0 Changes

- 2/16: Added typing rule for empty blocks

1 Types

The Xi type system uses a somewhat bigger set of types than can be expressed explicitly in the source language:

\[
\begin{align*}
\tau &::= \text{int} \mid \text{bool} \mid \tau[\,] \\
T &::= \tau \mid \text{unit} \mid (\tau_1, \tau_2, \ldots, \tau_n) \quad n \geq 2 \\
R &::= \text{unit} \mid \text{void} \\
\sigma &::= \text{var} \tau \\
&\quad \mid \text{ret} T \\
&\quad \mid \text{fn} T \to T'
\end{align*}
\]

Ordinary types expressible in the language are denoted by the metavariable \( \tau \), which can be int, bool, or an array type.

The metavariable \( T \) denotes an expanded notion of type that represents possible types in procedures, functions, and multiple assignments. It may be an ordinary type, unit, or a tuple type.

The type unit does not appear explicitly in the source language. This type is given to an element on the left-hand side of multiple assignments that uses the _ placeholder, allowing their handling be integrated directly into the type system. The unit type is also used to represent the result type of procedures.

Tuple types represent the parameter types of procedures and functions that take multiple arguments, or the return types of functions that return multiple results.

The metavariable \( R \) represents the outcome of evaluating a statement, which can be either unit or void. The unit type is the type of statements that might complete normally and permit the following statement to execute. The type void is the type of statements such as return that never pass control to the following statement. The void type should not be confused with the C/Java type void, which is actually closer to unit.

The set \( \sigma \) is used to represent typing-environment entries, which can either be normal variables (bound to var \( \tau \) for some type \( \tau \)), return types (bound to ret \( T \)), functions (bound to fn \( T \to T' \) where \( T' \neq \text{unit} \)), or procedures (bound to fn \( T \to \text{unit} \)), where the “result type” unit indicates that the procedure result contains no information other than that the procedure call terminated.

2 Subtyping

The subtyping relation on \( T \) is the least partial order consistent with this rule:

\[
\tau \leq \text{unit}
\]

For now, however, there is no subsumption rule, so subtyping matters only where it appears explicitly.

3 Type-checking expressions

To type-check expressions, we need to know what bound variables and functions are in scope; this is represented by the typing context \( \Gamma \), which maps names \( x \) to types \( \sigma \).

The judgment \( \Gamma \vdash e : T \) is the rule for the type of an expression; it states that with bindings \( \Gamma \) we can conclude that \( e \) has the type \( T \).
4 Type-checking statements

To type-check statements, we need all the information used to type-check expressions, plus the types of procedures, which are included in \( \Gamma \). In addition, we extend the domain of \( \Gamma \) a little to include a special symbol \( \rho \). To check the return statement we need to know what the return type of the current function is or if it is a procedure. Let this be denoted by \( \Gamma(\rho) = \text{ret } T \), where \( T \neq \text{unit} \) if the statement is part of a function, or \( T = \text{unit} \) if the statement is part of a procedure. Since statements include declarations, they can also produce new variable bindings, resulting in an updated typing context which we will denote as \( \Gamma' \). To update typing contexts, we write \( \Gamma[x \mapsto \text{var } \tau] \), which is an environment exactly like \( \Gamma \) except that it maps \( x \) to \( \text{var } \tau \). We use the metavariable \( s \) to denote a statement, so the main typing judgment for statements has the form \( \Gamma \vdash s : R, \Gamma' \).

Most of the statements are fairly straightforward and do not change \( \Gamma \):

\[
\frac{}{\Gamma \vdash \text{true} : \text{bool}} \quad \frac{}{\Gamma \vdash \text{false} : \text{bool}} \quad \frac{}{\Gamma \vdash \text{char} : \text{int}} \quad \frac{}{\Gamma \vdash n : \text{int}}
\]

\[
\frac{\Gamma(x) = \text{var } \tau}{\Gamma \vdash x : \tau} \quad \frac{\Gamma \vdash e_1 : \text{int} \quad \Gamma \vdash e_2 : \text{int} \quad \Theta \in \{+, -, *, ^{>, /}, \%\}} {\Gamma \vdash e_1 \Theta e_2 : \text{int}}
\]

\[
\frac{\Gamma \vdash e_1 : \text{int} \quad \Gamma \vdash e_2 : \text{int} \quad \Theta \in \{==, !=, <, <=, >, >=\}} {\Gamma \vdash e_1 \Theta e_2 : \text{bool}}
\]

\[
\frac{\Gamma \vdash e_1 : \text{bool} \quad \Gamma \vdash e_2 : \text{bool} \quad \Theta \in \{==, !=, \&, |\}} {\Gamma \vdash e_1 \Theta e_2 : \text{bool}}
\]

