

# **Assemblers, Linkers, and Loaders**

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**CS 3410**

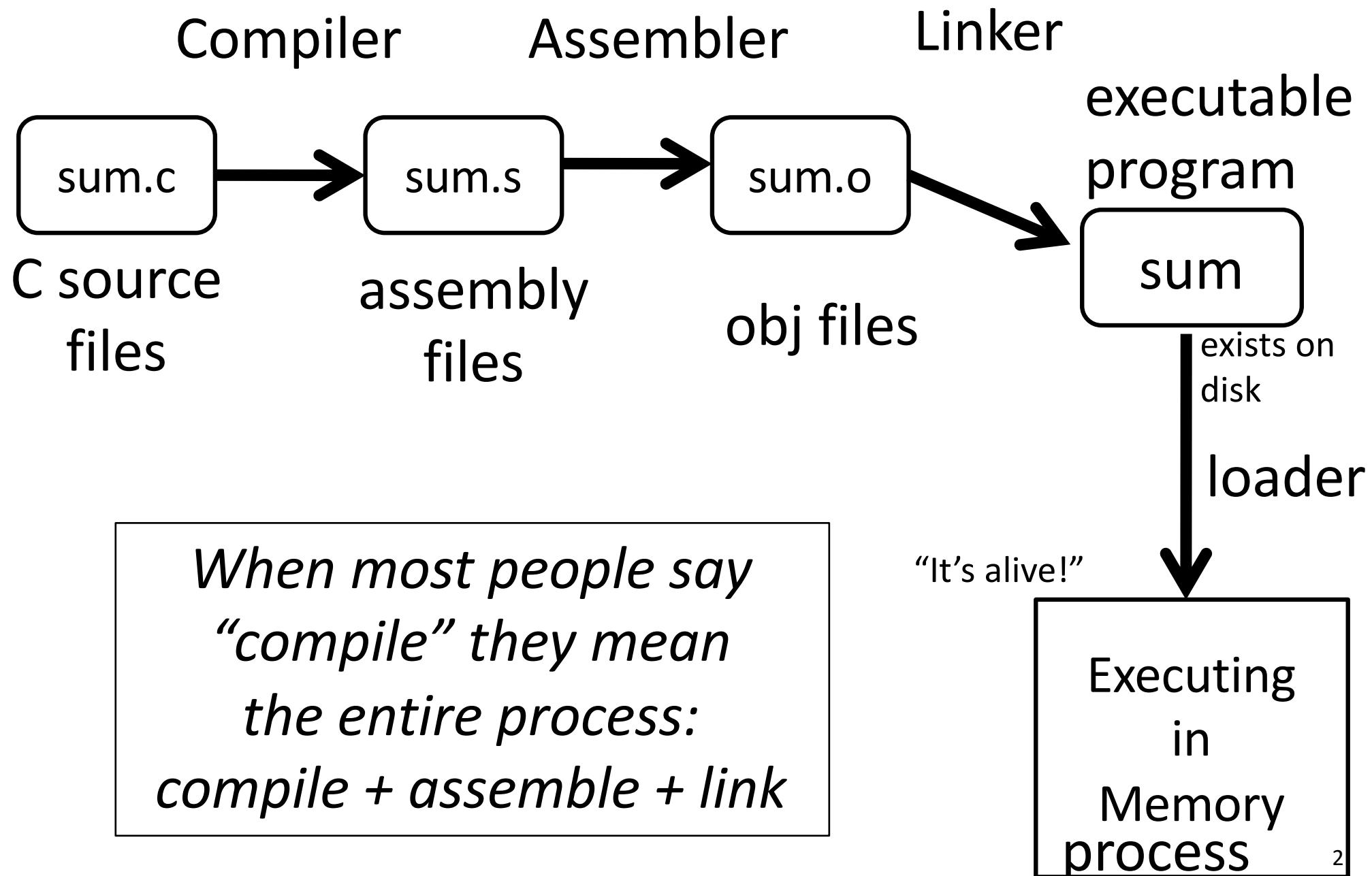
Computer Science

Cornell University

The slides are the product of many rounds of teaching CS 3410 by Professors Weatherspoon, Bala, Bracy, McKee, and Sirer.

