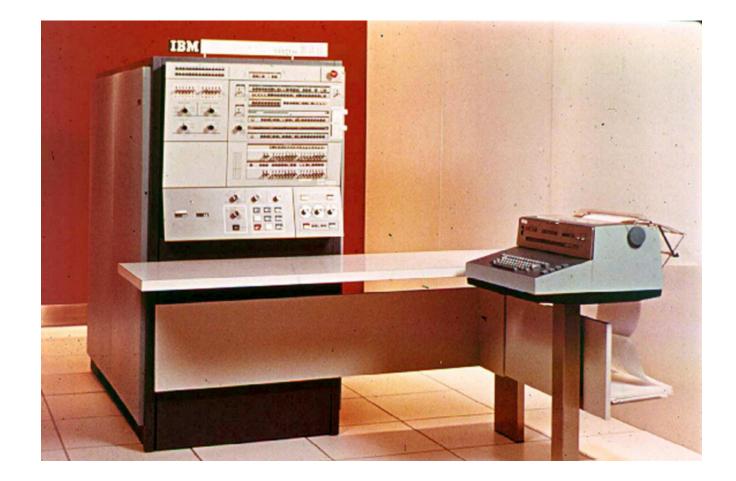
Review

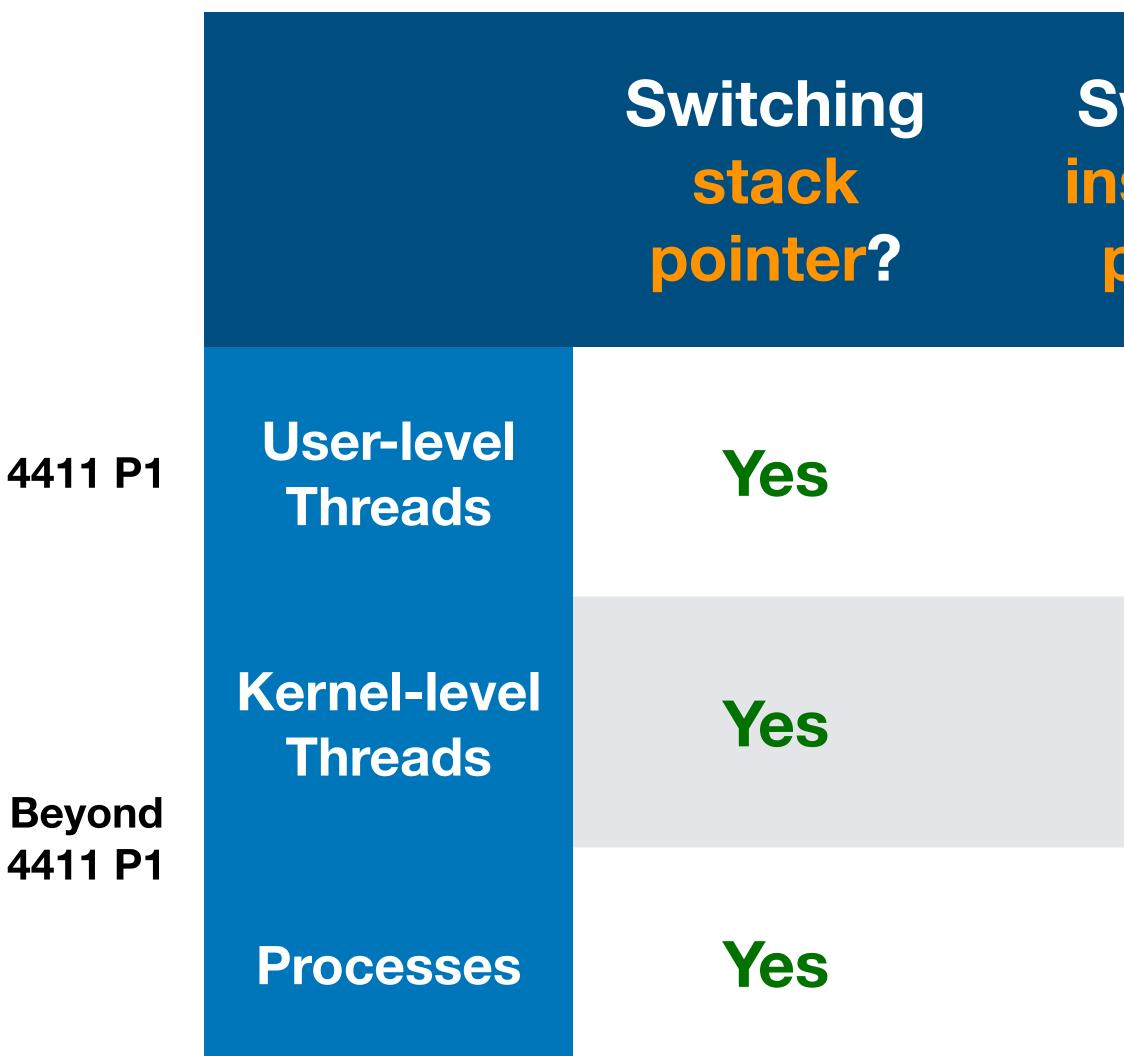
- Running a program requires the code & stack segments in memory.
- Context = memory address space + stack pointer + instruction pointer
 - A CPU is in the context of a program if its instruction pointer and stack pointer registers point to the code & stack segments of the program.
- Context-switch means switching the context of a CPU to different programs by modifying its stack pointer and instruction pointer.

Big picture of context-switch

- The initial goal of operating systems is multi-tasking.
 - A naive way of multi-tasking is batch processing.
 - The concept of context-switch enables time-sharing multi-tasking.
 - There are different implementations of context-switch: user-level threads, kernel-level threads, processes

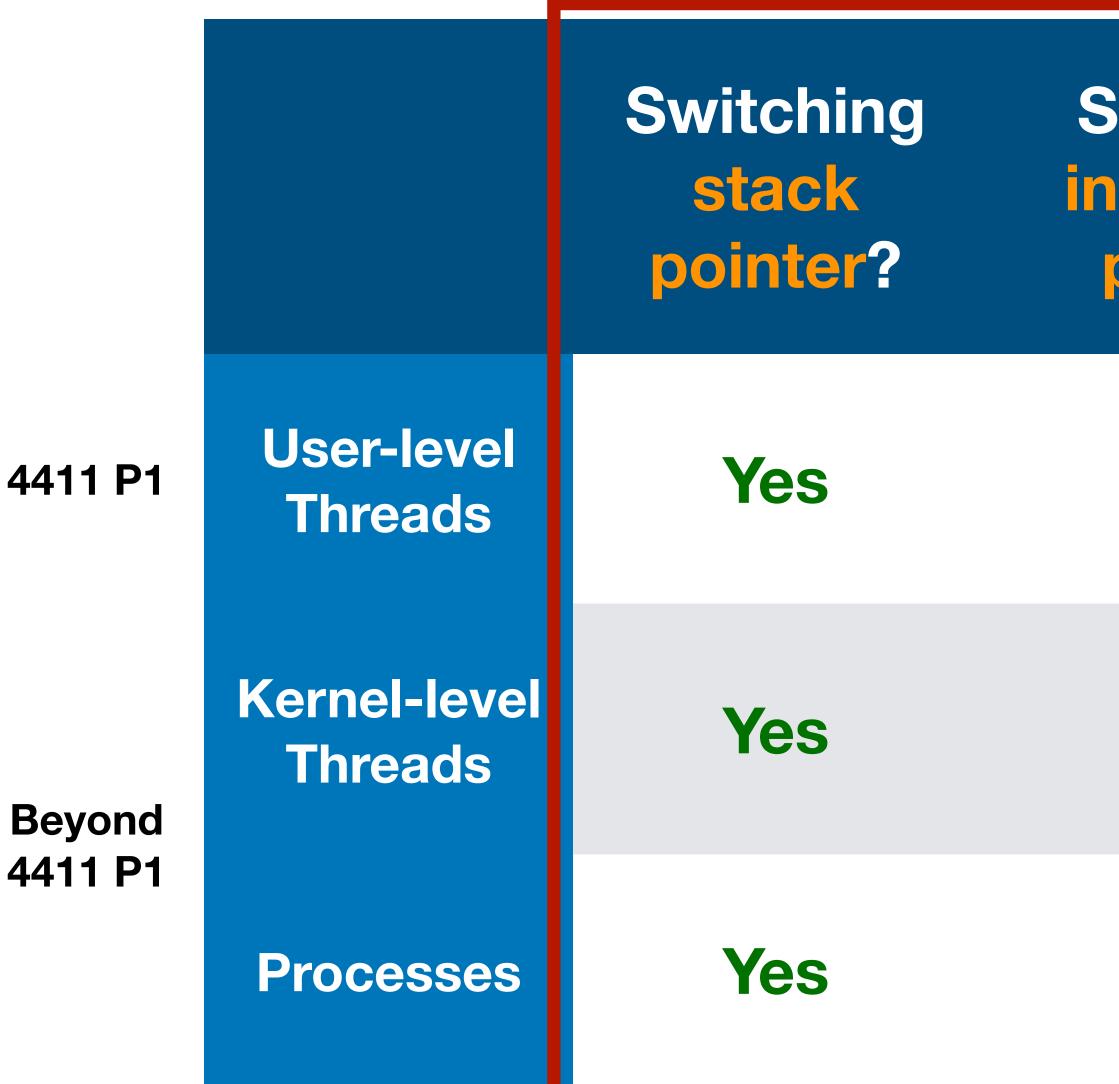
* Images from https://about.sourcegraph.com/blog/the-ibm-system-360-the-first-modular-general-purpose-computer/





