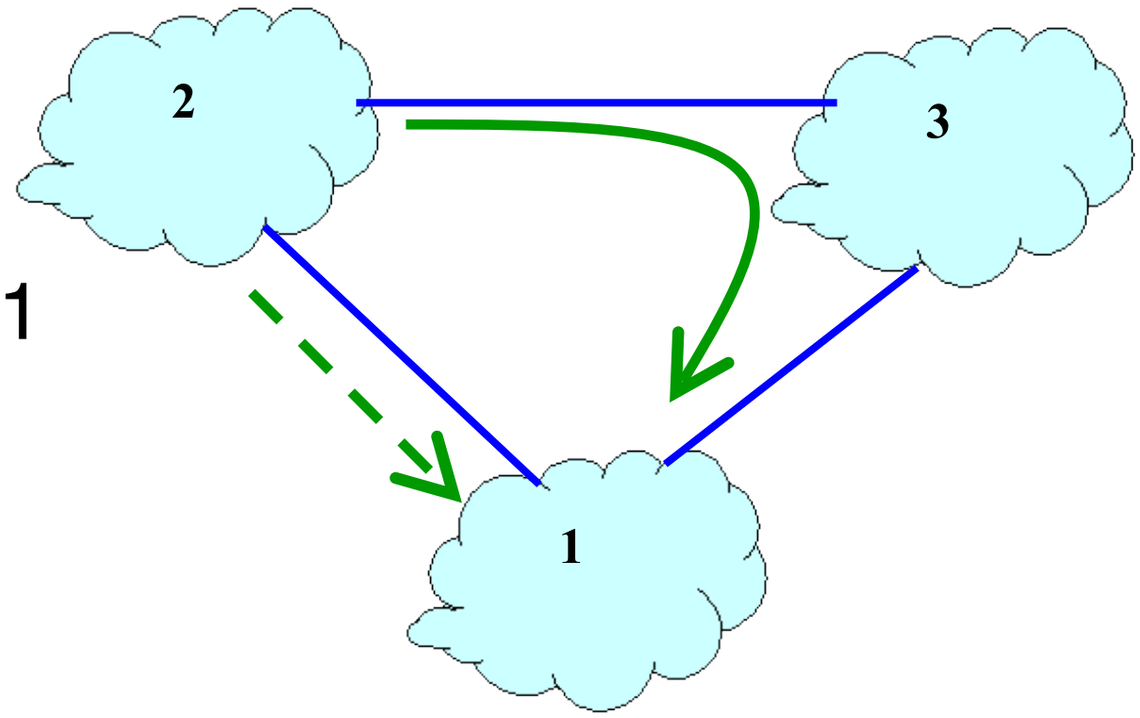
- Per-destination route advertisements
- No global sharing of network topology
- Iterative and distributed convergence on paths
- But, four key differences

BGP vs. DV

(1) BGP does not pick the shortest path routes!

- BGP selects route based on policy, not shortest distance/least cost

Node 2 may prefer 2, 3, 1 over 2, 1

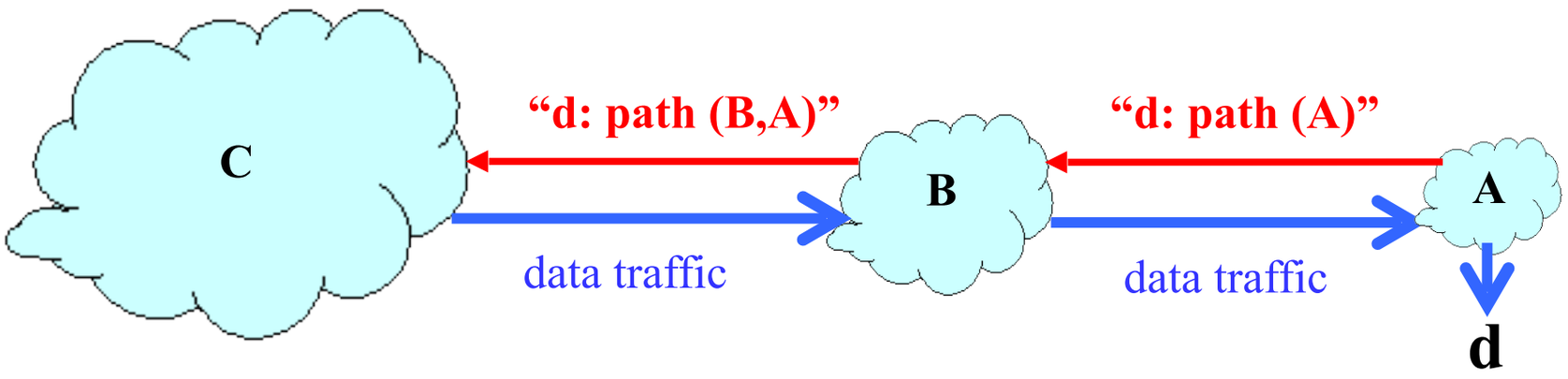


- How do we avoid loops?

BGP vs. DV

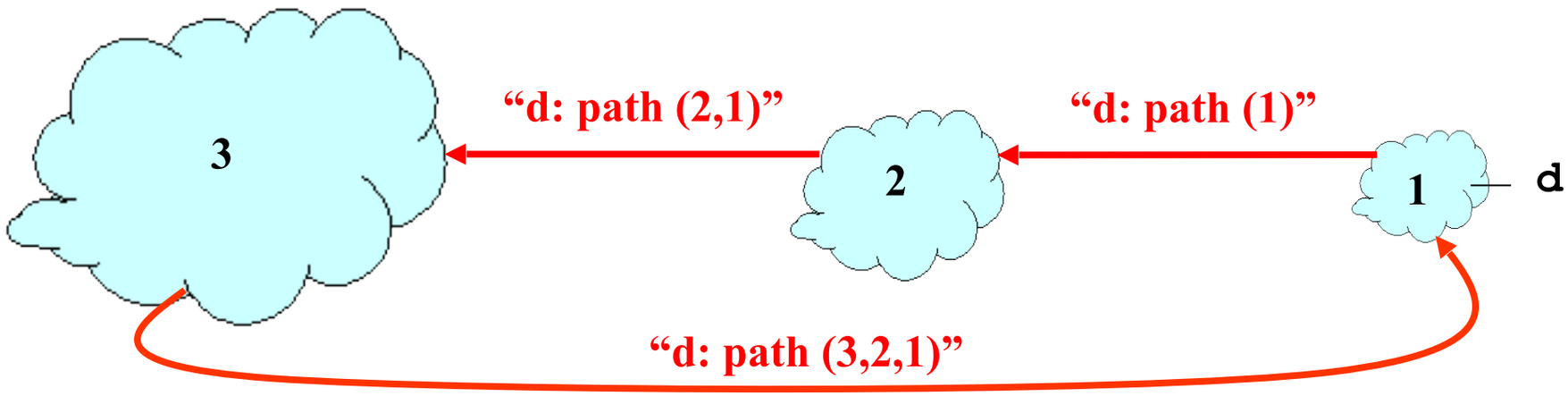
(2) Path-vector Routing

- Idea: advertise the entire path
  - Distance vector: send *distance metric* per dest. d
  - Path vector: send the *entire path* for each dest. d



Loop Detection with Path-Vector

- Node can easily detect a loop
  - Look for its own node identifier in the path
- Node can simply discard paths with loops
  - e.g. node 1 sees itself in the path 3, 2, 1



BGP vs. DV

(2) Path-vector Routing

- Idea: advertise the entire path
  - Distance vector: send *distance metric* per dest. d
  - Path vector: send the *entire path* for each dest. d

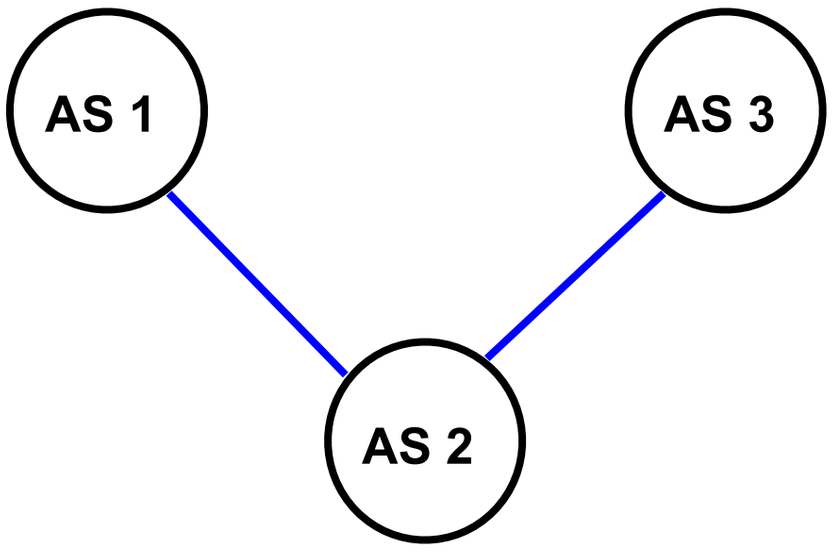
- Benefits
  - Loop avoidance is easy
  - Flexible policies based on entire path

