Project 3 Supplemental Lecture

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Today's Lecture

- Administrative Information
- Common mistakes on Project 2
- Project 3 FAQ
- Discussion

Administrative Information

- Project 2 is being graded
- See a TA if you need help fixing up your P2 before working on P3
- Project 3 deadline is October 19th, 11:59 PM.

Project 2 Common Errors

- Same semaphore used both with and without interrupts disabled (e.g. clean semaphore)
- Interrupt handler decreases alarm count for all alarms
- Only one alarm fired per clock interrupt, even if multiple are ready
- Sleep not implemented using semaphore
- A thread should run again immediately if it is the only thread of its level and it yields (and the level to schedule from hasn't changed)
- Memory leaks: sleep semaphores, alarm structs

Network header generation

- Having a common header format will be fun later on.
- Use the pack and unpack functions that we provide.
 - Do not use a simple byte copy of the network address, this
 is incorrect.
 - Port numbers may be stored as ints in your program but must be converted to unsigned shorts when packing.
 - Reminder: do not pack port numbers as ints!
- The protocol field will be useful in the next project.
 - Set the protocol char field to PROTOCOL_MINIDATAGRAM for this project.

Network header generation – the incorrect way

- Given the header specs:
 - 1 byte protocol type
 - 8 bytes source address
 - 2 bytes source port
 - 8 bytes destination address
 - 2 bytes source port
- Allocate a (1+8+2+8+2) byte buffer and manually fill in the contents.

Why is this a bad idea?

- Tedious to code.
 - Lots of memcpy() and sizeof() operations.
 - Code looks ugly.
- What if the header specs change later?
 - Must manually change all offsets.

Iteration #2

- Idea: use a struct to store all fields so they are arranged correctly in memory.
 - Compiler arranges a contiguous block of memory for the struct.
 - Memory layout of the struct follows the order declared by the struct.

```
struct header
{
    char protocol_type;
    network_address_t source_address;
    unsigned short source_port;
    network_address_t destination_address;
    unsigned short destination_port;
};
struct header hdr;
network_send_pkt((char*) &hdr, sizeof(hdr), ...);
```

Iteration #2: Close but not quite...

- Padding!
- Computers usually load in units of words.
 - If a multibyte variable spans 2 words, then 2 loads are needed.
 - Align the variable to some word boundary so it requires exactly 1 load.
- Padding is unpredictable and is a waste of resources to transmit.

Iteration #3

- Idea: Use a struct that cannot possibly have padding.
 - Chars require exactly 1 load no matter where they are located.
 - Therefore consecutive char fields in a struct are not padded.
 - Works regardless of compiler options for padding.

```
struct header
{
    char protocol_type;
    char source_address[8];
    char source_port[2];
    char destination_address[8];
    char destination_port[2];
};
struct header hdr;
network_send_pkt((char*) &hdr, sizeof(hdr), ...);
```

Use packing functions to convert and populate the char arrays.

Implementation Hints

- Use an array for your ports.
 - O(1) time when using unbound ports (since user specifies the port he wants).
 - O(1) time when creating bound ports before a wraparound; O(n) time afterwards is acceptable (since you need to check each port).
- Use semaphore_P and semaphore_V for blocking and unblocking threads.
 - Remember how we did this in project 2; consider places where you need to disable interrupts.

Implementation Hints

- Reuse your queue implementation.
 - This is useful for storing data in FIFO order.
- Perform sanity checks.
 - Is the protocol type correct?
 - Are you sure the received packet is meant for you?
 - Is the packet malformed (header too short, invalid port numbers, etc)?

Implementation Hints

- Consider semantics for unused ports.
 - Data sent to unused ports should not actually be transmitted.
 - Data received on an unused port should not be queued.
- Consider reuse semantics for unbound ports.
 - When an unbound port is destroyed and later recreated, any prior queued data should no longer be there.
 - Don't forget to reset the counting semaphore too.

Project 3 FAQ

- You may assume the specified port ranges will not change
 - No magic numbers, use #define
- Dynamic memory responsibilities.
 - Network interrupt handler passes you an network_interrupt_arg_t, which you have to eventually free.
 - The user-supplied buffer for both minimsg_send and minimsg_receive should not be freed by you.

Project 3 FAQ

- Mutexes and semaphores for unbound ports.
 - You will need a counting semaphore.
 - But you won't need a mutex. (Why?)

```
struct miniport {
    char port_type;
    int port_number;

union {
        struct {
            queue_t incoming_data;
            semaphore_t datagrams_ready;
        } unbound;

        struct {
            network_address_t remote_address;
            int remote_unbound_port;
        } bound;
};
```

Testing

- Occasional lost packets across machines.
 - This is normal.
 - Try re-executing your program again.
- Unable to communicate between two machines.
 - Make sure both machines can ping each other.
 - Try running on two machines in the CSUG lab.
 - Redrover is known to have problems with machine visibility.
- TAs will set up their solutions to test against in office hours this week.

Questions

Questions?