

Project-6: Unix-like File System Layer

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(original slides by Efe Gencer)

Project Scope

1. Implement a Unix-like file system layer: *ufsdisk*.
2. Use *free space bitmaps* to keep track of free and used blocks.
3. (Optional) Implement *a file system checker* (i.e. *fsck*) to check the integrity of your file system.

Recap: Intro

- File systems are built on one or more *block stores*.
- The block store abstraction provides:
 - a disk-like interface: read / write blocks
 - a sequence of blocks -- each typically a few kilobytes
- The block store abstraction doesn't deal with:
 - file naming
 - user permissions
 - distinguishing files from directories
 - ...

Recap: Block Store Abstraction

- `block_t block`
 - block of size `BLOCK_SIZE`
- `nblocks() -> integer`
 - returns size of the block store in `#blocks`
- `read(block number) -> block`
 - returns the contents of the given block number
- `write(block number, block)`
 - writes the block contents at the given block number
- `setsize(nblocks)`
 - sets the size of the block store

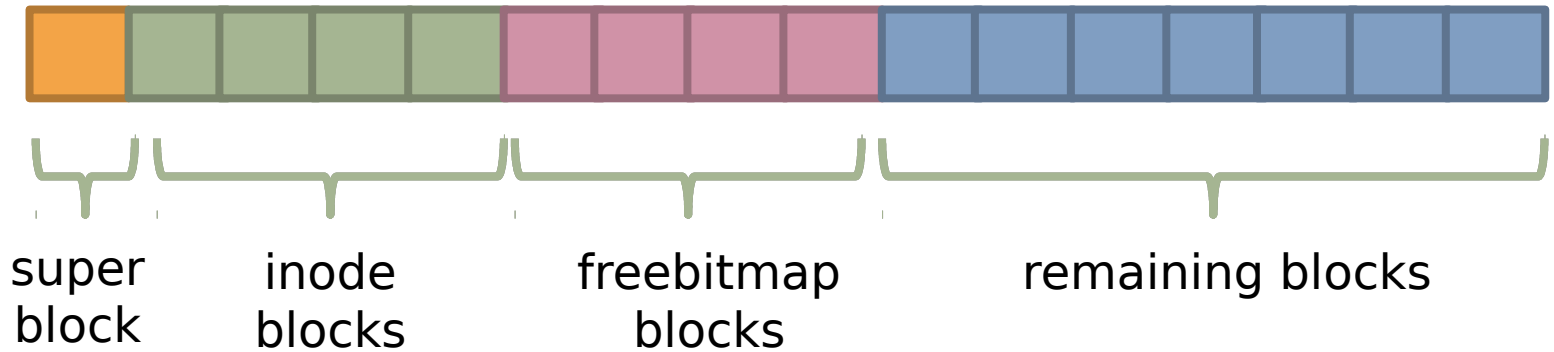
Recap: block_if.h

```
#define BLOCK_SIZE    512        // # bytes in a block
typedef unsigned int  block_no;   // index of a block

struct block { char bytes[BLOCK_SIZE]; };
typedef struct block block_t;

typedef struct block_if *block_if;
struct block_if {
    void *state;
    int (*nblocks)(block_if bif);
    int (*read)(block_if bif, block_no offset, block_t *block);
    int (*write)(block_if bif, block_no offset, block_t *block);
    int (*setsize)(block_if bif, block_no size);
    void (*destroy)(block_if bif);
};
```

Ufsdisk: layout



Ufsdisk: superblock (1 per underlying blockstore)

```
struct ufs_superblock {
    // magic number of ufsdisk
    // identifies the filesystem
    unsigned int magic_number;

    block_no n_inodeblocks;
    block_no n_freebitmapblocks;
};
```

Ufsdisk: inodeblock

```
#define INODESperBLOCK  
    (BLOCK_SIZE / sizeof(struct ufs_inode))  
  
struct ufs_inodeblock {  
    struct ufs_inode inodes[INODESperBLOCK];  
};
```

