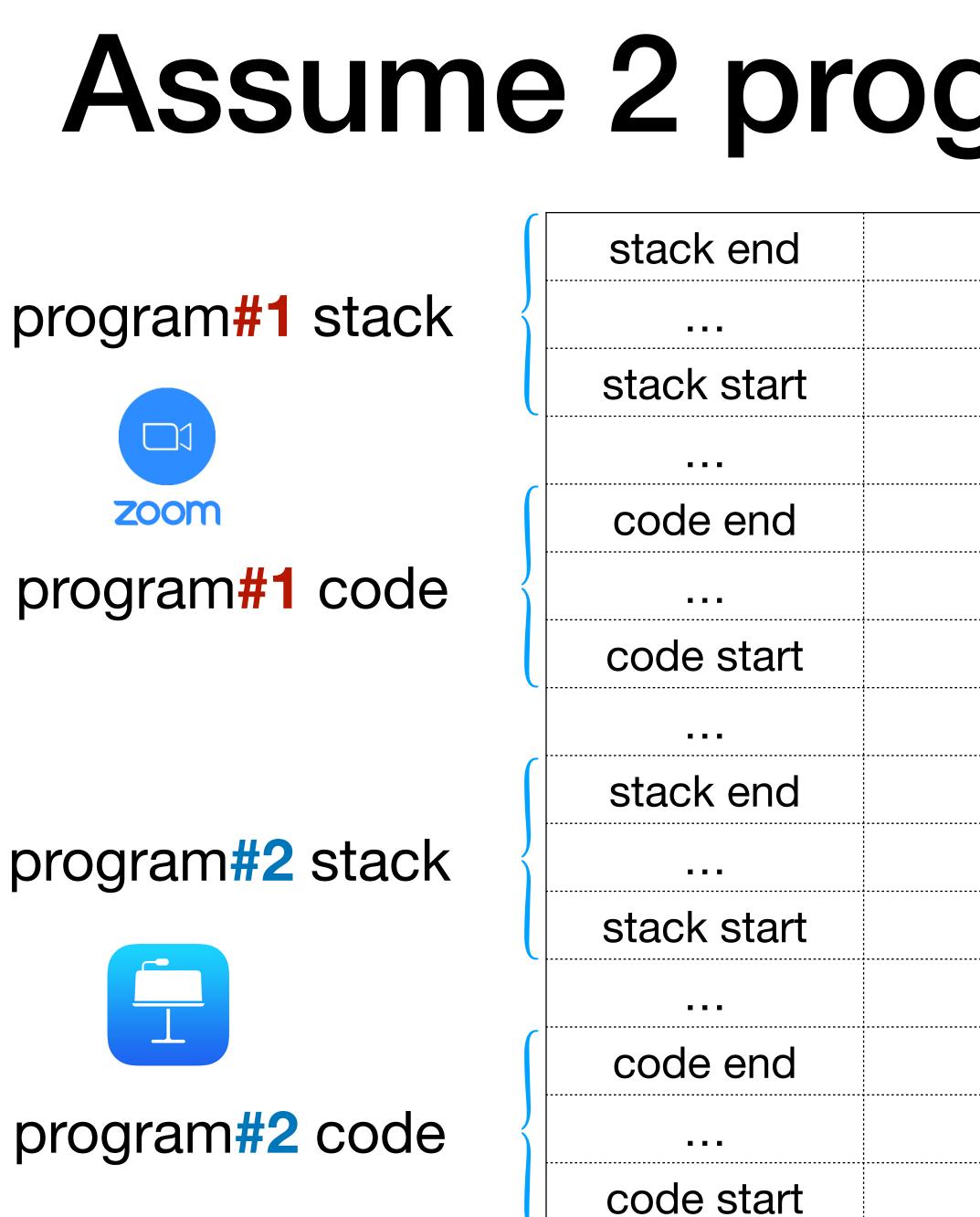
**Context-switch & Threads** 

### Goal of Today's Class

- Understand the concepts of context, context-switch and threads
- Understand the related functions in assignment P1
  - thread\_init, thread\_create, thread\_yield, thread\_exit
  - ctx\_entry, ctx\_start, ctx\_switch

**Review:** the minimal requirement of program execution is code & stack segments in memory address space.



### Assume 2 programs in memory

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

. . .

OS puts the code & stack of both programs in the memory so that they can take turns to execute.



### Review: context defines which program the CPU is executing; context = memory address space + stack pointer + instruction pointer

## CPU in the context of program #1

### CPU

#### Stack pointer register

#### Instruction pointer register

program <b>#1</b> stack end	
program #1 stack start	
program <b>#1</b> code end	
program <b>#1</b> code start	
program <b>#2</b> stack end	
program <b>#2</b> stack start	
program <b>#2</b> code end	
program <b>#2</b> code start	
	 program <b>#1</b> stack start  program <b>#1</b> code end  program <b>#1</b> code start  program <b>#2</b> stack end  program <b>#2</b> stack start 

## CPU in the context of program #2

#### CPU

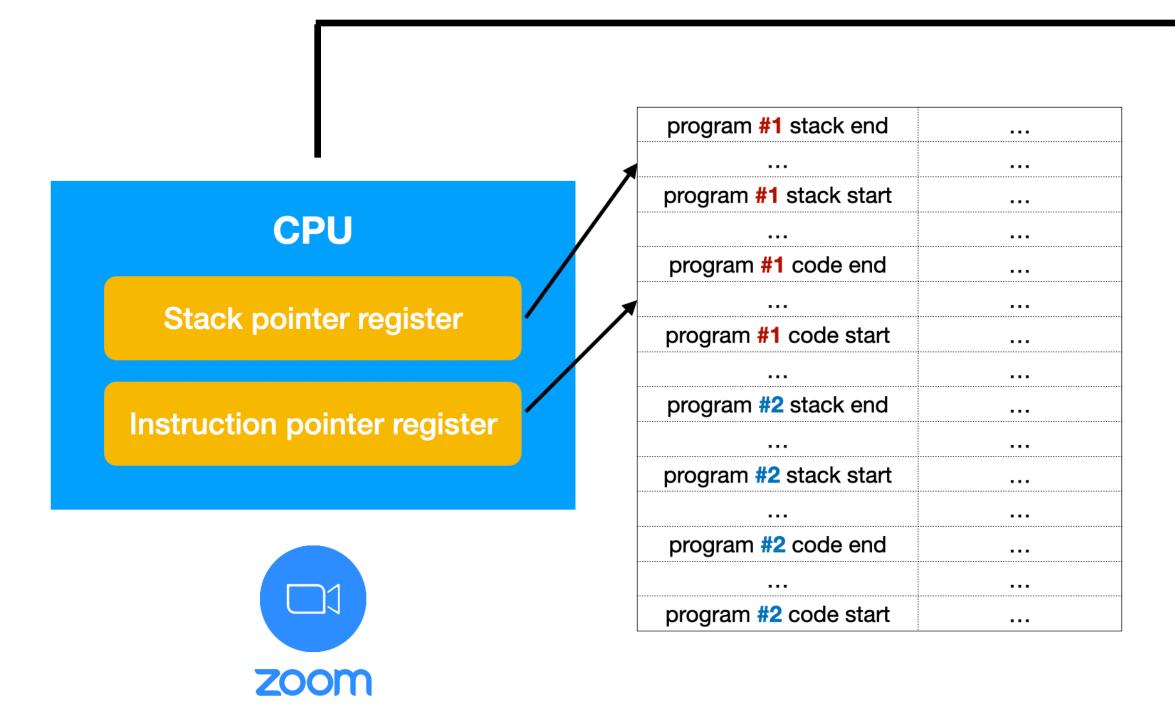
#### Stack pointer register

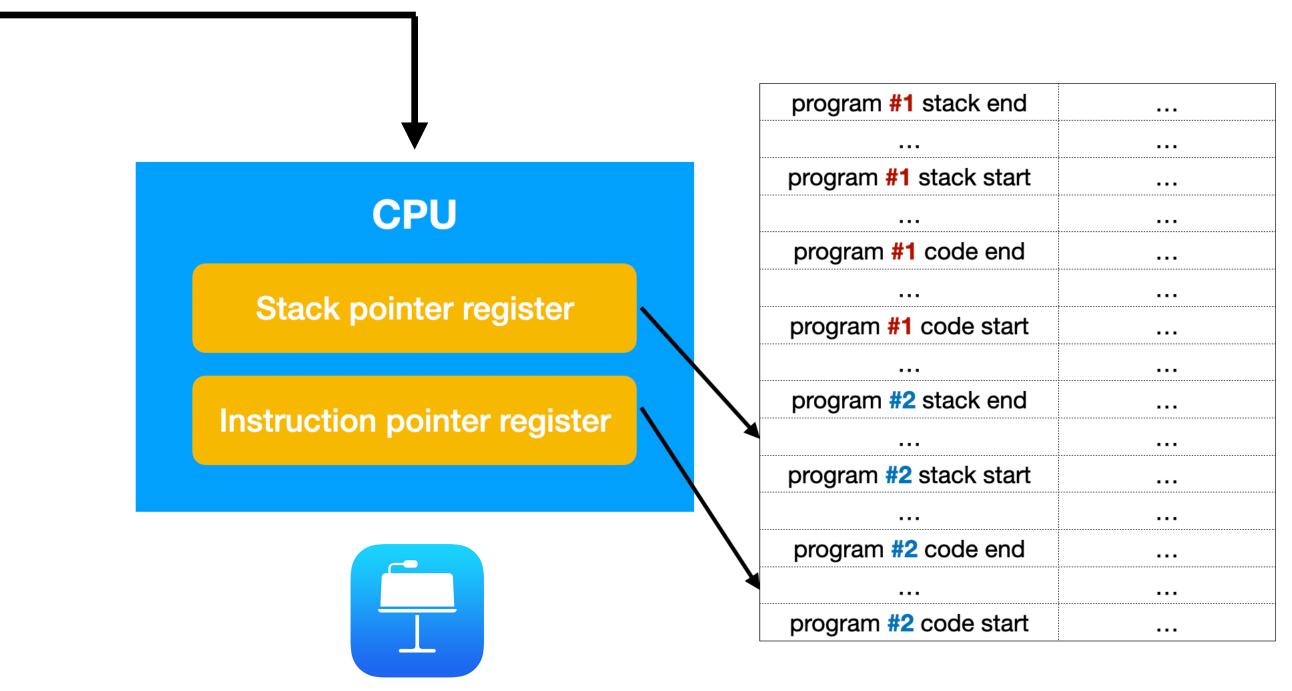
#### Instruction pointer register

_		
	program <b>#1</b> stack end	
	program <b>#1</b> stack start	
	program " otdok otdit	• • •
	program <b>#1</b> code end	
	program <b>#1</b> code start	
	program <b>#2</b> stack end	
	• • •	
	program <b>#2</b> stack start	
	program <b>#2</b> code end	
	• • •	
	program <b>#2</b> code start	

