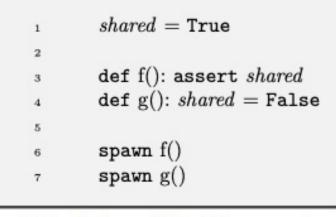
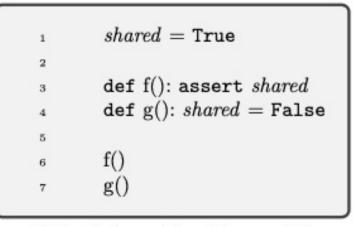


(a) [code/prog1.hny] Sequential

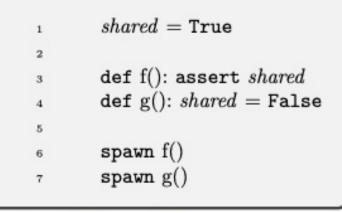


(b) [code/prog2.hny] Concurrent

Figure 3.1: A sequential and a concurrent program.



(a) [code/prog1.hny] Sequential



(b) [code/prog2.hny] Concurrent

Figure 3.1: A sequential and a concurrent program.

[Nemo:~/Documents/harmony] lorenzo% harmony code/prog1.hny	**Summary: something went wrong in an execution**
<pre>* Phase 1: compile Harmony program to bytecode * Phase 2: run the model checker (nworkers = 8)</pre>	Here is a summary of an execution that exhibits the issue:
<pre>* 2 states (time 0.00s, mem=0.000GB) * Phase 3: analysis</pre>	<pre>* Schedule thread T0:init() * Line 1: Initialize shared to True</pre>
* 2 components (0.00 seconds)	<pre>* **Thread terminated**</pre>
* Check for data races * **No issues found**	<pre>* Schedule thread T2: g() * Line 4: Set shared to False (was True)</pre>
<pre>* Phase 4: write results to code/prog1.hco * Phase 5: loading code/prog1.hco</pre>	<pre>* **Thread terminated** * Schedule thread T1: f()</pre>
	* Line 3: Harmony assertion failed

Two threads updating shared variable amount T_1 wants to decrement amount by \$10K \square T₂ wants to decrement amount by 50% T2 amount := amount - 10,000; amount := amount * 0.5; Memory amount 100,000

 T_1

What happens when T_1 and T_2 execute concurrently?

 T_2

Might execute like this:

r2 := load from amount r2 := 0.5 * r2 store r2 to amount

r1 := load from amount r1 := r1 - 10,000 store r1 to amount

 T_1

Memory

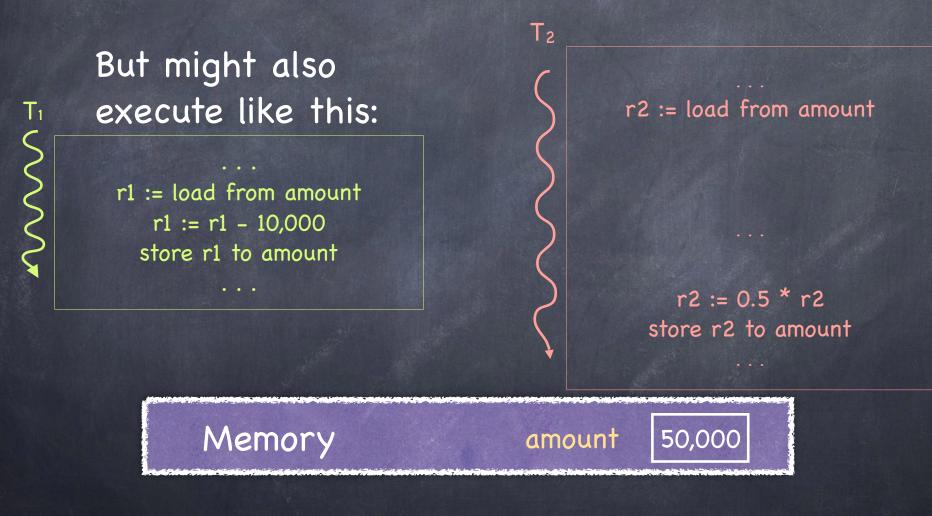




Or viceversa: T_1 and then T_2



Non-Atomicity



One update is lost!