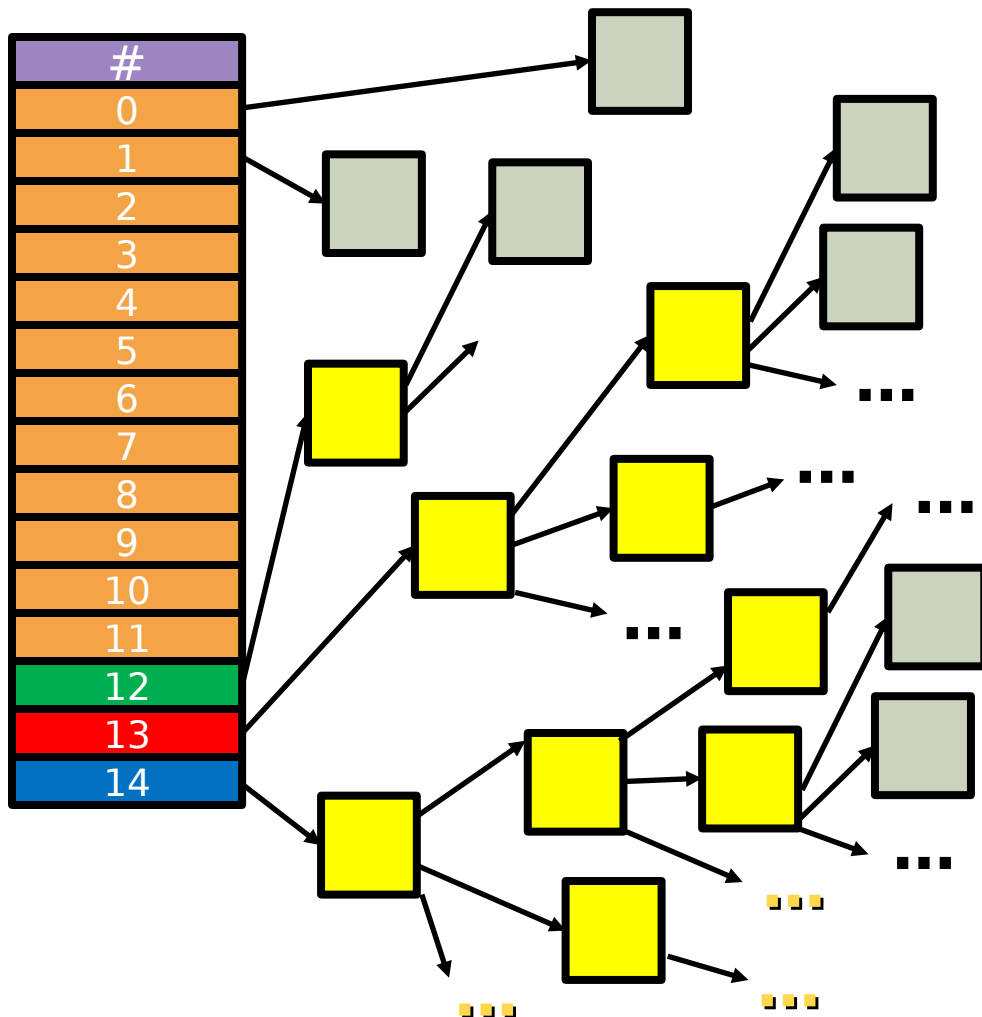

Ufsdisk: inode (1 per virtual blockstore)

```
#define REFS_PER_INODE    15

struct ufs_inode {
    // total size of the file
    block_no nblocks;
    block_no refs[REFS_PER_INODE];
};
```

Ufsdisk: inode

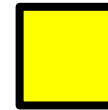
(1 per virtual blockstore)



