Outline

- Announcements
 - HWII due Today
 - HWIII due Monday
 - Wed. and Fri. in 484 Rhodes
- Parallel Computers
- Types of Parallelism
- Message Passing with MPI
- Parallelizing RAD1D

Basic Idea of Parallel Computing

- If you have N computers
 - Divide your problem into N pieces
 - Give each piece to its own computer
 - Solve your problem N-times faster

Parallel Computers

- 2-4 processor workstations are becoming common
- Sun, SGI make systems with up to 32 processors
- Research systems:
 - Up to 9632 processors!

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Parallel Computers

- Symmetric Mutli-Processors (SMP)
 - Several processors sharing same memory
 - Ex: 2x Dell Workstation, 32x SGI Origin
- Distributed Memory
 - Several independent machines linked by a network
 - Ex: "Beowulf" cluster
- Hybrid Systems
 - Several machines linked over network, each machine is SMP
 - Ex: Cornell's Velocity, ASCI White

Using Parallel Computers: Easy Way

- If your problem involves applying a simple model to several different data sets
 - Distribute data among machines
 - Run the program
 - If it takes T seconds to run on 1 machine
 - Your problem will run in T/N seconds plus time to distribute data (D)

Speed-Up

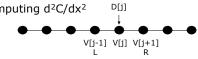
- Parallel performance is commonly measured as speed-up:
 - $-S(N)=T_1/T_N$
 - Ideal case: $T_N=T_1/N$, then S(N)=N
 - Easy parallelism
 - $T_N = T_1/N + D$
 - If D is small, S(N) is very close to N

Hard Parallelism

- You have a single very large problem
 - large in memory or time
- To parallelize:
 - you must divide your problem (data or program) into pieces
 - But, pieces are not independent
 - computers must share data
 - you must determine what data to share and when

Example:

• Computing d²C/dx²



- for(j=0;j<M;j++){
 - L=V[j-1] or V[M-1] R=V[j+1] or V[0]
- D[j]=(R-2V[j]+L)/dx/dx
- }

Computing d²C/dx²

• Divide domain into 2 pieces of length

M/2 D[j] - V[j-1] V[j] V[j+1] L R

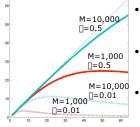
- for(j=0; j<M/2-1; j++){
 - L=V[j-1] or V[M/2-1] on neighbor
 - R=V[j+1] or V[0] on neighbor
 - D[j]=(R-2V[j]+L)/dx/dx
- }

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Computing d²C/dx²

- $T_N = T_1/N + 2*N*C$, where C=time to send a double
- $S(N)=T_1/(T_1/T_N+2NC)$
- $\bullet = N/(1+2N^2C/T_1)$
- Let T1=M\(\text{\textsuper}\)C, M=grid points, \(\text{\textsuper}\)
 =compute/communicate
- S(N)=N/(1+2N²/(M□))

Amdahl's Law of Parallel Processing



- Increasing N will reduce speed-up
 - There's a limit to parallel performance
- Increasing M will improve speed-up
 - Parallel is more attractive for big problems
- Increasing [] will improve speed-up
 - Hardware is critical

Message Passing with MPI

- We need a way to share data
- Message Passing Interface (MPI)
 - library for communication
 - device independent
 - same MPI code will run on cluster or SMP
 - well-specified, so very portable
 - versions available for all systems
 - MPICH, LAM are free

Using MPI

- First, set it up:
 - MPI_Init(&argc, &argv);
 - Starts MPI
 - MPI_Comm_rank(MPI_COMM_WORLD, &P
);
 - Gets "rank" of this job--processor number
 - MPI_Comm_size(MPI_COMM_WORLD, &N);
 - Gets number of jobs

Using MPI

- Point-to-Point Communication
 - MPI_Send(ovec, 10, MPI_DOUBLE, 1, tag, comm)
 - Sends 10 doubles in array ovec to process 1
 - tag is a message label
 - comm is a communicator--a subgroup of processes (MPI_COMM_WORLD is everyone)
 - MPI_Recv(ivec, 10, MPI_DOUBLE, 0, tag, comm, &status)
 - receives 10 doubles from process 0
 - status may contain error info

Point-to-Point Communication

- For every send, there must be a receive
- Send waits until Recv is posted & viceversa
 - Called blocking
- Allows for the possibility of deadlock
 - Recv's waiting for Sends that will never be posted

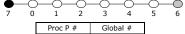
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Collective Communication

- Sometimes, you need to send data to all processes
- Ex: Computing sum of an array
 - Each processor computes the sum of its array
 - The partial sums must be added together and answer distributed to everyone
 - MPI_Allreduce(&locsum, &allsum, 1, MPI_DOUBLE, MPI_SUM, MPI_COMM_WORLD);

Parallelizing RAD1D

- Divide domain among N processors
 - Each proc gets M/N grid points



0	M/N*P+0
1	M/N*P+1
:	:
M/N-1	M/N*(P+1)-1
M/N	M/N*(P+1)
M/N+1	M/N*(P)-1

Need a subroutine to get last two entries from neighbors

Parallelizing RAD1D

- Each proc will get N and P
- Each proc will read inputs, taking what it needs
- Compute as usual (m=M/N) updating arrays with communication routine as needed: Update(C,m,P,N);

Update(C,m,P,N);
for(j=0;j<m;j++){
 left=j-1; if(j==0){left=m+1;}
 right=j+1;
 RHS[j]=C[j]+0.5*lambda[j]*(C[right]-C[left]);
}
Update(RHS,m,P,N);</pre>

Parallelizing RAD1D

Update(Arr, m, P, N);

if(N is odd){Error}

left=P-1; if(left<0){left=N-1;};

right=P+1; if(right==N){right=0;}

if(P is odd)

MP1_Send(&Arr[0], 1, MP1_DOUBLE, left, MP1_ANY_TAG,

MP1_COMM_WORLD);

MP1_Recv(&Arr[m+1], 1, MP1_DOUBLE, left, MP1_ANY_TAG,

MP1_COMM_WORLD, Satat);

MP1_Send(&Arr[m+1], 1, MP1_DOUBLE, right, MP1_ANY_TAG,

MP1_COMM_WORLD);

MP1_Recv(&Arr[m], MP1_DOUBLE, right, MP1_ANY_TAG,

MP1_COMM_WORLD, &Satat);

else

MP1_Recv(&Arr[m], 1, MP1_DOUBLE, right, MP1_ANY_TAG,

MP1_COMM_WORLD, &Satat);

MP1_Rend(&Arr[m], 1, MP1_DOUBLE, right, MP1_ANY_TAG,

MP1_COMM_WORLD), &Satat);

MP1_Recv(&Arr[m], 1, MP1_DOUBLE, left, MP1_ANY_TAG,

MP1_Recv(&Arr[m], 1, MP1_DOUBLE, left, MP1_ANY_TAG,

MP1_Recv(&Arr[m], 1, MP1_DOUBLE, left, MP1_ANY_TAG,

MP1_COMM_WORLD);

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