See: P&H Appendix A1-2, A.3-4 and 2.12

# From Writing to Running



# sum.c

```
#include <stdio.h>

int n = 100;

int main (int argc, char* argv[ ]) {
    int i;
    int m = n;
    int sum = 0;

    for (i = 1; i <= m; i++) {
        sum += i;
    }
    printf ("Sum 1 to %d is %d\n", n, sum);
}
```

csug03> mipsel-linux-gcc -S sum.c

export PATH=\${PATH}:/courses/cs3410/mipsel-linux/bin:/courses/cs3410/mips-sim/bin

or

setenv PATH \${PATH}:/courses/cs3410/mipsel-linux/bin:/courses/cs3410/mips-sim/bin

tells compiler to  
produce sum.s file

# From Writing to Running: Command Line

```
# Compile
```

```
csug01> mipsel-linux-gcc -S sum.c
```

```
# Assemble
```

```
csug01> mipsel-linux-gcc -c sum.s
```

```
# Link
```

```
csug01> mipsel-linux-gcc -o sum sum.o ${LINKFLAGS}
```

```
# -nostartfiles -nodefaultlibs
```

```
# -static -mno-xgot -mno-embedded-pic  
-mno-abicalls -G 0 -DMIPS -Wall
```

```
# Load
```

```
csug01> simulate sum
```

```
Sum 1 to 100 is 5050
```

```
MIPS program exits with status 0 (approx. 2007  
instructions in 143000 nsec at 14.14034 MHz)
```

# sum.s

	.data		\$L2:	lw	\$2,24(\$fp)
	.globl n			lw	\$3,28(\$fp)
	.align 2			slt	\$2,\$3,\$2
n:	.word 100			bne	\$2,\$0,\$L3
	.rdata			lw	\$3,32(\$fp)
	.align 2			lw	\$2,24(\$fp)
\$str0:	.asciiiz	"Sum 1 to %d is %d\n"		addu	\$2,\$3,\$2
	.text			sw	\$2,32(\$fp)
	.align 2			lw	\$2,24(\$fp)
	.globl main			addiu	\$2,\$2,1
main:	addiu \$sp,\$sp,-48			sw	\$2,24(\$fp)
	sw \$31,44(\$sp)			b	\$L2
	sw \$fp,40(\$sp)		\$L3:	la	\$4,\$str0
	move \$fp,\$sp			lw	\$5,28(\$fp)
	sw \$4,48(\$fp)			lw	\$6,32(\$fp)
	sw \$5,52(\$fp)			jal	printf
	la \$2,n			move	\$sp,\$fp
	lw \$2,0(\$2)			lw	\$31,44(\$sp)
	sw \$2,28(\$fp)			lw	\$fp,40(\$sp)
	sw \$0,32(\$fp)			addiu	\$sp,\$sp,48
	li \$2,1			j	\$31
	sw \$2,24(\$fp)				

# sum.s

	.data		\$L2:	lw	\$2,24(\$fp)	i=1
	.globl n			lw	\$3,28(\$fp)	m=100
	.align 2			slt	\$2,\$3,\$2	if(m < i)
n:	.word 100			bne	\$2,\$0,\$L3	100 < 1
	.rdata			lw	\$3,32(\$fp)	v1=0(sum)
	.align 2			lw	\$2,24(\$fp)	v0=1(i)
\$str0:	.asciiiz	"Sum 1 to %d is %d\n"		addu	\$2,\$3,\$2	v0=1(0+1)
	.text			sw	\$2,32(\$fp)	sum=1
	.align 2			lw	\$2,24(\$fp)	i=1
	.globl main			addiu	\$2,\$2,1	i=2 (1+1)
main:	addiu \$sp,\$sp,-48			sw	\$2,24(\$fp)	i=2
prologue	sw \$31,44(\$sp)			b	\$L2	
	sw \$fp,40(\$sp)		\$L3:	la	\$a0\$4,\$str0	str
	move \$fp,\$sp			call	\$a1\$5,28(\$fp)	m=100
	sw \$a0\$4,48(\$fp)		printf	1w	\$a2\$6,32(\$fp)	sum
	sw \$a1\$5,52(\$fp)			jal	printf	
	la \$v0\$2,n			move	\$sp,\$fp	
	lw \$2,0(\$2) \$v0=100		epilogue	1w	\$31,44(\$sp)	
	sw \$2,28(\$fp) m=100			1w	\$fp,40(\$sp)	
	sw \$0,32(\$fp) sum=0			addiu	\$sp,\$sp,48	
	li \$2,1			j	\$31	
	sw \$2,24(\$fp) i=1					

# Assembler

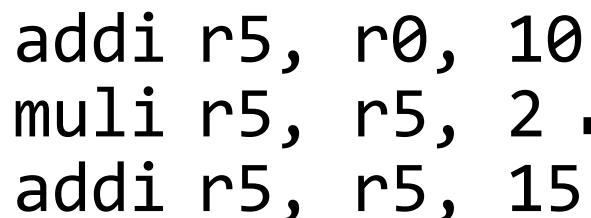
Input: Assembly File (.s)

- assembly instructions, pseudo-instructions
- program data (strings, variables), layout directives