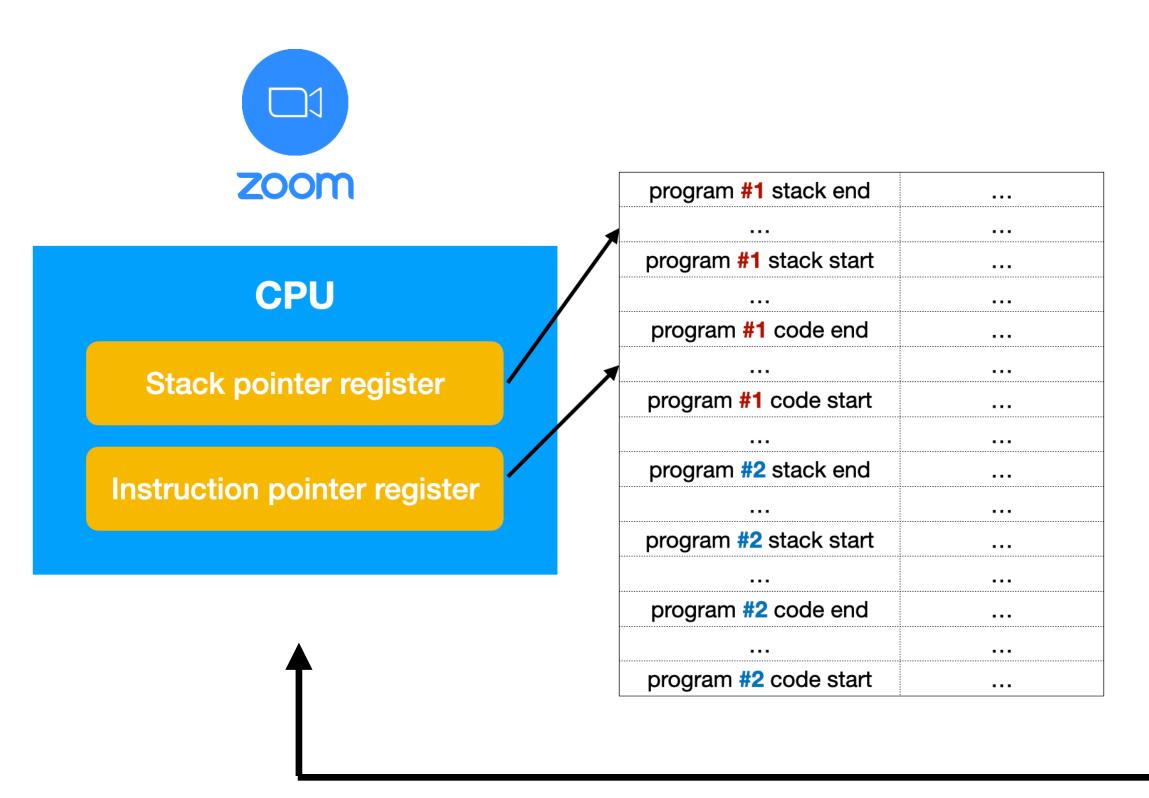
### **Context-switch**

#### CPU switches to the context of program #2

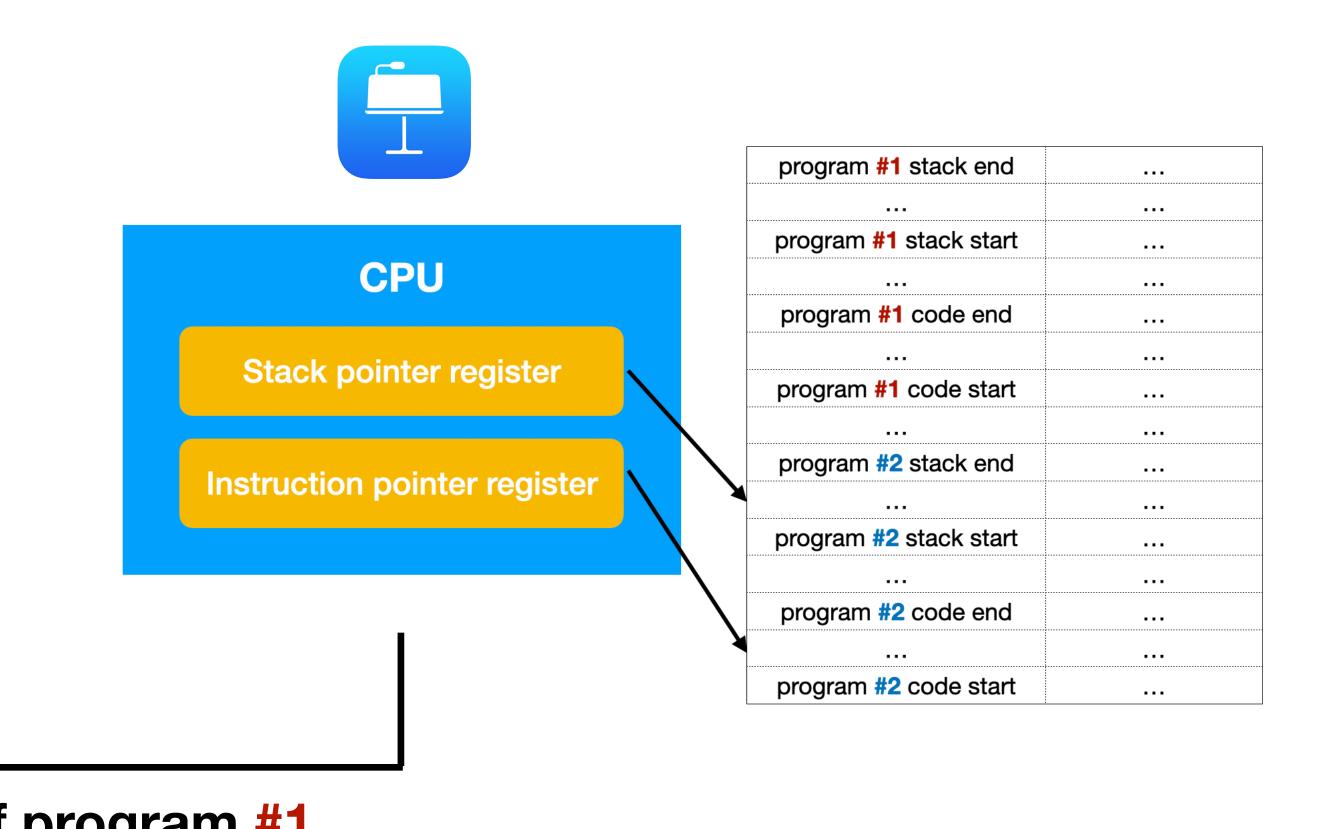




### **Context-switch**



#### CPU switches to the context of program #1



### Question: when does contextswitch happen?

### When does context-switch happen?

- Program terminates.
- Program calls yield system call. (next slide)
- CPU receives a timer interrupt. (later in assignment P2)
- CPU receives an I/O interrupt. (later in assignment P5)





#### int noble\_a() {

yield();

....

A program can occupy the CPU, but it decides to stop and let others to use the CPU first.

### yield is a noble behavior

A car can occupy the road, but it decides to stop and let others to use the road first.

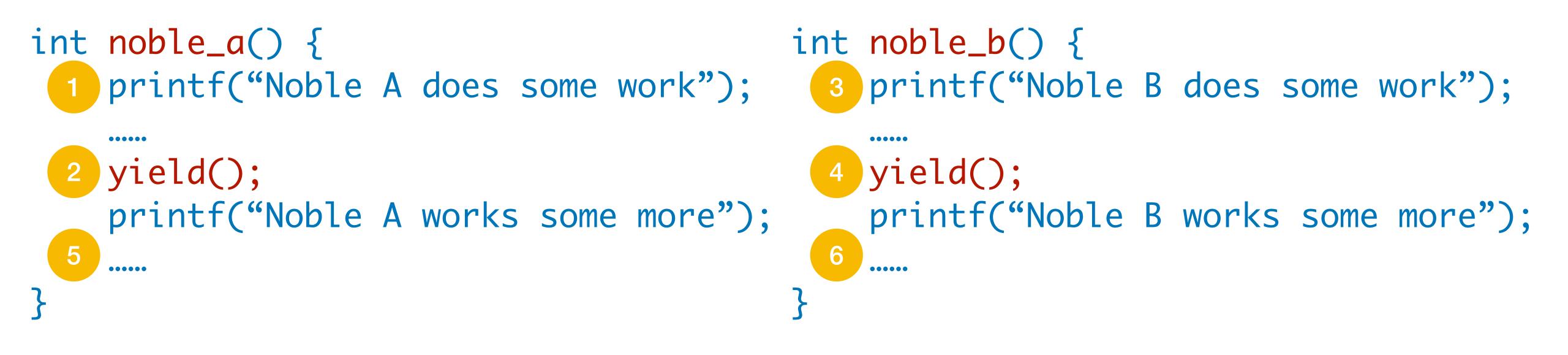
### Two noble functions

### int noble\_a() { •••••

int noble\_b() { printf("Noble A does some work"); printf("Noble B does some work"); ••••• yield(); yield(); printf("Noble A works some more"); printf("Noble B works some more"); ..... .....



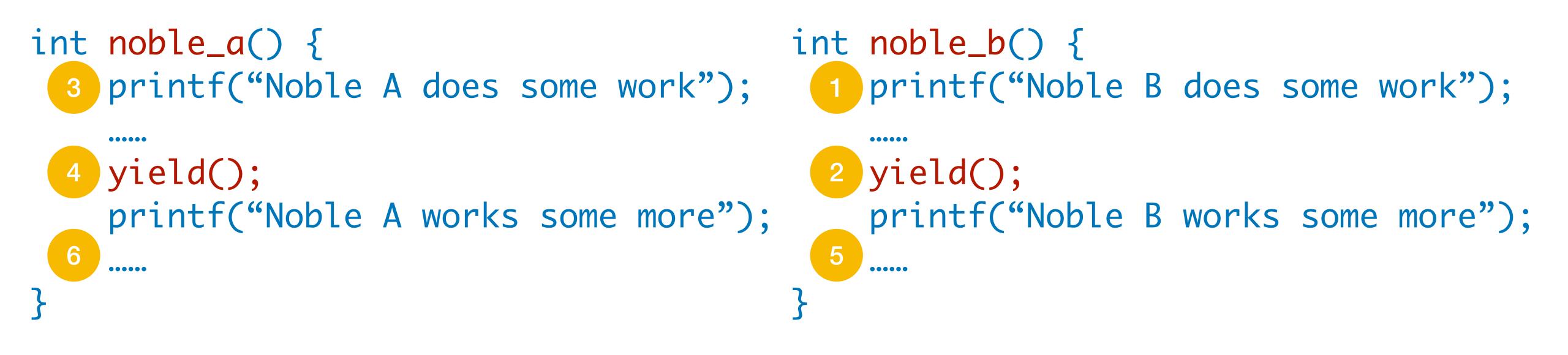
## One possible schedule



Output: Noble A does some work Noble B does some work Noble B works some more

- Noble A works some more

### Another possible schedule



Output: Noble B does some work Noble A does some work Noble A works some more

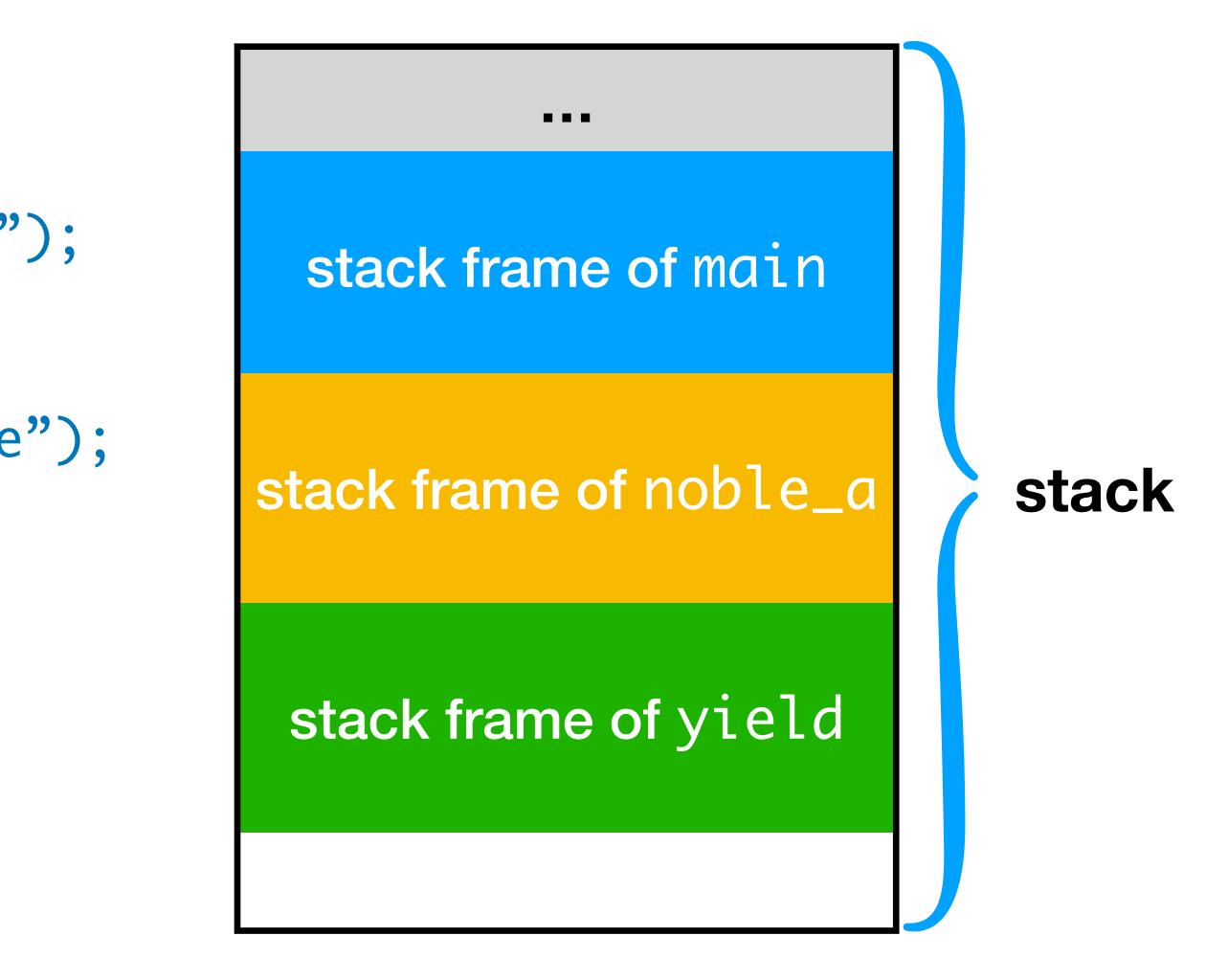
- Noble B works some more

# Question: how do we run two functions at the same time?

Let's review some knowledge of stack.