Wrong – and very hard to debug

Race Conditions

Timing dependent behaviors involving shared state

Behavior of race condition depends on how threads are scheduled!

a concurrent program can generate MANY "schedules" or "interleavings"

schedule: a total order of machine instructions

bug if any of them generates an undesirable behavior

All possible interleavings should be safe!

Race Conditions: Hard to Debug

- Only some interleavings may produce a bug
- But bad interleavings may happen very rarely
 program may run 100s of times without generating an unsafe interleaving
- Small changes to the program may hide bugs
 "The Case of the Print Statement"
- Compiler and processor hardware can reorder instructions

Students develop their code in Python or C, and test it by running it many times....

Testing can only prove the presence of bugs... not their absence!

True!

But there is testing and then testing... They submit their code, confident that it is correct...

and I test the code with my secret and evil methods...* ...and find that most submissions are broken!

> *uses homebrew líbrary that randomly samples from possíble ínterleavings ("fuzzíng")

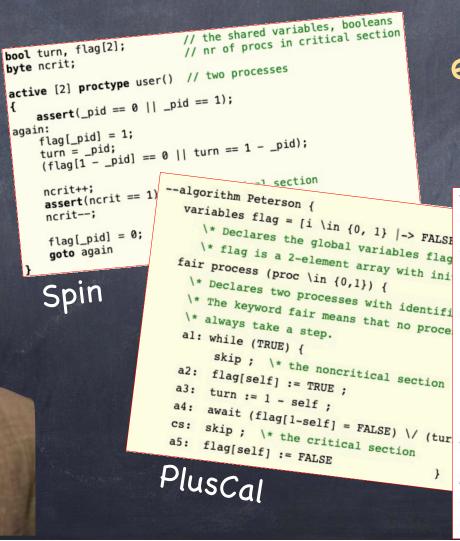
I am unhappy, and the students are unhappy!

Why is that?

- Studies show that heavily used code, implemented, reviewed and tested by expert programmers has lots of concurrency bugs
- Even professors who teach concurrency or write books or papers about concurrency get it wrong sometimes!

Hand-written proofs are just as likely to have bugs as programs... or even more likely, as you can't test hand-written proofs!

There are no mainstream tools to check concurrent algorithms... those that exist have a steep learning curve



Examples of existing tools

TLA+

VARIABLES flag, turn, pc vars $\stackrel{\Delta}{=} \langle flag, turn, pc \rangle$ Init $\stackrel{\Delta}{=} \wedge flag = [i \in \{0, 1\} \mapsto \text{FALSE}]$ $\wedge turn = 0$ $\wedge pc = [self \in \{0, 1\} \mapsto "a0"]$ $a3a(self) \stackrel{\Delta}{=}$ $\wedge pc[self] = "a3a"$ $\wedge \text{IF flag[Not(self)]}$ THEN pc' = [pc EXCEPT ![self] = "a3b"]ELSE pc' = [pc EXCEPT ![self] = "cs"] $\wedge \text{ UNCHANGED } \langle flag, turn \rangle$ $\backslash * \text{ remaining actions omitted}$ $proc(self) \stackrel{\Delta}{=} a0(self) \lor \ldots \lor a4(self)$ $Next \stackrel{\Delta}{=} \exists self \in \{0, 1\} : proc(self)$ $Spec \stackrel{\Delta}{=} Init \land \Box[Next]_{vars}$

Enter Harmony



A new concurrent programming language

heavily based on Python syntax to reduce learning curve for many

 A new underlying virtual machine, quite different from any other

it tries all possible executions of a program, until it finds a problem (if any)