BGP vs. DV

(3) Selective Route Advertisement

- For policy reasons, an AS may choose not to advertise a route to a destination

- As a result, reachability is not guaranteed even if the graph is connected

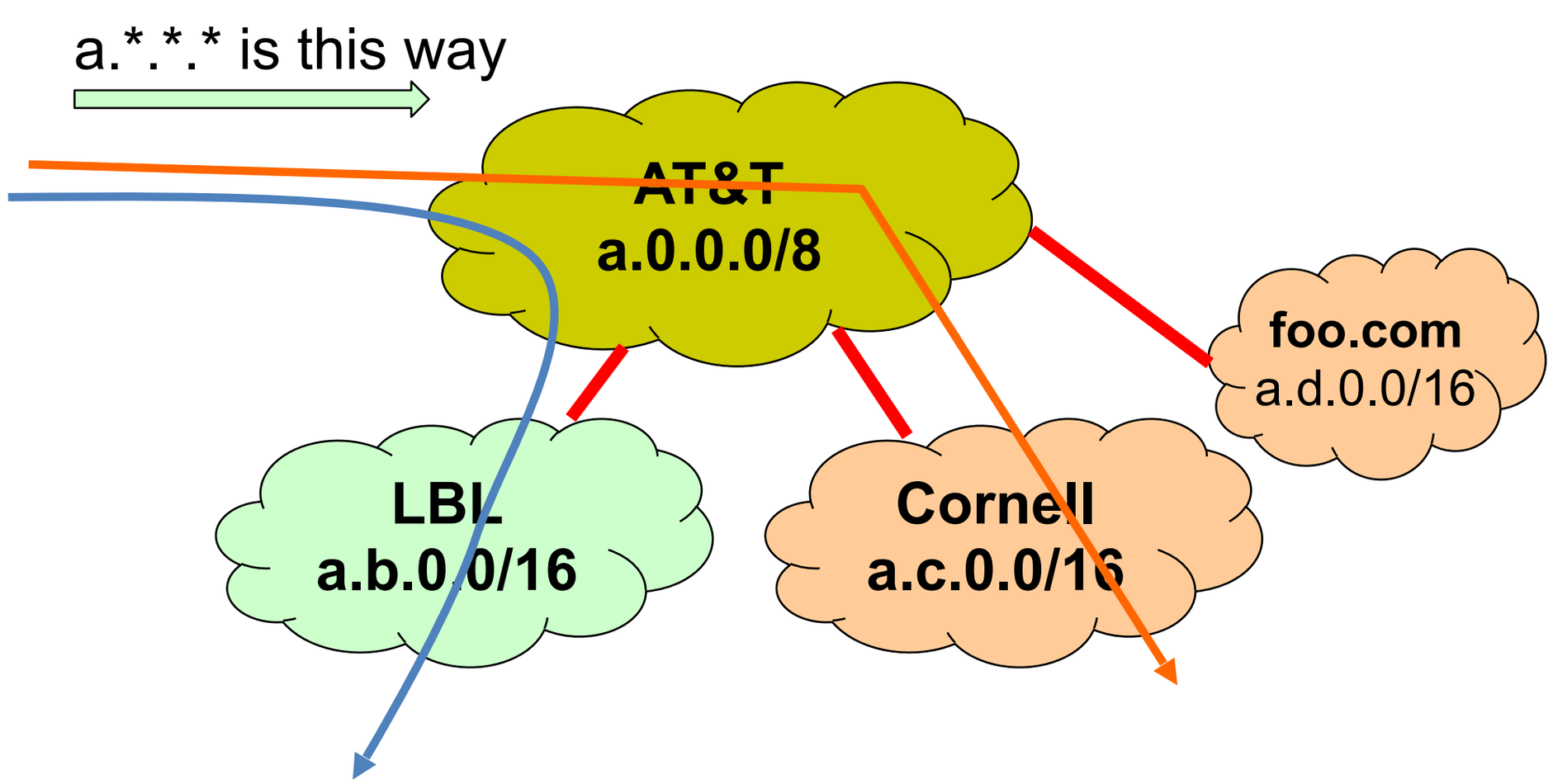


Example: AS#2 does not want to carry traffic between AS#1 and AS#3

BGP vs. DV

(4) BGP may aggregate routes

- For scalability, BGP may aggregate routes for different prefixes



BGP Outline

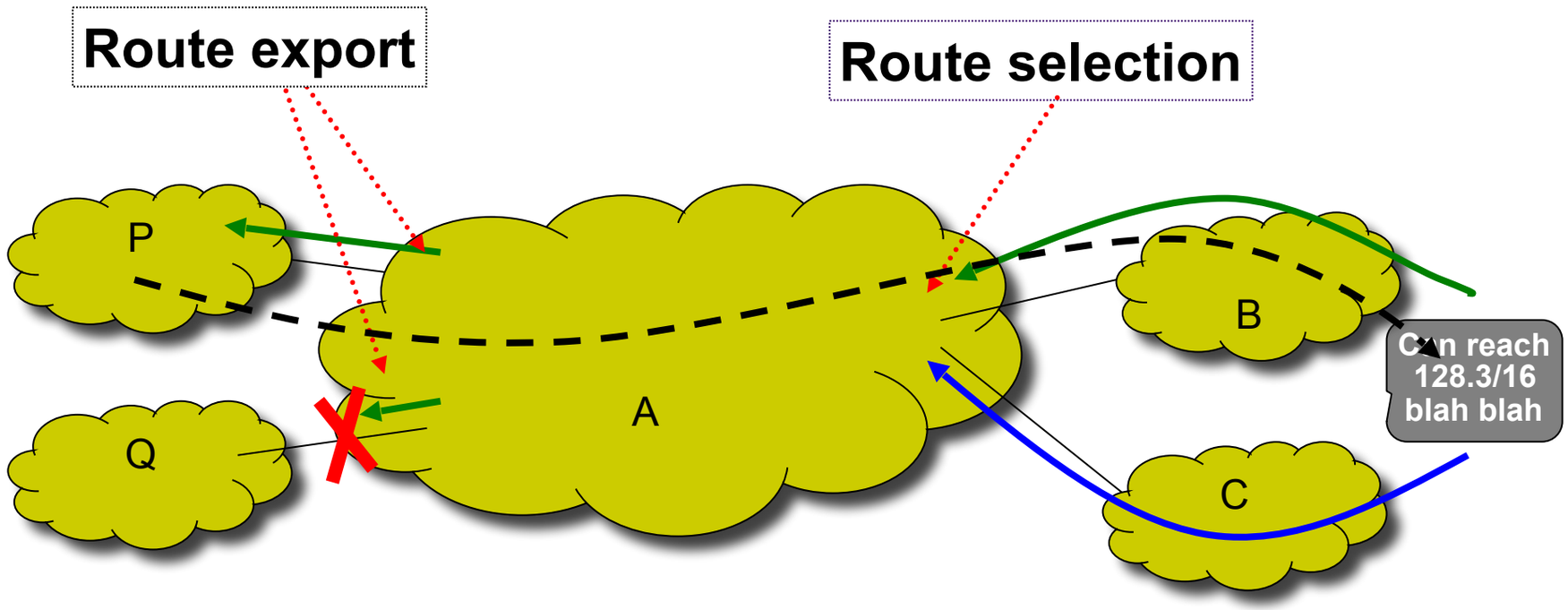
- BGP Policy
  - Typical policies and implementation

- BGP protocol details

- Issues with BGP

Policy:

Imposed in how routes are selected and exported



- **Selection**: Which path to use
  - Controls whether / how traffic leaves the network
- **Export**: Which path to advertise
  - Controls whether / how traffic enters the network

Typical Selection Policy

- In decreasing order of priority:
  1. Make or save money (send to customer > peer > provider)
  2. Maximize performance (smallest AS path length)
  3. Minimize use of my network bandwidth (“hot potato”)
  4. …

Typical Export Policy

Destination prefix advertised by…	Export route to…
Customer	Everyone (providers, peers, other customers)
Peer	Customers
Provider	Customers

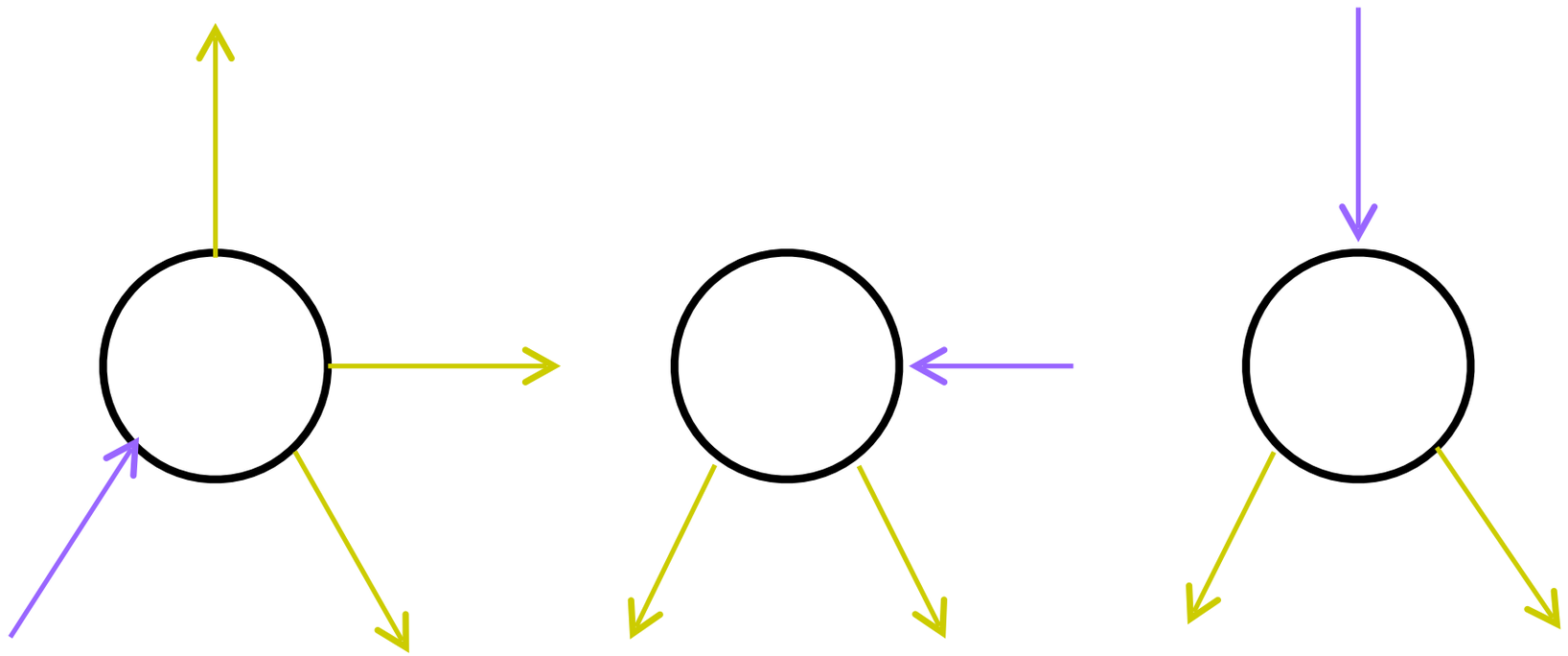
Known as the “Gao-Rexford” rules
Capture common (but not required!) practice

Gao-Rexford

providers

peers

customers



With Gao-Rexford, the AS policy graph is a
DAG (directed acyclic graph) and routes are “valley free”

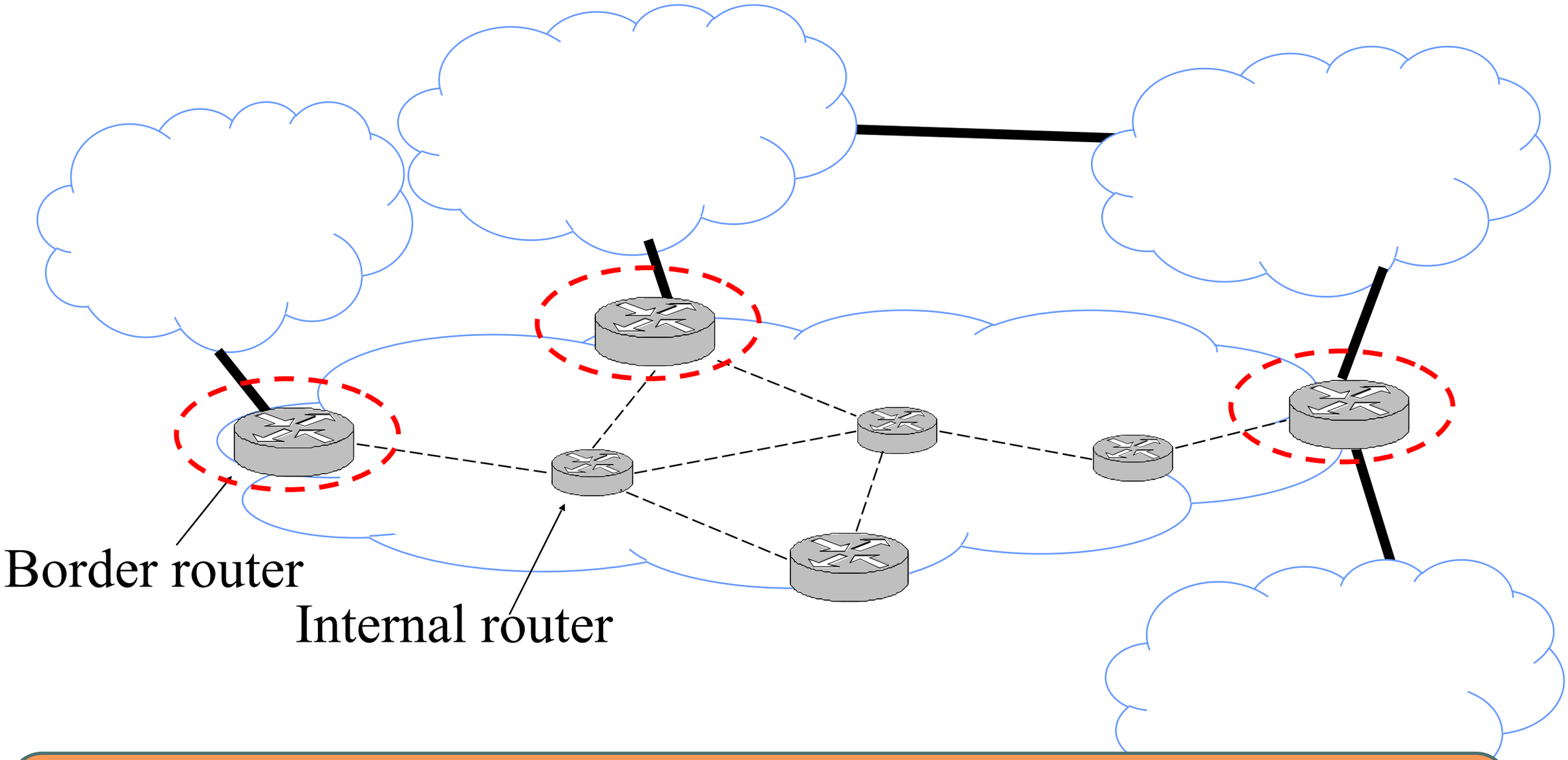
BGP Outline

- BGP Policy
  - Typical policies and implementation

- **BGP protocol details**

- Issues with BGP

Who speaks BGP?



Border routers at an Autonomous System

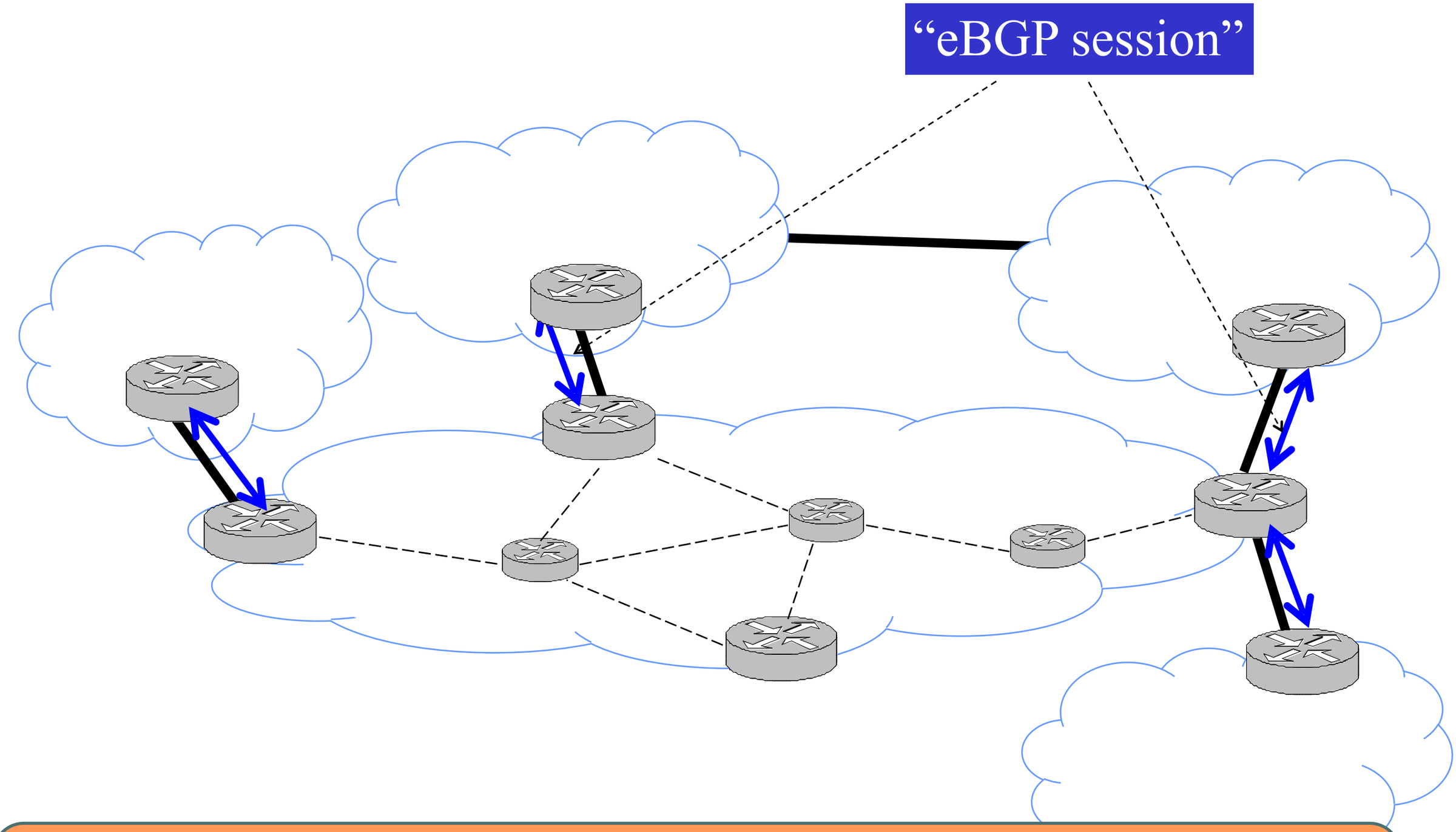
What Does “speak BGP” Mean?