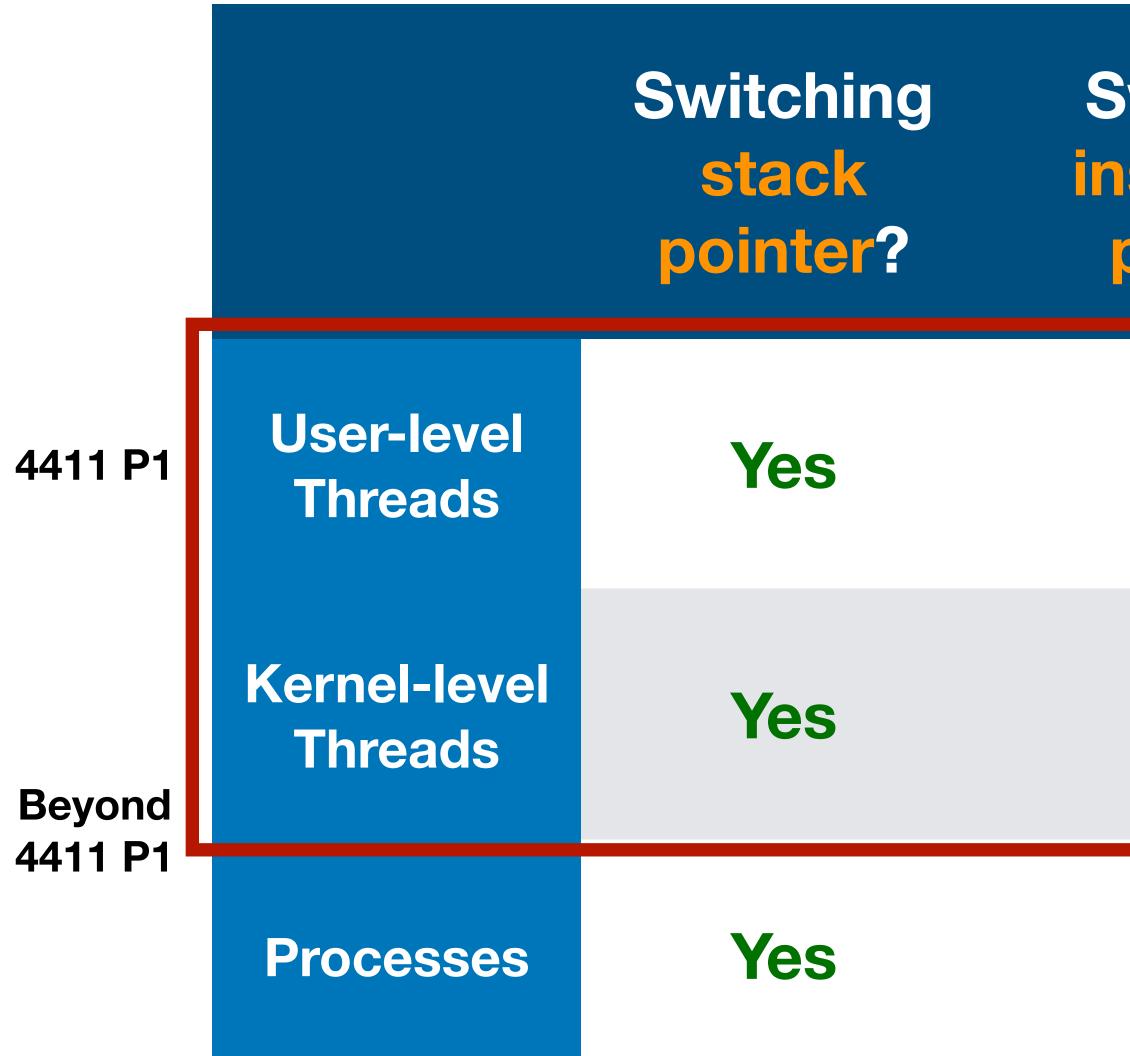
Source of the structure	Switchig memory address space?	Switching kernel/user mode?	
No	No	No	
Yes	No	Yes	
Yes	Yes	Yes	





Switching struction	Switching memory address	Switching kernel/user	
pointer?	space?	mode?	
No	No	No	
Yes	No	Yes	
Yes	Yes	Yes	





Switching struction pointer?	Switchig memory address space?	Switching kernel/user mode?	
No	No	No	
Yes	No	Yes	
Yes	Yes	Yes	



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4411 P1	User-level Threads	Yes	No	No	No
Beyond 4411 P1	Kernel-level Threads	Yes	Yes	No	Yes
	Processes	Yes	Yes	Yes	Yes





Context-switch solves the problem of time-sharing multi-tasking.

Processes and threads implement the concept of context-switch.

Context-switch solves the problem of time-sharing multi-tasking.

Processes and threads implement the concept of context-switch.

Next problem: how do different processes/threads communicate?

Next problem: how do different processes/threads communicate?

- user interface, one for microphone and one for camera.
 - microphone thread to stop recording.
 - interface thread.

• For example, say there are 3 threads running my zoom together: one for

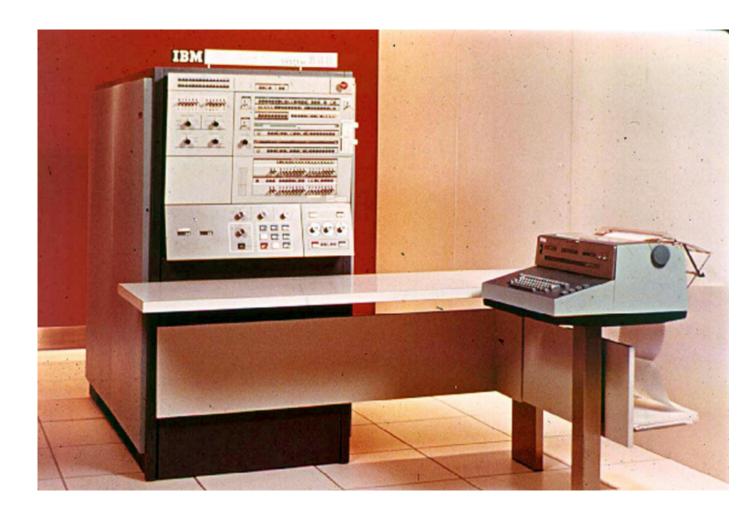
When I click the "mute" button, the user interface thread should tell the

The camera thread should continuously transfer video data to the user

Interprocess communication (IPC)

- The terminology of this problem is IPC which is extensively studied in the operating systems literature. Performance is the key!
 - If IPC has poor performance, the camera thread cannot transfer video data to the user interface thread in time, leading to poor experience.
- Note: we will use the general term IPC to represent both communications among processes and communications among threads.

Historical representatives



• IBM 360 is a representative of time-sharing multi-tasking with context-switch.

1960s

• AT&T UNIX System V is a representative of interprocess communication (IPC).

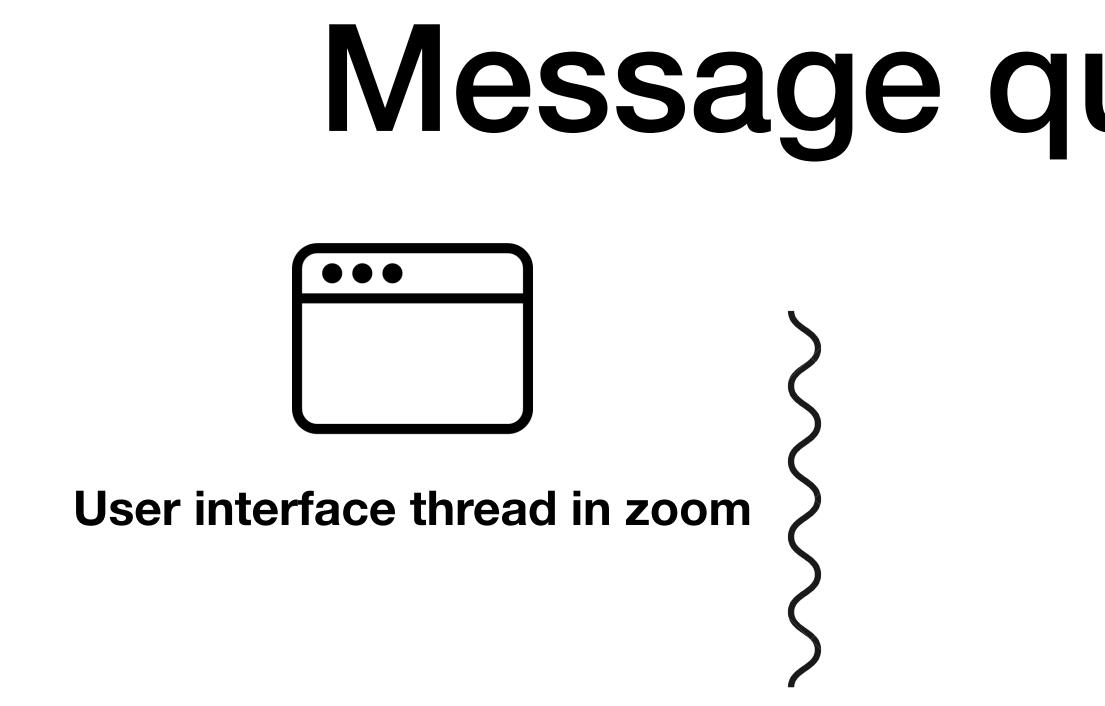
1980s

UNIX System V IPC

System V IPC is the name given to three interprocess communication mechanisms that are widely available on UNIX systems: message queues, semaphore, and shared memory.