(this is called "model checking")

def T1():

amount -= 10000 done1 = **True**

def T2():

amount /= 2 done2 = **True**

def T1():

amount -= 10000 done1 = **True**

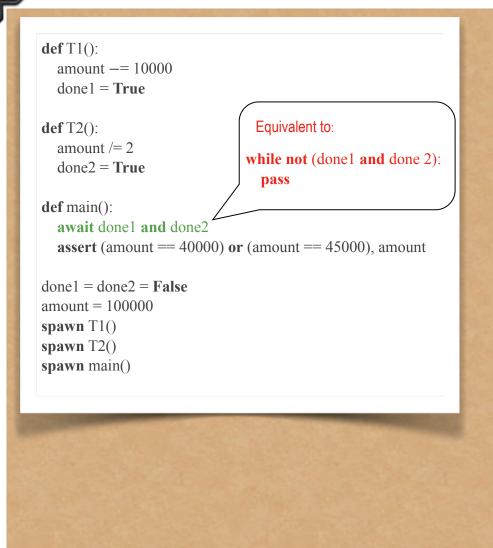
def T2(): amount /= 2

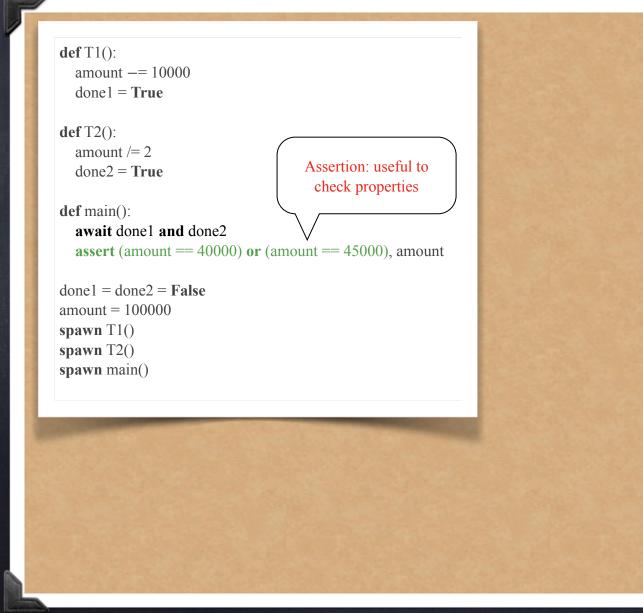
done2 = True

def main():

await done1 and done2
assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main()





def T1(): amount -= 10000done1 = True def T2(): amount $\neq 2$ Output amount if done2 = **True** assertion fails def main(): await done1 and done2 assert (amount == 40000) or (amount == 45000), amount done1 = done2 = Falseamount = 100000spawn T1() spawn T2() spawn main()

An important note on assertions

An assertion is not part of your algorithm
 Semantically an assertion is a no-op

 it is never expected to fail because it is supposed to state a fact

That said...

Assertions are super-useful

© @label: assert P is a type of invariant:

▶ $pc = label \Rightarrow P$

Output Use them liberally

in C, Java, ..., they are automatically removed in production code — or automatically optimized out if you have a really good compiler

They are great for testing

They are executable documentation

comments tend to get outdated over time

That said...

- Comment them out before submitting a programming assignment
 - vou don't want your assertions to fail while we are testing your code...

def T1():

amount -= 10000 done1 = **True**

def T2():

amount /= 2 done2 = **True**

def main():

await done1 and done2 assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main()



def T1():

amount -= 10000 done1 = **True**

def T2():

amount /= 2 done2 = **True**

def main():

await done1 and done2 assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main() Spawn main()

def T1():

amount -= 10000 done1 = **True**

def T2():
 amount /= 2
 done2 = True

def main():

await done1 and done2 assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main()