Output: Object File in binary machine code

MIPS instructions in executable form  
(.o file in Unix, .obj in Windows)

```
addi r5, r0, 10
muli r5, r5, 2
addi r5, r5, 15
```



```
00100000000010100000000000001010
000000000000101001010001000000
00100001010010100000000000001111
```

# MIPS Assembly Instructions

## Arithmetic/Logical

- ADD, ADDU, SUB, SUBU, AND, OR, XOR, NOR, SLT, SLTU
- ADDI, ADDIU, ANDI, ORI, XORI, LUI, SLL, SRL, SLLV, SRLV, SRAV, SLTI, SLTIU
- MULT, DIV, MFLO, MTLO, MFHI, MTHI

## Memory Access

- LW, LH, LB, LHU, LBU, LWL, LWR
- SW, SH, SB, SWL, SWR

## Control flow

- BEQ, BNE, BLEZ, BLTZ, BGEZ, BGTZ
- J, JR, JAL, JALR, BEQL, BNEL, BLEZL, BGTZL

## Special

- LL, SC, SYSCALL, BREAK, SYNC, COPROC

# Pseudo-Instructions

Assembly shorthand, technically not machine instructions, but easily converted into 1+ instructions that are

Pseudo-Insns	Actual Insns	Functionality
NOP	SLL r0, r0, 0	# do nothing
MOVE reg, reg	ADD r2, r0, r1	# copy between regs
LI reg, 0x45678	LUI reg, 0x4 ORI reg, reg, 0x5678	#load immediate
BLT reg, reg, label	SLT r1, rA, rB BNE r1, r0, label	# branch less than

*+ a few more...*

# Symbols and References

```
int pi = 3;  
int e = 2;  
  
static int randomval = 7;  
  
(external == defined in another file)  
extern char *username;  
extern int printf(char *str, ...);  
  
int square(int x) { ... }  
static int is_prime(int x) { ... }  
int pick_prime() { ... }  
int pick_random() {  
    return randomval;  
}
```

Global labels: Externally visible “exported” symbols

- Can be referenced from other object files
- Exported functions, global variables
- Examples: pi, e, username, printf, pick\_prime, pick\_random

Local labels: Internally visible only symbols

- Only used within this object file
- static functions, static variables, loop labels, ...
- Examples: randomval, is\_prime

# Handling forward references

## Example:

bne \$1, \$2, L

# *Looking for L*

sll \$0, \$0, 0

L: addiu \$2, \$3, 0x2

## *Found L*

The assembler will change this to

bne \$1, \$2, +1

sll \$0, \$0, 0

addiu \$2, \$3, \$0x2

# Final machine code

0X14220001 # bne

0x00000000 # sll

0x24620002 # addiu

# Object file

## Header

- Size and position of pieces of file

## Text Segment

- instructions

## Data Segment

- static data (local/global vars, strings, constants)

## Debugging Information

- line number → code address map, *etc.*

## Symbol Table

- External (exported) references
- Unresolved (imported) references

# Object File Formats

## Unix

- a.out
- COFF: Common Object File Format
- ELF: Executable and Linking Format
- ...

## Windows

- PE: Portable Executable

All support both executable and object files

# Objdump disassembly

```
csug01> mipsel-linux-objdump --disassemble math.o
math.o:      file format elf32-tradlittlemips
Disassembly of section .text:
```

00000000 <pick\_random>:

0:	27bdfff8	addiu	sp,sp,-8	<i>prologue</i>	unresolved symbol (see symbol table next slide)
4:	afbe0000	sw	s8,0(sp)		
8:	03a0f021	move	s8,sp		
c:	3c020000	lui	v0,0x0		
10:	8c420008	lw	v0,8(v0)	<i>body</i>	
14:	03c0e821	move	sp,s8		
18:	8fbe0000	lw	s8,0(sp)		
1c:	27bd0008	addiu	sp,sp,8		
20:	03e00008	jr	ra	<i>epilogue</i>	
24:	00000000	nop			

```
static int randomval = 7;
int pick_random() { return randomval; }
```