what you need to implement in 4411 P1

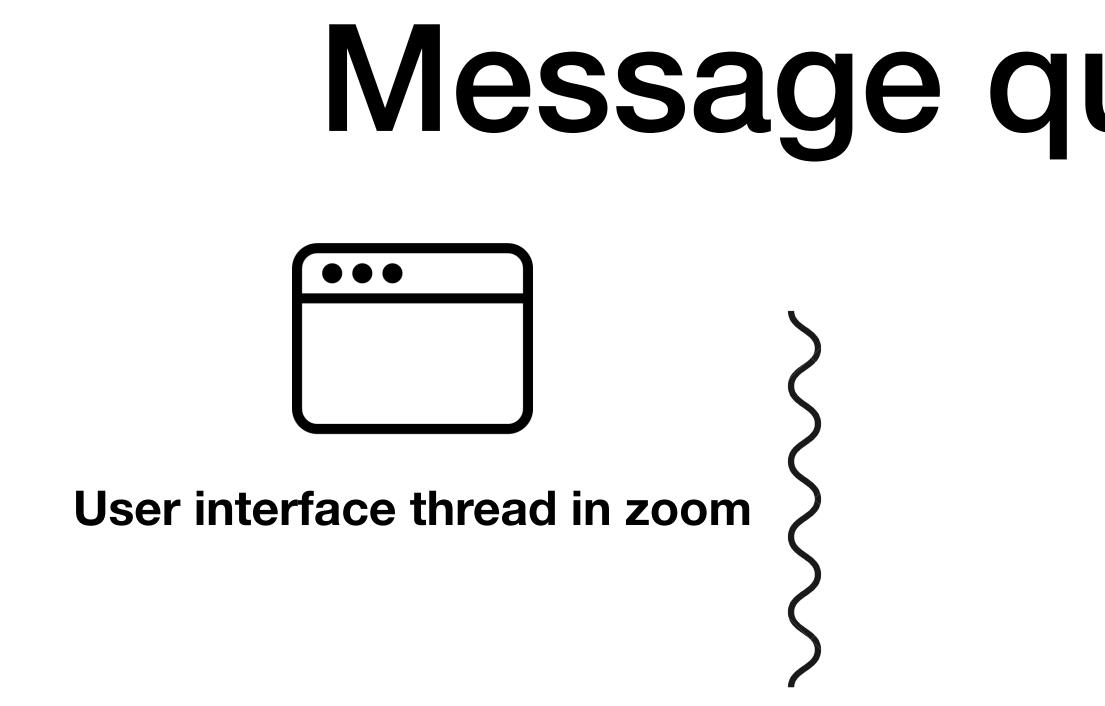
https://man7.org/linux/man-pages/man7/svipc.7.html



User-level

Kernel-level

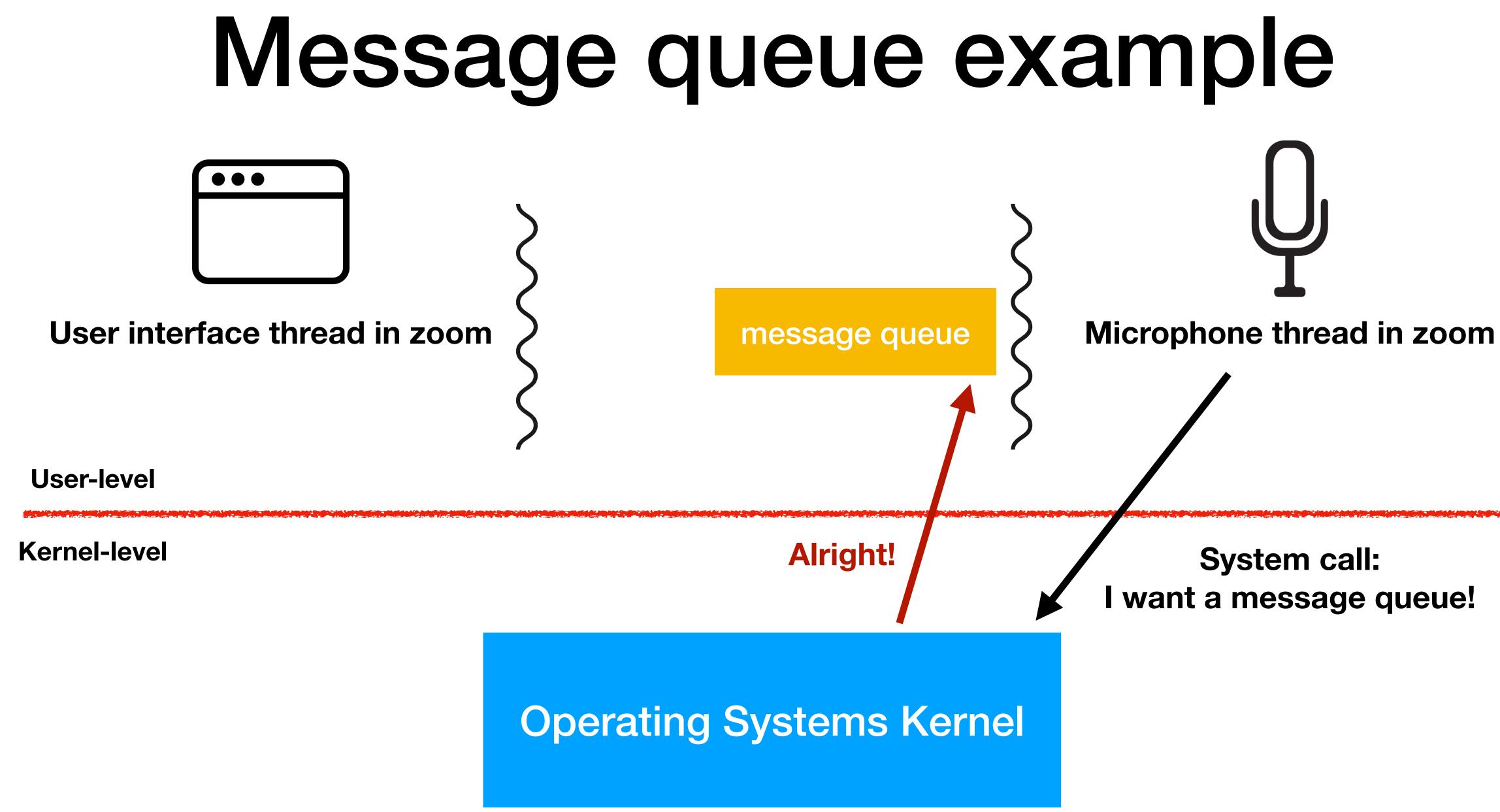
Message queue example **Microphone thread in zoom**