[Nemo:~/Documents/harmony/mycode] lorenzo% harmony example.hny * Phase 1: compile Harmony program to bytecode * Phase 2: run the model checker (nworkers = 8) * 103 states (time 0.00s, mem=0.000GB) * Phase 3: analysis * **Safety Violation** * Phase 4: write results to example.hco * Phase 5: loading example.hco **Summary: something went wrong in an execution** Here is a summary of an execution that exhibits the issue: * Schedule thread T0: __init__() * Line 13: Initialize done2 to False * Line 13: Initialize done1 to False * Line 14: Initialize amount to 100000 * **Thread terminated** * Schedule thread T2: T2() * Preempted in T2() about to store 50000 into amount in line 6 * Schedule thread T1: T1() * Line 2: Set amount to 90000 (was 100000) * Line 3: Set done1 to True (was False) * **Thread terminated** * Schedule thread T2: T2() * Line 6: Set amount to 50000 (was 90000) * Line 7: Set done2 to True (was False) * **Thread terminated** * Schedule thread T3: main() * Line 11: Harmony assertion failed: 50000

def T1():

amount -= 10000 done1 = **True**

def T2(): amount /= 2 done2 = True

def main():

await done1 and done2 assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main()



Summary: something went wrong in an execution

Here is a summary of an execution that exhibits the issue:

* Schedule thread T0: __init__()

- * Line 13: Initialize done2 to False
- * Line 13: Initialize done1 to False
- * Line 14: Initialize amount to 100000
- * **Thread terminated**

* Schedule thread T1: T1()

```
* Preempted in T1()
```

about to store 90000 into amount in line 2 * Schedule thread T2: T2()

* Line 6: Set amount to 50000 (was 100000)

* Line 7: Set done2 to True (was False)

* **Thread terminated**

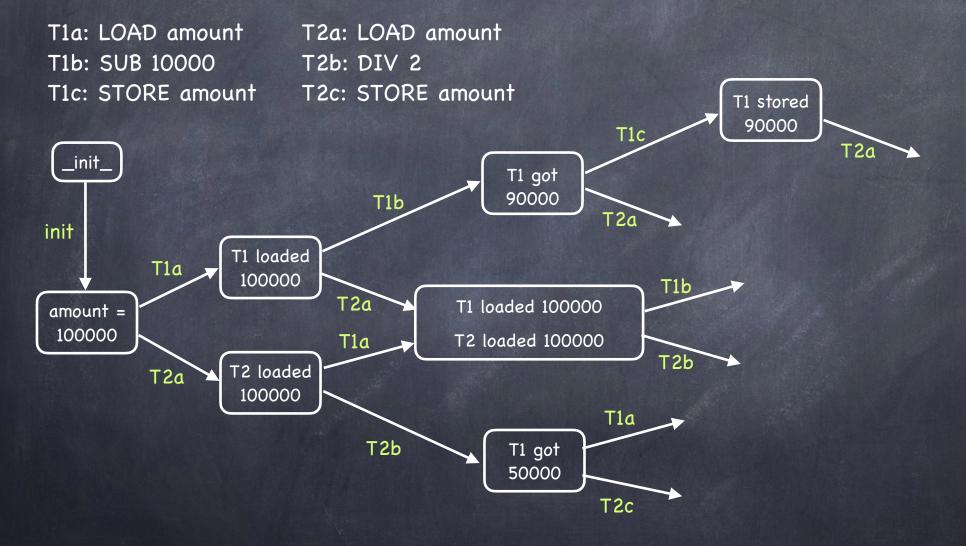
* Schedule thread T1: T1()

- * Line 2: Set amount to 90000 (was 50000)
- * Line 3: Set done1 to True (was False)

* **Thread terminated**

- Schedule thread T3: main()
- * Line 11: Harmony assertion failed: 90000

Simplified model (ignoring main)



def T1():

amount -= 10000 done1 = **True**

def T2():
 amount /= 2
 done2 = True

def main():

await done1 and done2 assert (amount == 40000) or (amount == 45000), amount

done1 = done2 = False amount = 100000 spawn T1() spawn T2() spawn main()



