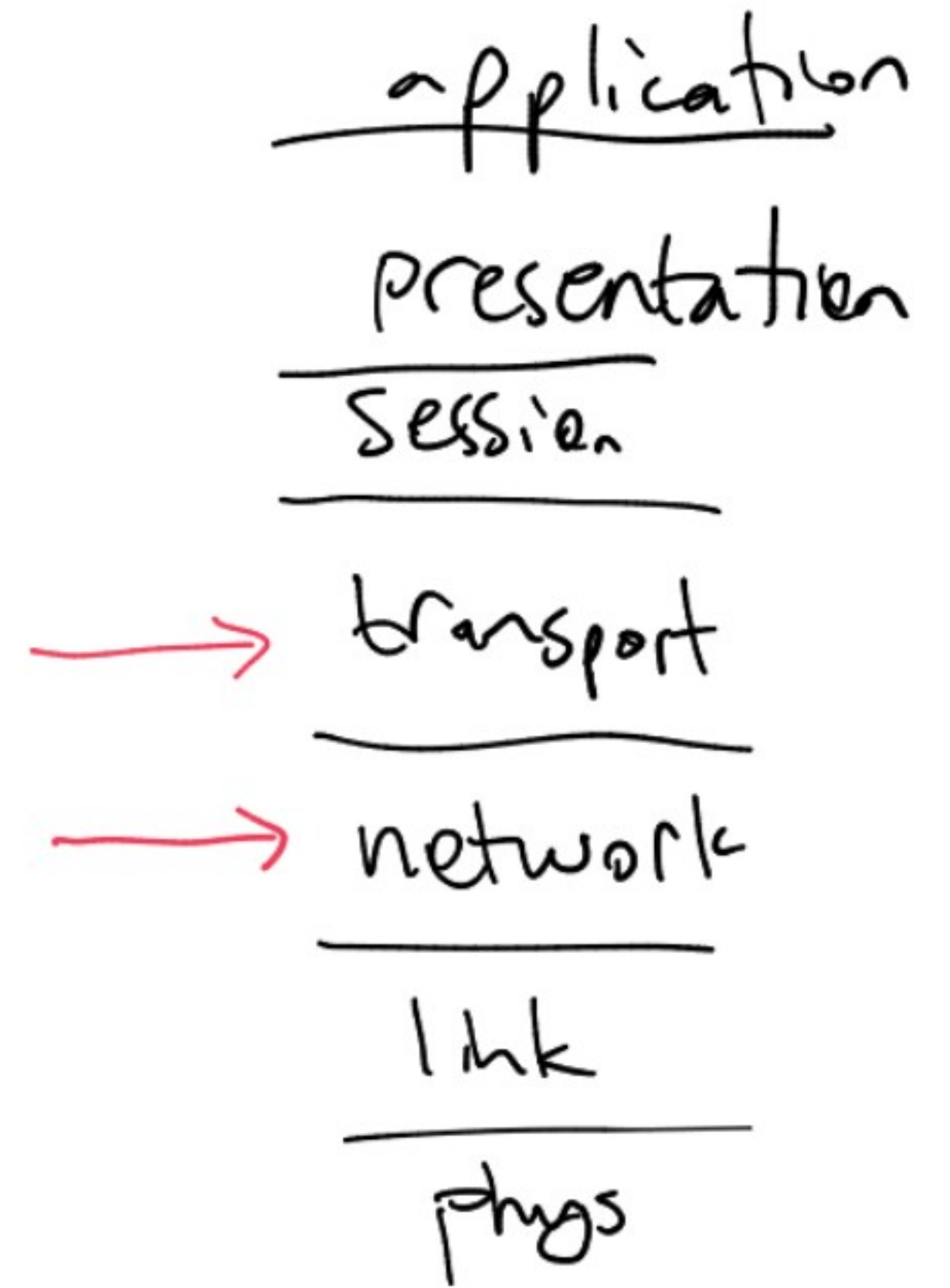


# Lec 22: routing & transport

- Routing
  - source vs. path routing, BGP
- transport - layer