\[
\frac{\Gamma \vdash e : \tau[\cdot]} {\Gamma \vdash \text{length}(e) : \text{int}}
\]

\[
\frac{\Gamma \vdash e_i : \tau \quad \ldots \quad \Gamma \vdash e_n : \tau \quad n \geq 0} {\Gamma \vdash \{e_1, \ldots, e_n\} : \tau[\cdot]} \quad \frac{\Gamma \vdash e_1 : \tau[\cdot] \quad \Gamma \vdash e_2 : \text{int} \quad \Theta \in \{==, !=\}} {\Gamma \vdash e_1 \Theta e_2 : \text{bool}}
\]

\[
\frac{\Gamma(f) = \text{fn } \text{unit} \rightarrow T' \quad T' \neq \text{unit}} {\Gamma \vdash f() : T'}
\]

\[
\frac{\Gamma(f) = \text{fn } \tau \rightarrow T' \quad T' \neq \text{unit}} {\Gamma \vdash e : \tau}
\]

\[
\frac{\Gamma(f) = \text{fn } (\tau_1, \ldots, \tau_n) \rightarrow T' \quad T' \neq \text{unit} \quad \Gamma \vdash e_i : \tau_j \quad (\forall j \in 1..n) \quad n \geq 2} {\Gamma \vdash f(e_1, \ldots, e_n) : T'}
\]
The function $\text{hub}$ is defined as follows:

$$\text{hub}(R, R) = R \quad \text{hub}(\text{unit}, R) = \text{hub}(R, \text{unit}) = \text{unit}$$

Therefore, the type of an if is void only if all branches have that type.

Assignments require checking the left-hand side to make sure it is assignable:

$$\Gamma(x) = \text{var} \tau \quad \Gamma \vdash e : \tau \quad (\text{ASSIGN})$$

$$\Gamma \vdash x = e : \tau, \Gamma$$

$$\Gamma \vdash e_1 : \tau[i] \quad \Gamma \vdash e_2 : \text{int} \quad \Gamma \vdash e_3 : \tau \quad (\text{ARRASSIGN})$$

$$\Gamma \vdash e_1[e_2] = e_3 : \tau, \Gamma$$

Declarations are the source of new bindings. Three kinds of declarations can appear in the source language: variable declarations, multiple assignments, and function/procedure declarations. We are only concerned with the first two kinds within a function body. To handle multiple assignments, we define a declaration $d$ that can appear within a multiple assignment:

$$d ::= x : \tau \mid _{-}$$

and define functions $\text{typeof}(d)$ and $\text{varsof}(d)$ as follows:

$$\text{typeof}(x : \tau) = \tau \quad \text{typeof}(_{-}) = \text{unit}$$

$$\text{varsof}(x : \tau) = \{x\} \quad \text{varsof}(_{-}) = \emptyset$$

Using these notations, we have the following rules:

$$\Gamma \vdash x : \tau, \Gamma[x \mapsto \text{var} \tau] \quad (\text{VARDECL})$$

$$\Gamma \vdash x : \tau, \Gamma[x \mapsto \text{var} \tau] \quad (\text{VARINIT})$$

$$\Gamma \vdash e : \tau \quad (\text{EXPRSTMT})$$

$$\Gamma \vdash x : \tau, \Gamma[x \mapsto \tau[i] \ldots [j] : \text{unit}, \Gamma[x \mapsto \text{var} \tau[i] \ldots [j]]] \quad (\text{ARRAYDECL})$$

$$\Gamma \vdash e : (\tau_1, \ldots, \tau_n) \quad \tau_i \leq \text{typeof}(d_i) \quad (\forall i \in 1..n)$$

$$\text{dom}(\Gamma) \cap \text{varsof}(d_i) = \emptyset \quad (\forall i \in 1..n)$$

$$\text{varsof}(d_i) \cap \text{varsof}(d_j) = \emptyset \quad (\forall i, j \in 1..n, i \neq j)$$

$$\Gamma \vdash d_1, \ldots, d_n = e : \text{unit}, \Gamma[x_i \mapsto \text{var} \text{typeof}(d_i)] \quad (\forall i \in 1..n \forall x_i, \text{varsof}(d_i) = \{x_i\}) \quad (\text{MULTIASSIGN})$$

The final premise in rule $\text{MULTIASSIGN}$ prevents shadowing by ensuring that $\text{dom}(\Gamma)$ and all of the $\text{varsof}(d_i)$’s are disjoint from each other.
5 Checking top-level declarations

