File System Implementation

File-System Structure

- · File structure
 - Logical storage unit
 - Collection of related information
- File system resides on secondary storage (disks)
- · File system organized into layers
- File control block storage structure consisting of information about a file

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A Typical File Control Block file permissions file dates (create, access, write) file owner, group, ACL file size file data blocks or pointers to file data blocks

In-Memory File System Structures

open (file name) directory structure file-control block secondary storage
(a) secondary storage
(b) file-control block
secondary storage
(b)

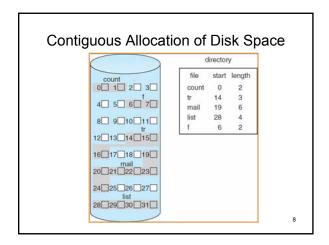
File System Layout File System is stored on disks Disk is divided into 1 or more partitions Sector 0 of disk called Master Boot Record End of MBR has partition table (start & end address of partitions) First block of each partition has boot block Loaded by MBR and executed on boot Fartition table Disk partition Boot block Boot block Super block Free space mgmt I -nodes Root dir Files and directories

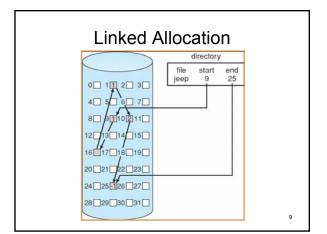
Allocation Methods

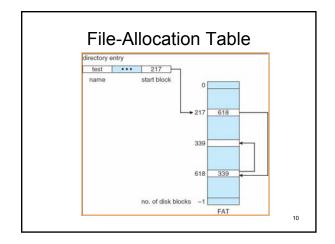
- An allocation method refers to how disk blocks are allocated for files:
- · Contiguous allocation
- · Linked allocation
- Indexed allocation

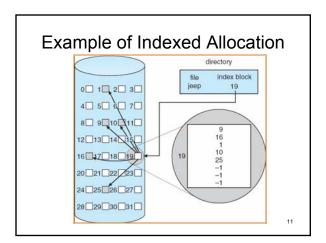
Contiguous Allocation

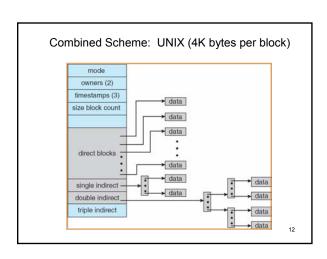
- Each file occupies a set of contiguous blocks on the disk
- Simple only starting location (block #) and length (number of blocks) are required
- Random access
- Wasteful of space (dynamic storage-allocation problem)
- · Files cannot grow





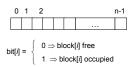






Free-Space Management

Bit vector (n blocks)



Block number calculation

(number of bits per word) *
(number of 0-value words) +
offset of first 1 bit

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Free-Space Management (Cont.)

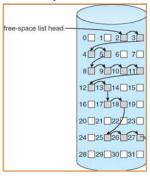
- · Bit map requires extra space
 - Example:

block size = 2^{12} bytes disk size = 2^{30} bytes (1 gigabyte) $n = 2^{30}/2^{12} = 2^{18}$ bits (or 32K bytes)

- · Easy to get contiguous files
- Linked list (free list)
 - Cannot get contiguous space easily
 - No waste of space
- · Grouping (many links per block)
- · Counting (aggregate continuous blocks)

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Linked Free Space List on Disk



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Directory Implementation

- · Each directory is a file
- Linear list of file names with pointer to the data blocks
 - simple to program
 - time-consuming to execute
- · Hash Table linear list with hash data structure
 - decreases directory search time
 - collisions situations where two file names hash to the same location
 - fixed size

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Efficiency and Performance

- · Efficiency dependent on:
 - disk allocation and directory algorithms
 - types of data kept in file's directory entry
- Performance
 - disk cache separate section of main memory for frequently used blocks
 - free-behind and read-ahead techniques to optimize sequential access
 - improve PC performance by dedicating section of memory as virtual disk, or RAM disk

Page Cache

- A page cache caches pages rather than disk blocks using virtual memory techniques
- · Memory-mapped I/O uses a page cache
- Routine I/O through the file system uses the buffer (disk) cache
- · Sync, Async, flush

Recovery

- Consistency checking compares data in directory structure with data blocks on disk, and tries to fix inconsistencies
- Use system programs to back up data from disk to another storage device (floppy disk, magnetic tape, other magnetic disk, optical)
- Recover lost file or disk by restoring data from backup

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Log Structured File Systems

- Log structured (or journaling) file systems record each update to the file system as a transaction
- All transactions are written to a log
 - A transaction is considered **committed** once it is written to the log
 - However, the file system may not yet be updated
- The transactions in the log are asynchronously written to the file system
 - When the file system is modified, the transaction is removed from the log
- If the file system crashes, all remaining transactions in the log must still be performed