  - UDP & TCP

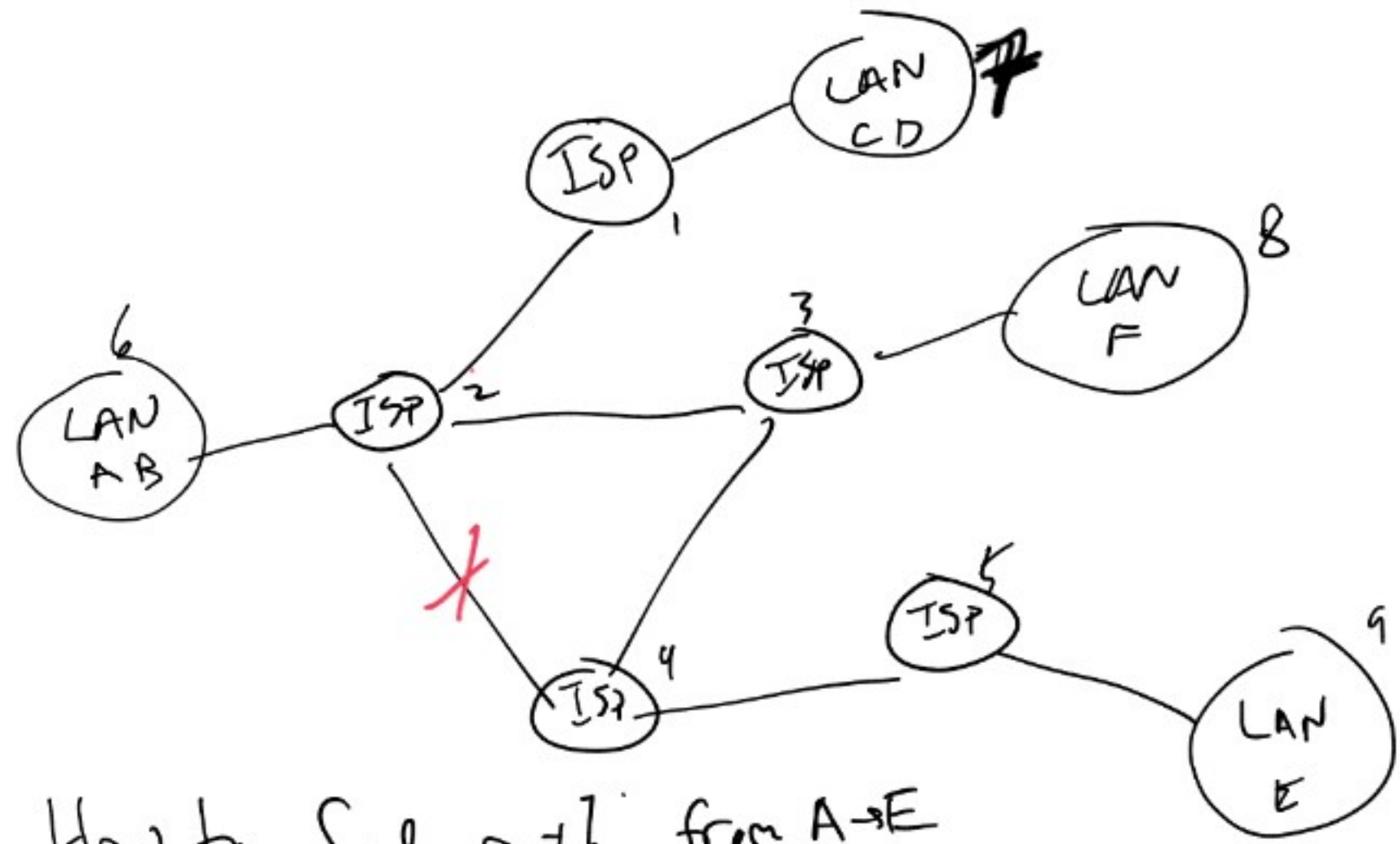


Reality: routing table entries correspond to subnetworks of different sizes.