- Implement the BGP Protocol Standard
  - Internet Engineering Task Force (IETF) RFC 4271

- Specifies what messages to exchange with other BGP “speakers”
  - Message types (e.g. route advertisements, updates)
  - Message syntax

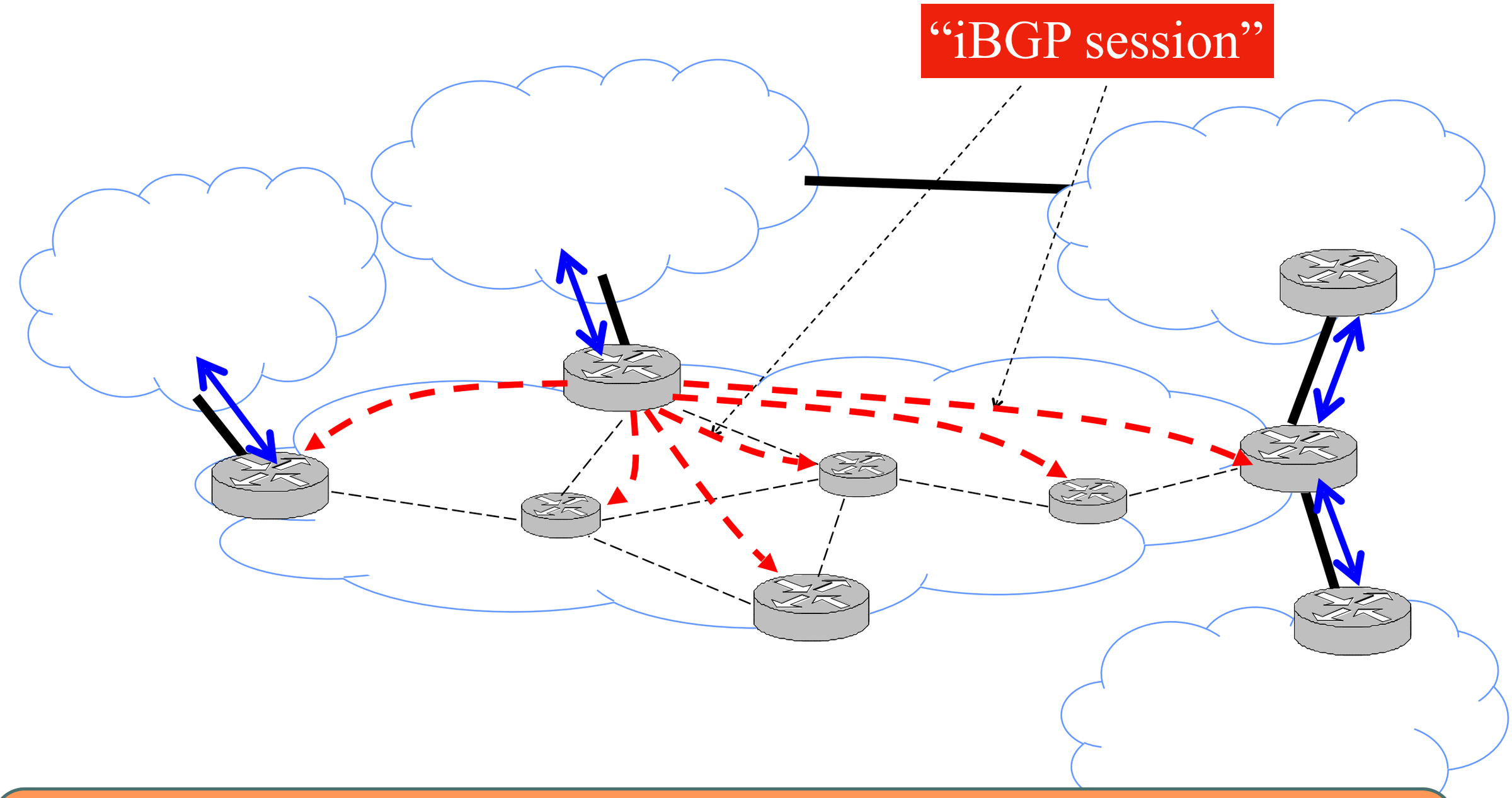
- Specifies how to process these messages
  - When you receive a BGP update, do x
  - Follows BGP state machine in the protocol spec and policy decisions, etc.

BGP Sessions



A border router speaks BGP with
border routers in other ASes

BGP Sessions



A border router speaks BGP with other (interior and border) routers in its own AS

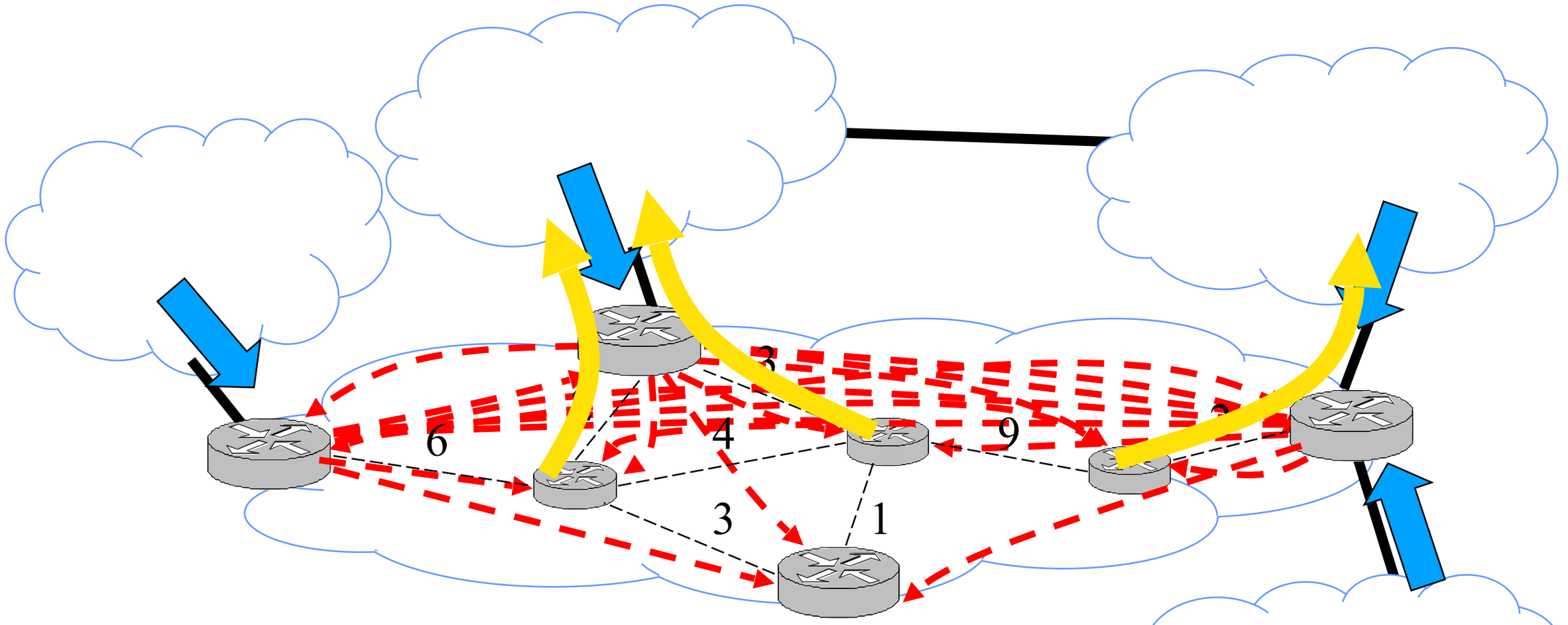
eBGP, iBGP, IGP

- eBGP: BGP sessions between border routers in different ASes
  - Learn routes to external destinations

- iBGP: BGP sessions between border routers and other routers within the same AS
  - Distribute externally learned routes internally

- IGP: Interior Gateway Protocol = Intradomain routing protocol
  - Provides internal reachability
  - e.g. OSPF, RIP

Putting the Pieces Together



1. Provide internal reachability (**IGP**) - - - - - - -
2. Learn routes to external destinations (**eBGP**)
3. Distribute externally learned routes internally (**iBGP**)
4. Travel shortest path to egress (IGP)

Basic Messages in BGP

- **Open**
  - Establishes BGP session
  - BGP uses TCP
- **Notification**
  - Report unusual conditions
- **Update**
  - Inform neighbor of new routes
  - Inform neighbor of old routes that become inactive
- **Keepalive**
  - Inform neighbor that connection is still viable

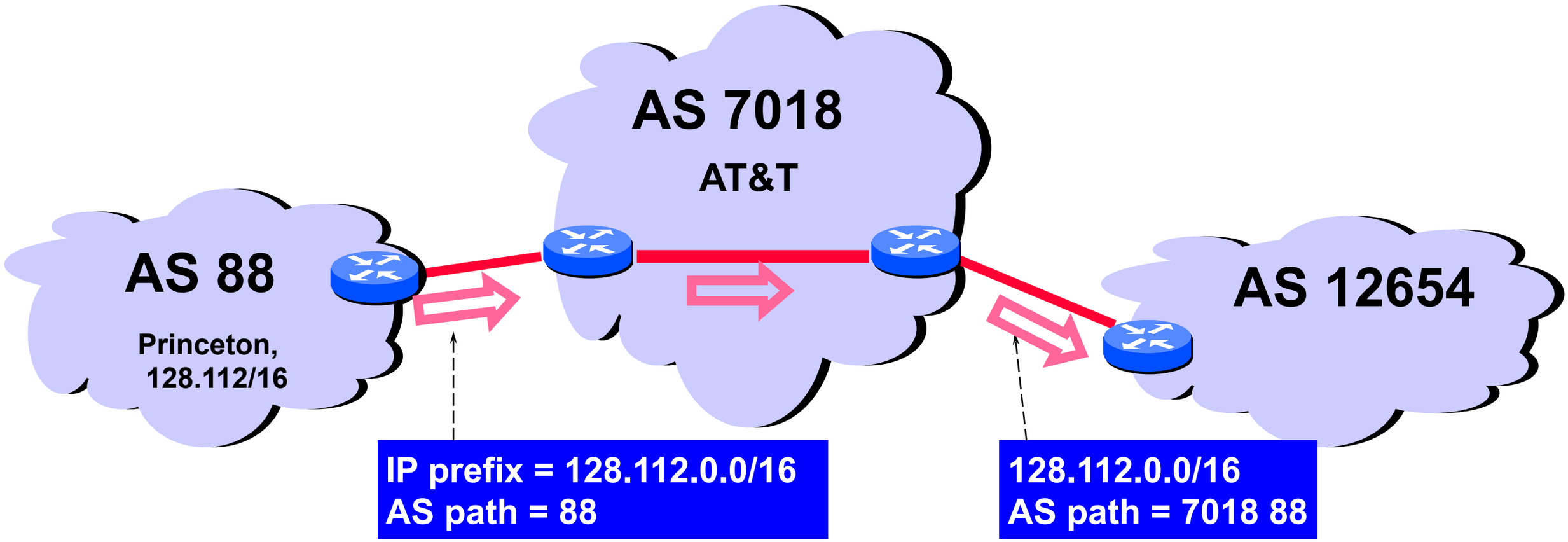
Route Updates

- Format: *<IP prefix: route attributes>*
- Two kinds of updates:
  - Announcements: new routes or changes to existing routes
  - Withdrawals: remove routes that no longer exist

- Route Attributes
  - Describe routes, used in selection/export decisions
  - Some attributes are local
    - i.e. private within an AS, not included in announcements
  - Some attributes are propagated with eBGP route announcements
  - Many standardized attributes in BGP

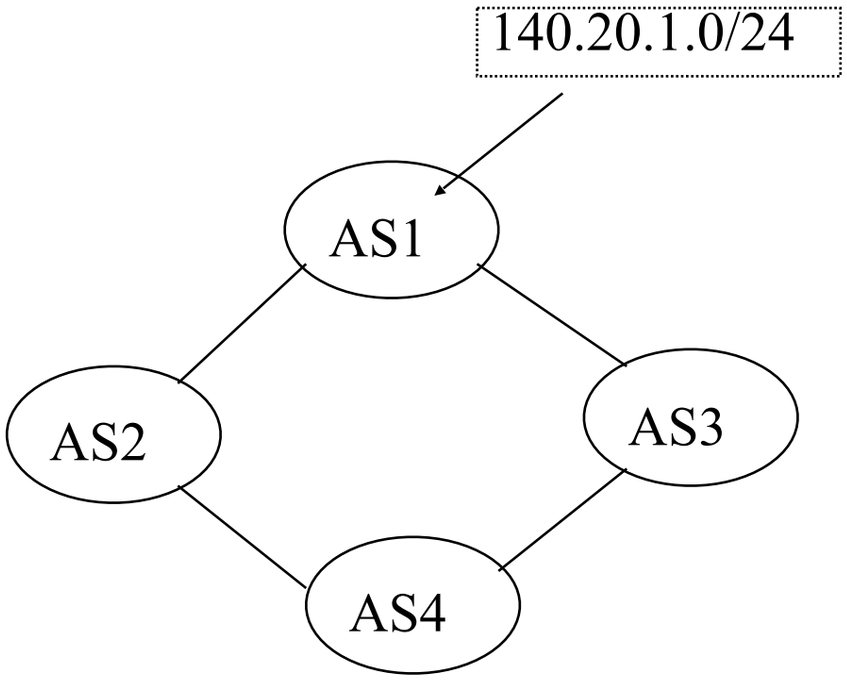
Route Attributes (1): ASPATH

- Carried in route announcements
- Vector that lists all the ASes a route advertisement has traversed (in reverse order)



Route Attributes (2): LOCAL PREF

- "Local Preference"
- Used to choose between different AS paths
- The higher the value, the more preferred
- Local to an AS; carried only in iBGP messages