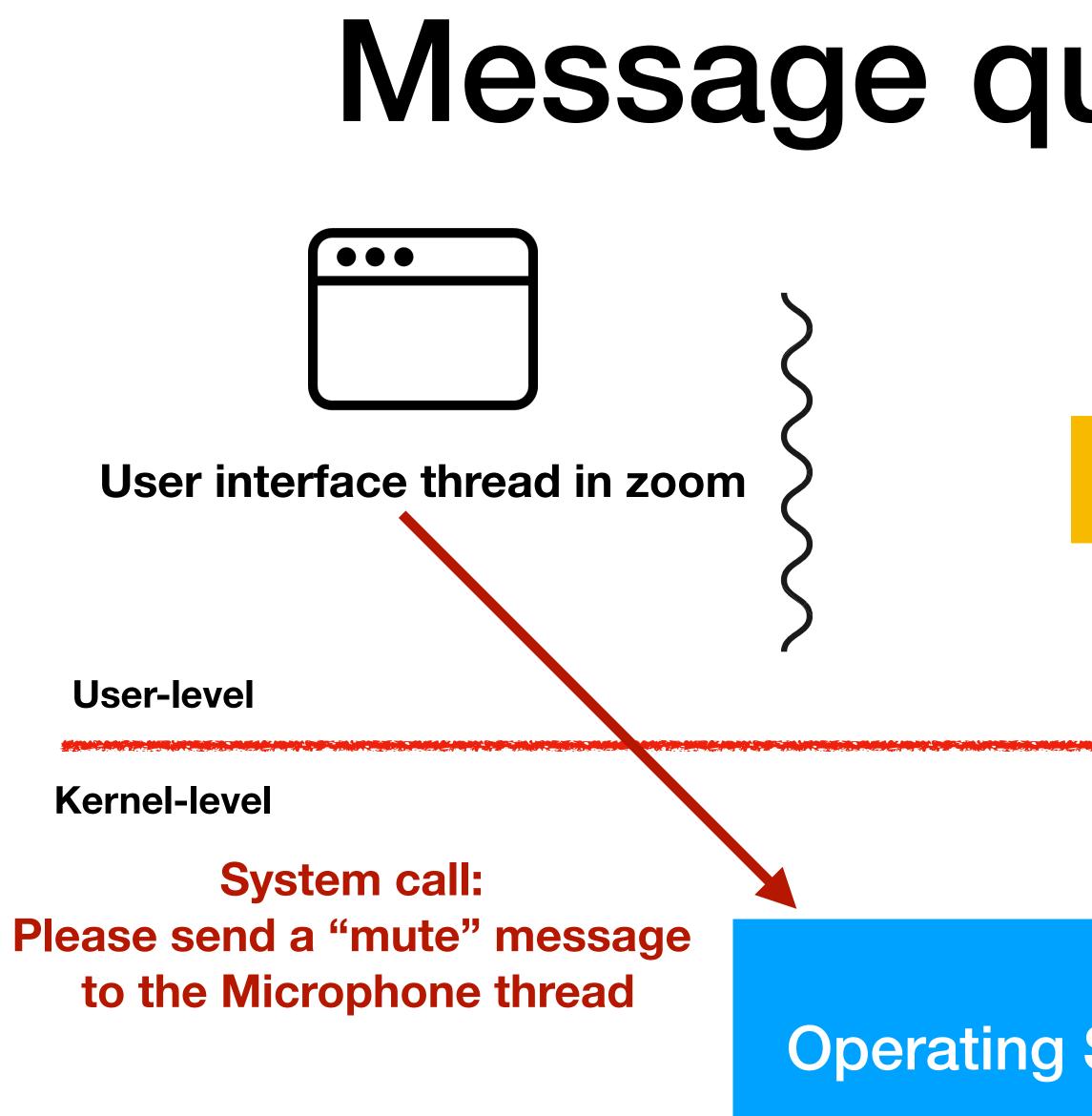


User-level

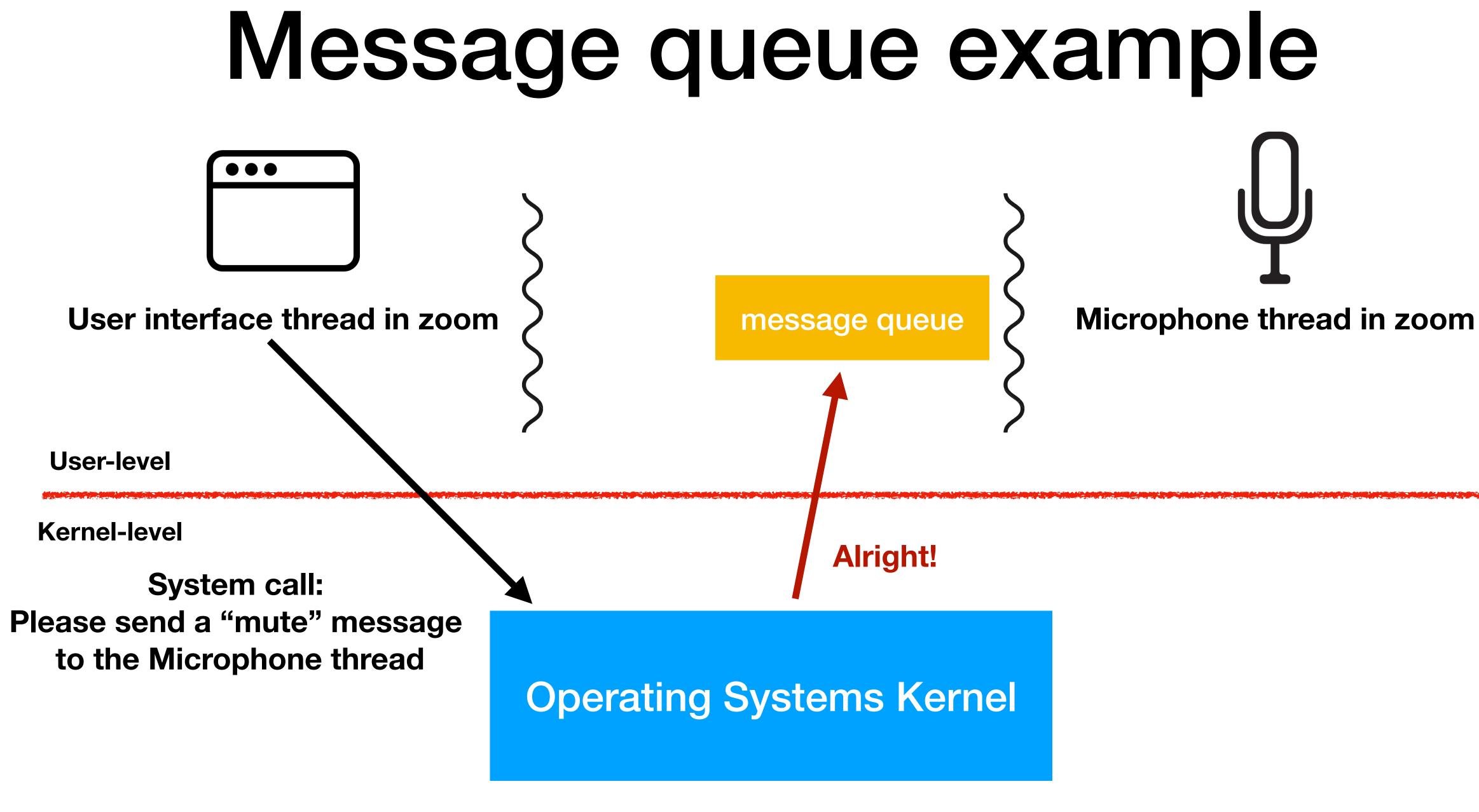
Kernel-level

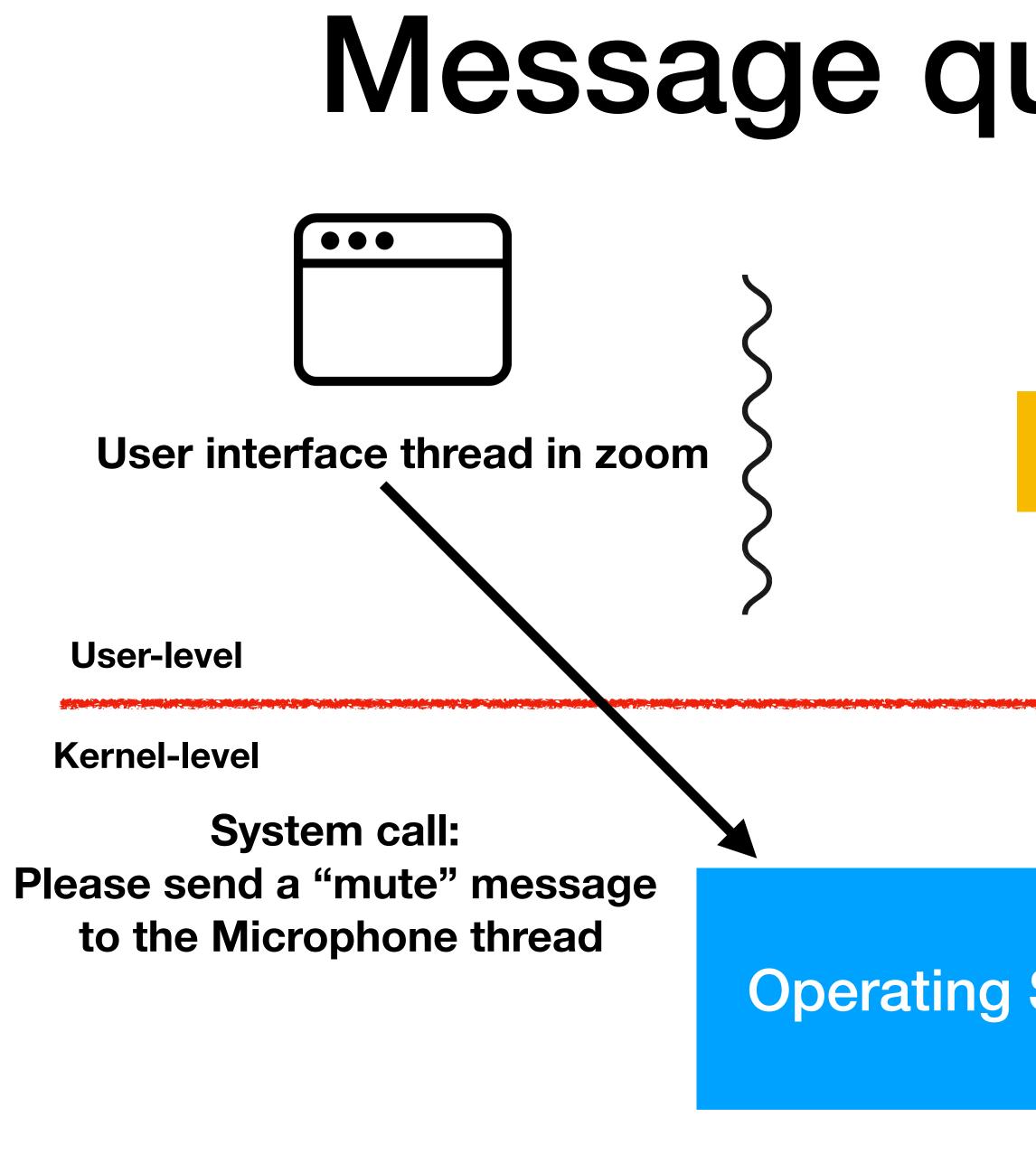
Message queue example **Microphone thread in zoom System call:** I want a message queue! **Operating Systems Kernel**