Subnetwork is identified by IP address + mask

e.g. 1.2.x.x  
represented as  $\underbrace{1.2.0.0}_{\text{IP}} / \underbrace{255.255.0.0}_{\substack{\text{bits in first} \\ \text{2 digits set}}} \\ \text{clear}$

to see if address is in a network:  
(logical) and address w/ netmask,  
compare to network IP.



How to find path from A → E

- A computes shortest path through graph, writes path into packet, each ISP on path forwards to next.

- Eg: 6 → 2 → 4 → 5 → 9

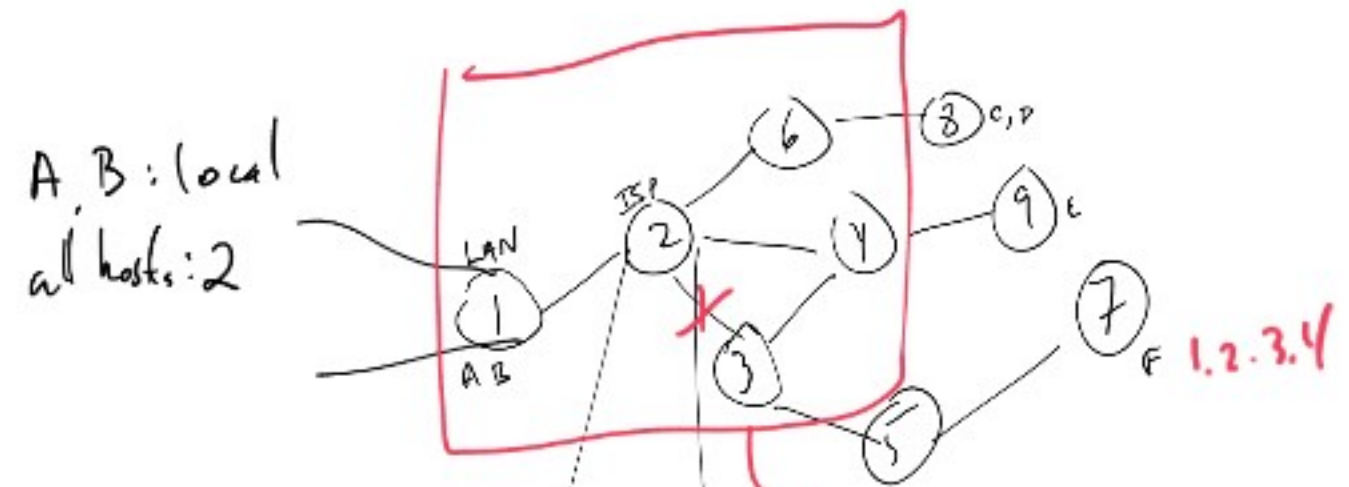
Alice is in 6, forward to 2, 2 looks @ path, forward to 4, ...

Source routing:  
source msg (A) finds route.

Pro: • shortest path

Con:

- data can get dropped if topology changes
- each packet has a path in it: overhead.
- endpoints need to know entire graph



A, B: local  
all hosts: 2

Send from A → F  
 A → (gateway) 1 → 2  
 2 → 3  
 3 → 5  
 5 → 7  
 7 (gateway) → F

routing tbl for 2

if destined for	forward it to
A, B	1
C, D	6
E	4
F	3

routing tbl for 3:

A, B, C, D: 2  
 E: 4  
 F: 5

routing tbl for 5

F: 7  
 anyone else: 3

prob: → ISPs can dynamically adjust to respond to traffic.  
 - topology can be private.