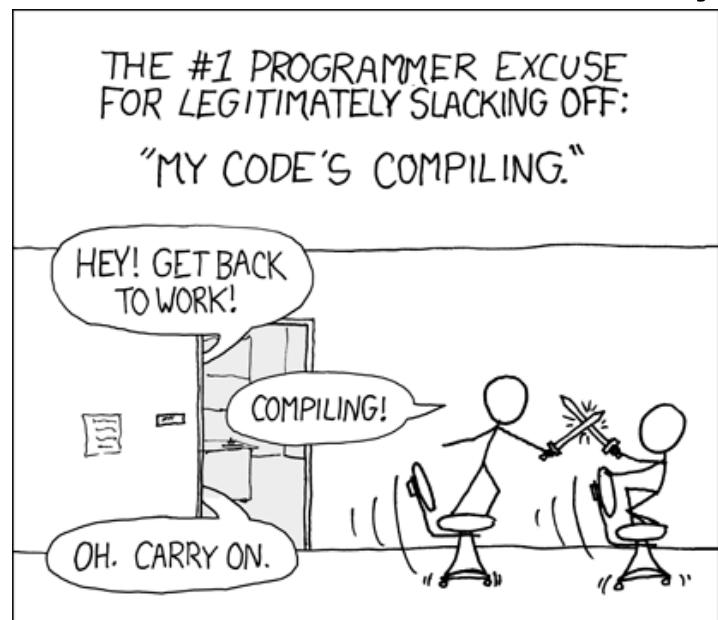
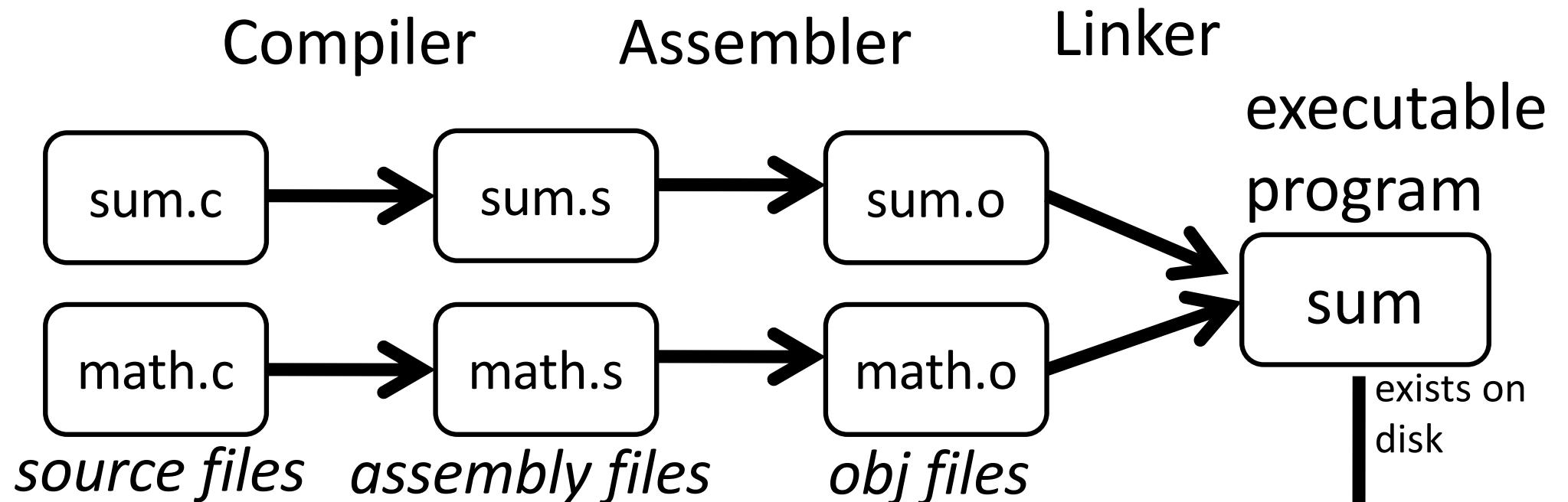
# Objdump symbols

csug01 ~\$ mipsel-linux-objdump --syms math.o  
 math.o: file format elf32-tradlittlemips

<u>SYMBOL TABLE:</u>	<i>segment</i>	<i>segment size</i>	[F]unction [O]bject [1]ocal [g]lobal
00000000 l	df *ABS*	00000000	math.c
00000000 l	d .text	00000000	.text
00000000 l	d .data	00000000	.data
00000000 l	d .bss	00000000	.bss
00000000 l	d .mdebug.abi32	00000000	.mdebug.abi32
00000008 l	O .data	00000004	randomval
00000060 l	F .text	00000028	is_prime ← static local
00000000 l	d .rodata	00000000	.rodata      function @
00000000 l	d .comment	00000000	.comment      address 0x60
00000000 g	O .data	00000004	pi
00000004 g	O .data	00000004	e      Size = x28 bytes
00000000 g	F .text	00000028	pick_random
00000028 g	F .text	00000038	square
00000088 g	F .text	0000004c	pick_prime
00000000	*UND*	00000000	username
00000000	*UND*	00000000	printf

*external references (undefined)*

# Separate Compilation & Assembly



small change ?  
→ recompile one  
module only

Executing  
in  
Memory  
process <sub>16</sub>

# Linkers

Linker combines object files into an executable file

- Resolve as-yet-unresolved symbols
- Each has illusion of own address space
  - Relocate each object's text and data segments
- Record top-level entry point in executable file