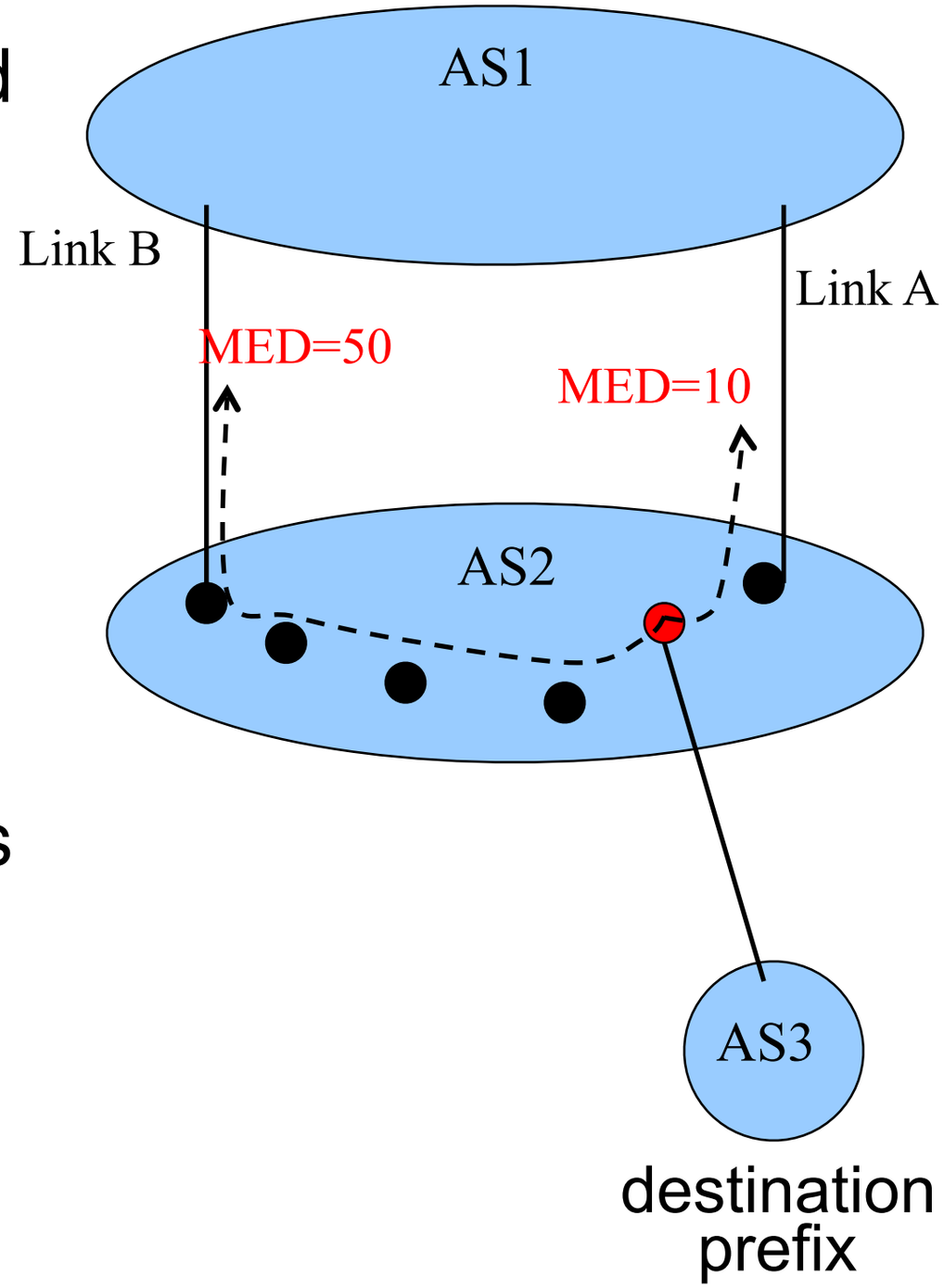


BGP table at AS4:

Destination	AS Path	Local Pref
140.20.1.0/24	**AS3 AS1**	**300**
140.20.1.0/24	**AS2 AS1**	**100**

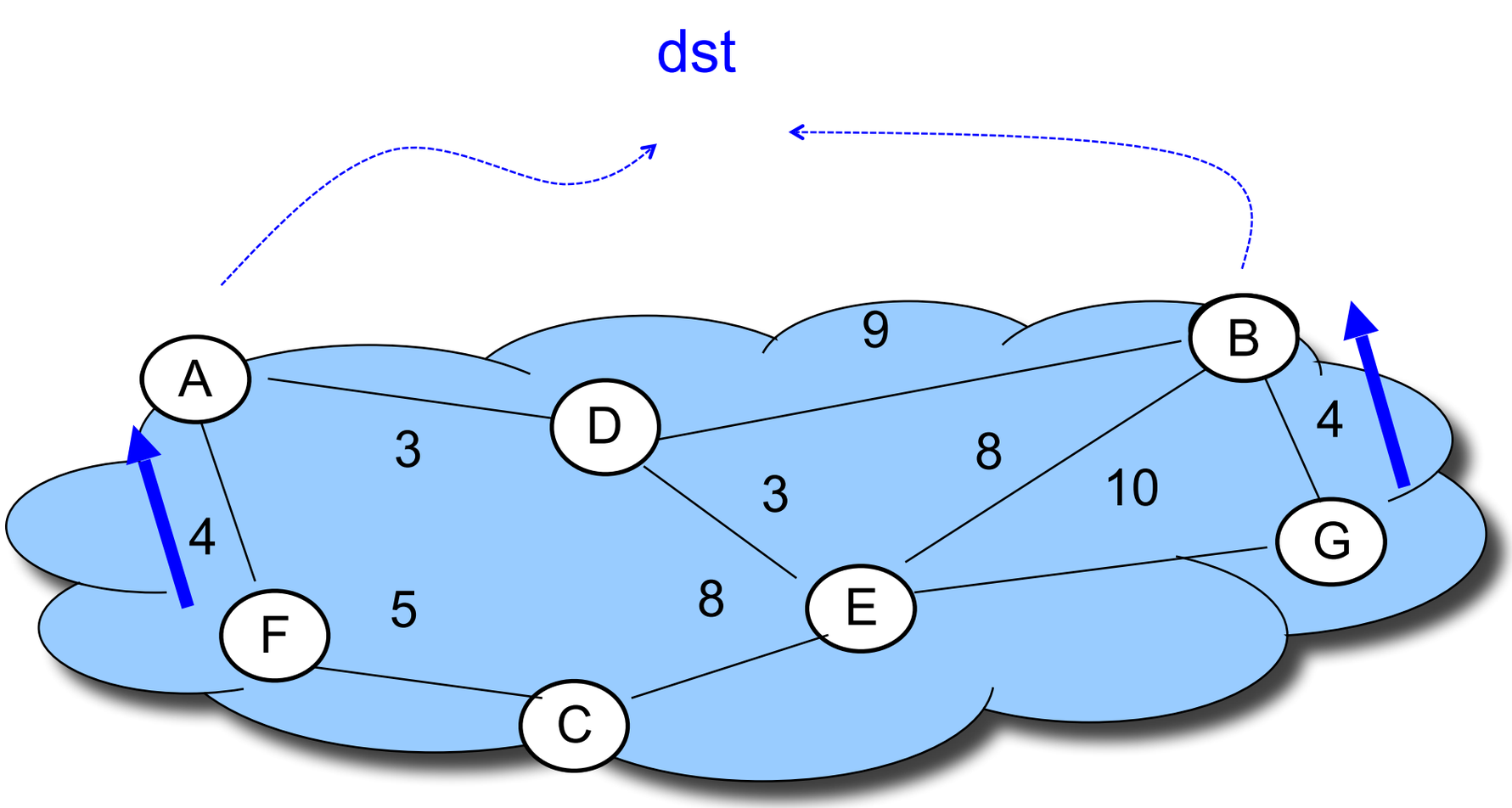
Route Attributes (3) : MED

- "Multi-Exit Discriminator"
- Used when ASes are interconnected via two or more links
  - Specifies how close a prefix is to the link it is announced on
- Lower is better
- AS announcing prefix sets MED
- AS receiving prefix (optionally!) uses MED to select link



Route Attributes (4): IGP Cost

- Used for hot-potato routing
  - Each router selects the closest egress point based on the path cost in intra-domain protocol



Using Attributes

- Rules for route selection in priority order

1. Make or save money (send to customer > peer > provider)
2. Maximize performance (smallest AS path length)
3. Minimize use of my network bandwidth ("hot potato")
4. …

Using Attributes

- Rules for route selection in priority order

Priority	Rule	Remarks
1	`LOCAL PREF`	Pick highest `LOCAL PREF`
2	`ASPATH`	Pick shortest `ASPATH` length
3	`MED`	Lowest `MED` preferred
4	eBGP > iBGP	Did AS learn route via eBGP (preferred) or iBGP?
5	iBGP path	Lowest IGP cost to next hop (egress router)
6	Router ID	Smallest next-hop router's IP address as tie-breaker

BGP Update Processing

Open ended programming.
Constrained only by vendor configuration language

Control plane

BGP Updates

Apply Import Policies

Best Route Selection

Best Route Table

Apply Export Policies

BGP Updates

Data plane

forwarding Entries

Data packets

IP Forwarding Table

Data packets

BGP Outline

- BGP Policy
  - Typical policies and implementation

- BGP protocol details

- **Issues with BGP**

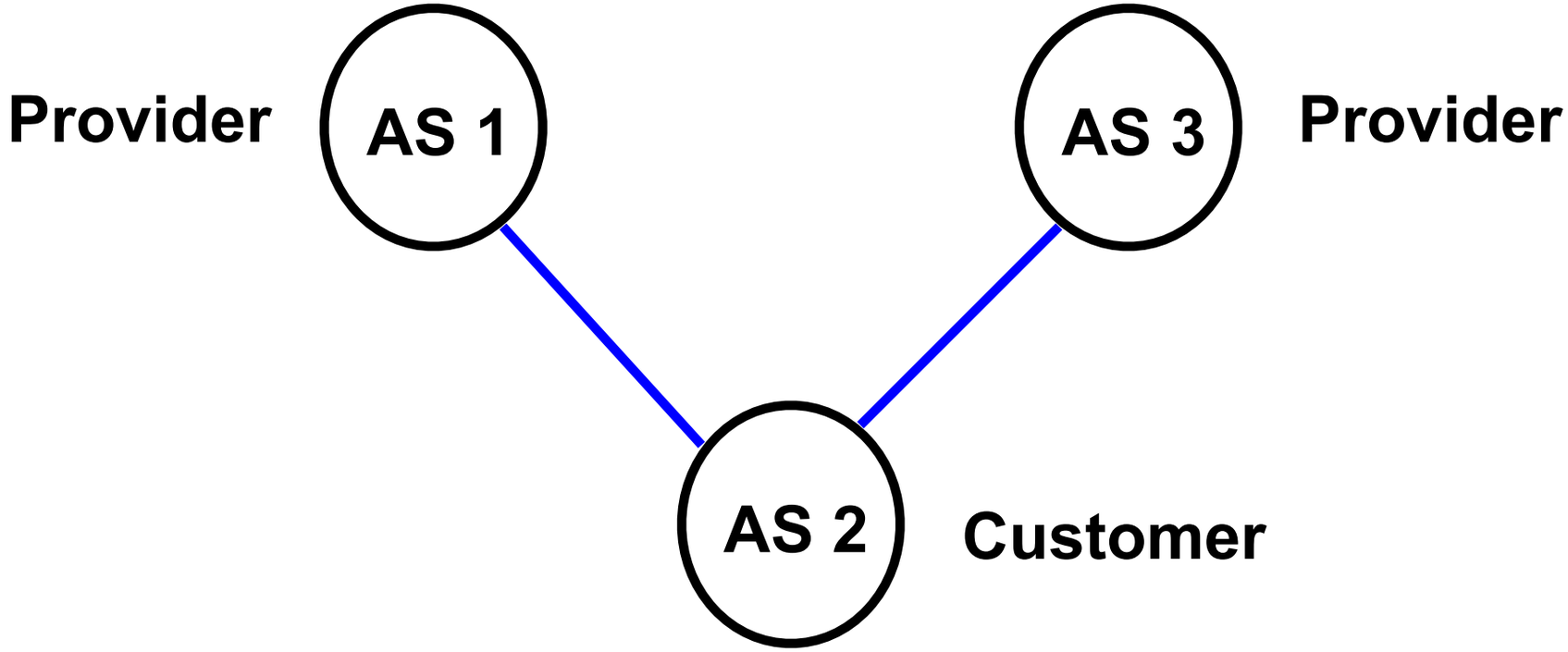
BGP: Issues

- Reachability
- Security
- Convergence
- Performance
- Anomalies

Reachability

- In normal routing, if graph is connected then reachability is assured

- With policy routing, this doesn't always hold



Security

- An AS can claim to serve a prefix that they actually don't have a route to (blackholing traffic)
  - Problem not specific to policy or path vector
  - Important because of AS autonomy
  - *Fixable: make ASes prove they have a path*

- But…
  - AS may forward packets along a route different from what is advertised
    - Tell customers about a fictitious short path…
    - Much harder to fix!