### Review of stack

```
int noble_a() {
    printf("Noble A does some work");
    ....
    yield();
    printf("Noble A works some more");
    ....
int main() {
    noble_a();
    return 0;
```



main() calls noble\_a() calls yield()



### Review of stack

<b>Only continue when</b> noble_a <b>returns</b>	stack fra
<b>Only continue when</b> yield <b>returns</b>	stack fran
<b>Currently running function</b>	stack fra

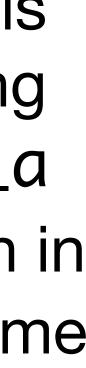
main() calls noble\_a() calls yield()

rame of main

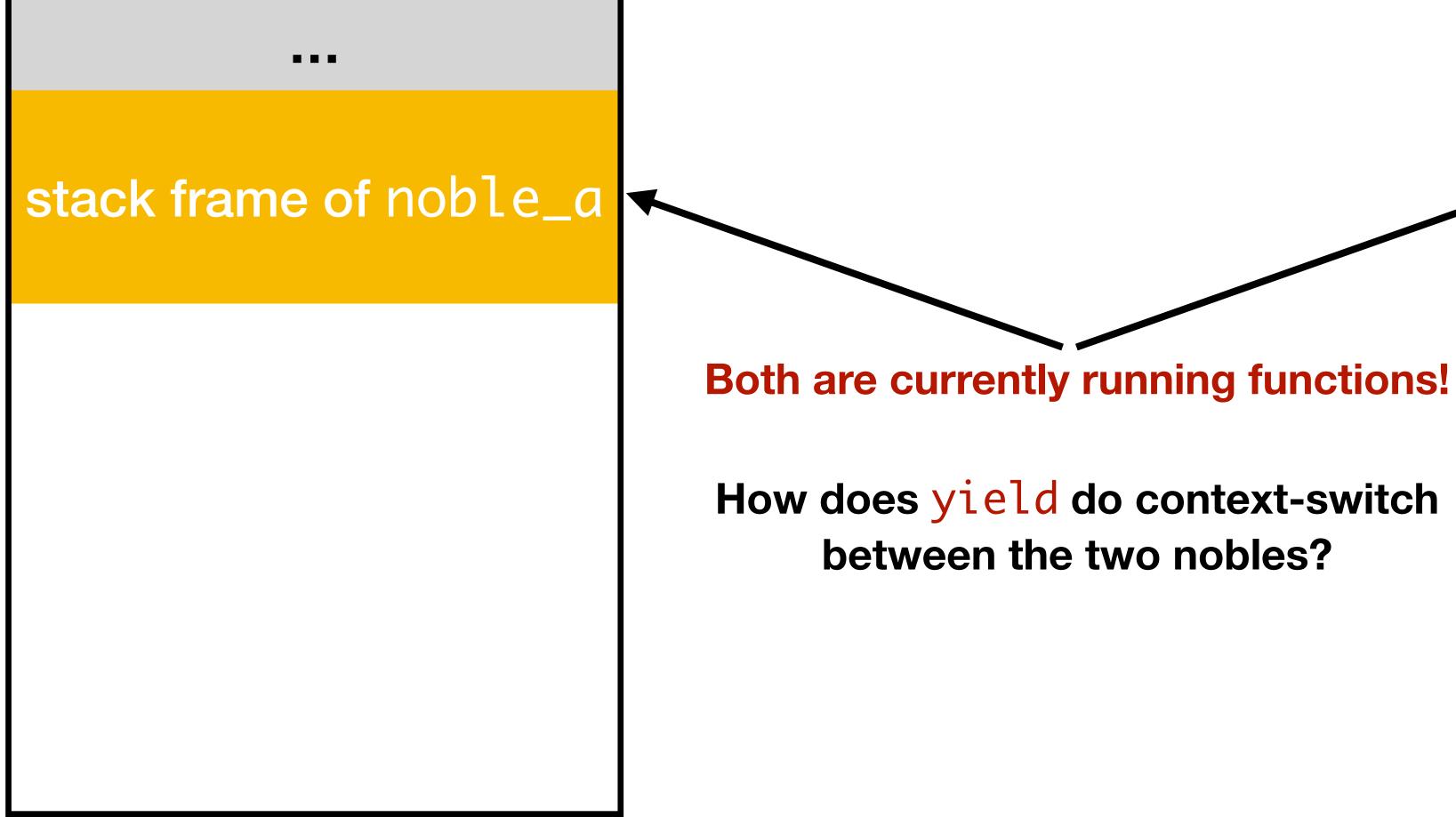
me of noble\_a

ame of yield

For a single stack, there is only one currently running function, so that noble\_a and noble\_b cannot run in the same stack at the same time.



### Two stacks for two nobles



#### stack frame of noble\_b



## Noble A does some work

#### int noble\_a() { 1 printf("Noble A does some work"); •••••

.....

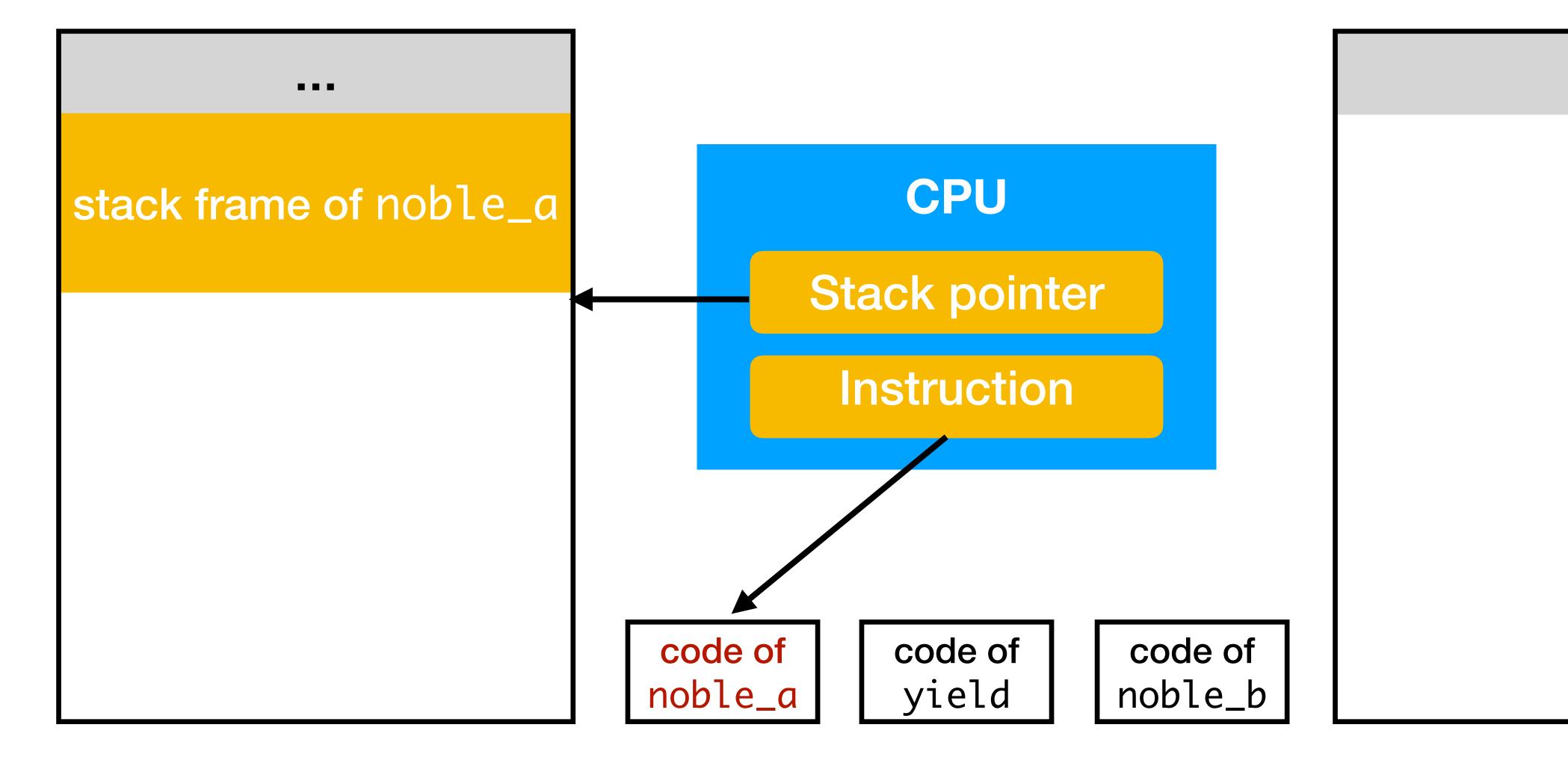
yield(); printf("Noble A works some more");

> Output: Noble A does some work

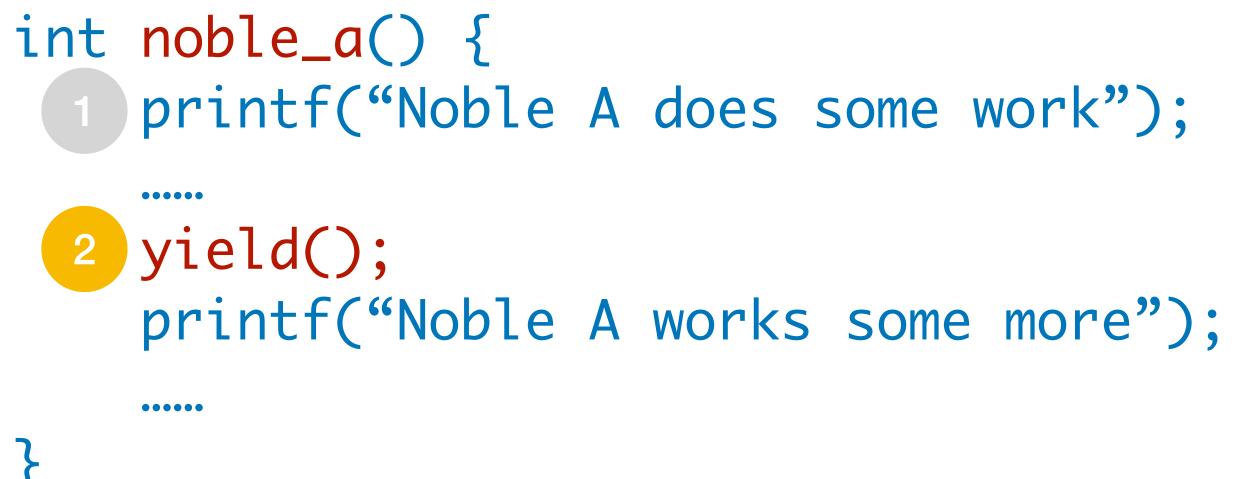
```
int noble_b() {
    printf("Noble B does some work");
    •••••
    yield();
    printf("Noble B works some more");
    .....
```