End result: a program on disk, ready to execute

- E.g.    ./sum                              Linux  
              ./sum.exe                          Windows  
              simulate sum                      Class MIPS simulator

# Static Libraries

*Static Library:* Collection of object files  
(think: like a zip archive)

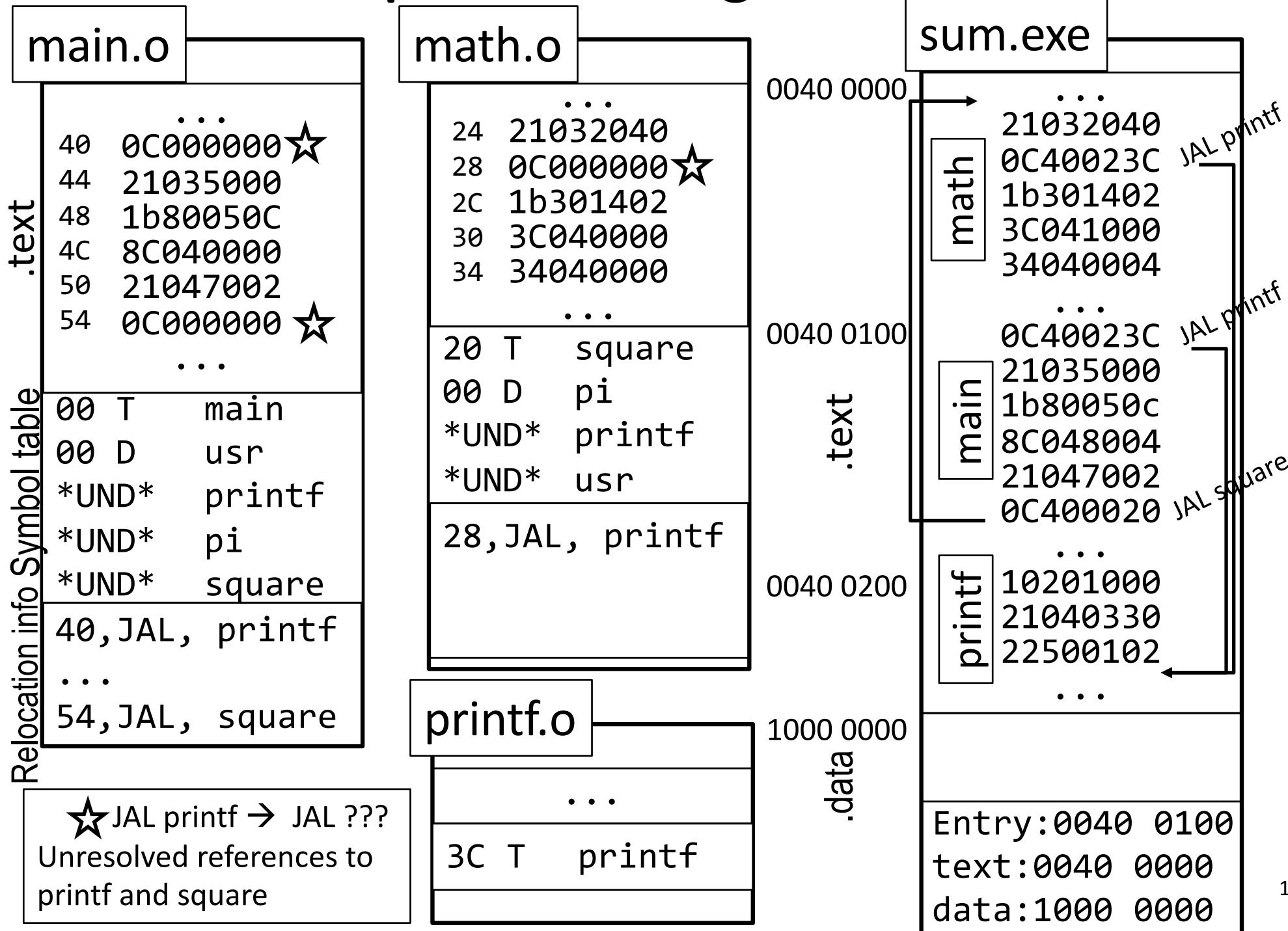
Q: Every program contains the entire library?!?

A: No, Linker picks only object files needed to resolve undefined references at link time

e.g. libc.a contains many objects:

- printf.o, fprintf.o, vprintf.o, sprintf.o, snprintf.o, ...
- read.o, write.o, open.o, close.o, mkdir.o, readdir.o, ...
- rand.o, exit.o, sleep.o, time.o, ....