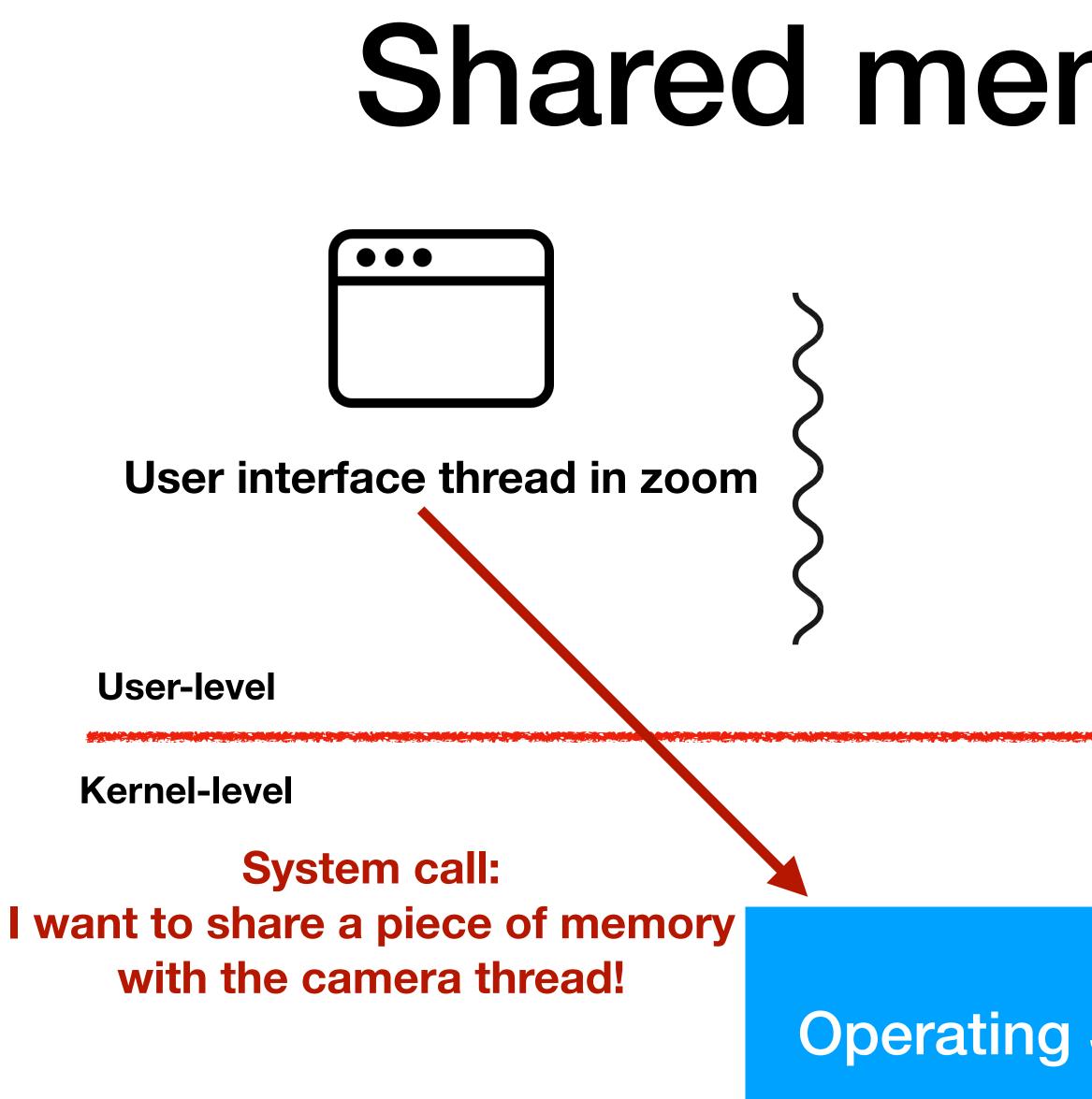
Message queue example



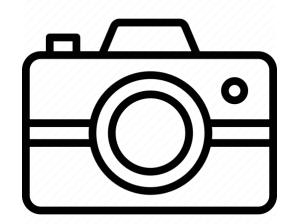


Message queue example **Microphone thread in zoom** message queue **OK**, I'll stop recording

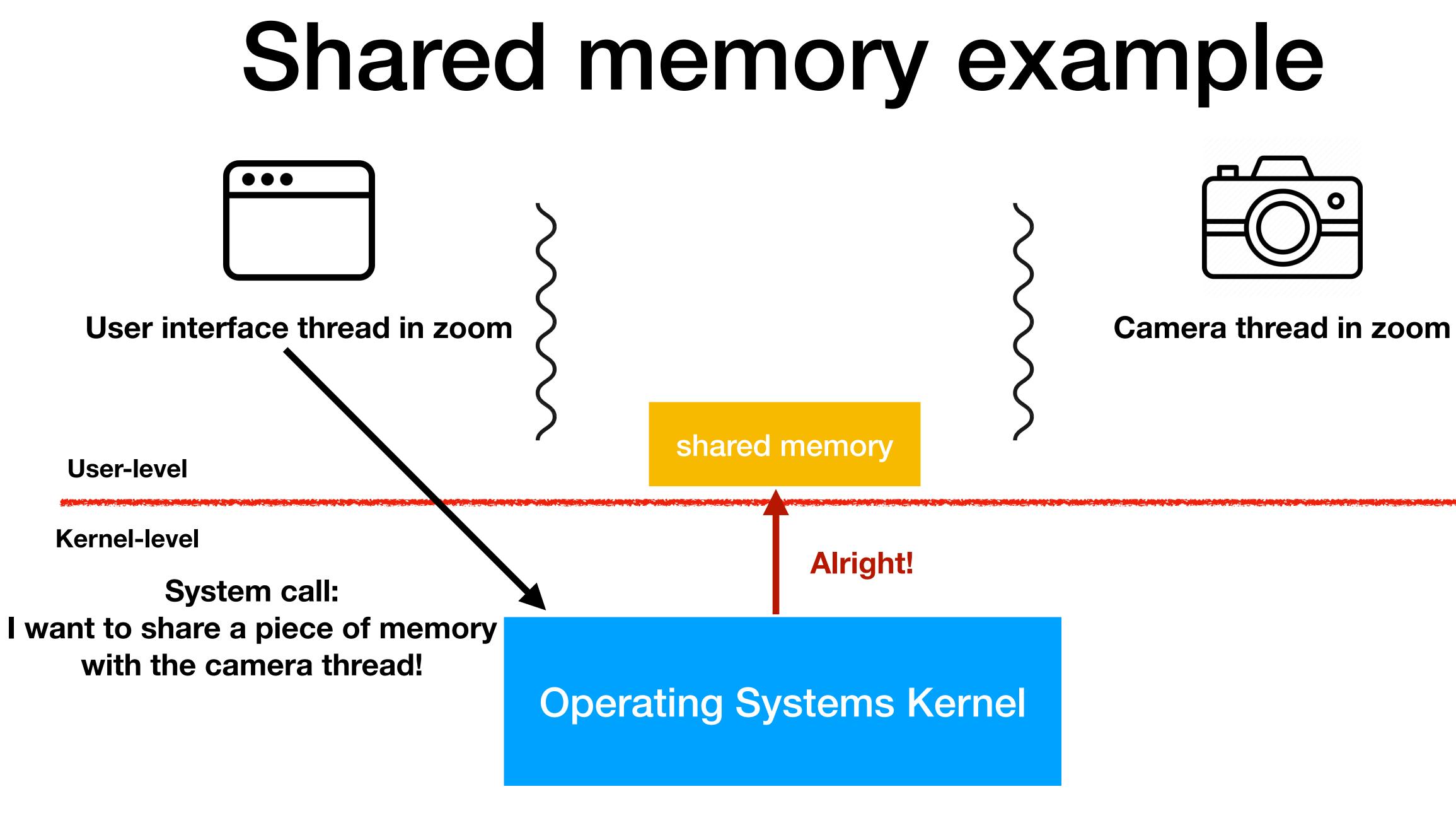
Alright!