Legend



:Data block



:Indirect block



:Number of blocks



:Direct pointer



:Indirect pointer



:Double Indirect pointer



:Triple Indirect pointer

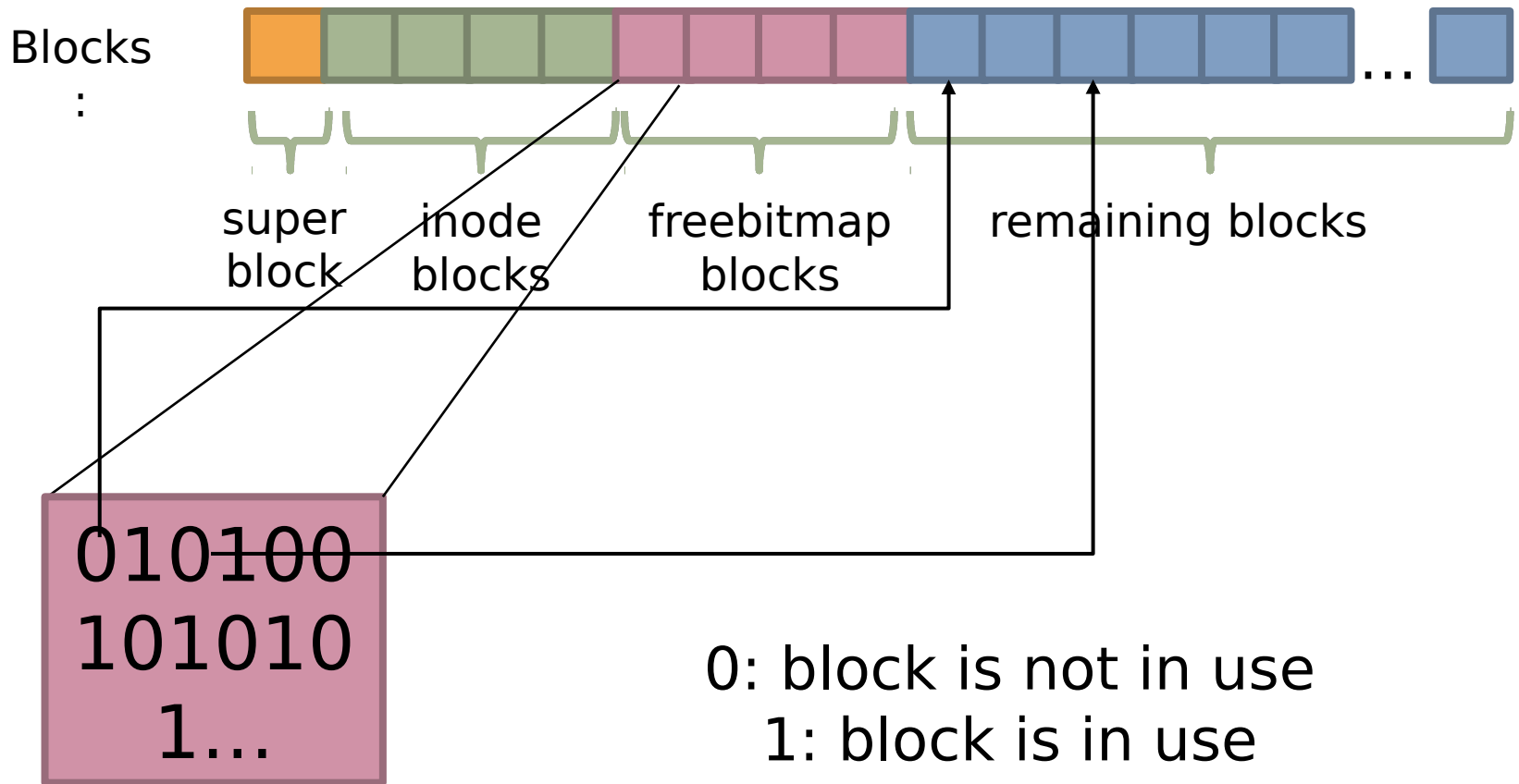


:Other data blocks



:Other indirect blocks

Free space bitmaps



Ufsdisk: freebitmap blocks

- Each freebitmap block is a list of bits
- How many freebitmap blocks, f , do I need?

One superblock

How many inode blocks are there?

$$f = \left\lceil \frac{nblocks - 1 - \left\lfloor \frac{n_inodes}{INODES_PER_BLOCK} \right\rfloor}{1 + BLOCK_SIZE * 2^3} \right\rceil$$

Free bitmap blocks take up space too

Each freebitmap block holds $BLOCK_SIZE * 8$ bits.

File system checker

- Verifies the consistency of your filesystem – e.g. try fsck (UNIX), chkdsk (Windows).
- If system crashes, filesystem may be corrupted.
- Checks filesystems -- and repairs fixable issues if broken.
 - a datablock is in use but marked as free.
 - a particular block is both an indirblock and datablock.
 - a particular datablock has been used more than once.
 - other issues...

Questions?

- This is the **last** project!
- Begin early so you have time to study for the finals.