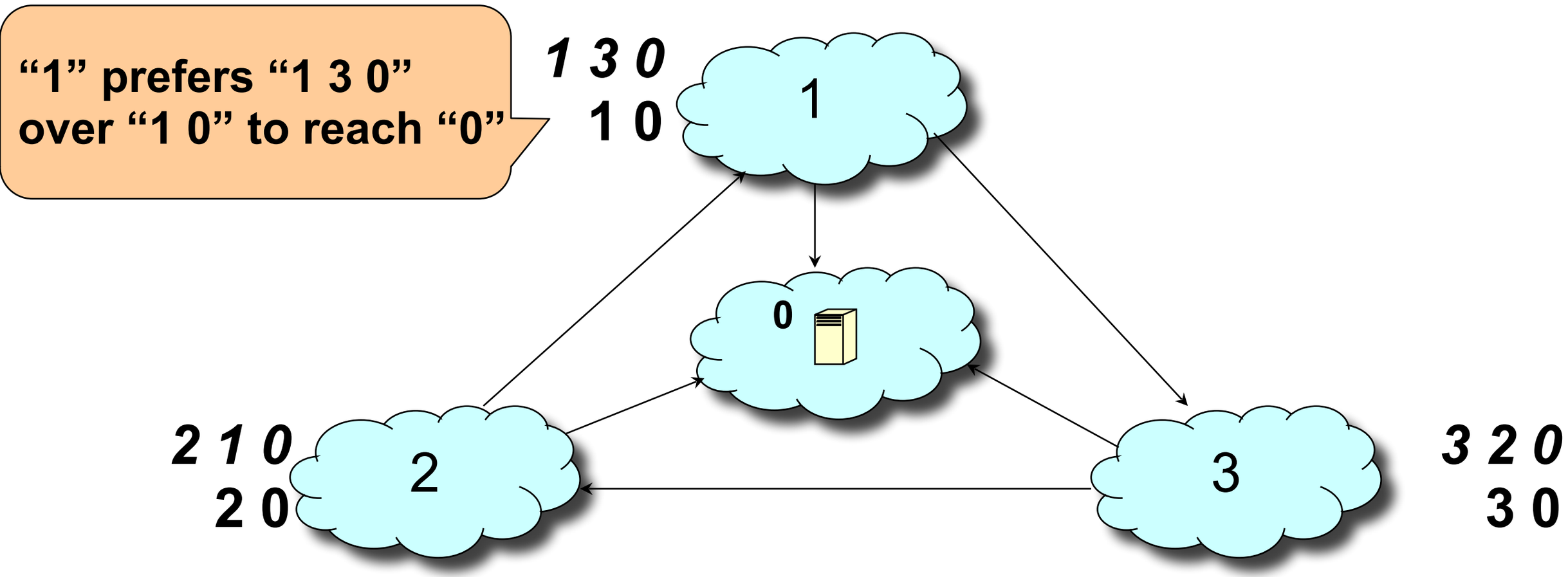
Convergence

- If all AS policies follow Gao-Rexford rules,
  - Then BGP is guaranteed to converge (safety)

- For arbitrary policies, BGP may fail to converge!

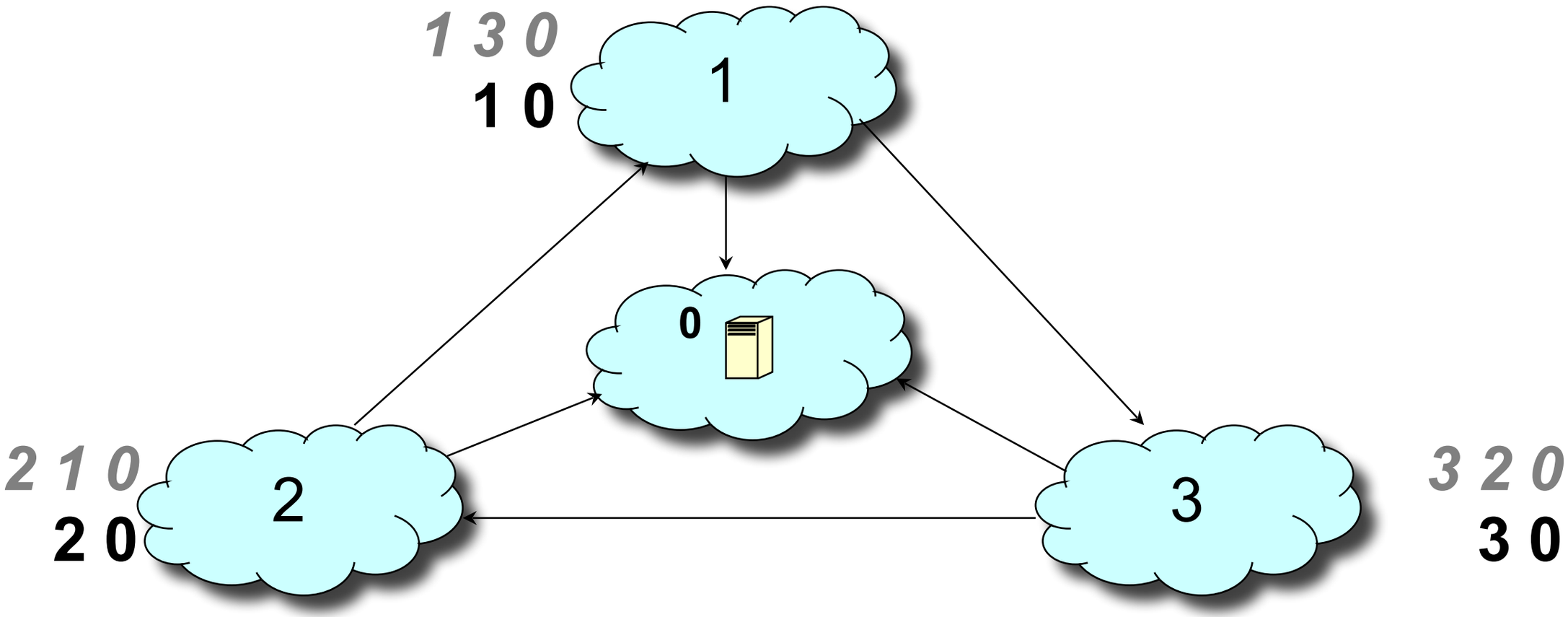
Example of Policy Oscillation

"1" prefers "1 3 0"
over "1 0" to reach "0"