Summary: something went wrong in an execution

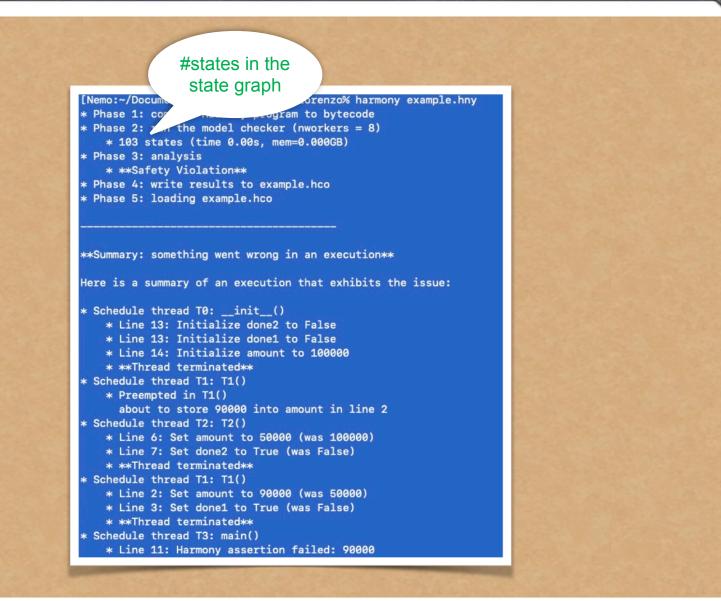
Here is a summary of an execution that exhibits the issue:

* Schedule thread T0: __init__()

- * Line 13: Initialize done2 to False
- * Line 13: Initialize done1 to False
- * Line 14: Initialize amount to 100000
- * **Thread terminated**
- * Schedule thread T1: T1()
 - * Preempted in T1()
- about to store 90000 into amount in line 2 * Schedule thread T2: T2()
 - * Line 6: Set amount to 50000 (was 100000)
 - * Line 7: Set done2 to True (was False)
 - * **Thread terminated**

* Schedule thread T1: T1()

- * Line 2: Set amount to 90000 (was 50000)
- * Line 3: Set done1 to True (was False)
- * **Thread terminated**
- * Schedule thread T3: main()
 - * Line 11: Harmony assertion failed: 90000







Harmony's VM State

Three parts:
code (which never changes)
values of shared variables
states of each of the running threads
a.k.a. "contexts"

State represents one vertex in the graph model

Context (State of a Process)

- Method name and parameters
- PC (program counter)
- local variables
 - parameters (a.k.a. arguments)
 - result
 - there is no return statement
 - local variables
 - declared in var, let, and for statements

Harmony != Python

Harmony	Python
tries all possible executions	executes just one
() == [] ==	1 != [1] != (1)
1, == [1,] == (1,) != (1) == [1] == 1	[1,] == [1] != (1) == 1 != (1,)
f(1) == f 1 == f[1]	f 1 and f[1] are illegal (if f is method)
{ } is empty set	{ } is empty dictionary
few operator precedence rules use parentheses often	many operator precedence rules
variables global unless declared otherwise	depends Sometimes must be explicitly declared global
no return , break , continue	various flow control escapes
no classes	object-oriented

. . .

 $\mathbf{x}_{i} \in \mathbf{x}_{i}$

I/O in Harmony

Input

□ choose expression

▶ $x = choose(\{1,2,3\})$

allows Harmony to know all possible inputs

□ const expression

• const x = 3

▶ can be overridden with "-c x = 4" to Harmony

Ø Output

 \square print x + y

 \square assert x + y < 10, (x, y)