## CPU in context of noble\_a





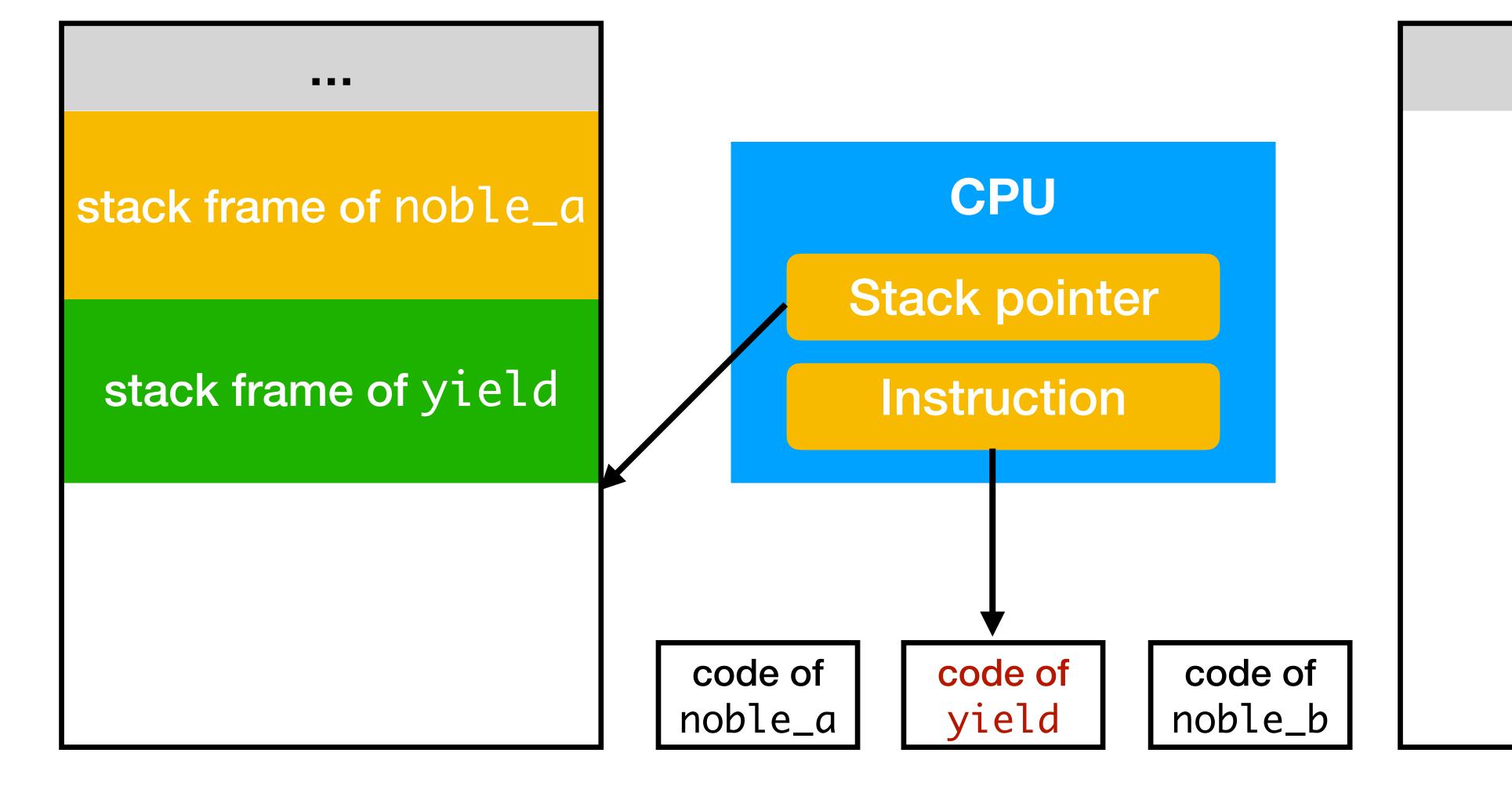


Output: Noble A does some work

## Noble A yields

```
int noble_b() {
    printf("Noble B does some work");
    ....
    yield();
    printf("Noble B works some more");
    •••••
```





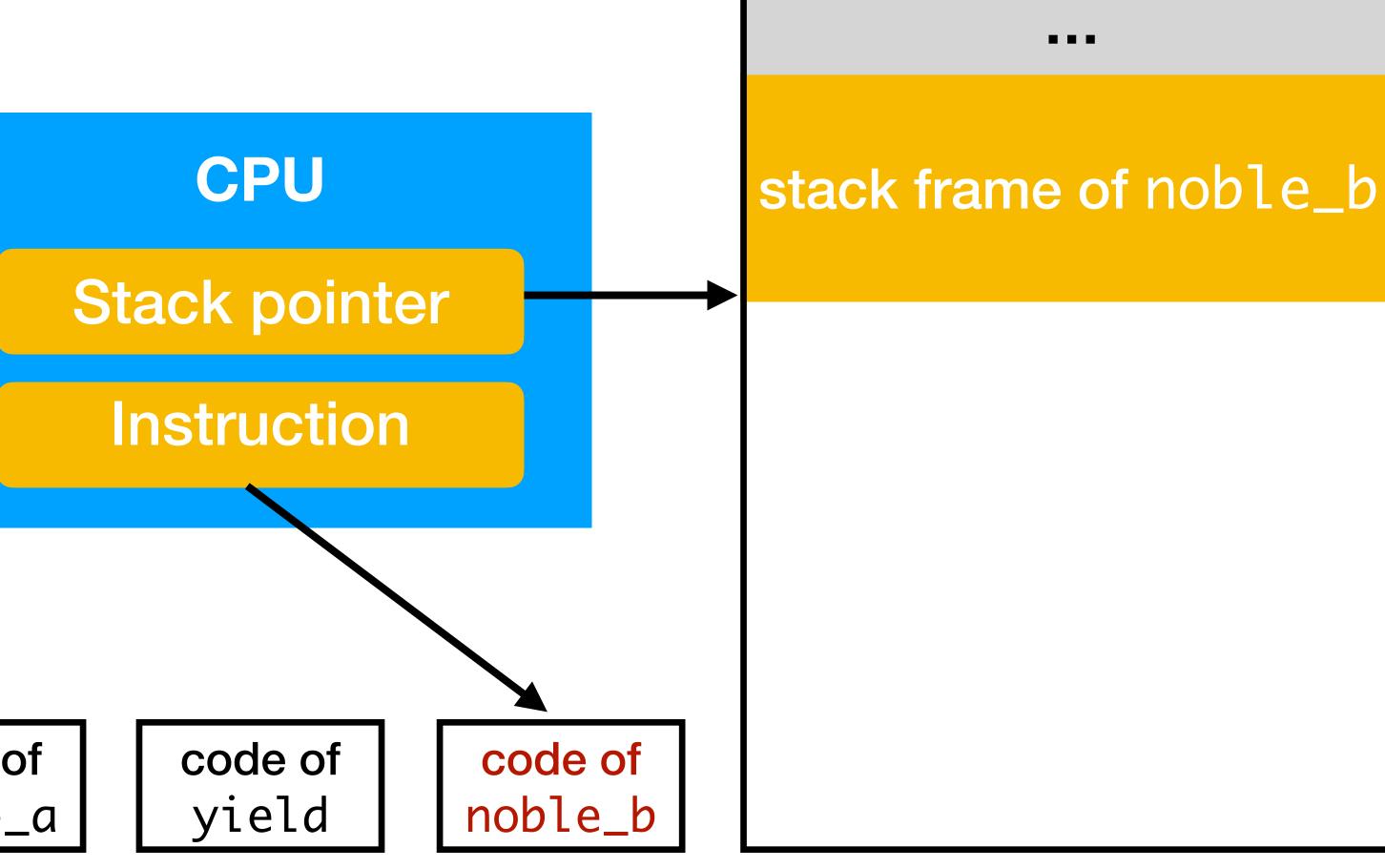
## noble\_a calls yield



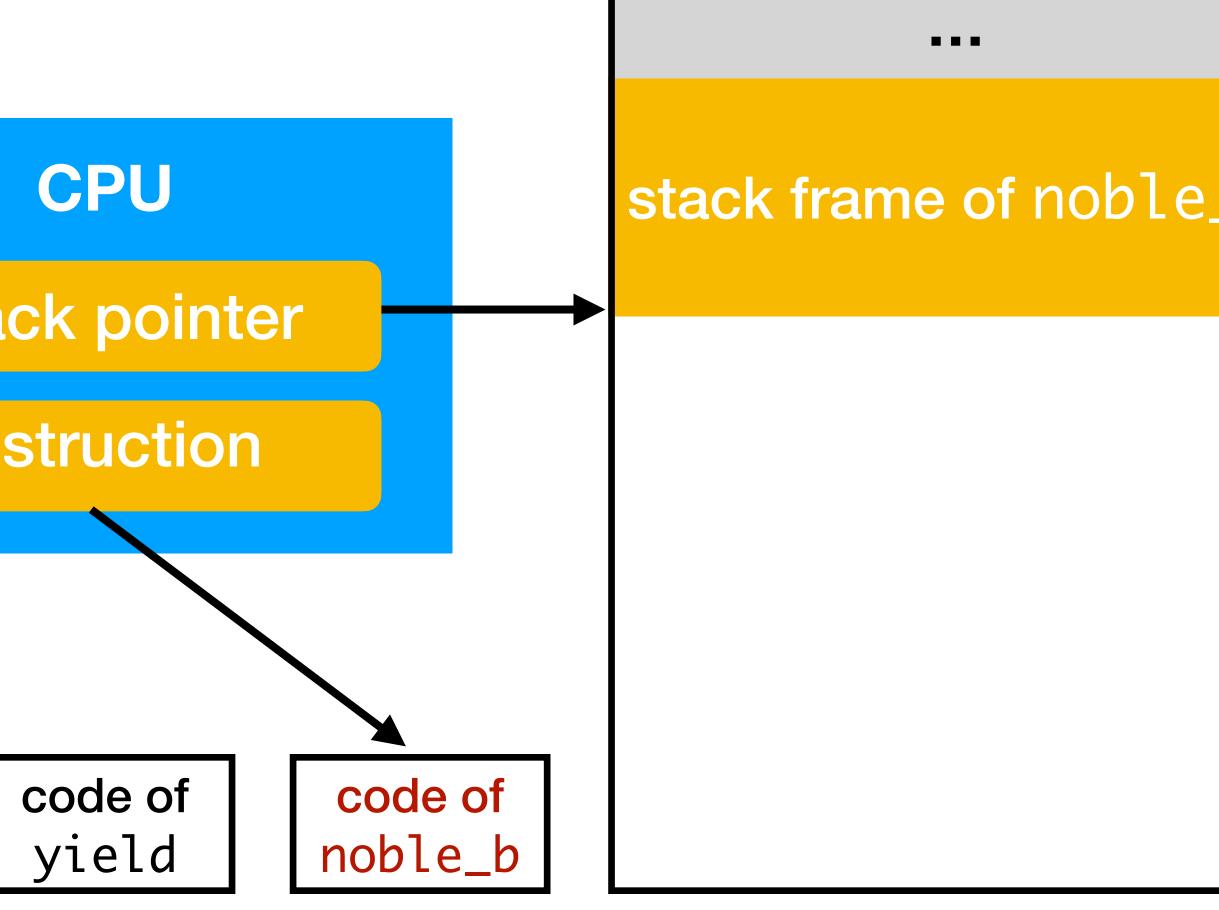
### yield switches context to noble\_b

#### stack frame of noble\_a

#### stack frame of yield

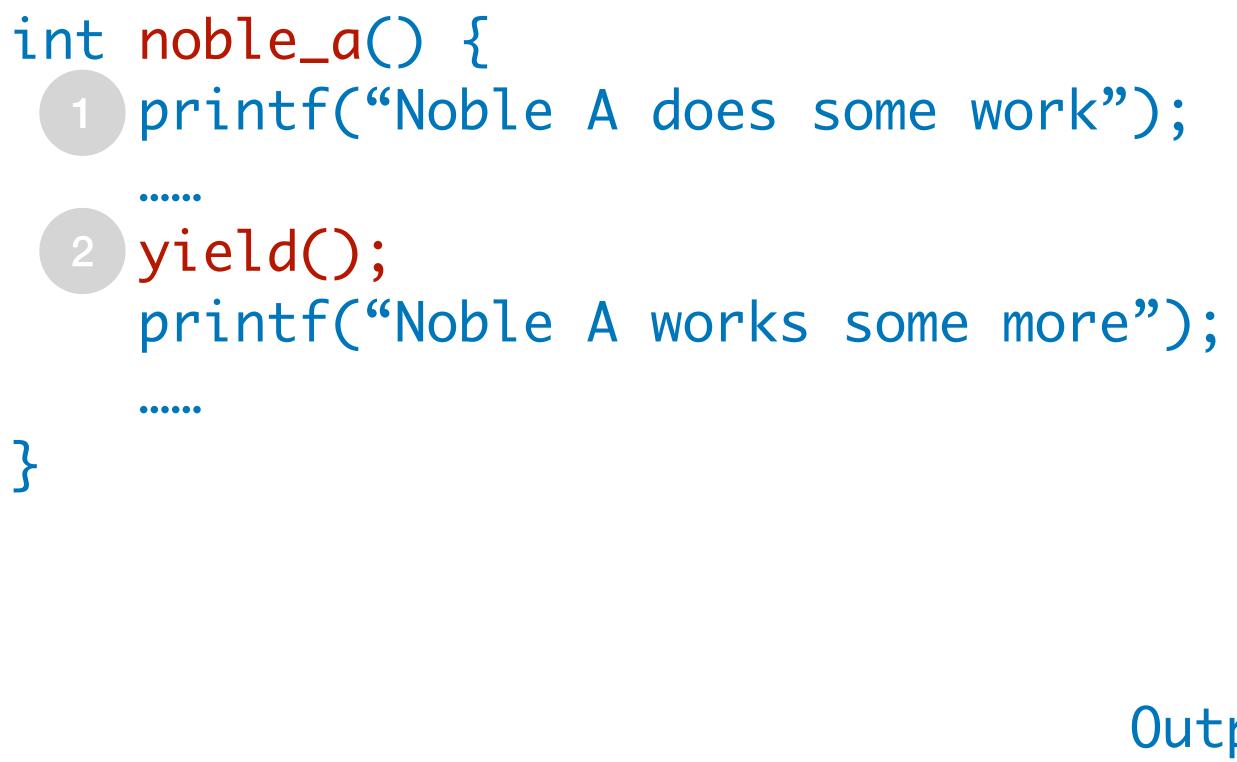


code of noble\_a





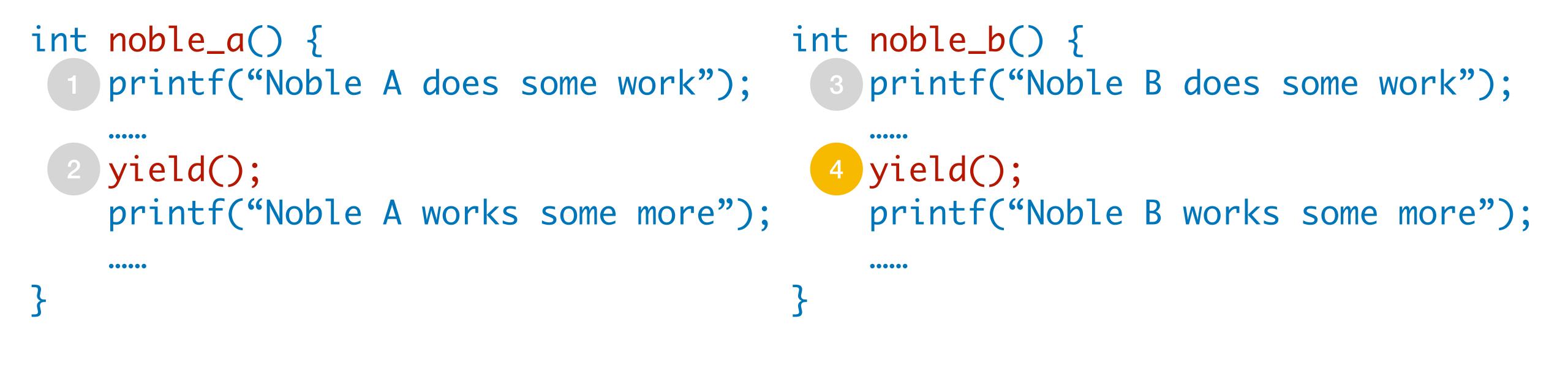
## Noble B does some work



```
int noble_b() {
 3 printf("Noble B does some work");
    ....
    yield();
    printf("Noble B works some more");
    .....
```