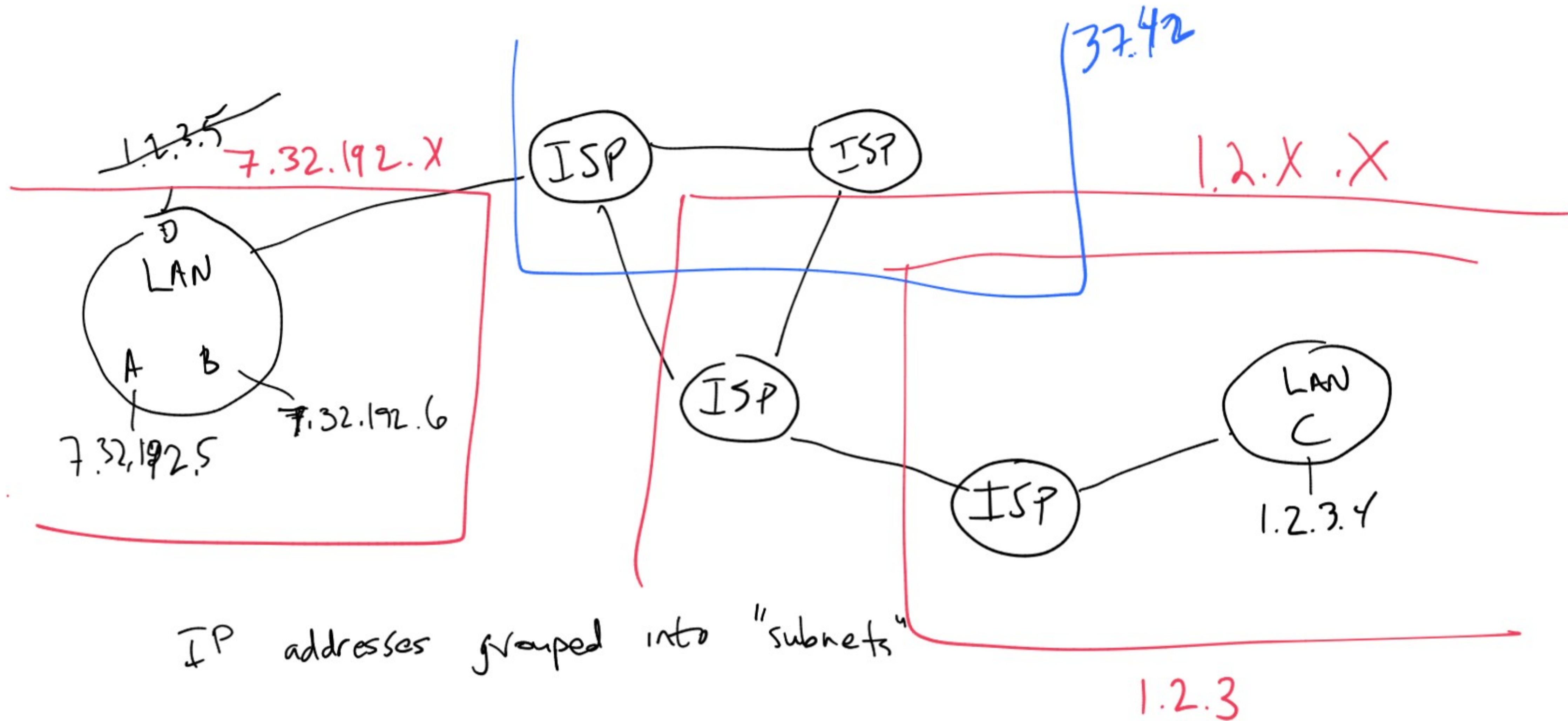
Con: = each ISP needs routing tbl:  
 lot of storage?

Path routing:  
 each node in path has table telling it how to forward packets. (based on dest addr).

- need to update routing tables when network changes.

