

# CS 3110

## Formal Methods

Nate Foster

Spring 2019

Today's music: Theme from *Downton Abbey* by John Lunn

# Review

## Previously in 3110:

- Functional programming
- Modular programming
- Data structures
- Interpreters

Next unit of course: [formal methods](#)

## Today:

- Proof assistants
- Functional programming in Coq
- Propositional logic
- Simple proofs about programs

# Approaches to validation [lec 11]

- Social
  - Code reviews
  - Extreme/Pair programming
- Methodological
  - Design patterns
  - Test-driven development
  - Version control
  - Bug tracking
- Technological
  - Static analysis (“lint” tools, FindBugs, ...)
  - Fuzzers
- Mathematical
  - Sound type systems
  - “Formal” verification



Less formal: Techniques may miss problems in programs

All of these methods should be used!

Even the most formal can still have holes:

- did you prove the right thing?
- do your assumptions match reality?

More formal: eliminate *with certainty* as many problems as possible.

# Verification

- In the 1970s, scaled to about tens of LOC
- Now, research projects scale to real software:
  - **CompCert**: verified C compiler
  - **seL4**: verified microkernel OS
  - **Ynot**: verified DBMS, web services
  - **NetCore**: software-defined network controller
- In another 40 years?

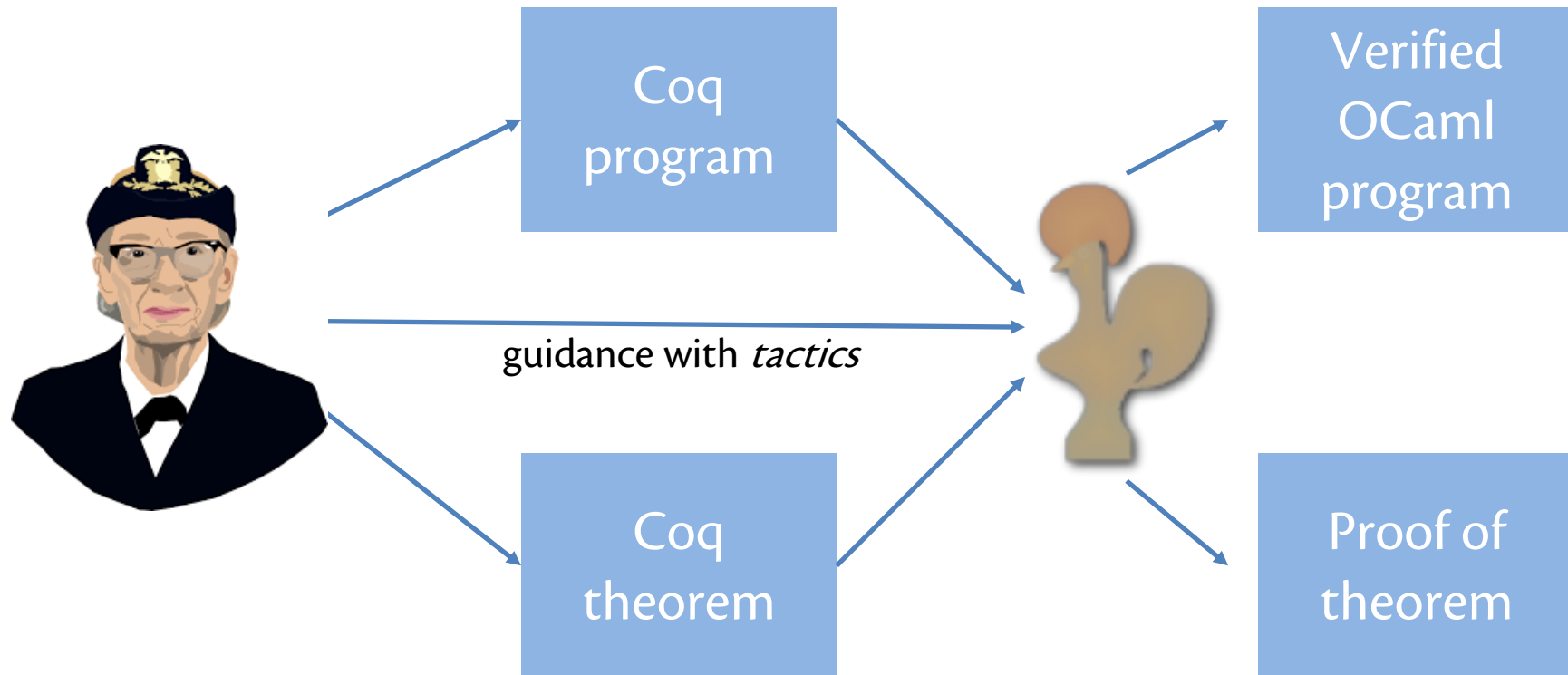
# Coq

- 1984: Coquand and Huet implement Coq based on *calculus of inductive constructions*
- 1992: Coq ported to Caml
- Now implemented in OCaml