I/O in Harmony

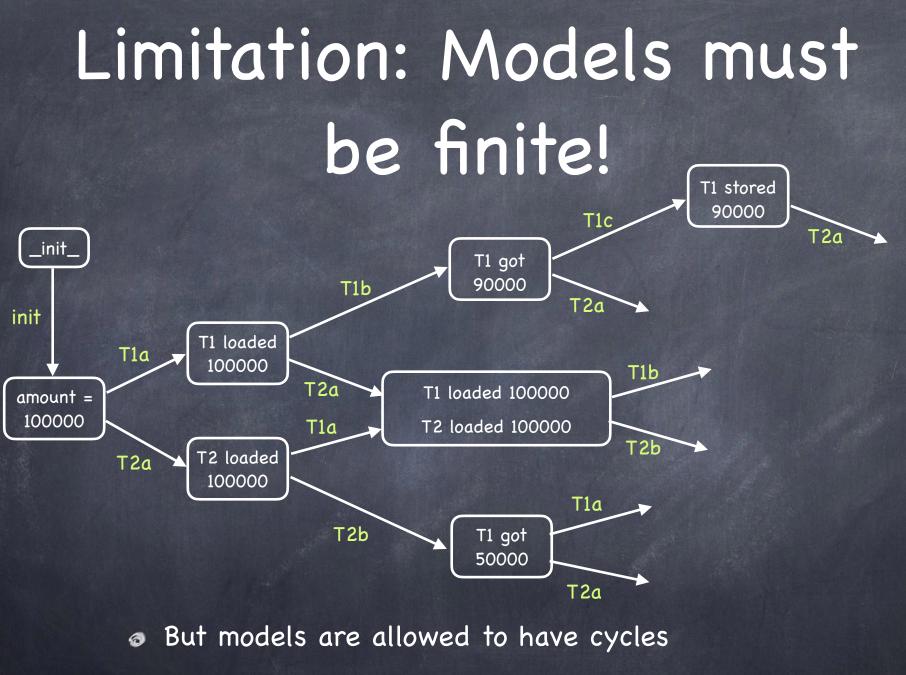
Input

No open(), read(), or input() statements

Output
print x + y
assert x + y < 10, (x, y)

Non-determinism in Harmony

Three sources
 choose expressions
 thread interleavings
 interrupts



- Securiors are allowed to be unbounded
- Harmony checks for the possibility of termination

Back to our problem...

Two threads updating shared variable amount
T₁ wants to decrement amount by \$10K
T₂ wants to decrement amount by 50%

amount := amount - 10,000;

Memory

 T_1

amount := amount * 0.5;

100,000

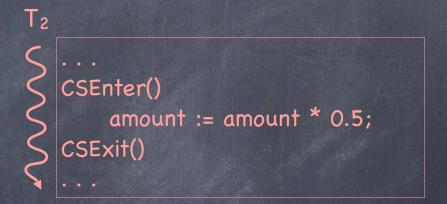
amount

How to "serialize" these executions?

Critical Section

Shared memory access: must be serialized

T₁ Second states and the second states of the se



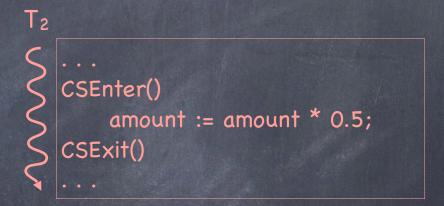
Goals

Mutual exclusion: at most 1 thread in CS at any time
 Progress: all threads wanting to enter CS eventually do
 Fairness: equal chances to get into CS (uncommon in practice)