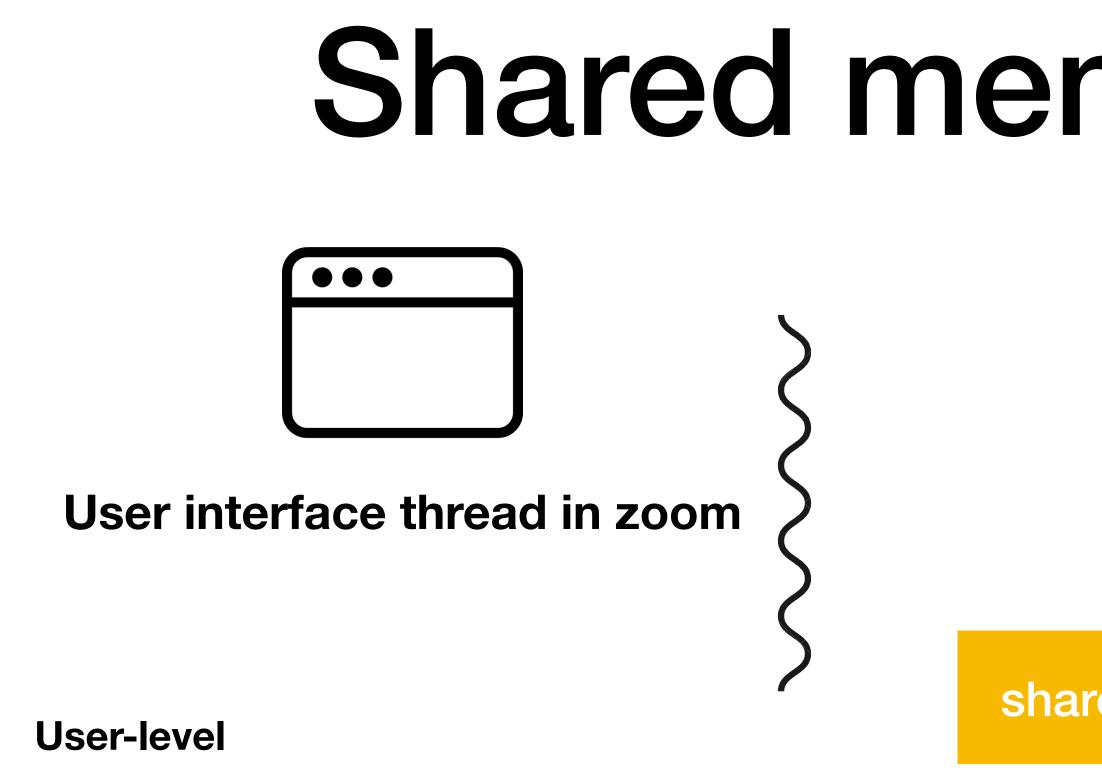


Shared memory example



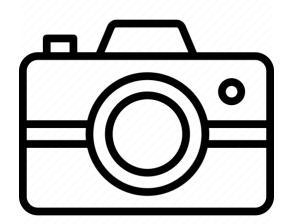
Camera thread in zoom





Kernel-level

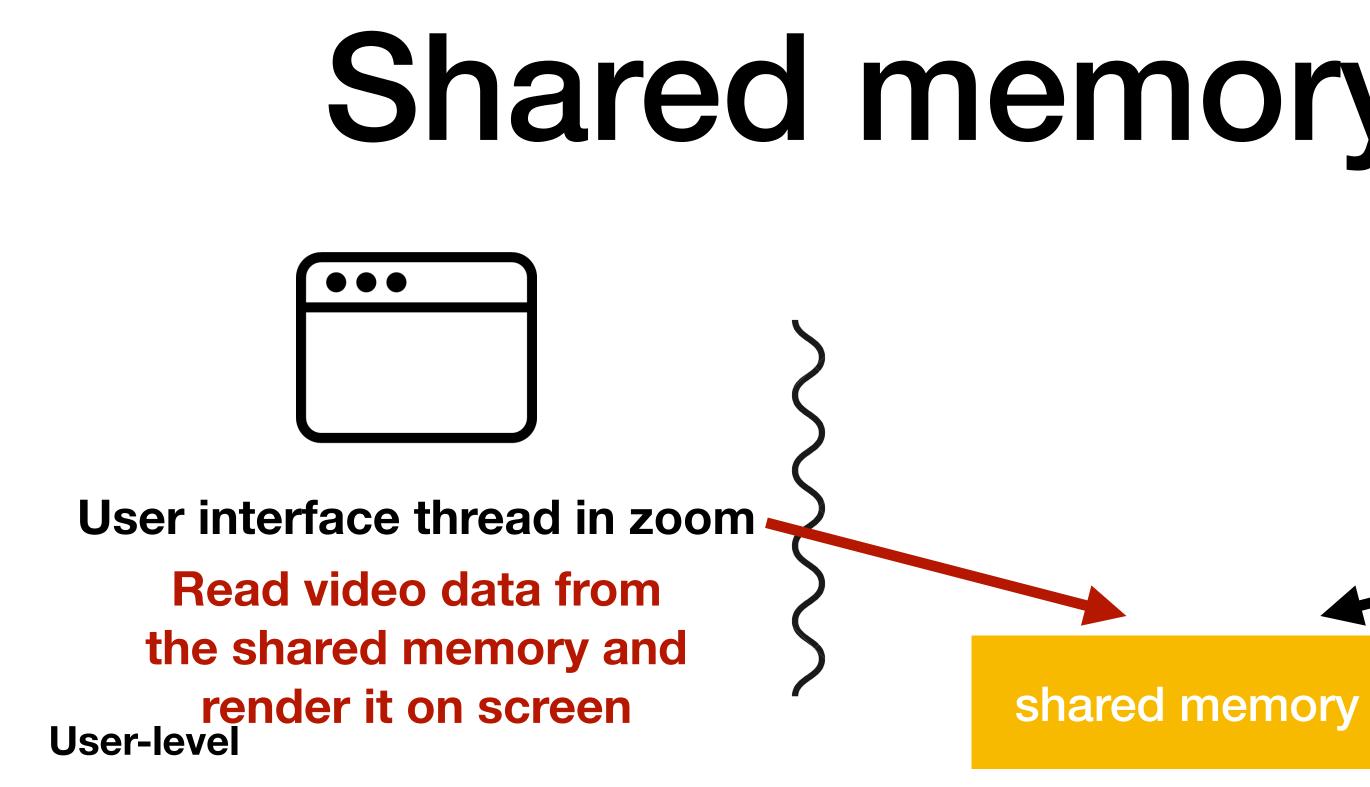
Shared memory example



shared memory

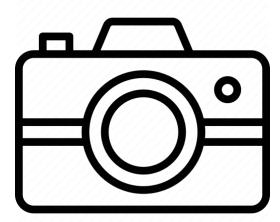
Camera thread in zoom

Write video data to the shared memory



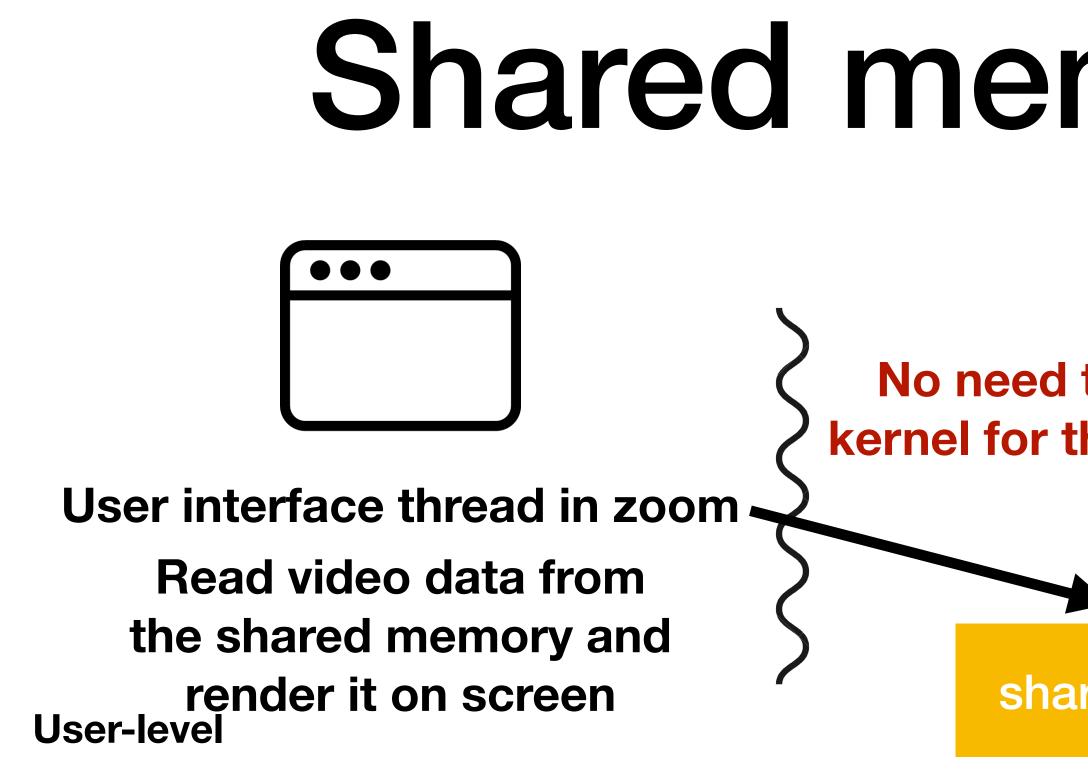
Kernel-level

Shared memory example



Camera thread in zoom

Write video data to the shared memory



Kernel-level

Shared memory example No need to go through the kernel for this communication! **Camera thread in zoom** Write video data to the shared memory shared memory

Lesson: shared memory has better performance than message queues because communications get around the kernel.

The third IPC mechanism: semaphores

System V IPC is the name given to three interprocess communication mechanisms that are widely available on UNIX systems: message queues, semaphore, and shared memory.

https://man7.org/linux/man-pages/man7/svipc.7.html

what you need to implement in 4411 P1

The producer-consumer problem

- There are two types of threads (or processes): producer and consumer.
 - Producer produces some kind of resources (e.g., HTTP web request) and consumer consume the resources (e.g., process the request).
- Goal: consumer should only be scheduled when some resource produced by the producer is available (i.e., has not been consumed).
- The core of semaphore is a counter of such available resources. If counter is greater than 0, a consumer thread will be scheduled.