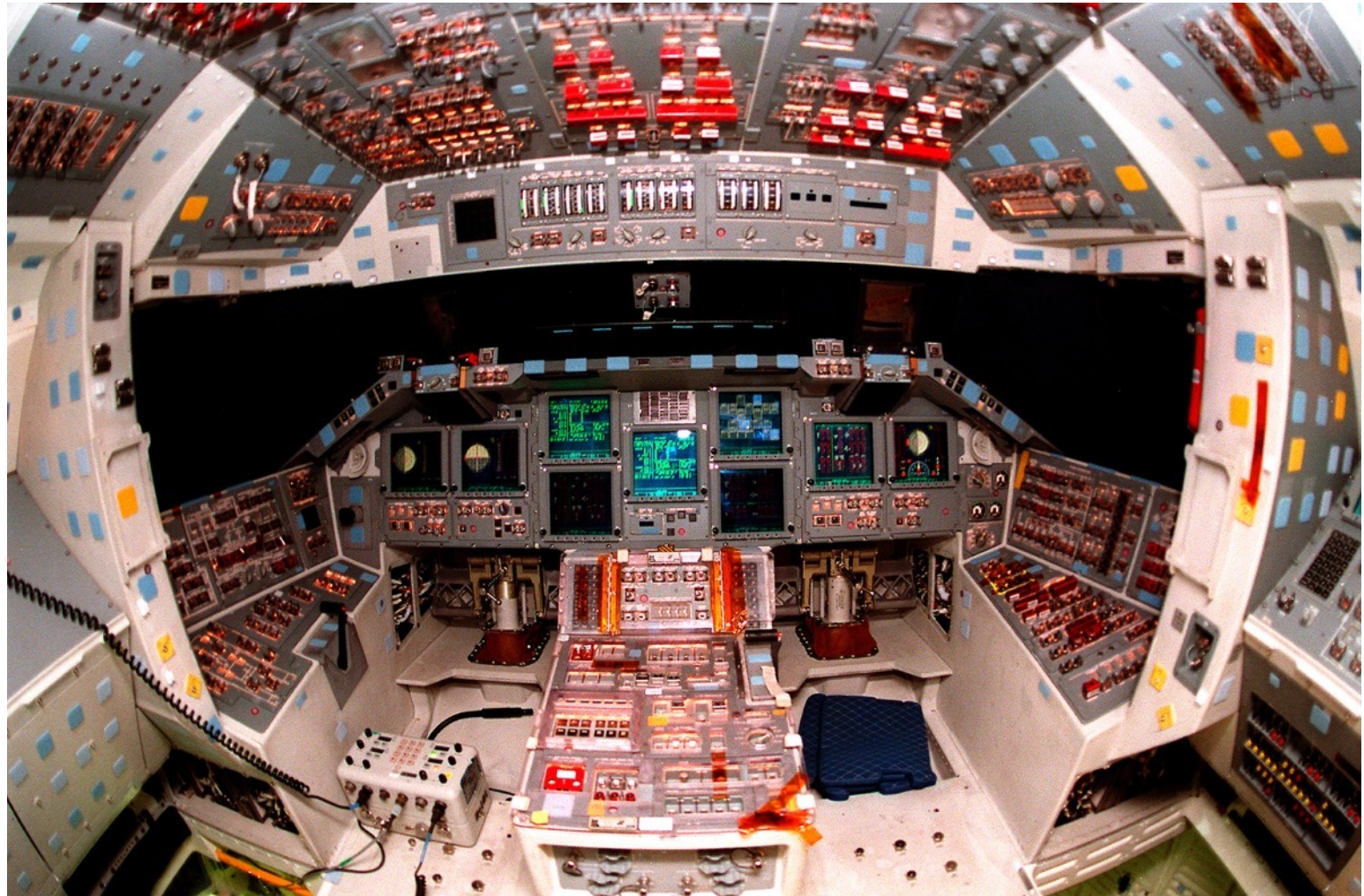


Thierry Coquand  
1961 –

# Coq for program verification



# Coq's full system



# Subset of Coq we'll use





# Our goals

- Write **basic functional programs** in Coq
  - no side effects, mutability, I/O
- Prove **simple theorems** in Coq
  - CS 3110 programs: lists, options, trees
  - CS 2800 mathematics: induction, logic
- **Non goal:** full verification of large programs
- Rather:
  - help you understand what verification involves
  - expose you to the future of functional programming
  - solidify concepts about proof and induction by developing machine-checked proofs

Definitions and Functions

Lists

# **FUNCTIONAL PROGRAMMING IN COQ**

Demo

# INDUCTION

# Structure of inductive proof

**Theorem:**

for all natural numbers  $n$ ,  $P(n)$ .

**Proof:** by induction on  $n$

**Case:**  $n = 0$

**Show:**  $P(0)$

**Case:**  $n = k+1$

**IH:**  $P(k)$

**Show:**  $P(k+1)$

**QED**

# Sum to n

```
let rec sum_to n =  
  if n=0 then 0  
  else n + sum_to (n-1)
```

$$\sum_{i=0}^n i$$

**Theorem:**

for all natural numbers n,  
sum\_to n = n \* (n+1) / 2.

**Proof:** by induction on n

**Discussion:** What is P? Base case? Inductive case? Inductive hypothesis?

# Proof

```