At the top level of the program, we need to figure out the types of procedures and functions, and make sure their bodies are well-typed. Since mutual recursion is supported, this needs to be done in two passes. First, we use the judgment $\Gamma \vdash fd : \Gamma'$ to state that the function or procedure declaration $fd$ extends top-level bindings $\Gamma$ to $\Gamma'$:

$$f \notin \text{dom}(\Gamma) \quad f \notin \text{dom}(\Gamma)$$

$$\Gamma \vdash f(\tau) \quad s : \Gamma[f \mapsto \text{fn } \tau \rightarrow \text{unit}] \\
\Gamma \vdash f(x) \quad s : \Gamma[f \mapsto \text{fn } \tau \rightarrow \text{unit}]$$

$$f \notin \text{dom}(\Gamma) \quad \Gamma' = \Gamma[f \mapsto \text{fn } (\tau_1, \ldots, \tau_n) \rightarrow \text{unit}] \quad n \geq 2$$

$$\Gamma \vdash f(x_1 : \tau_1, \ldots, x_n : \tau_n) \quad s : \Gamma'$$

$$f \notin \text{dom}(\Gamma) \quad \Gamma' = \Gamma[f \mapsto \text{fn } (\tau_1, \ldots, \tau_n) \rightarrow \tau'] \quad n \geq 2$$

$$\Gamma \vdash f(x_1 : \tau_1, \ldots, x_n : \tau_n) : \tau' \quad s : \Gamma'$$

$$f \notin \text{dom}(\Gamma) \quad \Gamma' = \Gamma[f \mapsto \text{fn } (\tau_1', \ldots, \tau_m') \rightarrow \text{unit}] \quad m \geq 2$$

$$\Gamma \vdash f(x_1 : \tau_1', \ldots, x_n : \tau_n) \quad s : \Gamma'$$

$$f \notin \text{dom}(\Gamma) \quad \Gamma' = \Gamma[f \mapsto \text{fn } (\tau_1', \ldots, \tau_m') \rightarrow (\tau_1'', \ldots, \tau_m'')] \quad m \geq 2$$

$$\Gamma \vdash f(x_1 : \tau_1', \ldots, x_n : \tau_n) : \tau_1'', \ldots, \tau_m''' \quad s : \Gamma'$$

The second pass over the program is captured by the judgment $\Gamma \vdash fd \; \text{def}$, which defines how to check well-formedness of each function definition against a top-level environment $\Gamma$, ensuring that parameters do not shadow anything and that the body is well-typed. We treat procedures just like functions that return the unit type. The body of a procedure definition may have any type, but the body of a function definition must have type void, which ensures that the function body does not fall off the end without returning.

$$\Gamma[\rho \mapsto \text{ret unit}] \vdash s : R, \Gamma'$$

$$\Gamma \vdash f(\tau) \quad s \; \text{def}$$

$$x \notin \text{dom}(\Gamma) \quad \Gamma[\tau \mapsto \text{var } \rho \mapsto \text{ret unit}] \vdash s : R, \Gamma'$$

$$\Gamma \vdash f(x : \tau) \quad s \; \text{def}$$

$$|\text{dom}(\Gamma) \cup \{x_1, \ldots, x_n\}| = |\text{dom}(\Gamma)| + n \quad n \geq 2$$

$$\Gamma[x_1 \mapsto \text{var } \tau_1, \ldots, x_n \mapsto \text{var } \tau_n, \rho \mapsto \text{ret unit}] \vdash s : R, \Gamma'$$

$$\Gamma \vdash f(x_1 : \tau_1, \ldots, x_n : \tau_n) \quad s \; \text{def}$$

$$\Gamma[\rho \mapsto \text{ret } \tau'] \vdash s : \text{void}, \Gamma'$$

$$\Gamma \vdash f(\tau') \quad s \; \text{def}$$

$$x \notin \text{dom}(\Gamma) \quad \Gamma[\tau \mapsto \text{var } \rho \mapsto \text{ret } \tau'] \vdash s : \text{void}, \Gamma'$$

$$\Gamma \vdash f(x : \tau) \quad s \; \text{def}$$
Using the previous judgments, we can define when an entire program $fd_1\; fd_2\; \ldots\; fd_n$ that does not contain a use declaration is well-formed, written $\vdash fd_1\; fd_2\; \ldots\; fd_n$. For brevity, the rules for adding declarations appearing in interfaces are omitted. These rules are slightly different from those of the form $\Gamma \vdash fd : \Gamma'$ in Section 5, where $f \not\in \text{dom}(\Gamma)$ is replaced with appropriate conditions. Once added, these declarations also permits declarations in the source file of identical signature. See Section 8 of the Xi Language Specification.