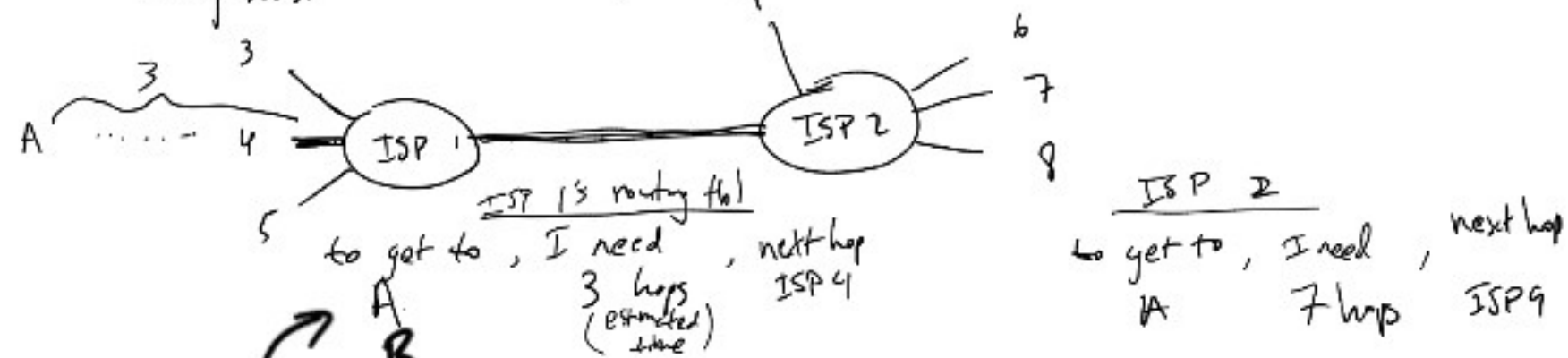
other concerns:  
 what about traffic on each link?

network IDs,  
 not host IDs.



## How to build/maintain routing tables?

each ISP will periodically share routing table with neighbors.



distance vector

ISP 1 sends to ISP 2, who notices it can forward packets to A using only 4 hops (better than before) by forwarding packets to ISP 1.

- Distributed solution to all-pairs shortest path algorithm.

- if topology changes, ISPs closest to change will notice first; eventually data will propagate to rest of network.

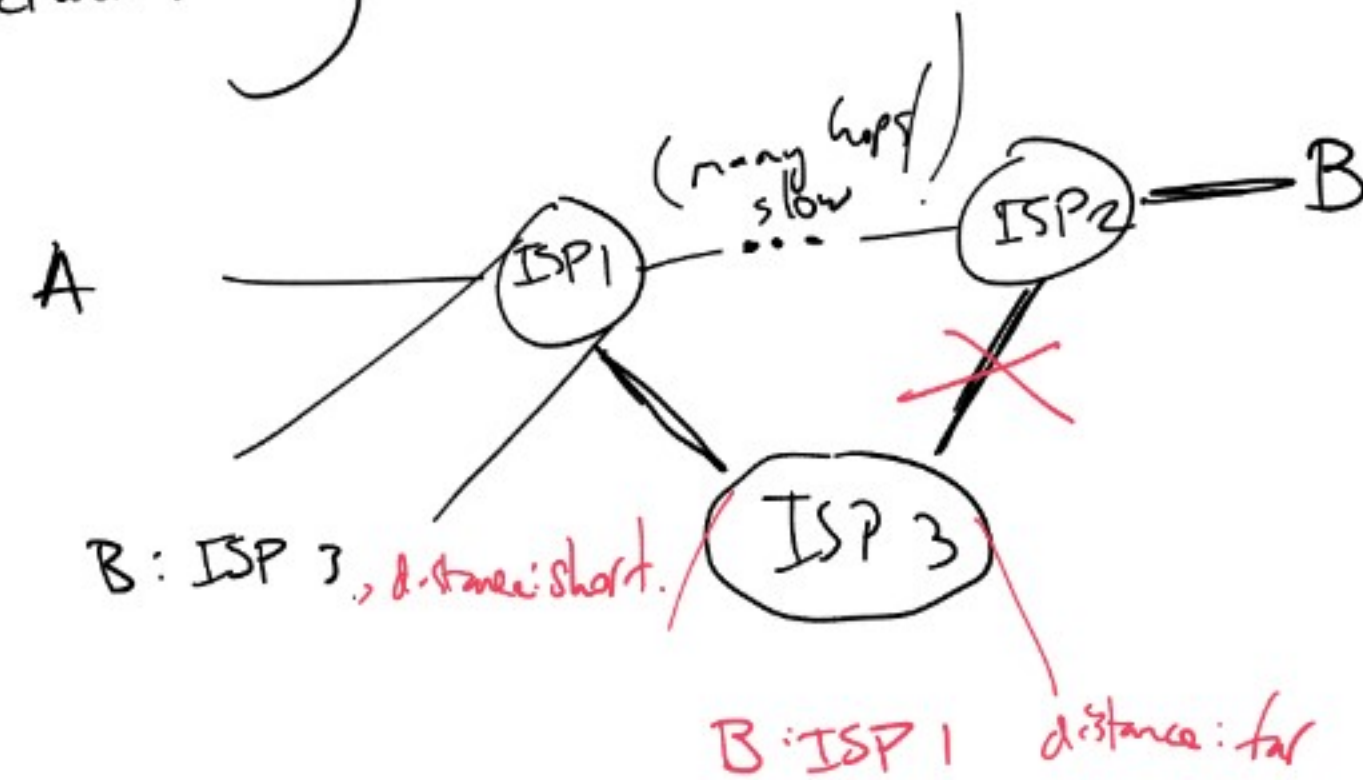
Suboptimal in mean time.