Critical Section

Shared memory access: must be serialized

T₁ Second states and the second states of the se



Goals

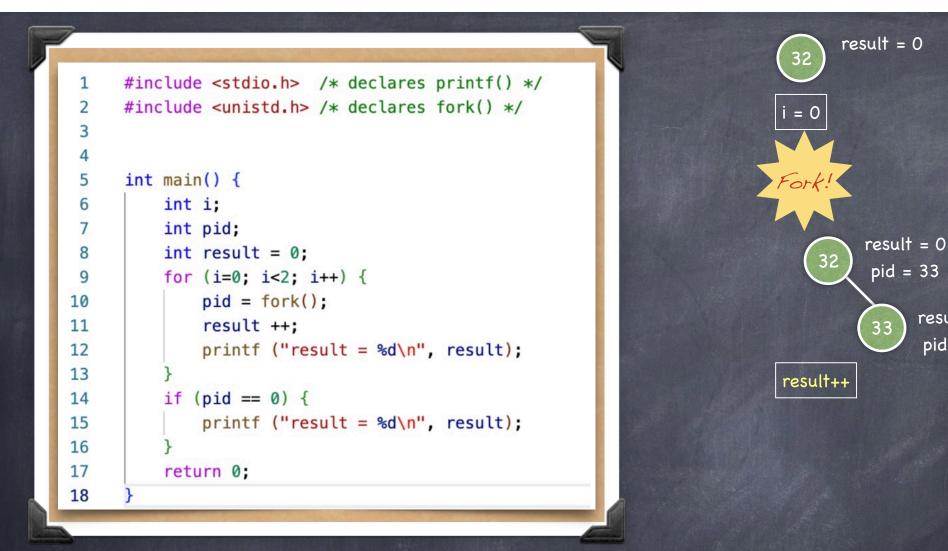
D Mutual exclusion: at most 1 thread in CS at any time

Progress: if any threads want to enter the CS, at least one does

What makes the Critical Section problem hard?

- Mutual exclusion?
- Progress?
- It is the combination!
 - both properties, on their own, are trivial to achieve
 - \square there is much more to this...