# Linker Example: Resolving an External Fn Call



# iClicker Question

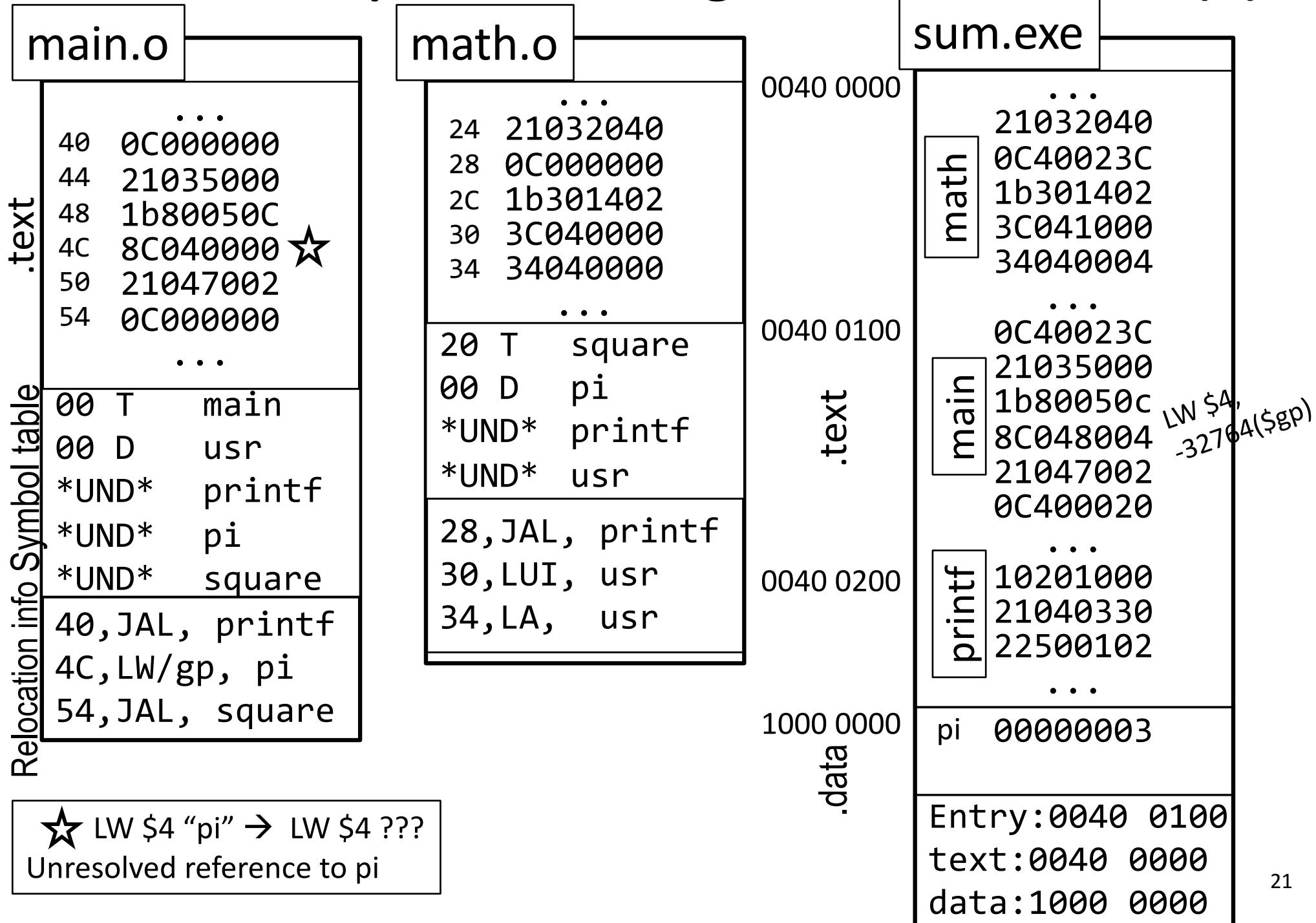
main.o		
Relocation info	Symbol table	.text
		40 0C000000 ...
		44 21035000
		48 1b80050C
		4C 8C040000
		50 21047002
		54 0C000000
		...
		00 T main
		00 D usr
		*UND* printf
		*UND* pi
		*UND* square
		40, JAL, printf
		4C, LW/gp, pi
		54, JAL, square

math.o		
Relocation info	Symbol table	.text
		24 21032040 ...
		28 0C000000
		2C 1b301402
		30 3C040000
		34 34040000
		...
		20 T square
		00 D pi
		*UND* printf
		*UND* usr
		28, JAL, printf
		30, LUI, usr
		34, LA, usr

Which symbols are undefined according to both main.o and math.o's symbol table?