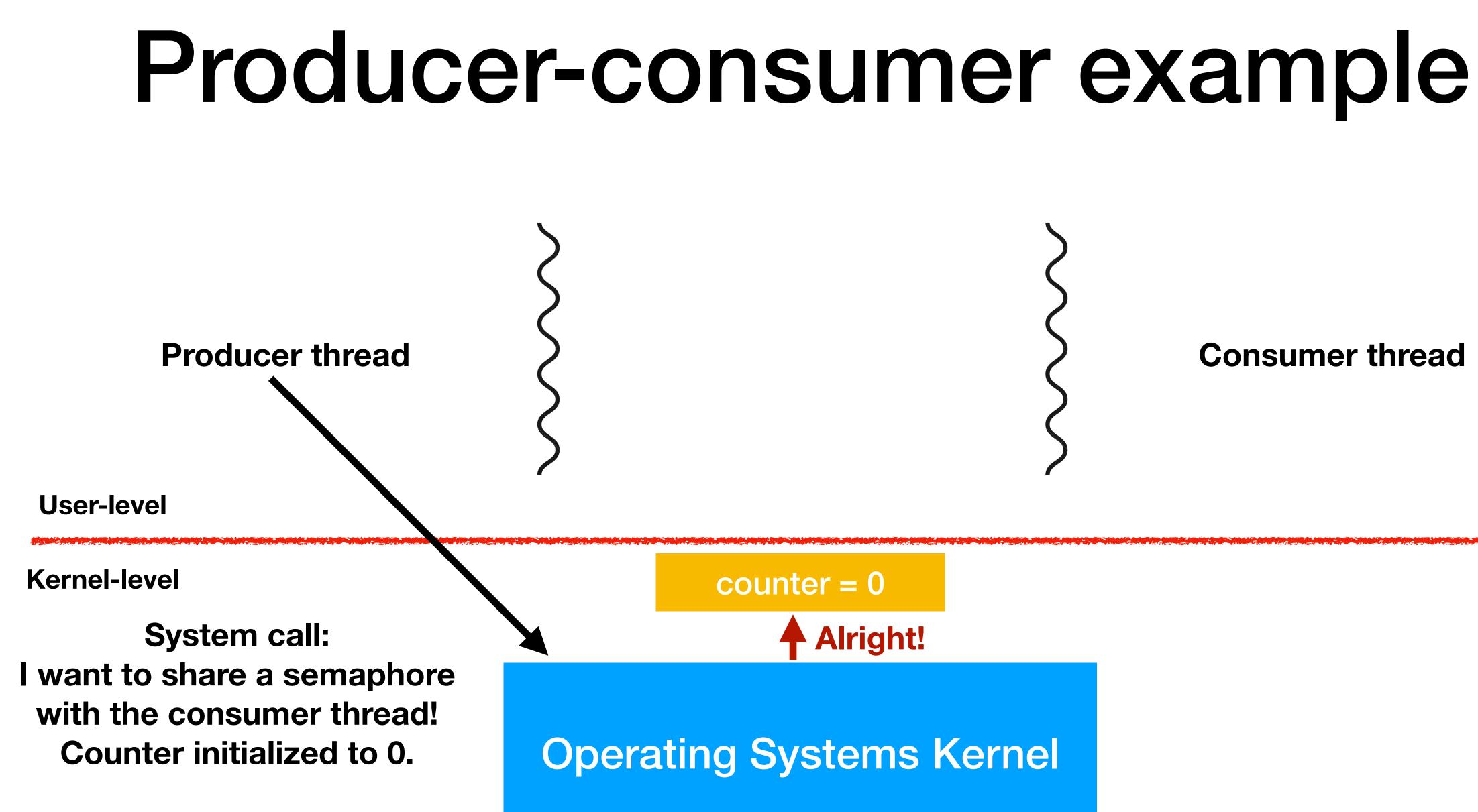
Producer thread

User-level

Kernel-level

System call: I want to share a semaphore with the consumer thread! **Counter initialized to 0.**