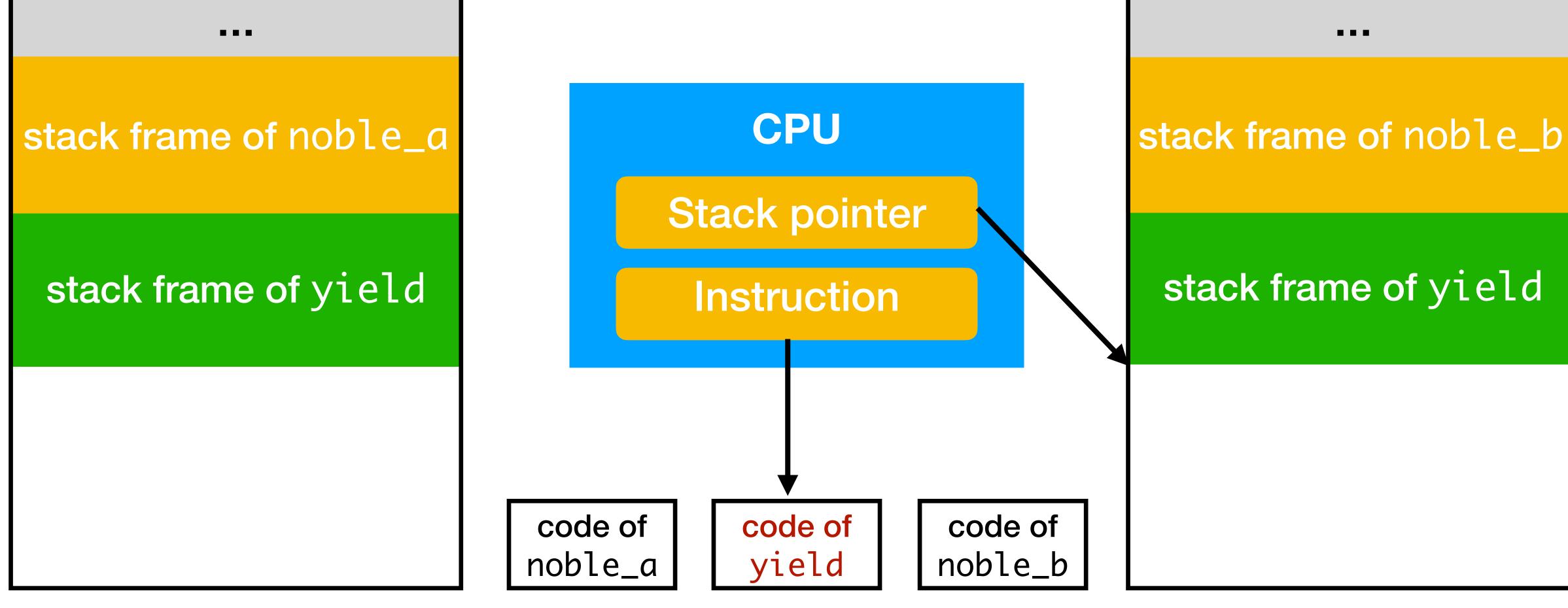
Output: Noble A does some work Noble B does some work





## Noble B yields

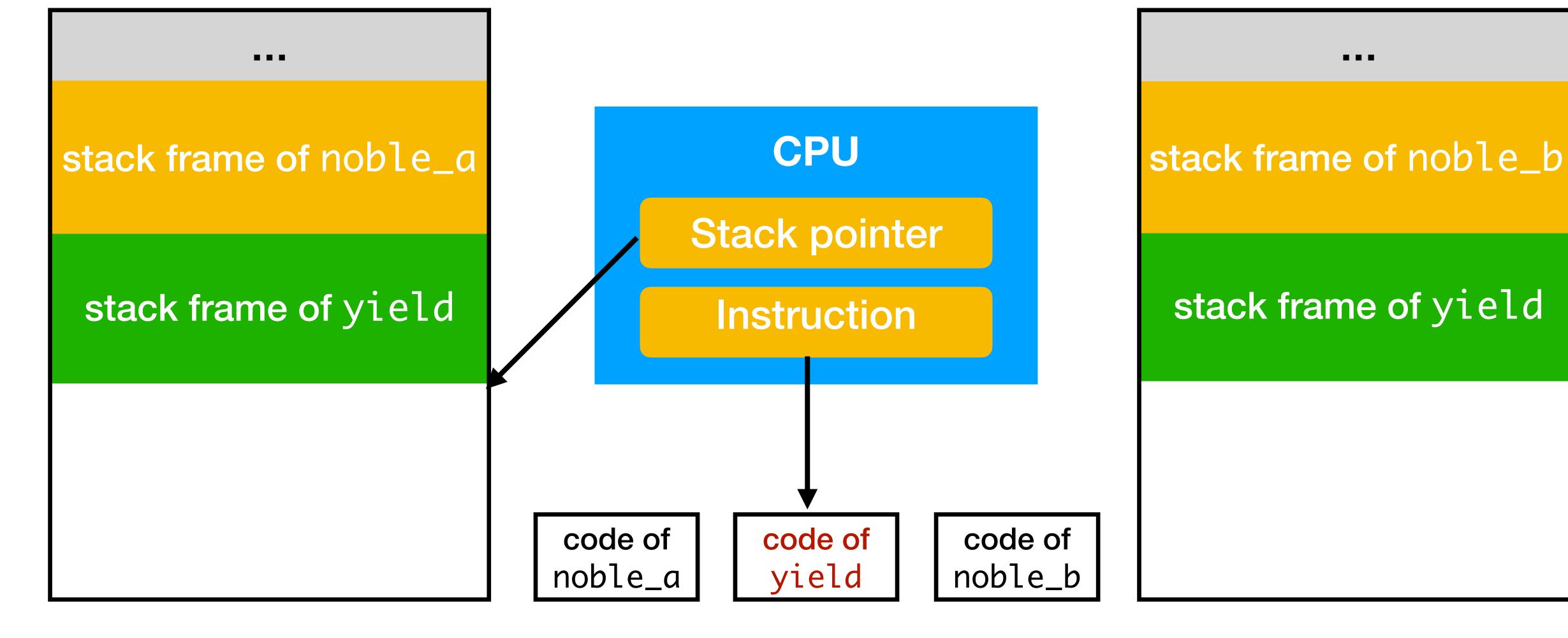
Output: Noble A does some work Noble B does some work