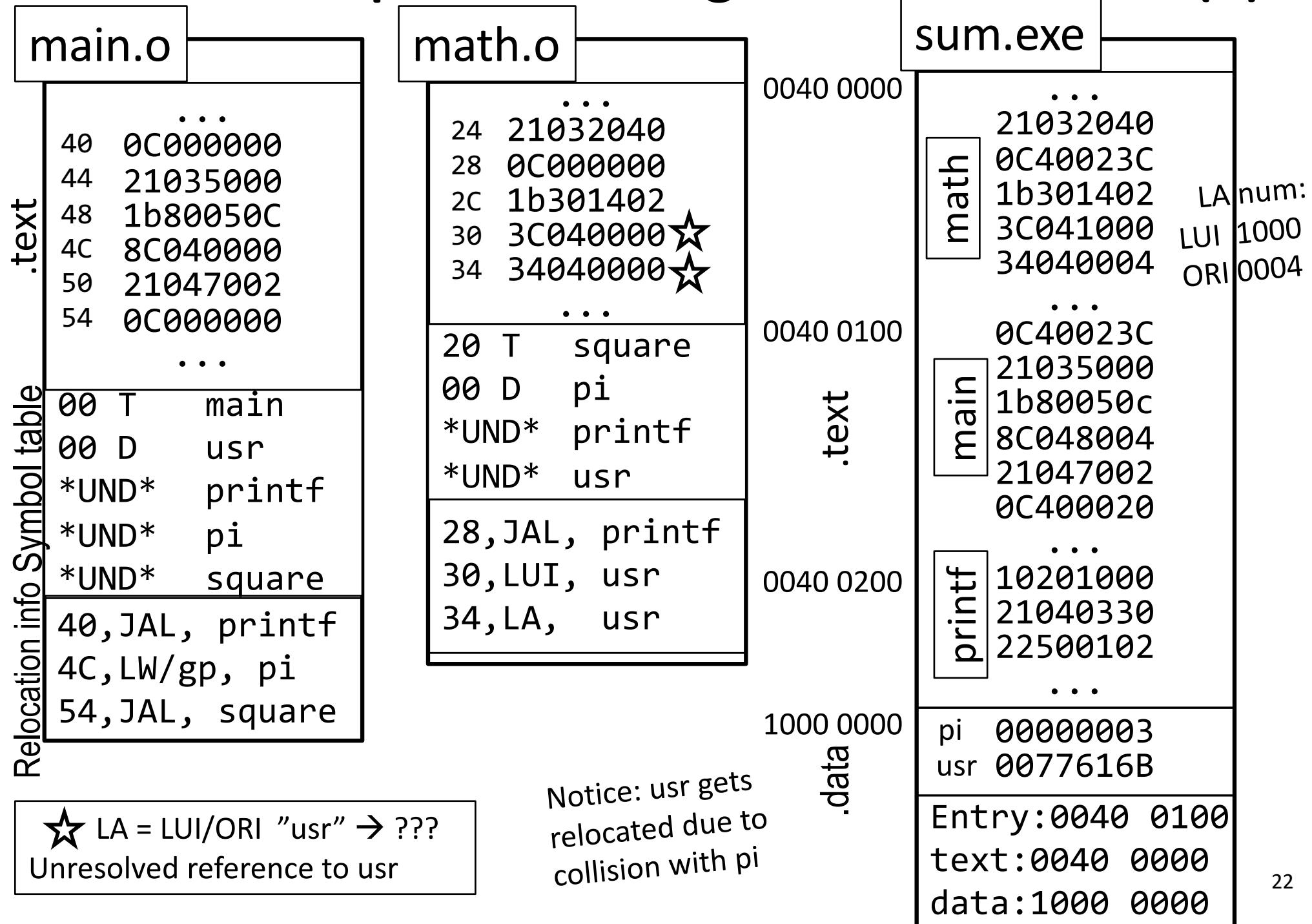
- A) printf
- B) pi
- C) square
- D) usr
- E) printf & pi