(Could even loop!)

## Border gateway protocol (BGP)

= maintain "distance vector", distance to end end network

- exchange with neighbors periodically.



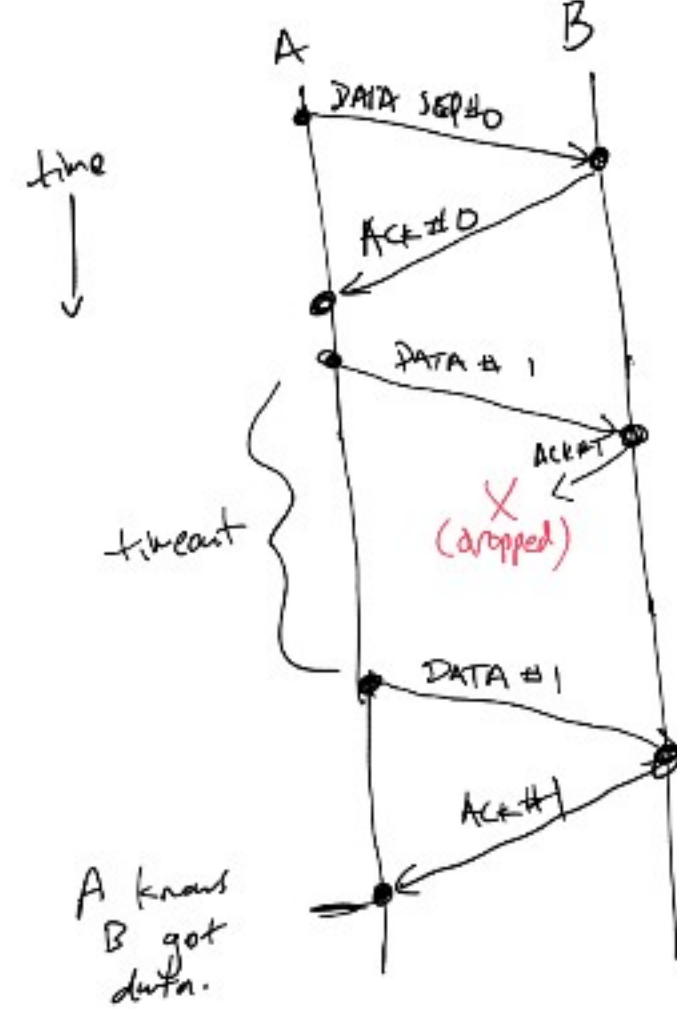
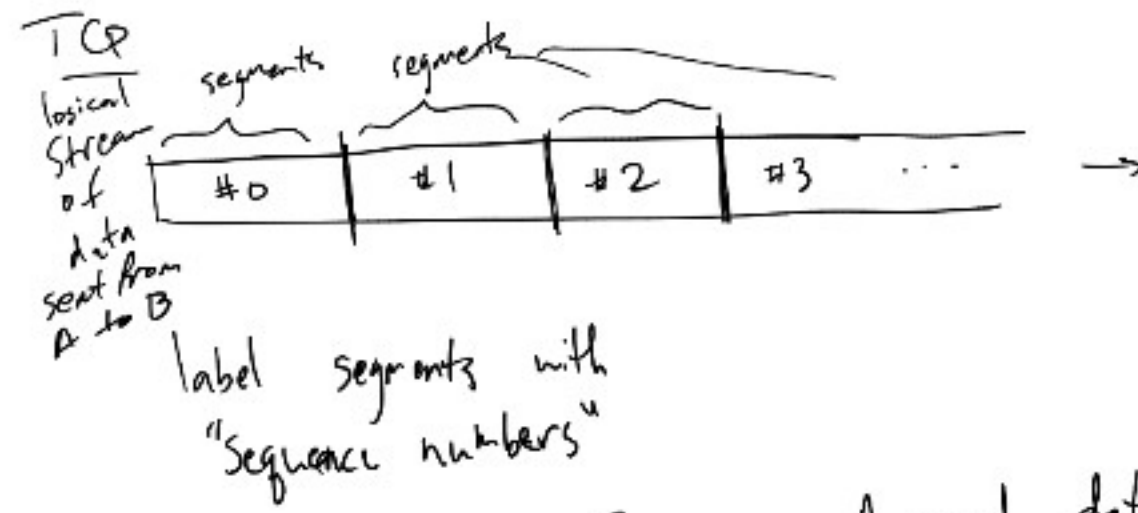
- Transmission Control protocol (TCP)