## noble\_b calls yield

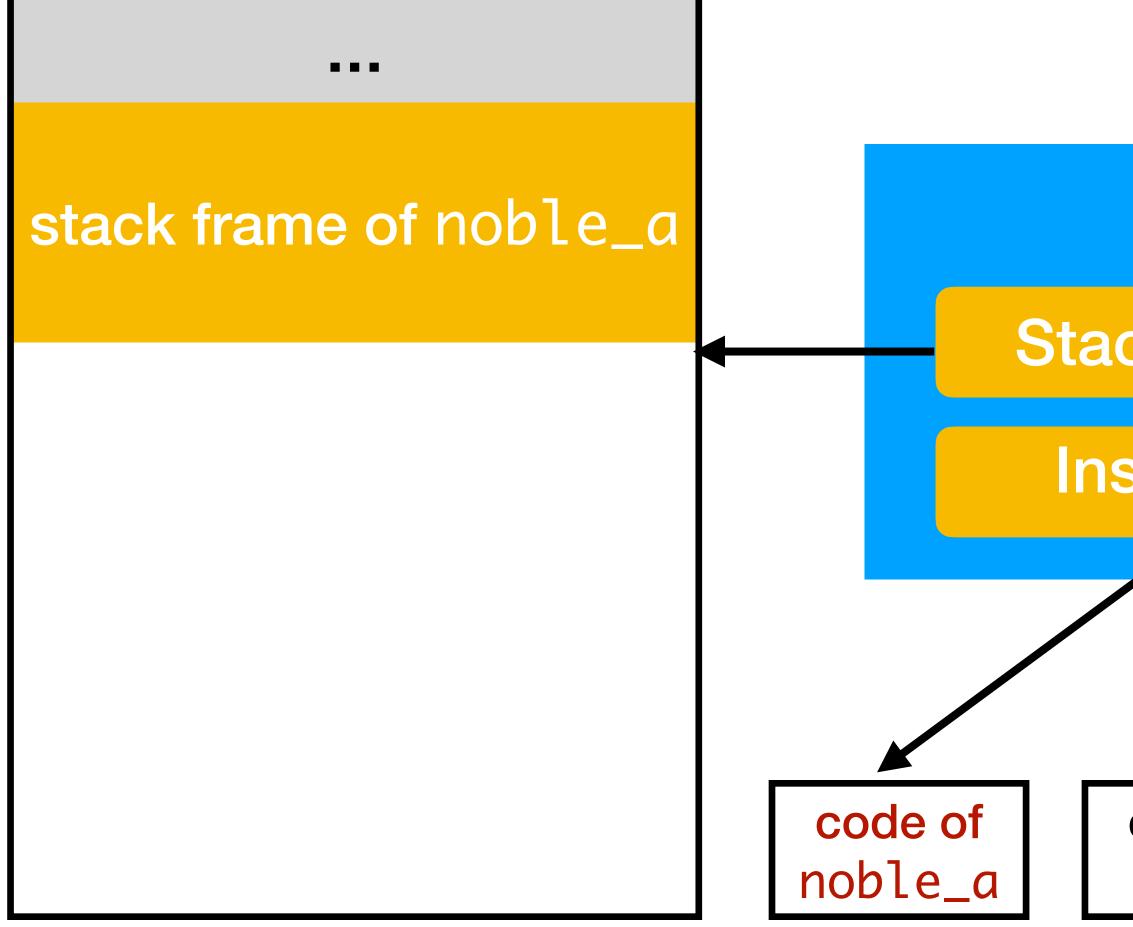


### yield switches context to noble\_a





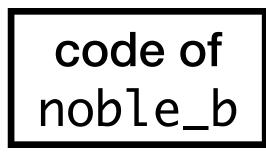
## yield returns to noble\_a



### CPU Stack pointer

Instruction

code of yield

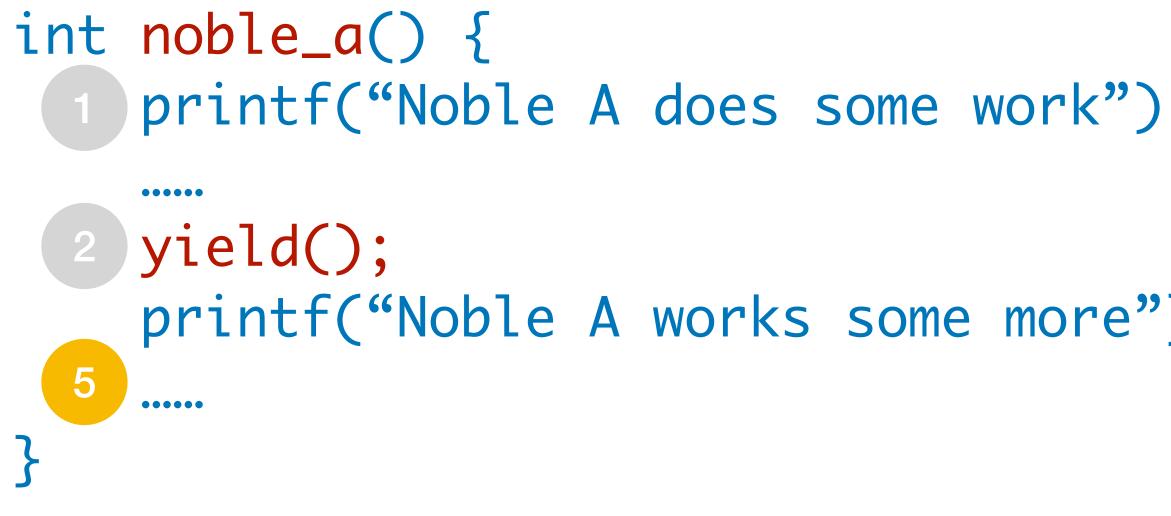


#### stack frame of noble\_b

#### stack frame of yield



### Noble A works some more

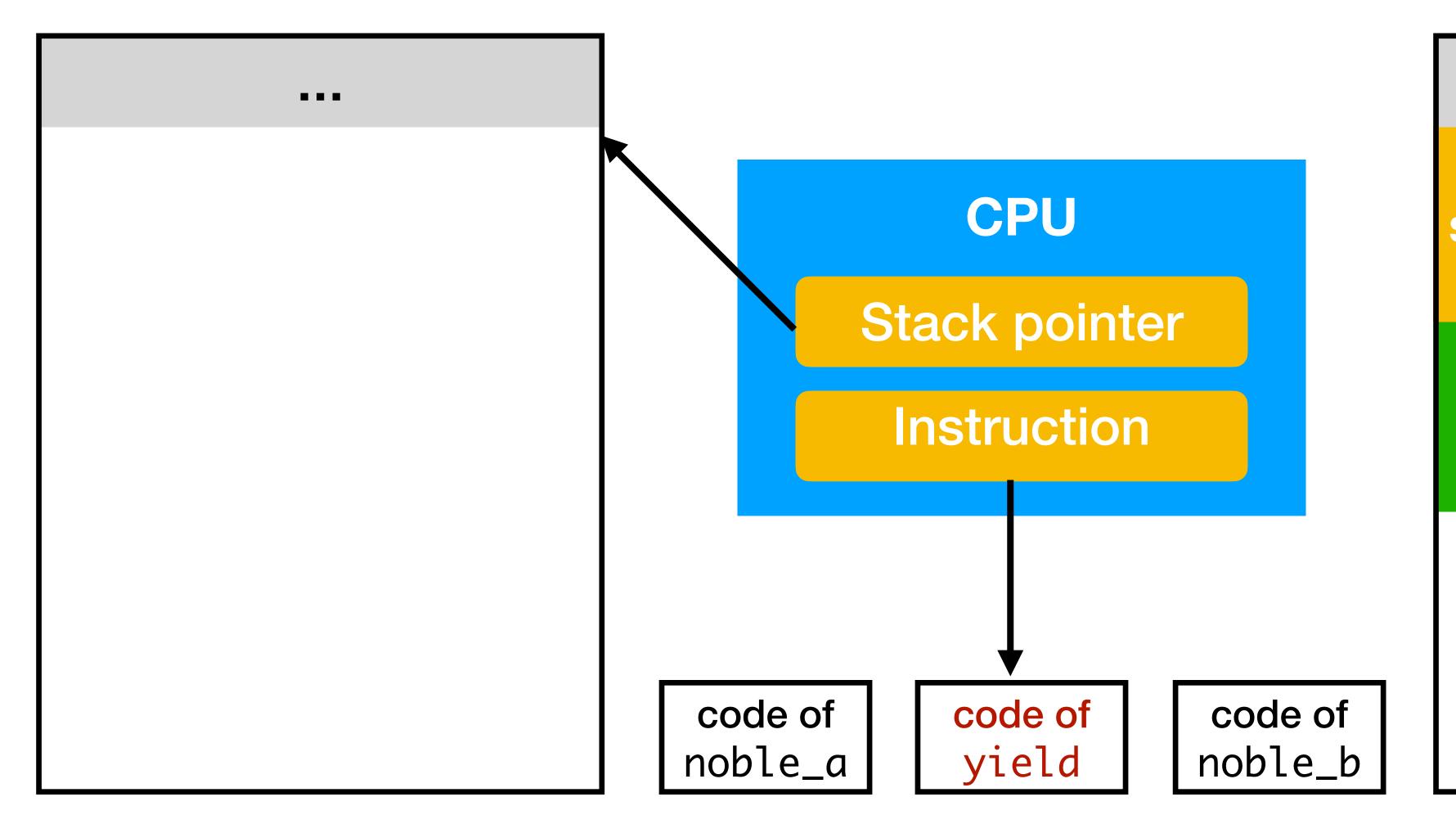


Output: Noble A does some work Noble B does some work Noble A works some more

•		<pre>noble_b() { printf("Noble</pre>	В	does	some	work"
);	4	<pre>yield(); printf("Noble</pre>	B	works	5 SOMe	e more