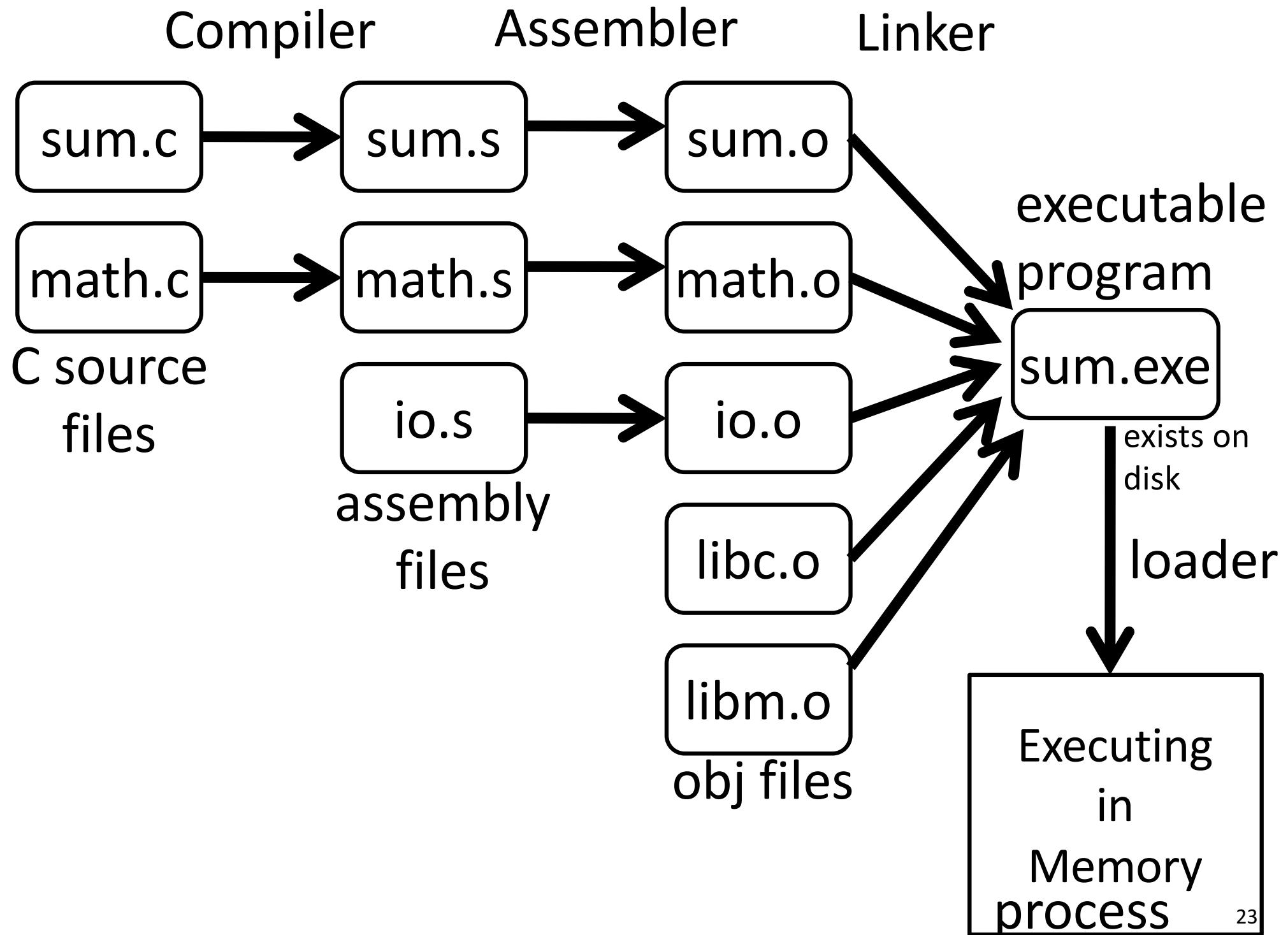
# Linker Example: Resolving Global Variables (1)



★ LW \$4 "pi" → LW \$4 ???  
Unresolved reference to pi

# Linker Example: Resolving Global Variables (2)





# Loaders

*Loader* reads executable from disk into memory

- Initializes registers, stack, arguments to first function
- Jumps to entry-point

Part of the Operating System (OS)

# Shared Libraries

Q: Every program contains parts of same library?!?

A: No, they can use shared libraries

- Executables all point to single *shared library* on disk
- final linking (and relocations) done by the loader

Optimizations:

- Library compiled at fixed non-zero address
- Jump table in each program instead of relocations
- Can even patch jumps on-the-fly

# Static and Dynamic Linking

## Static linking

- Big executable files (all/most of needed libraries inside)
- Don't benefit from updates to library
- No load-time linking

## Dynamic linking

- Small executable files (just point to shared library)
- Library update benefits all programs that use it
- Load-time cost to do final linking
  - But dll code is probably already in memory
  - And can do the linking incrementally, on-demand

# Takeaway

Compiler produces assembly files

- (contain MIPS assembly, pseudo-instructions, directives, etc.)

Assembler produces object files

- (contain MIPS machine code, missing symbols, some layout information, etc.)

Linker joins object files into one executable file

- (contains MIPS machine code, no missing symbols, some layout information)

Loader puts program into memory, jumps to 1<sup>st</sup> insn,  
and starts executing a *process*

- (machine code)