Prelim Interlude



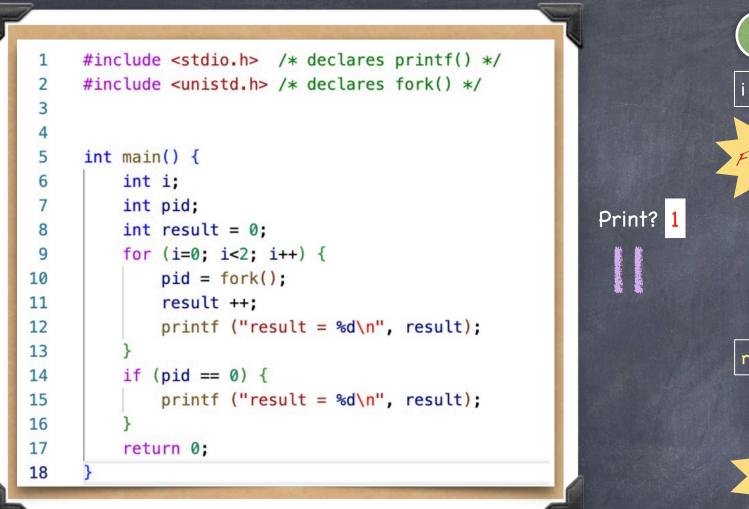
pid = 33

33

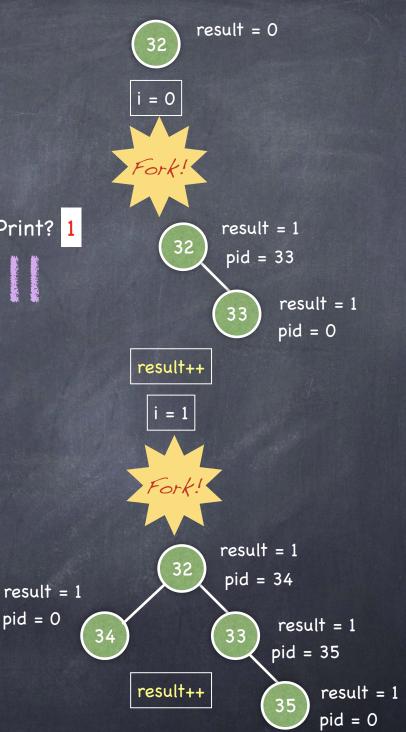
result = 0

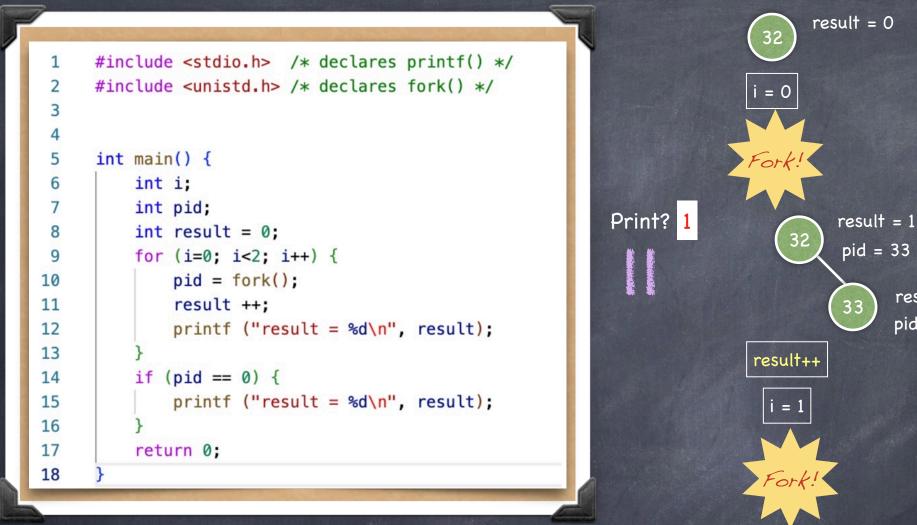
pid = 0

How many times will the value of result be printed? First value(s)? Last value(s)?

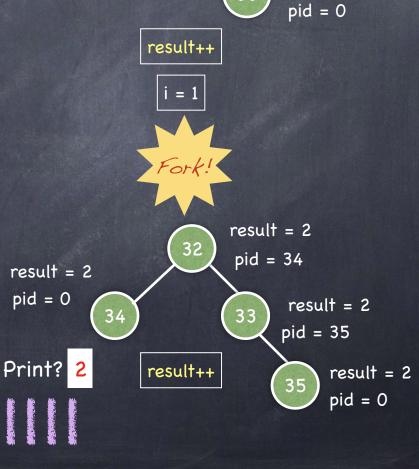


How many times will the value of result be printed? First value(s)? Last value(s)?

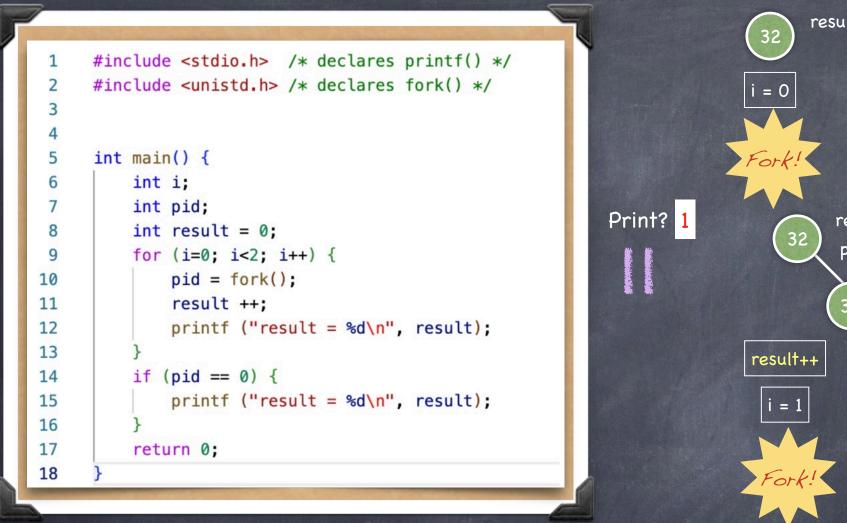




How many times will the value of result be printed? First value(s)? Last value(s)?



result = 1



How many times will the value of result be printed? First value(s)? Last value(s)?