### noble\_a terminates (implicit yield)

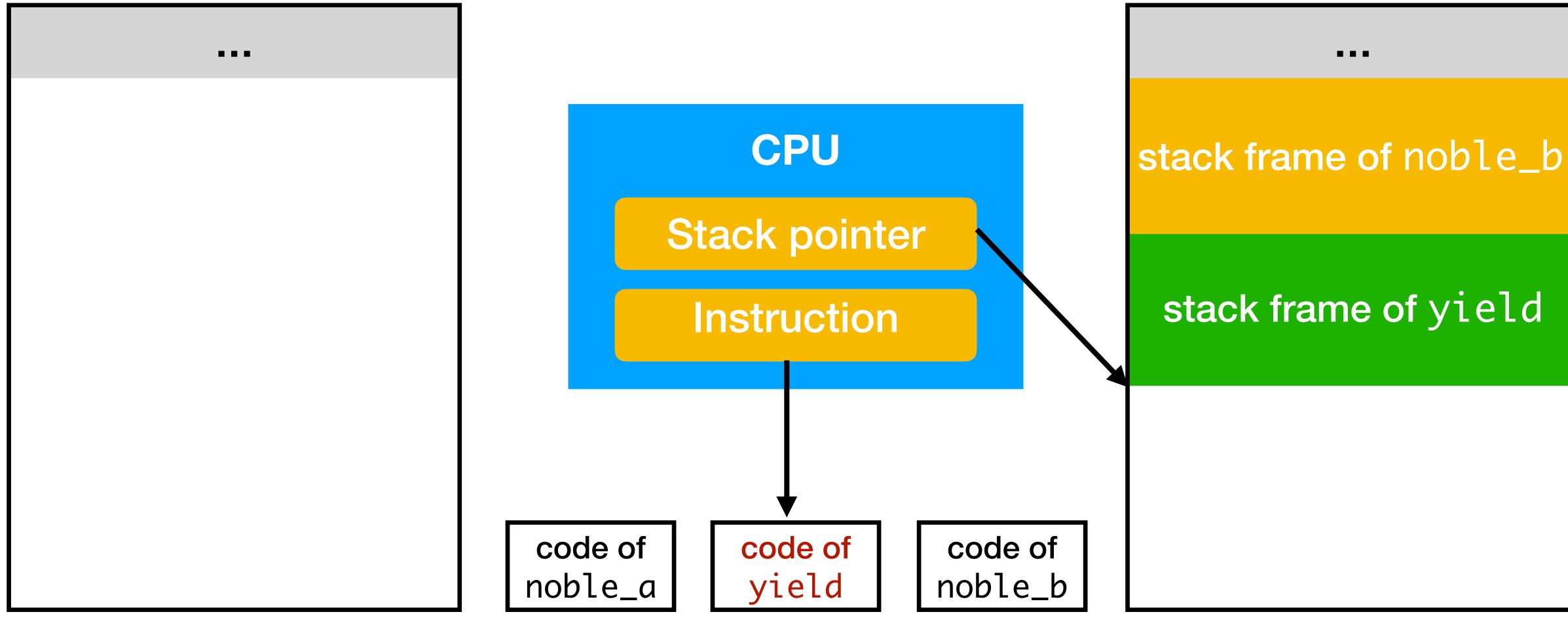


stack frame of noble\_b

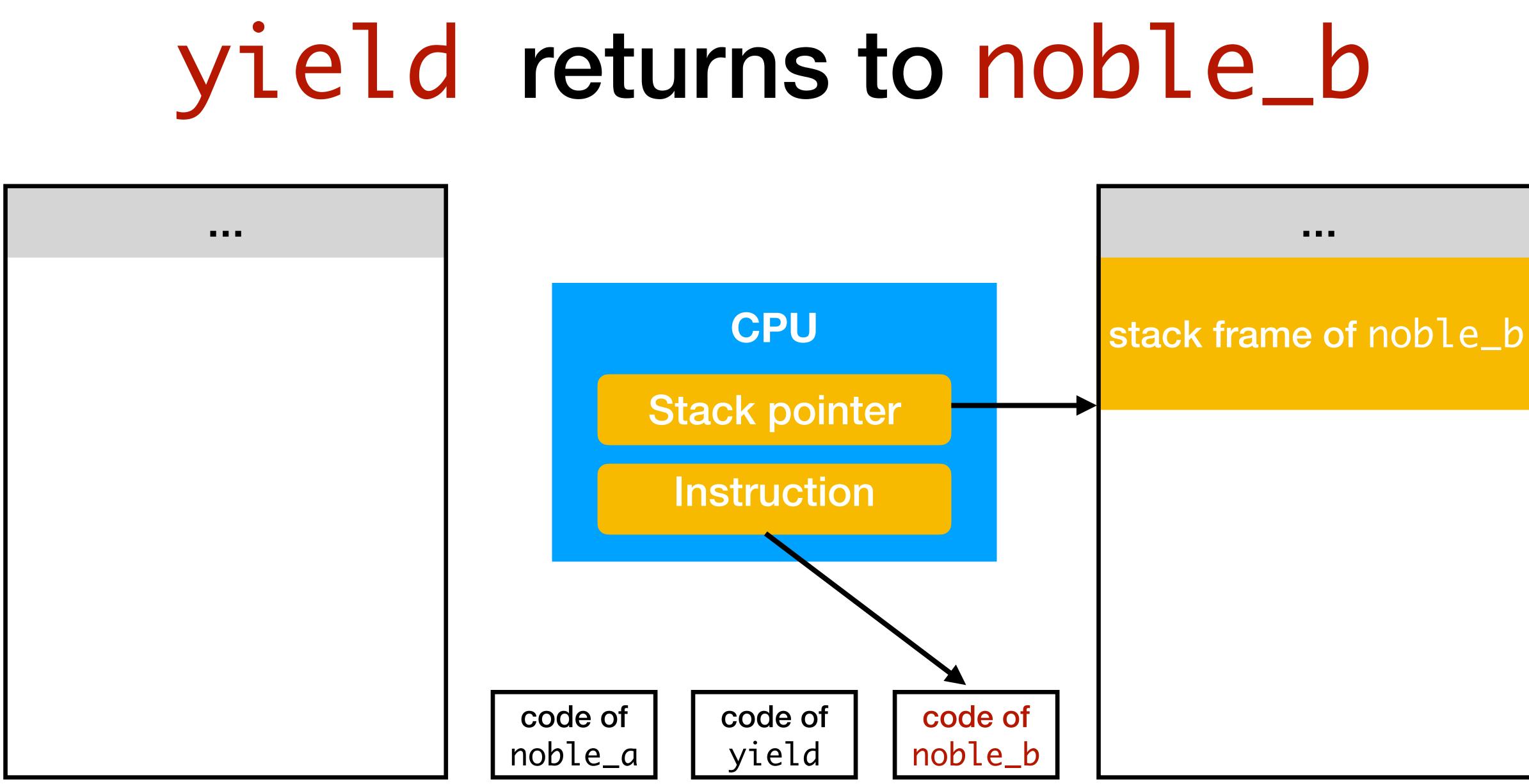
#### stack frame of yield



### yield switches context to noble\_b

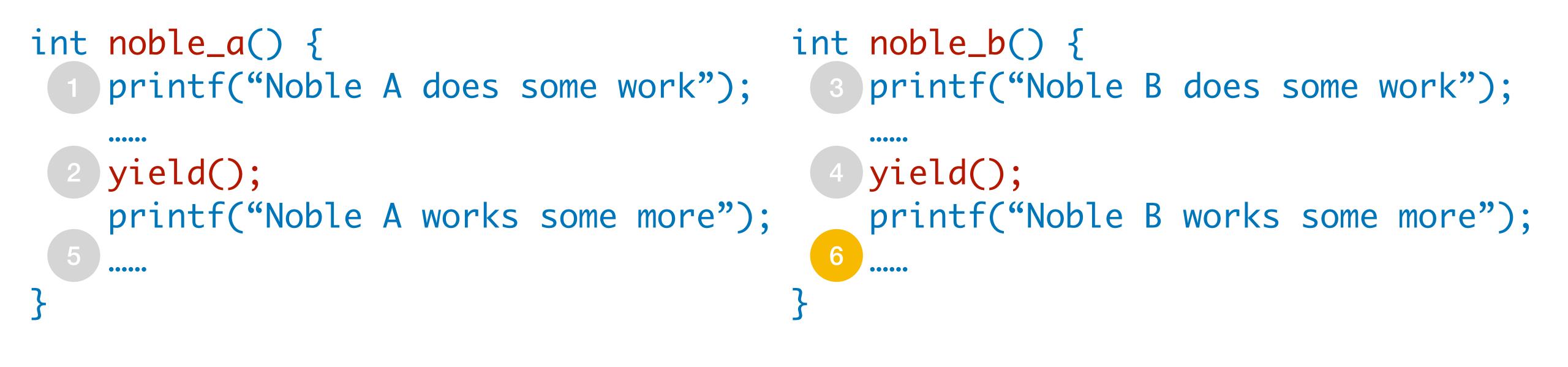






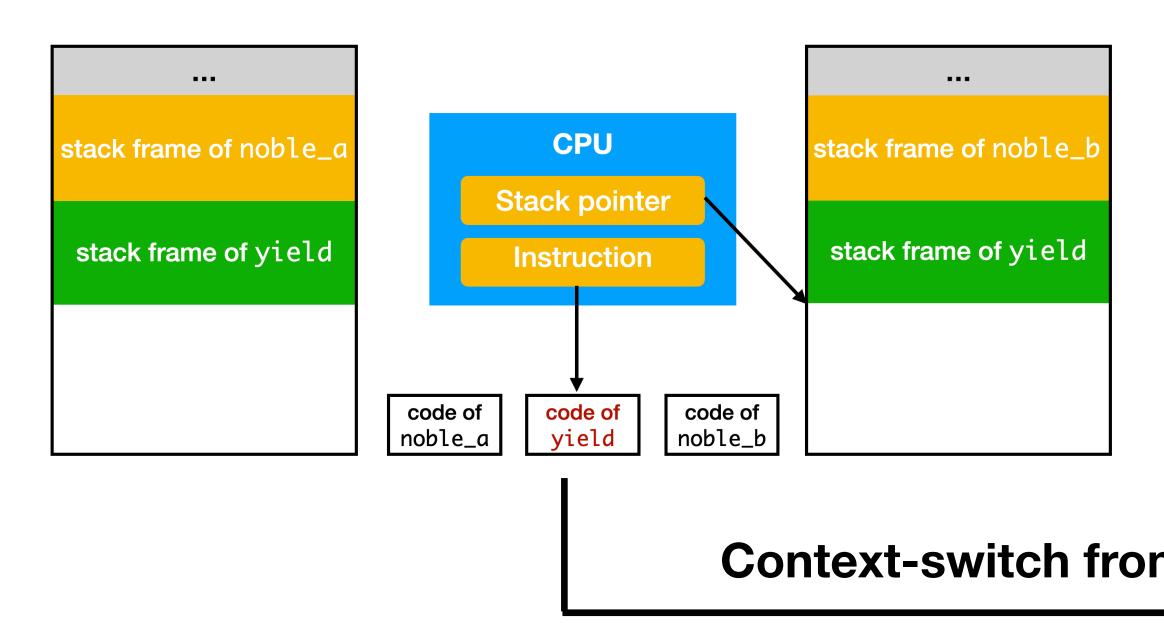


### Noble B works some more

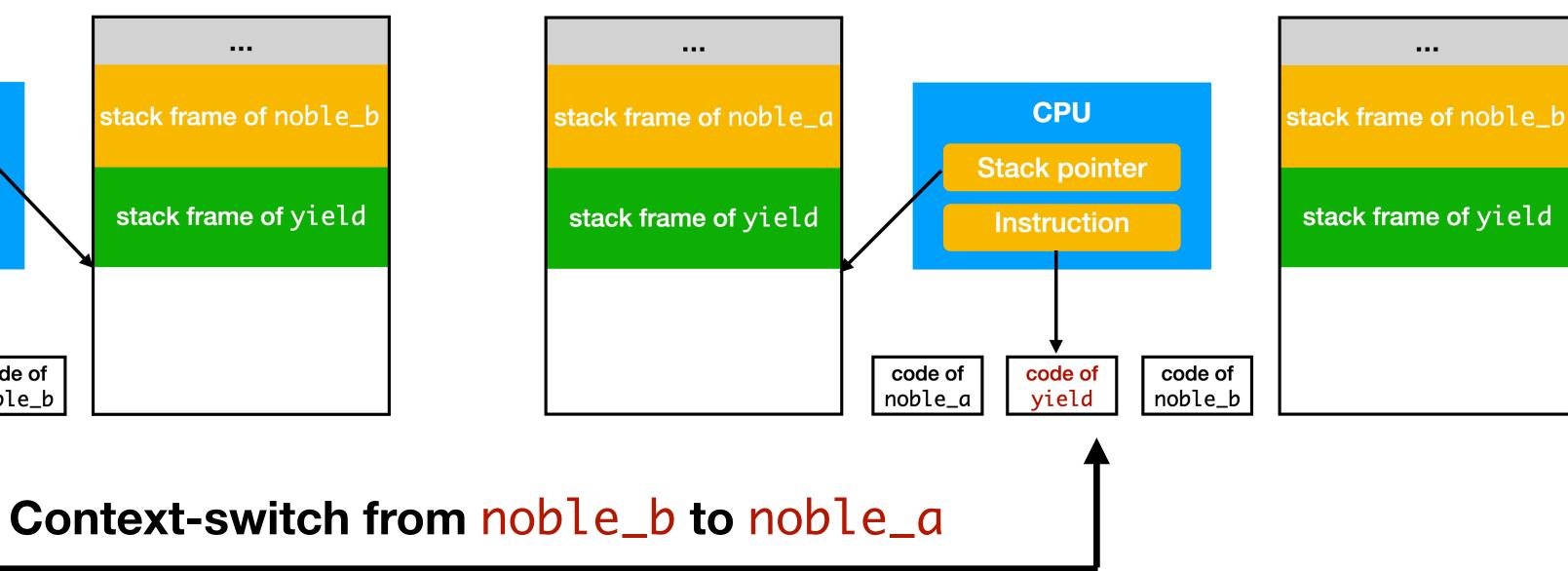


Output: Noble A does some work Noble B does some work Noble A works some more Noble B works some more

### How does yield do context-switch?



- The answer is simple: change the stack pointer!
  - when switching from noble\_a to noble\_b, the yield function needs to record the stack pointer of noble\_a
  - when switching back to noble\_a, the yield function restores the stack pointer of noble\_a





### **Recall: when does context-switch happen?**

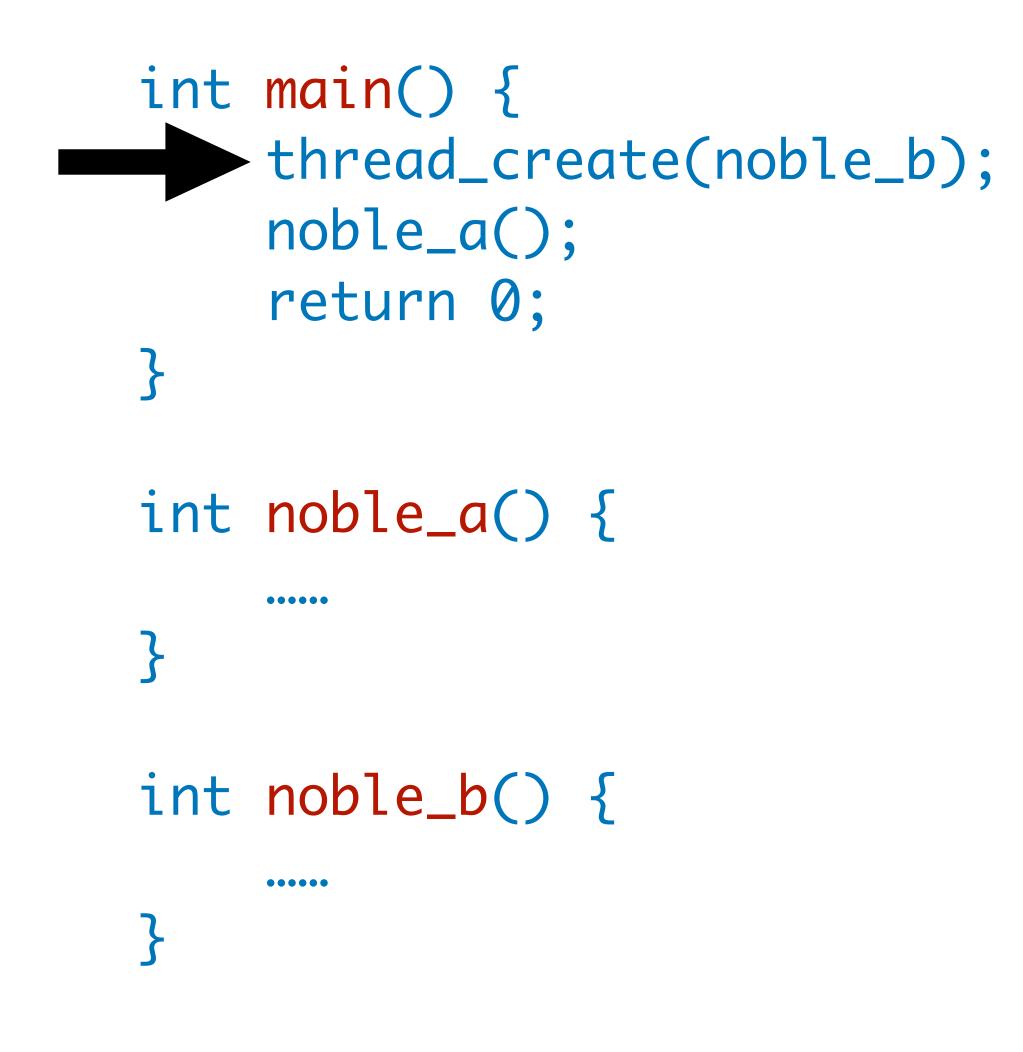


Program calls yield system call.

- CPU receives a timer interrupt. (later in assignment P2)
- CPU receives an I/O interrupt. (later in assignment P5)