- multiplexing

- "stream" abstraction.

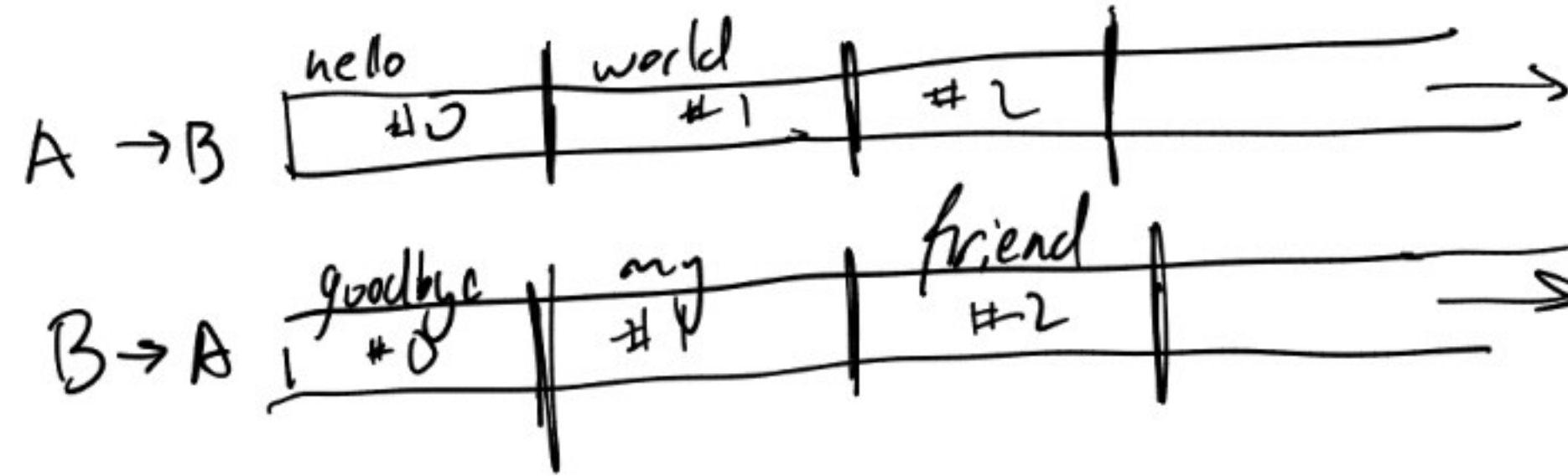
- two-way stream of data,  
each endpoint can write into  
stream, read from stream.

- In-order, guaranteed delivery.



- A send data, with sequence #
- B reply with acknowledgement (ACK) also containing seq. #.
- when A receives ACK, knows B received data.
  - if it doesn't, re-send data.

actually 2 logical streams in a TCP connection



Can combine DATA & ACK messages to send 1 packet instead of 2.

DATA / ACK

"piggybacking"