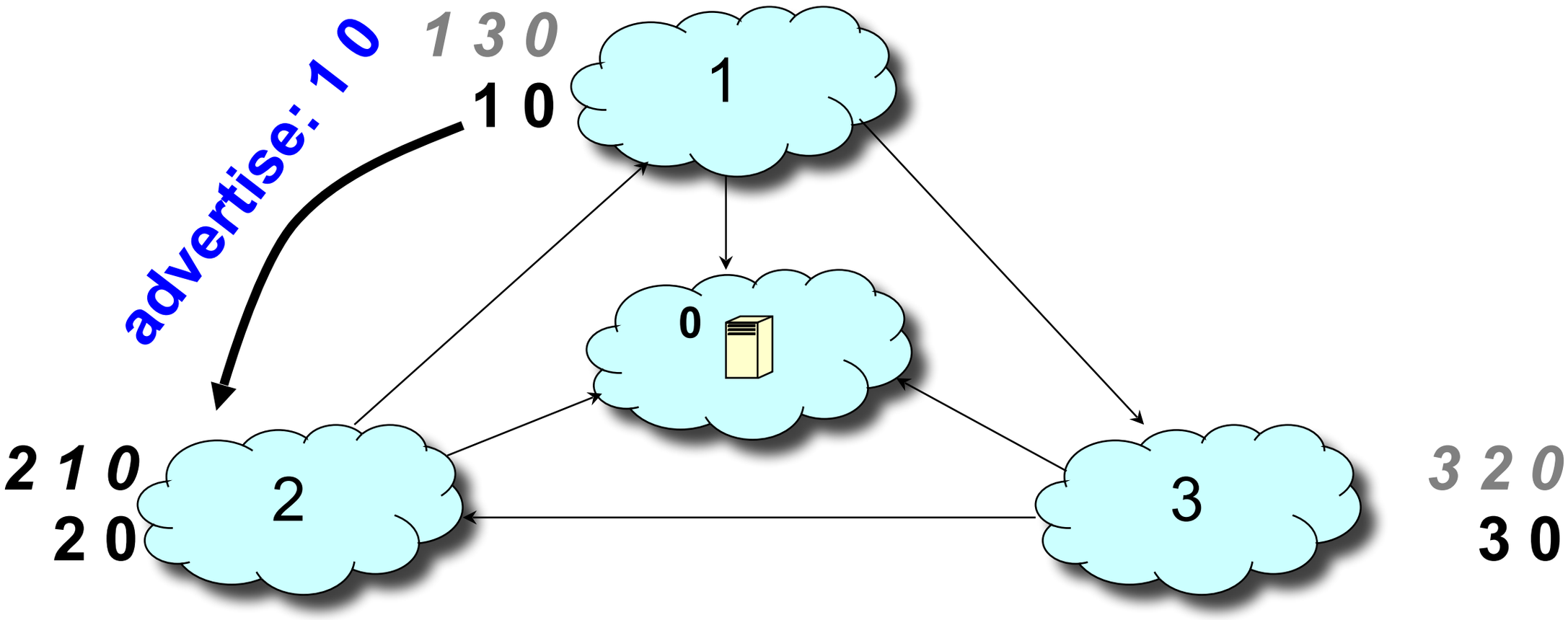
Step-by-step Policy Oscillation

Initially: nodes 1, 2, 3 know only shortest path to 0

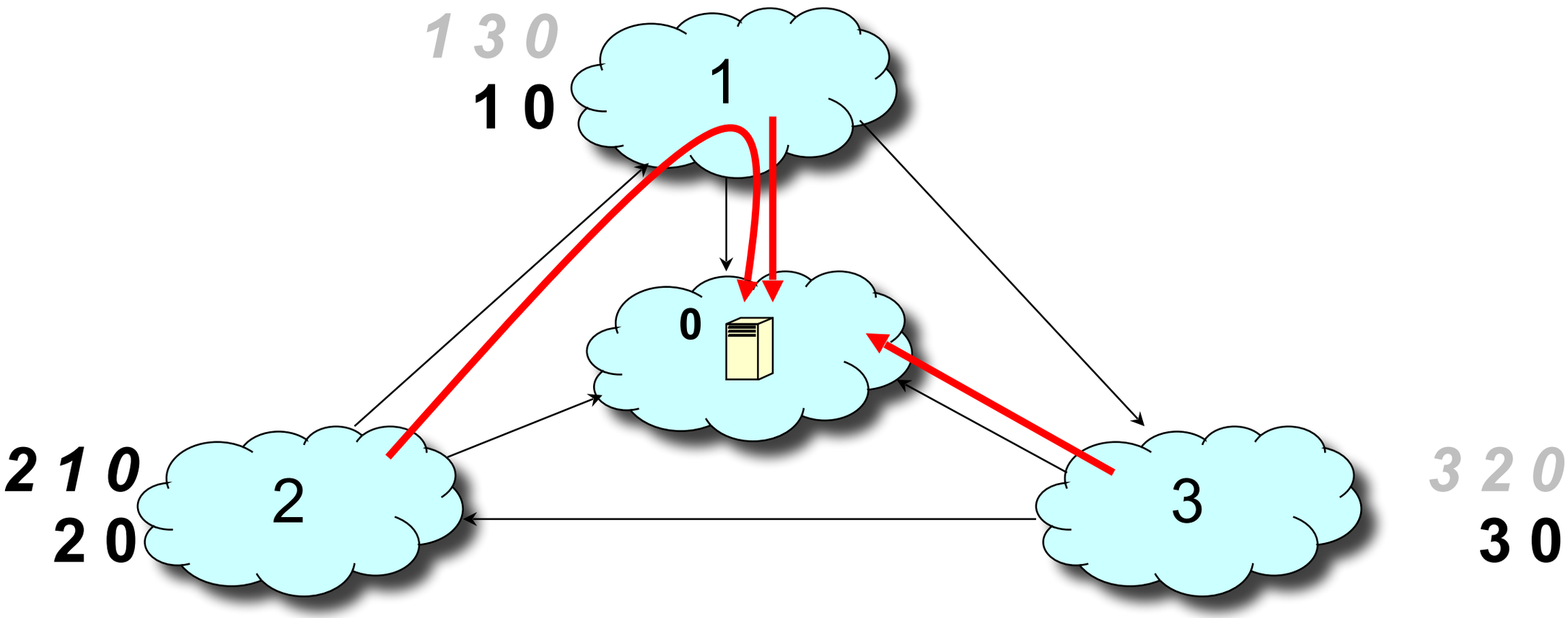


Step-by-step Policy Oscillation

1 **advertises** its path 1 0 to 2

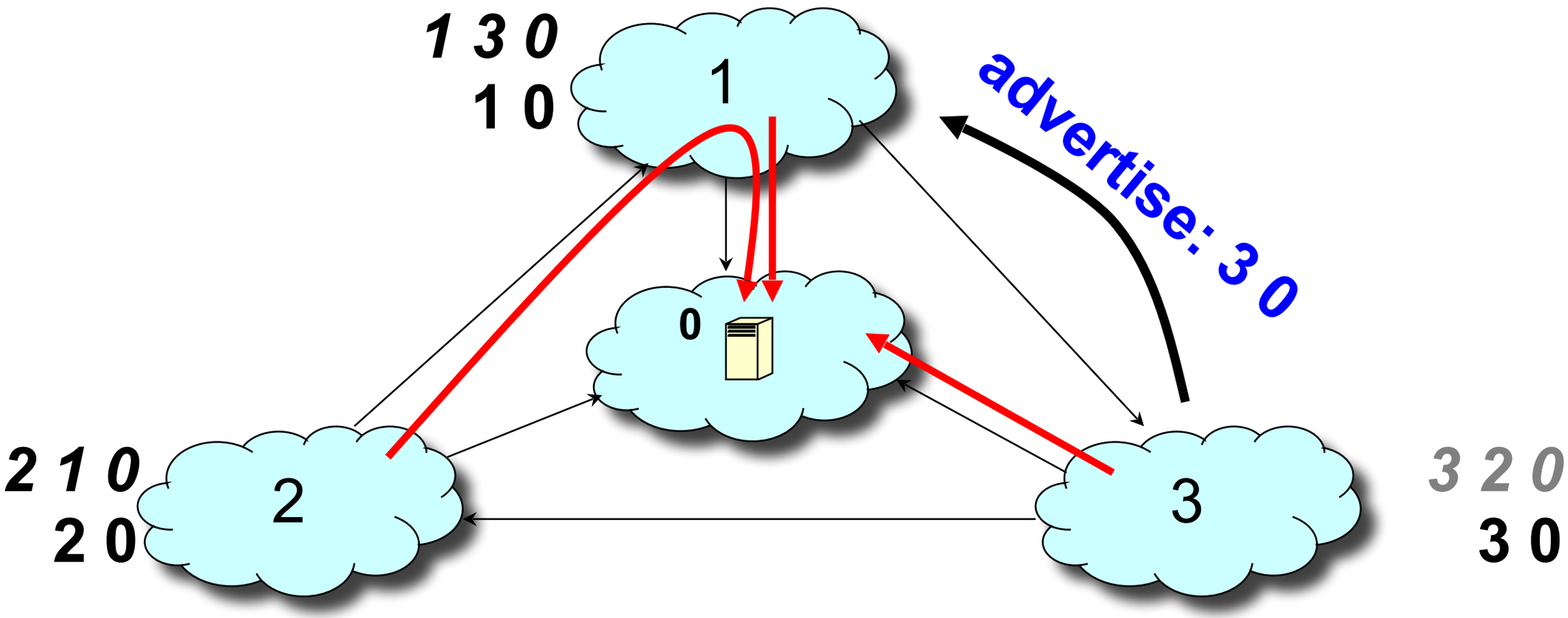


Step-by-step Policy Oscillation

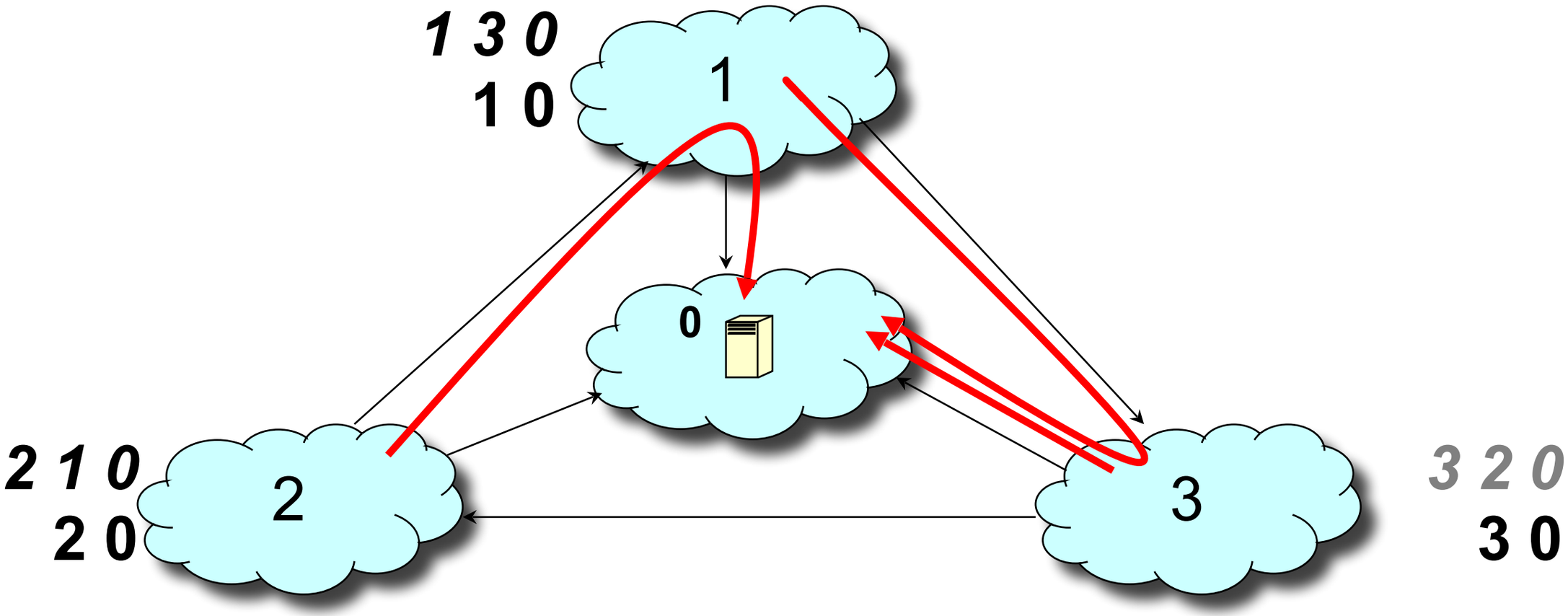


Step-by-step Policy Oscillation

3 **advertises** its path 3 0 to 1

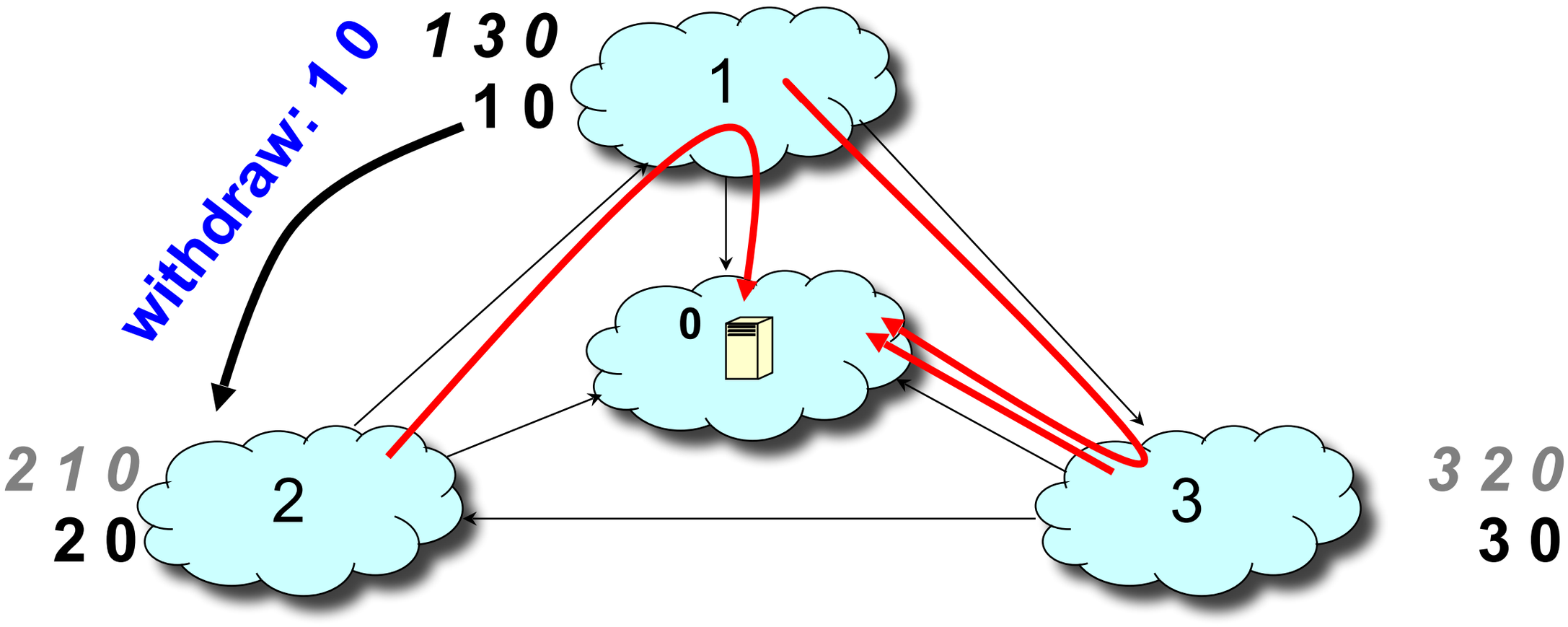


Step-by-step Policy Oscillation

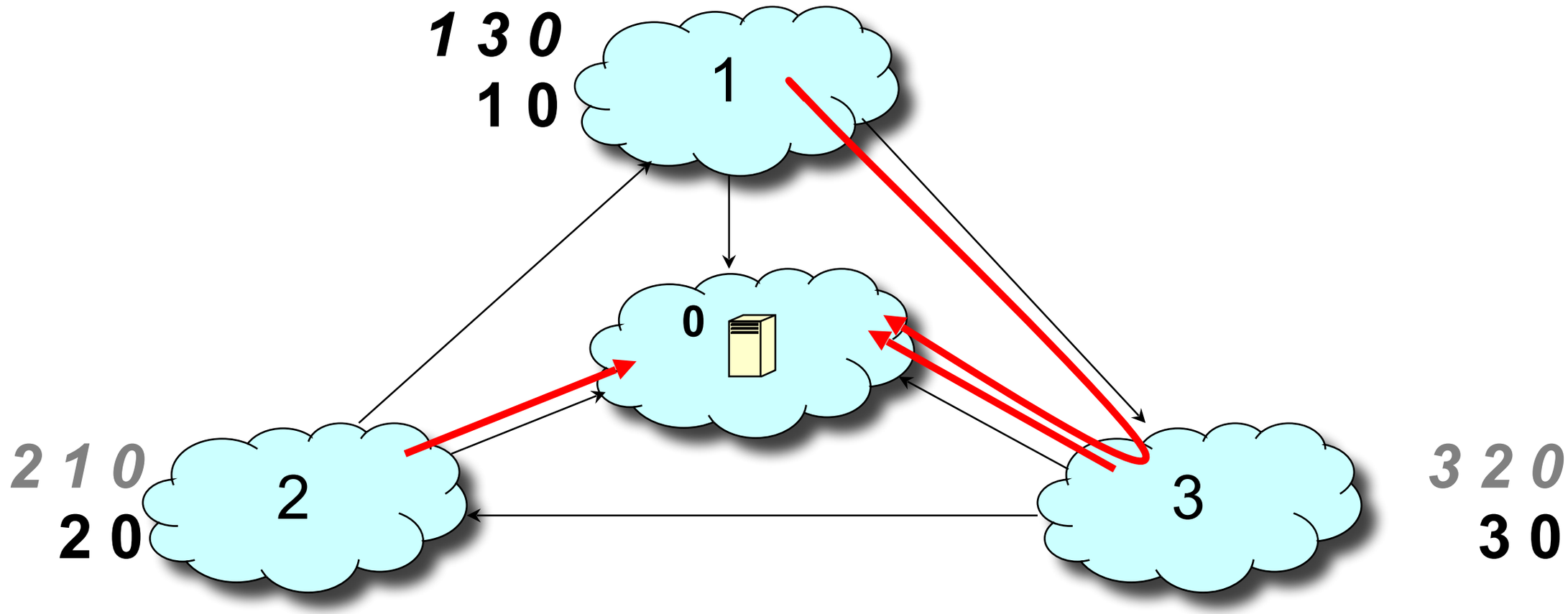


Step-by-step Policy Oscillation

1 **withdraws** its path 1 0 from 2

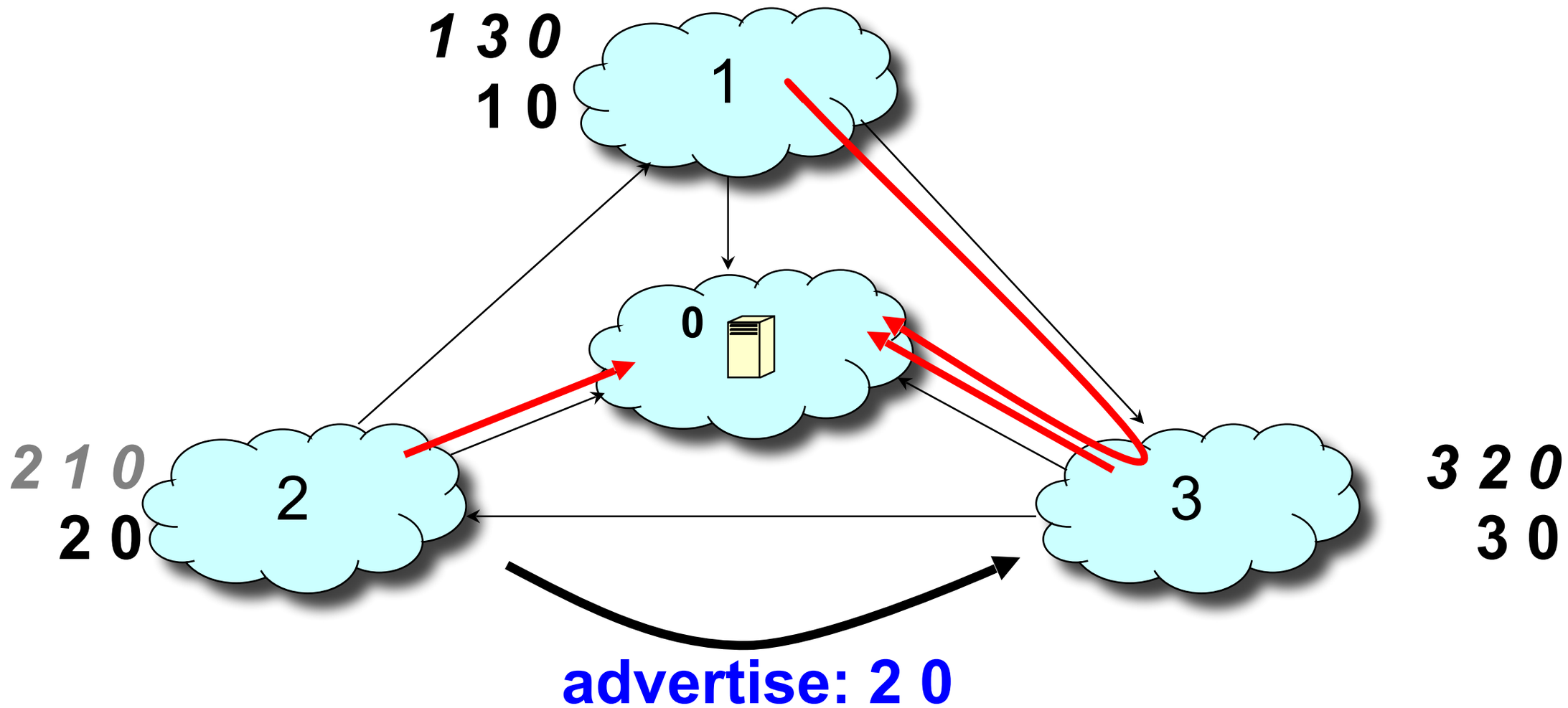


Step-by-step Policy Oscillation

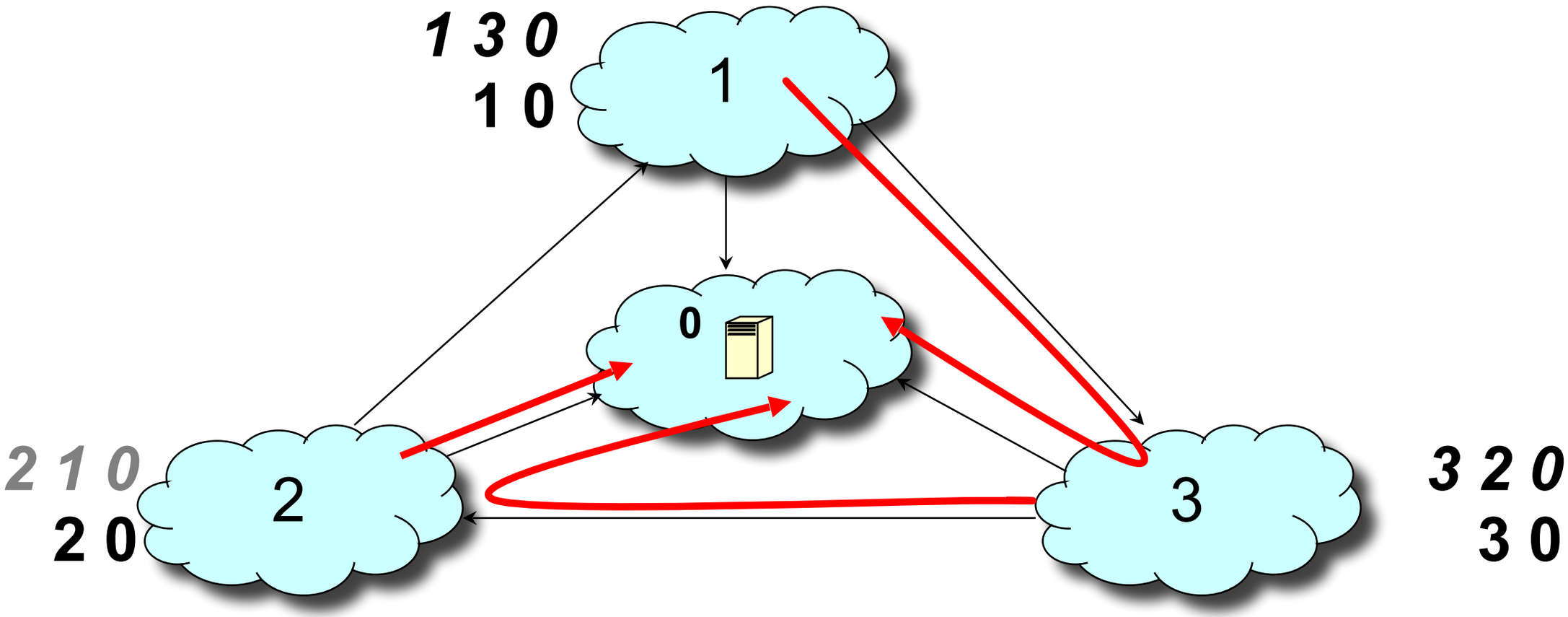


Step-by-step Policy Oscillation

2 **advertises** its path 2 0 to 3

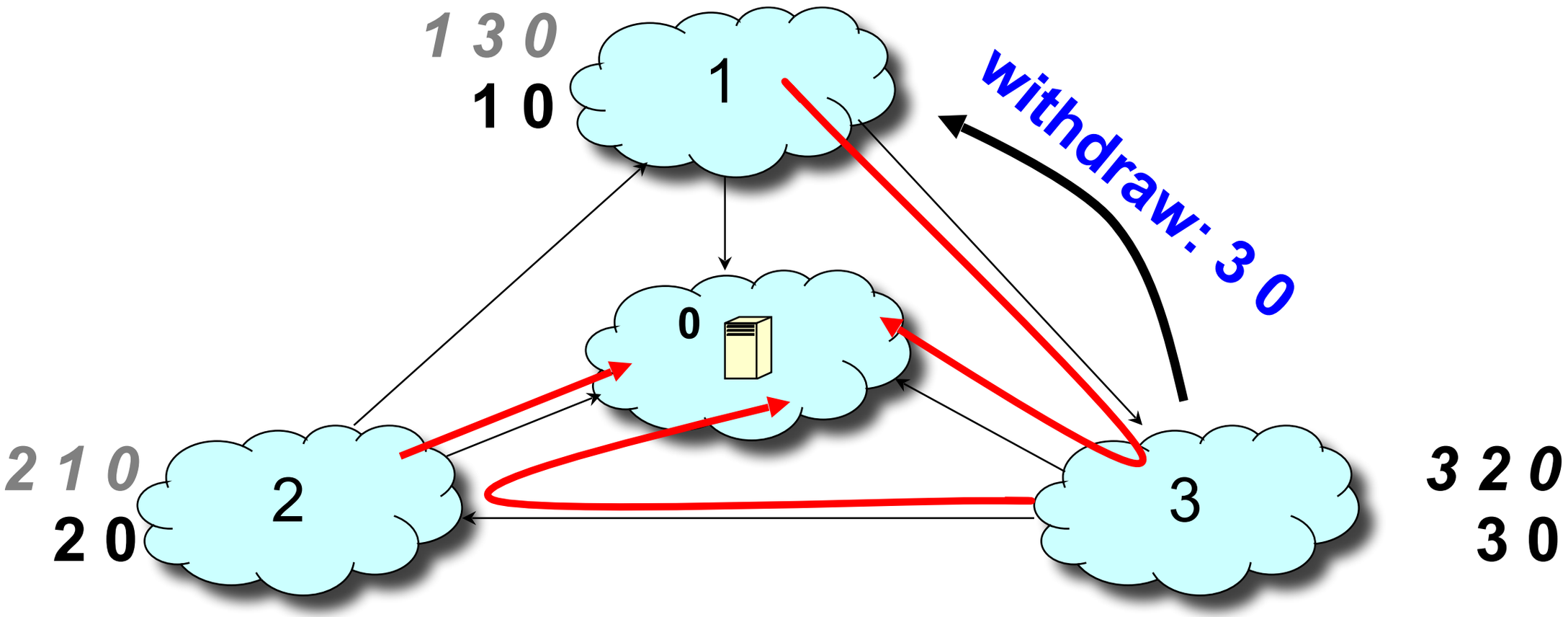


Step-by-step Policy Oscillation

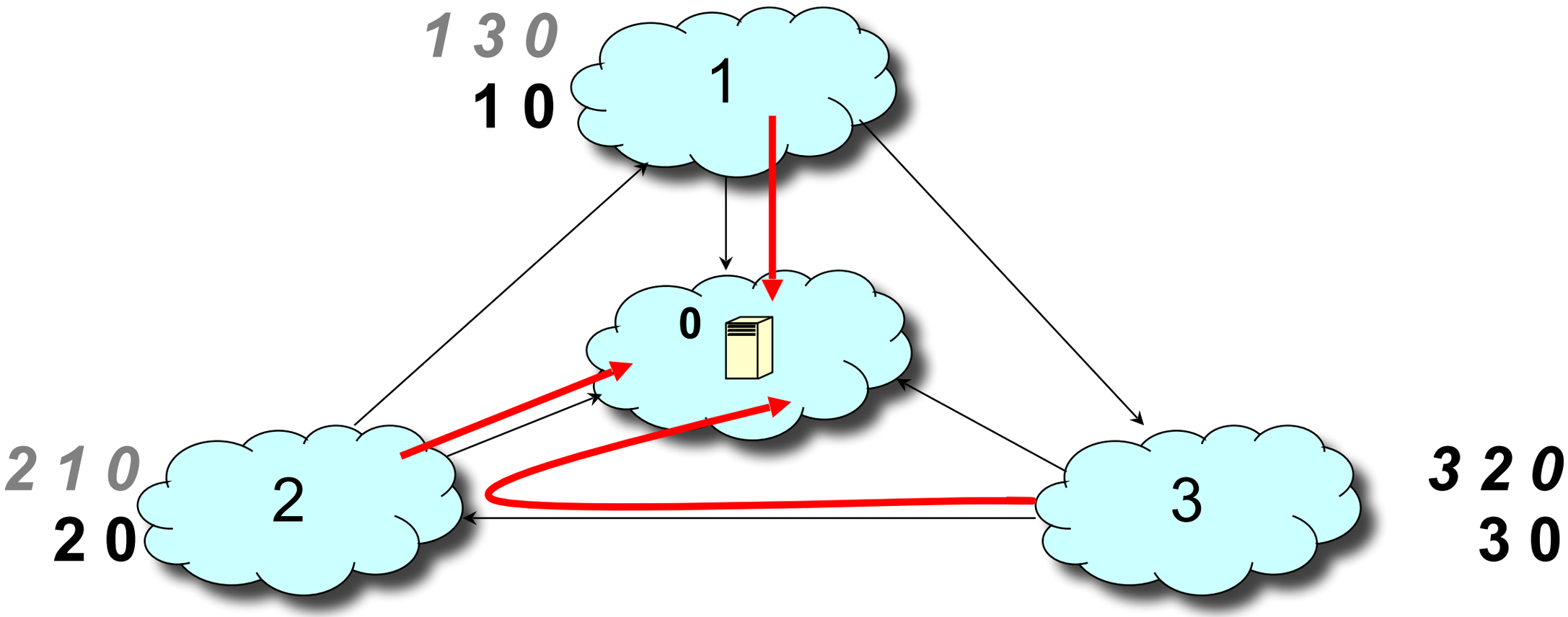


Step-by-step Policy Oscillation

3 **withdraws** its path 3 0 from 1

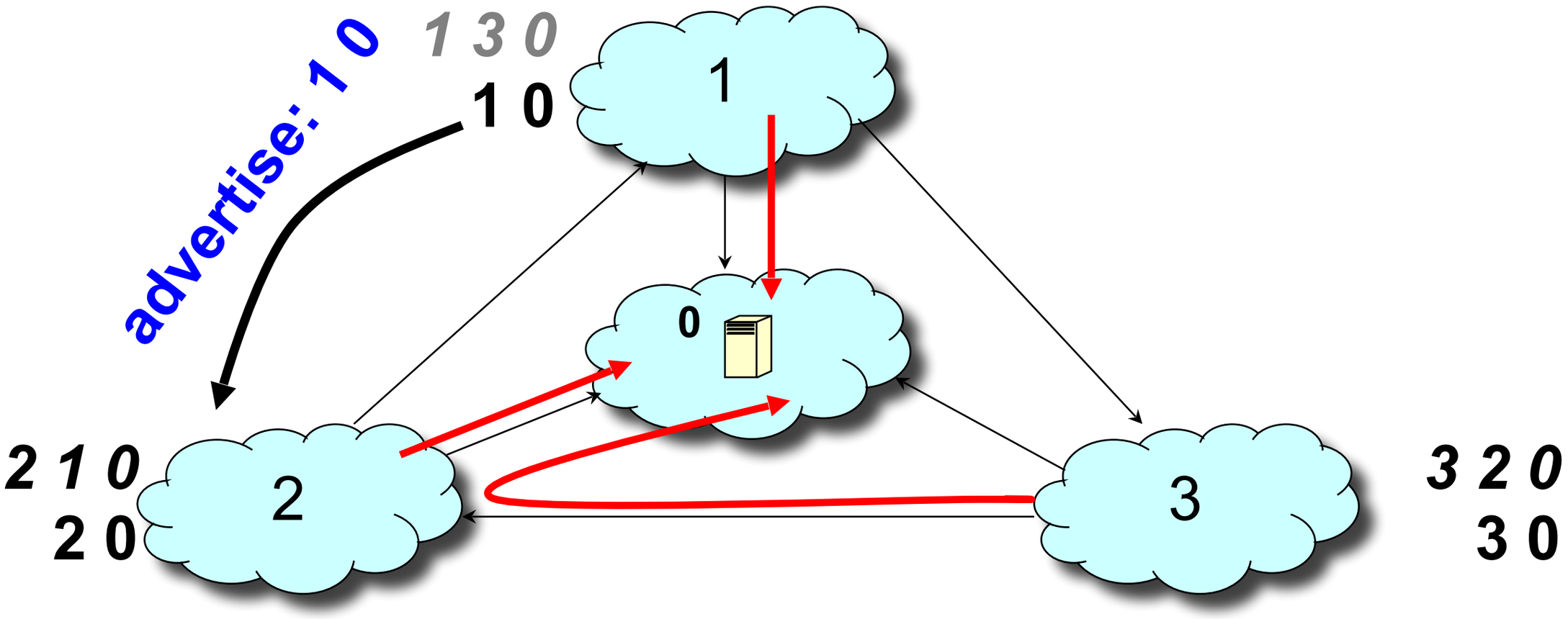


Step-by-step Policy Oscillation

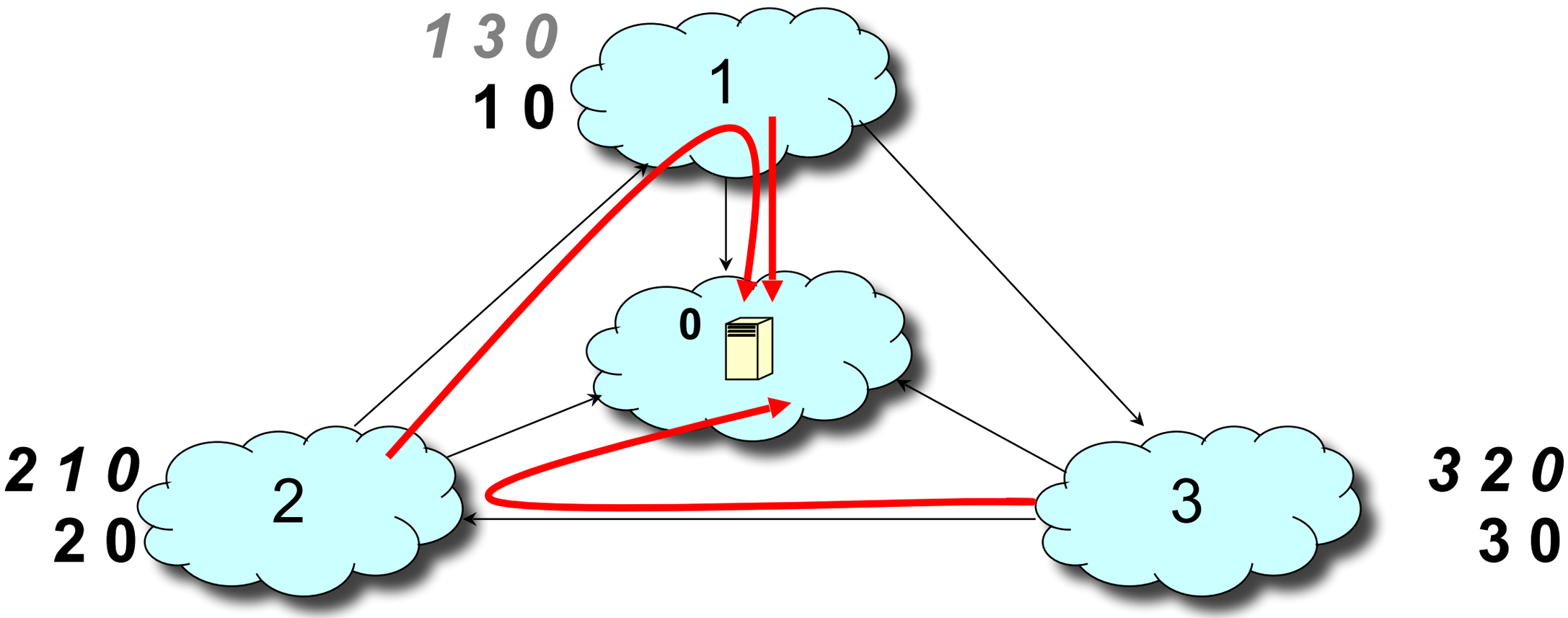


Step-by-step Policy Oscillation

1 **advertises** its path 1 0 to 2