Consumer thread





Consumer thread

Producer thread



Kernel-level

Consumer thread

counter = 0

System call: want to consume a resource please decrement the counter



Producer thread



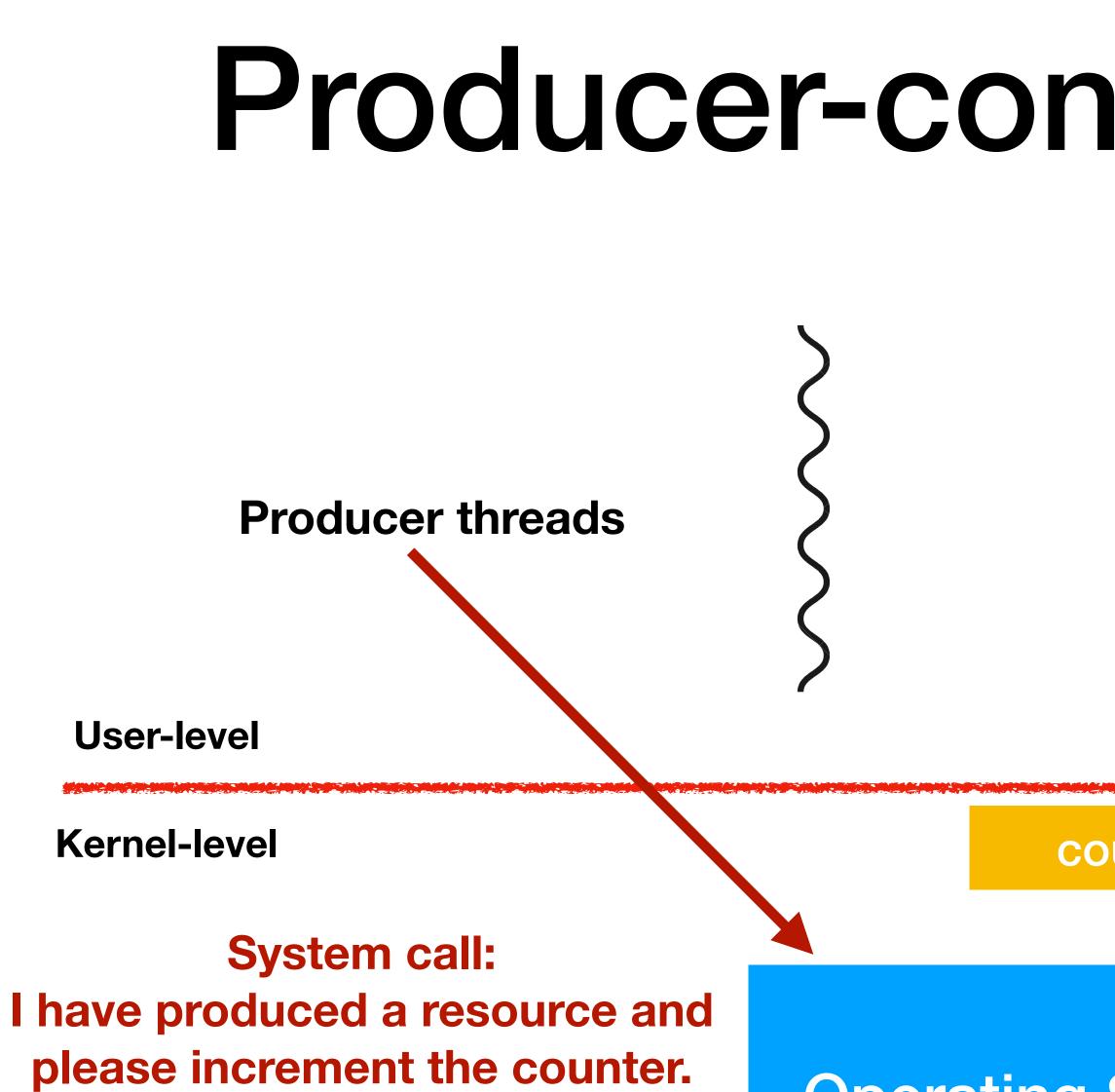
Kernel-level



Operating Systems Kernel

No resource available, suspend the thread until further notice

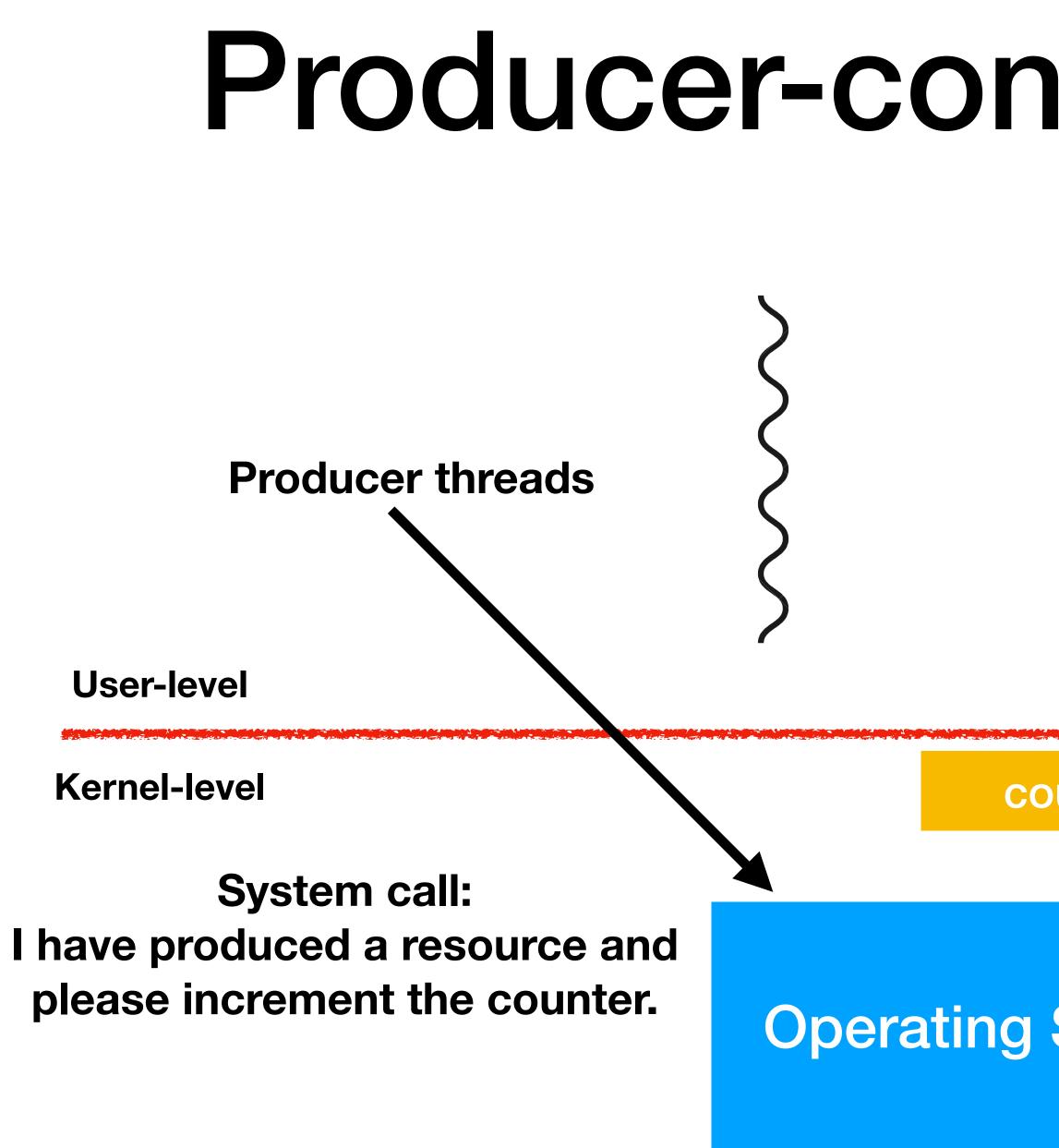




counter = 0

Operating Systems Kernel

Consumer thread (suspended)



counter = 1



Operating Systems Kernel

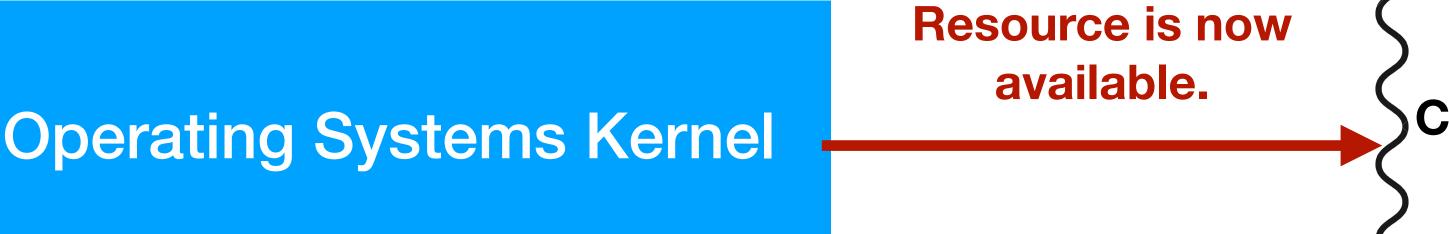
Consumer thread (suspended)

Producer threads



Kernel-level

counter = 1

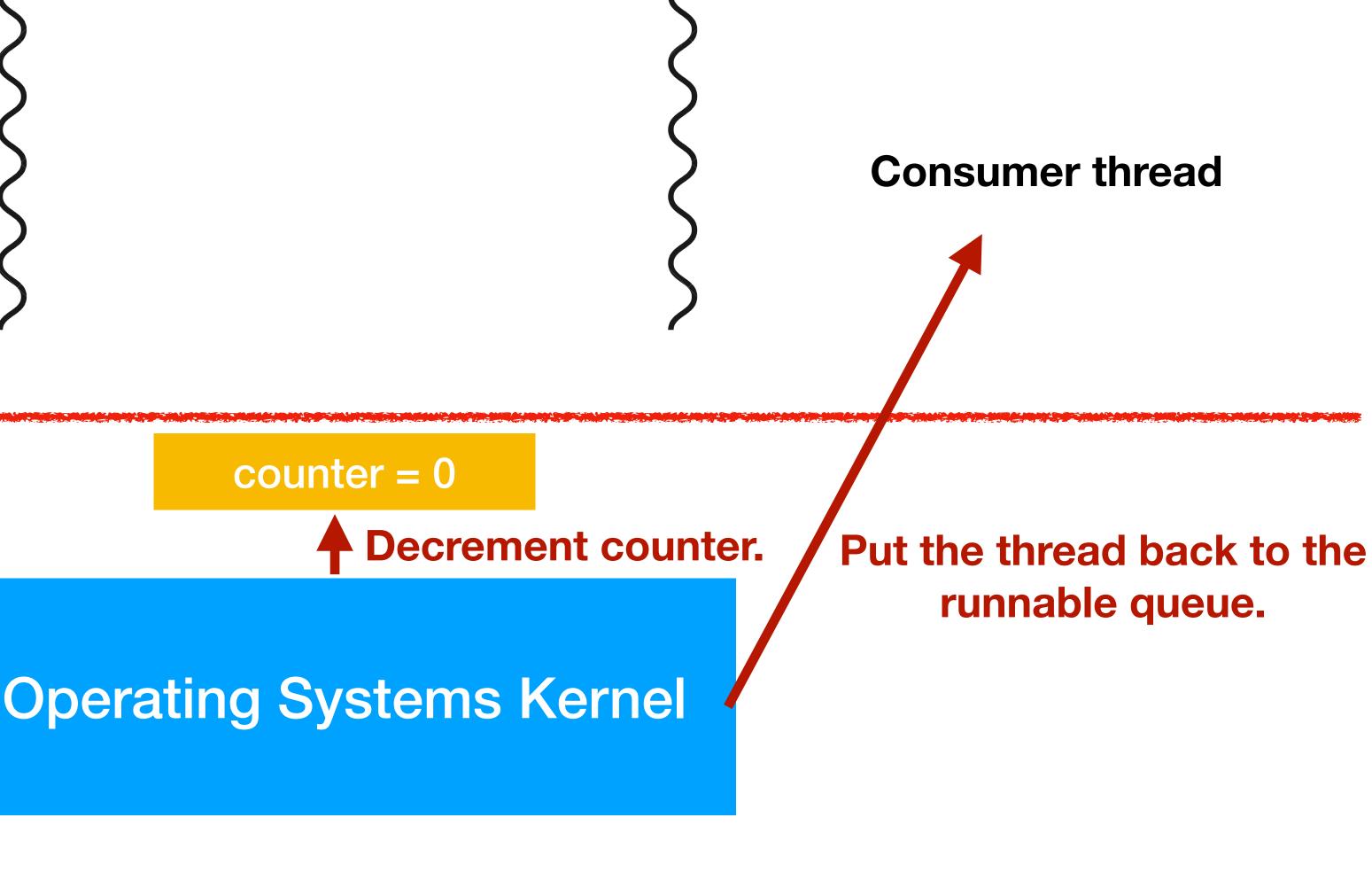




Producer threads



Kernel-level





Variants of producer-consumer

- There can be multiple producer threads and consumer threads.

• Bounded buffer: a producer can only produce if the number of available resources (the value of the counter) is not greater to a given number.

Semaphores in P1

struct sema { // counter // queue of threads that are put to sleep // feel free to add other fields that you need };

// initialize a semaphore void sema_init(struct sema *sema, unsigned int count)

// produce a resource by incrementing the semaphore void sema_inc(struct sema *sema)

// consume a resource by decrementing the semaphore void sema_dec(struct sema *sema)

Lesson: semaphore is easy to implement, but it is not very useful and one should try to avoid using it.

Refer to "12 Commandments of Synchronization" by Emin Gün Sirer

Homework

- P1: implement semaphores and test your semaphore with producer-consumer and other synchronization problems.
- Read the Linux manual of System V IPC: https:// man7.org/linux/man-pages/man7/svipc.7.html
- Next lecture: introduce the concepts of timer interrupt and scheduling for P2.

Concepts in the real-world (Unix/Linux)

- User-level threads
 - getcontext, setcontext, ...
- Kernel-level threads
 - pthread_create, pthread_join, ...
- Processes

• fork, …

Concepts in the real-world (Unix/Linux)

- Message queues
 - msgget, msgsnd, msgrcv, ...
- Semaphores
 - semget, semop, ...
- Shared memory

shmget, shmdt, ...