# Lesson: A thread owns a stack running a given function.

### A thread owns a stack running a given function



#### stack of noble\_b thread

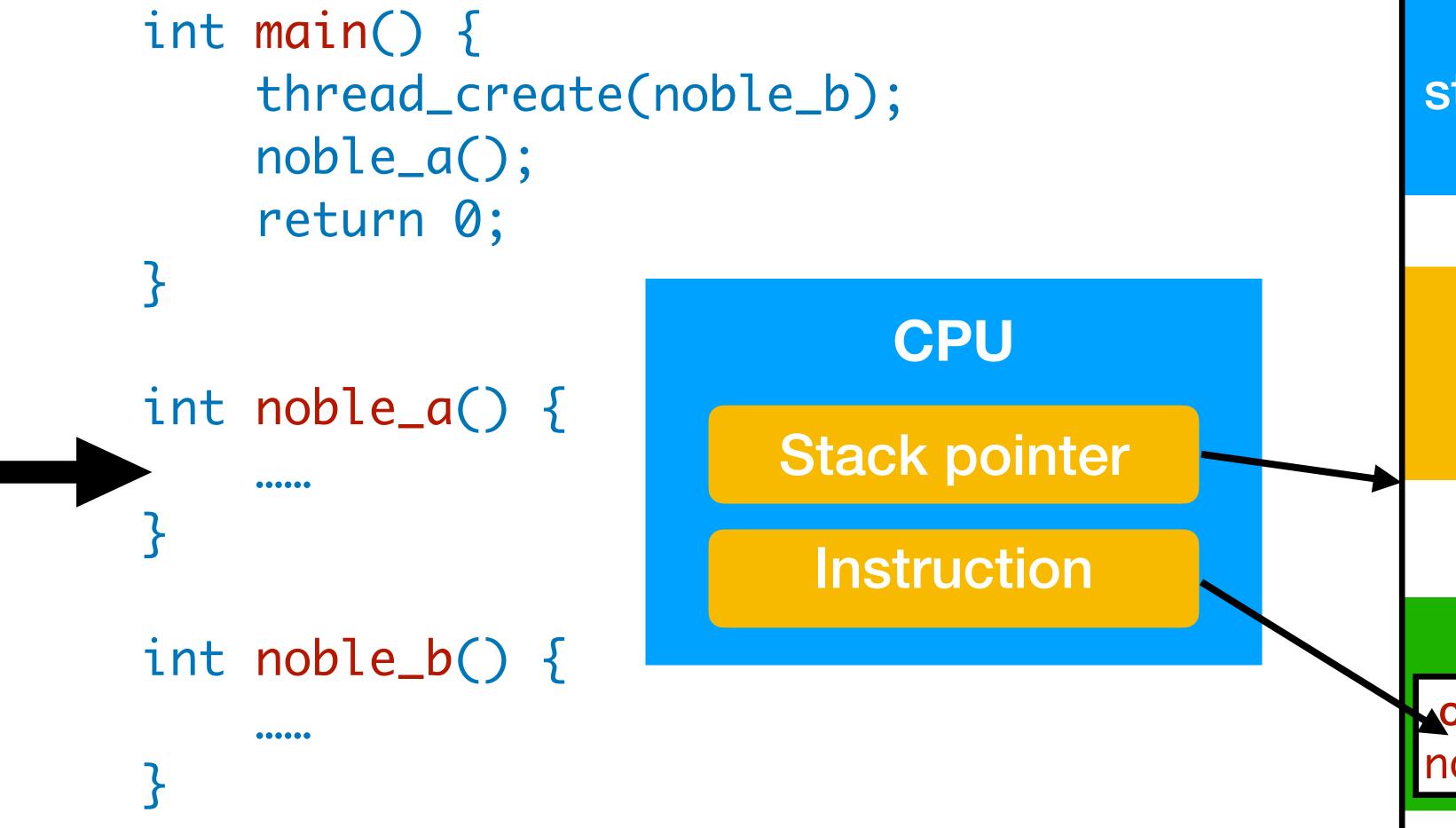
#### stack of main thread

#### code

code ofcode ofnoble\_ayieldnoble\_b



### **Context-switch between threads**



stack of noble\_b thread

#### stack of main thread

#### code

**L**code of noble\_a yield noble\_b

code of



## **Context-switch between threads**

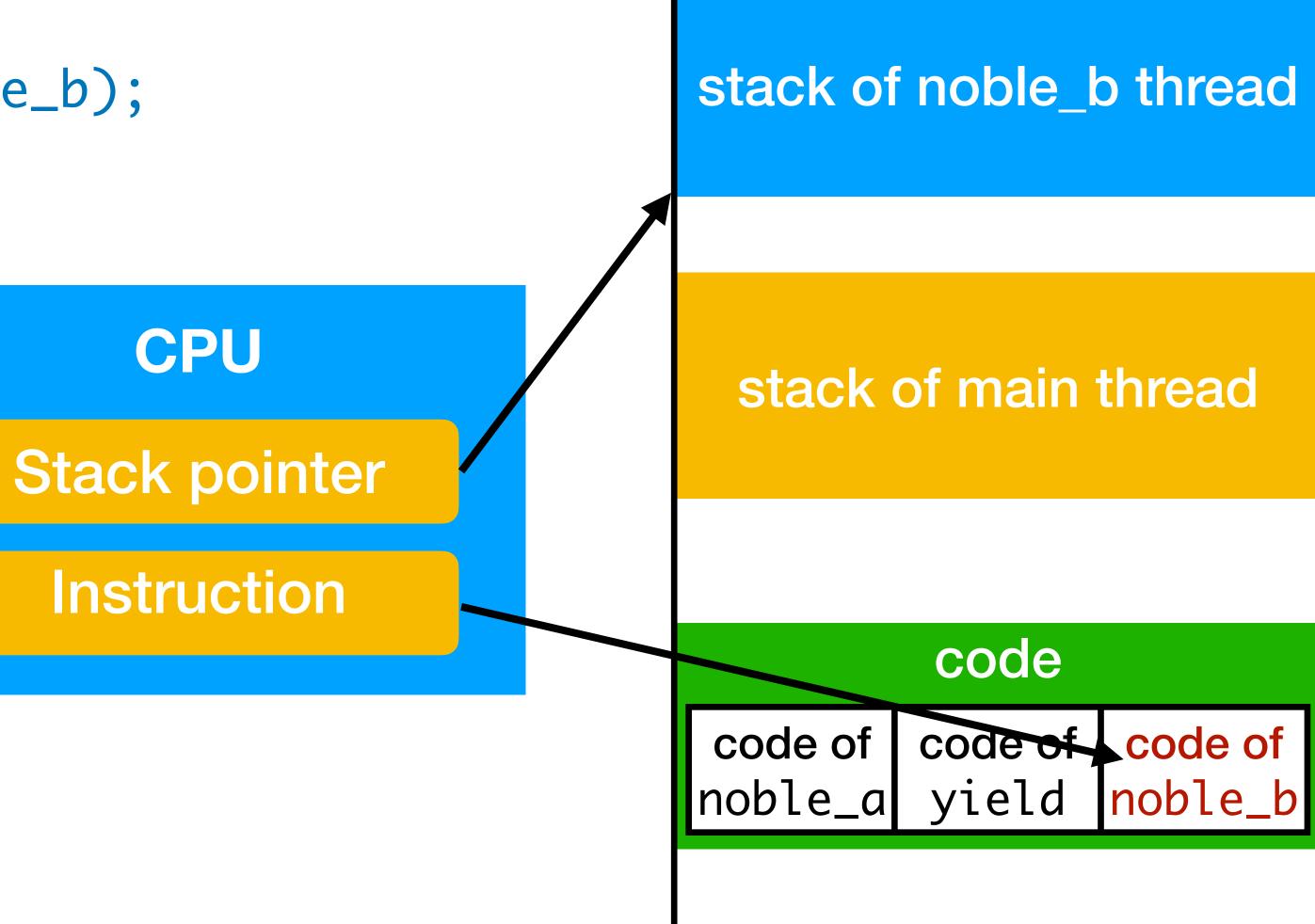
int main() { thread\_create(noble\_b); noble\_a(); return 0;

int noble\_a() {

....

....

int noble\_b() {





### Goal of Today's Class



- Understand the related functions in assignment P1
  - thread\_init, thread\_create, thread\_yield, thread\_exit
  - ctx\_entry, ctx\_start, ctx\_switch

## thread\_init

struct thread {

- // stored stack pointer for yield // \*func to call
- // \*arg to the function
- // ..... feel free to add new fields };
- struct thread current\_thread; struct queue\_t runnable\_threads;

int main() { // initialize the two global variables thread\_init(); return 0;

They live here so shared by threads.

#### stack of main thread

#### data code

code of code of code of noble\_a yield noble\_b



### thread\_create

struct thread current\_thread; struct queue\_t runnable\_threads;

void noble\_b(void\* arg) {

....

int main() { // initialize the two global variables thread\_init(); // create a thread by modifying // the global variables thread\_create(noble\_b, NULL, 16 \* 1024); return 0;

#### stack of main thread

#### data code code of code of code of noble\_a yield noble\_b



### thread\_create

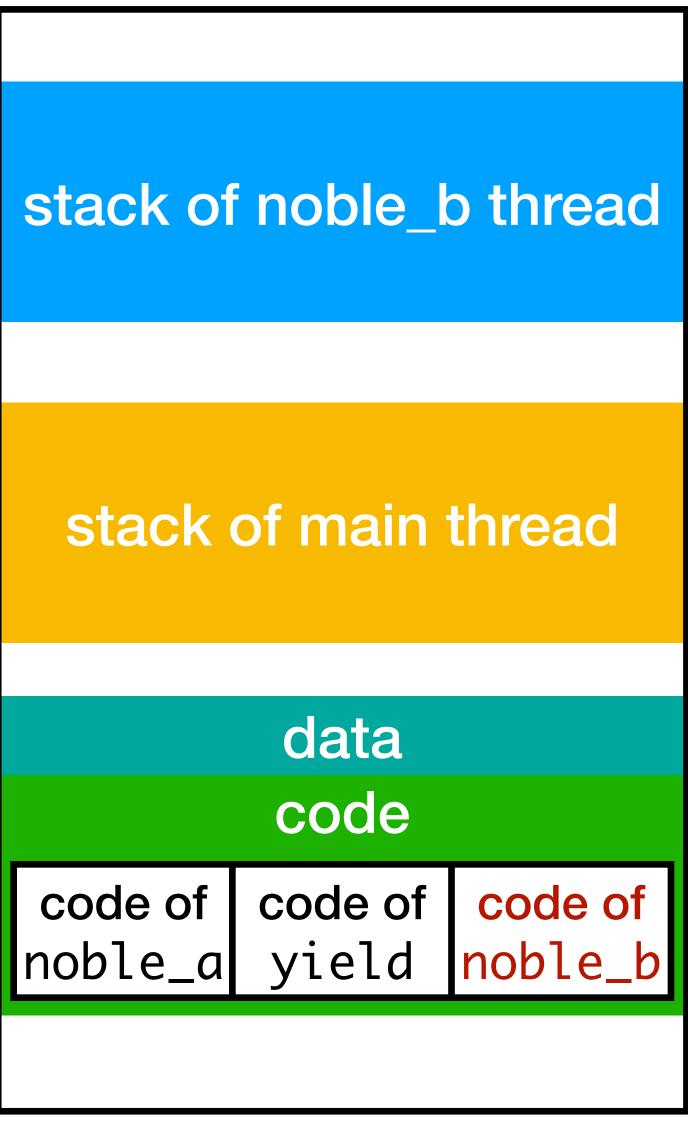
struct thread current\_thread; struct queue\_t runnable\_threads; .....

•••••

void thread\_create(void (\*f)(void \*), void \*arg,

unsigned int stack\_size){

// allocate a stack for noble\_b // modify global variables // call ctx\_start to run function f // ctx\_start is defined in /src/lib/\*.s

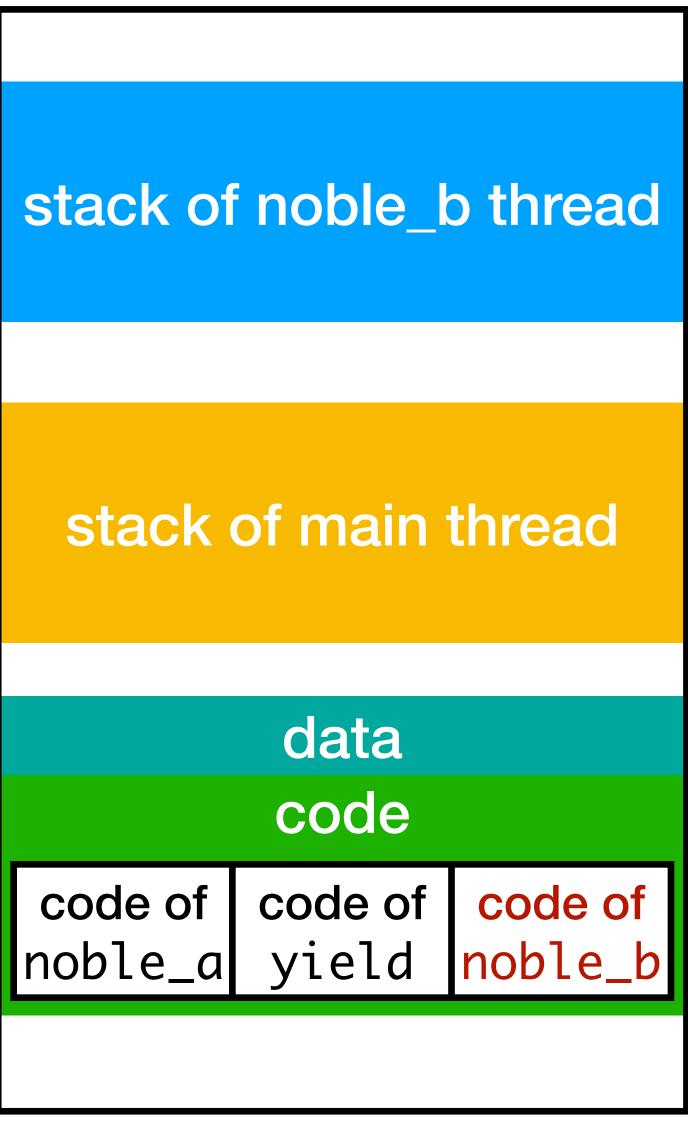


## thread\_yield

struct thread current\_thread; struct queue\_t runnable\_threads; .....

•••••

void thread\_yield(){ // choose next thread to run // call ctx\_switch to run the next thread // ctx\_switch requires previously stored // stack pointers // ctx\_switch is defined in /src/lib/\*.s



## Goal of Today's Class

#### Understand the concepts of context, context-switch and threads

Understand the related functions in assignment P1

- thread\_init, thread\_create, thread\_yield, thread\_exit
- ctx\_entry, ctx\_start, ctx\_switch
- It's your job to explore the details of how to implement and use these functions. ;-) Try to understand yourself before coming to office hours.

### Homework

- P1 is due on Oct 2. Start early.
- Implement the concepts of thread, context-switch and synchronization of threads (next lecture).