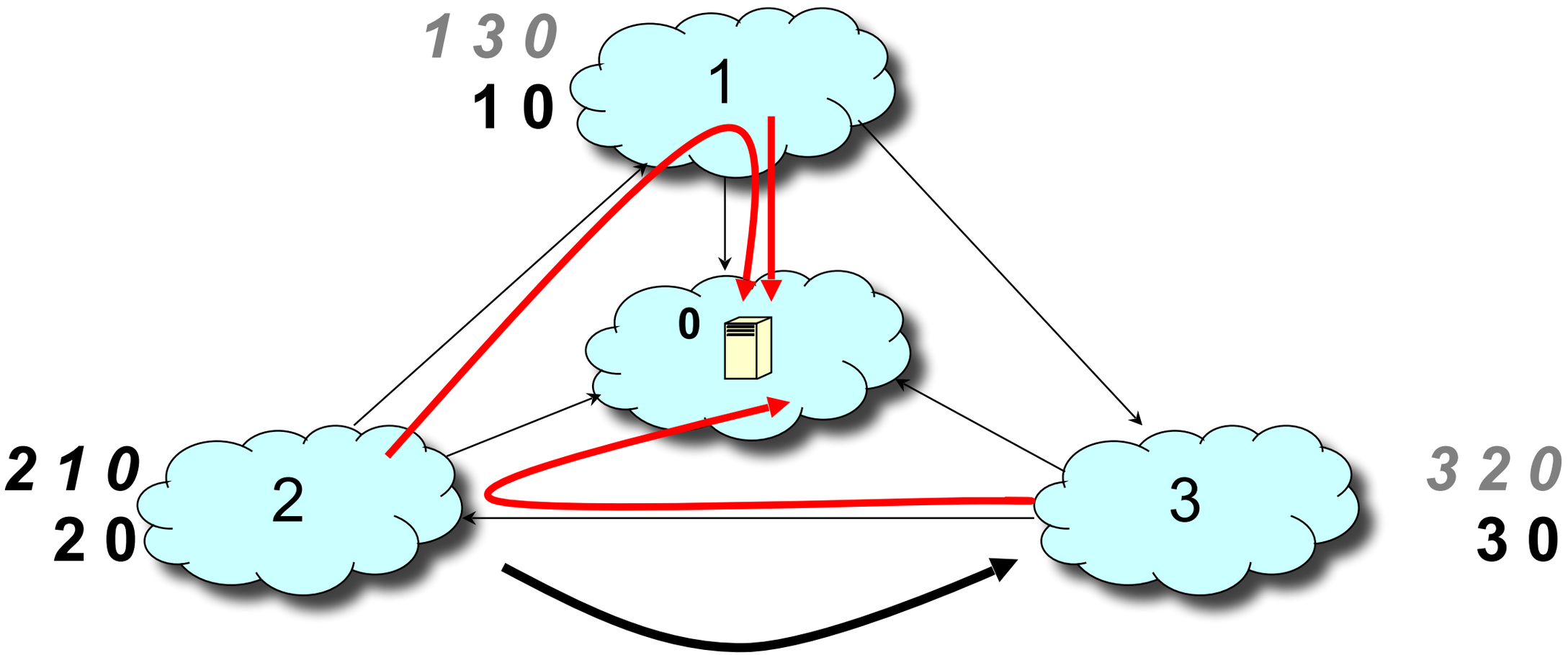


Step-by-step Policy Oscillation

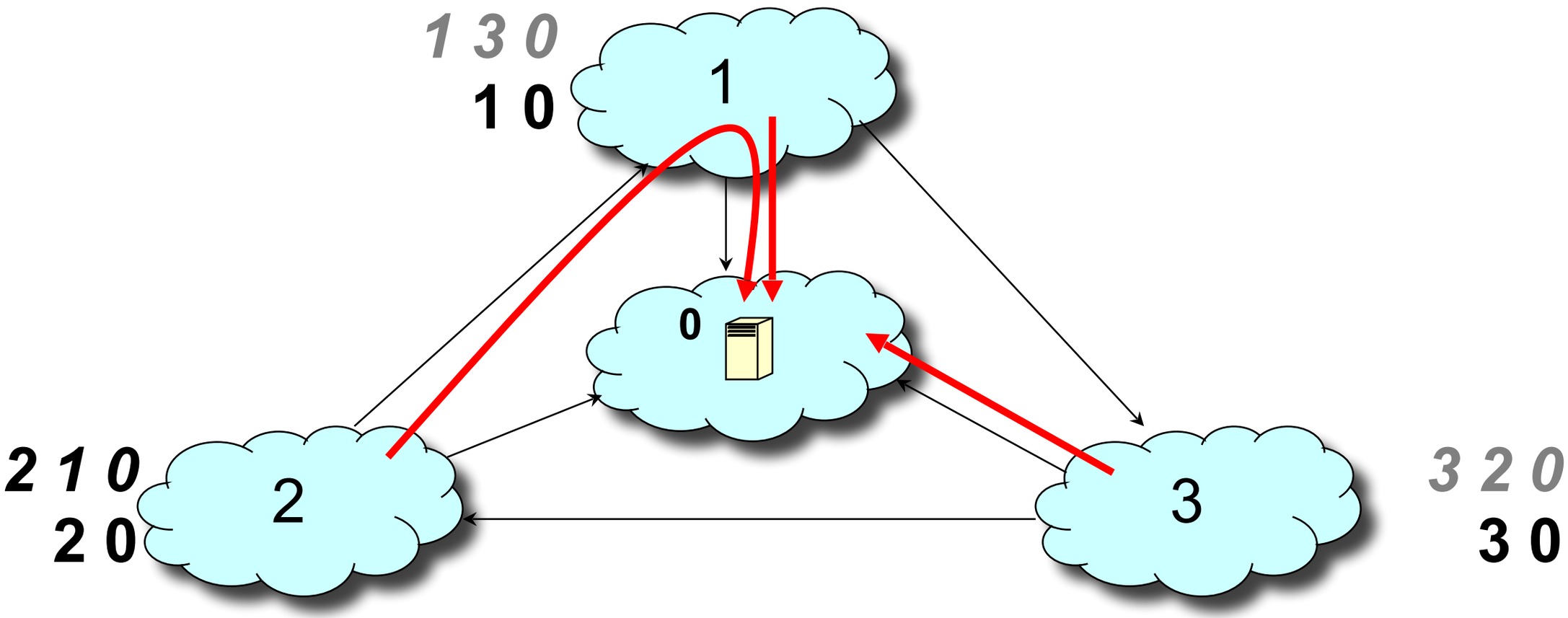


Step-by-step Policy Oscillation

2 **withdraws** its path 2 0 from 3



Step-by-step Policy Oscillation



We are back to where we started!

Convergence

- If all AS policies follow Gao-Rexford rules,
  - Then BGP is guaranteed to converge (safety)

- For arbitrary policies, BGP may fail to converge!

- Why should this trouble us?

Performance Non-Issues

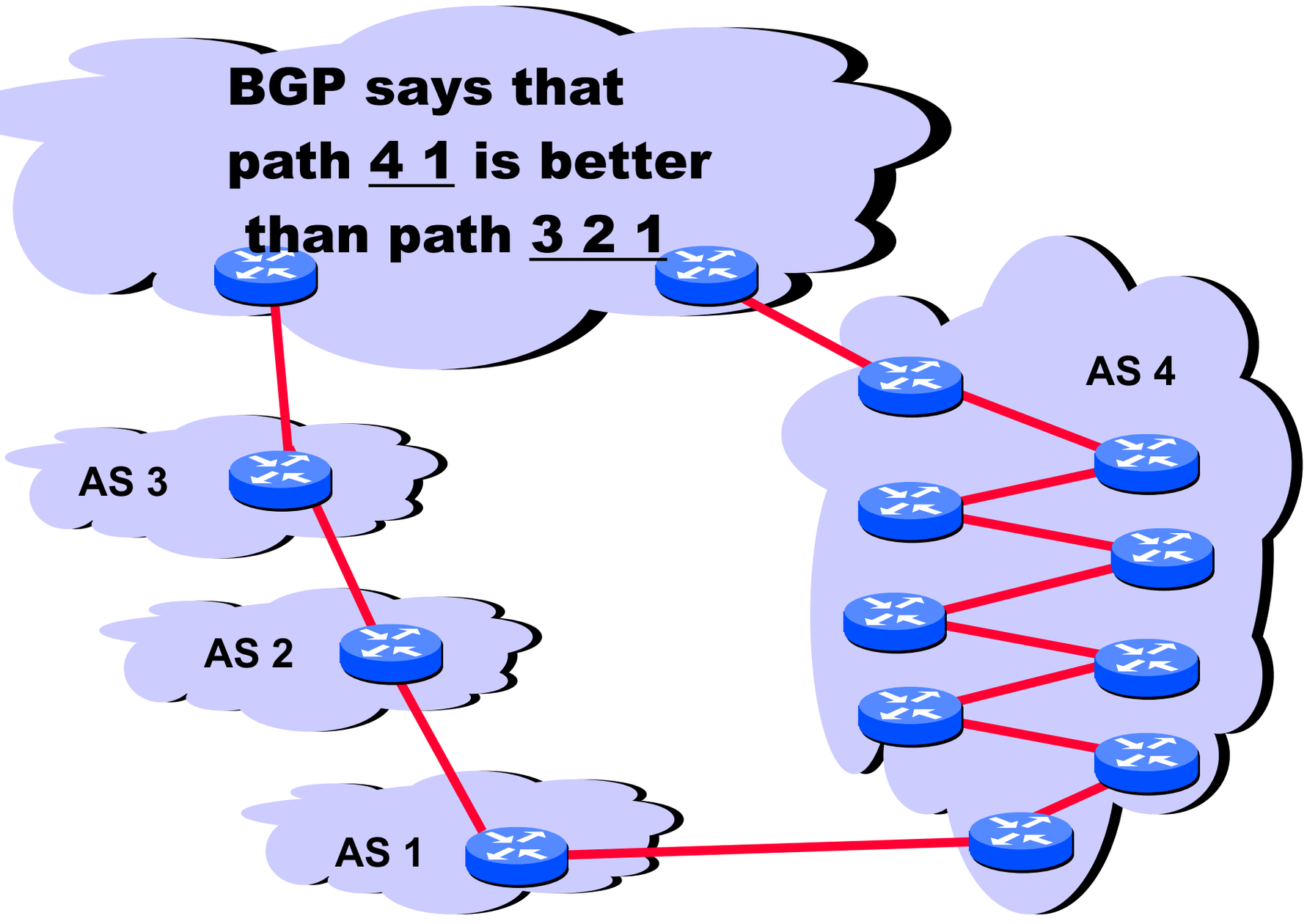
- Internal Routing
  - Domains typically use “hot potato” routing
  - Not always optimal, but economically expedient

- Policy not about performance
  - So policy-chosen paths aren’t shortest

- AS path length can be misleading
  - 20% of paths inflated by at least 5 router hops

Performance (example)

- AS path length can be misleading
  - An AS may have many router-level hops



Performance: Real Issue

Slow Convergence

- BGP outages are biggest source of Internet problems

- Labovitz et al. *SIGCOMM'97*
  - 10% of routes available less than 95% of the time
  - Less than 35% of routes available 99.99% of the time

- Labovitz et al. *SIGCOMM 2000*
  - 40% of path outages take 30+ minutes to repair

- But most popular paths are very stable

BGP Misconfigurations

- BGP protocol is both bloated and underspecified
  - Lots of attributes
  - Lots of leeway in how to set and interpret attributes
  - Necessary to allow autonomy, diverse policies
  - … But also gives operators plenty of rope

- Much of this configuration is manual and *ad hoc*

- And the core abstraction is fundamentally flawed
  - Disjoint per-router configuration to effect AS-wide policy
  - Now strong industry interest in changing this!

BGP: How did we get here?

- BGP was designed for a different time
  - Before commercial ISPs and their needs
  - Before address aggregation
  - Before multi-homing

- **1989 : BGP-1 [RFC 1105]**
  - **Replacement for EGP (1984, RFC 904)**
- **1990 : BGP-2 [RFC 1163]**
- **1991 : BGP-3 [RFC 1267]**
- **1995 : BGP-4 [RFC 1771]**
  - **Support for Classless Interdomain Routing (CIDR)**

- We don’t get a second chance: ‘clean slate’ designs virtually impossible to deplay
- Thought experiment: how would you design a policy-driven interdomain routing solution?
  - How would you deploy it?