let rec sum_to n =  
  if n=0 then 0  
  else n + sum_to (n-1)
```

$$P(n) \equiv (\text{sum\_to } n = n * (n+1) / 2)$$

**Case:**  $n = 0$

**Show:**

$$P(0)$$

**Case:**  $n = k+1$

**IH:**  $P(k) \equiv \text{sum\_to } k = k * (k+1) / 2$

**Show:**

$$P(k+1)$$

**QED**

# INDUCTION ON LISTS

# Structure of inductive proof

**Theorem:**

for all natural numbers  $n$ ,  $P(n)$ .

**Proof:** by induction on  $n$

**Case:**  $n = 0$

**Show:**  $P(0)$

**Case:**  $n = k+1$

**IH:**  $P(k)$

**Show:**  $P(k+1)$

**QED**



# Structure of inductive proof

**Theorem:**

for all `lists lst`,  $P(\text{lst})$ .

**Proof:** by induction on `lst`

**Case:** `lst = []`

**Show:**  $P([])$

**Case:** `lst = h::t`

**IH:**  $P(t)$

**Show:**  $P(h::t)$

**QED**

# Append nil

```
let rec (@) lst1 lst2 =  
  match lst1 with  
  | []      -> lst2  
  | h::t    -> h :: (t @ lst2)
```

## Theorem:

for all lists  $lst$ ,  $lst @ [] = lst$ .

**Proof:** by induction on  $lst$

**Discussion:** What is P? Base case? Inductive case? Inductive hypothesis?

```

let rec (@) lst1 lst2 =
  match lst1 with
  | [] -> lst2
  | h::t -> h :: (t @ lst2)

```

## Base case

$$P(\text{lst}) \equiv \text{lst} @ [] = \text{lst}$$

**Case:**  $\text{lst} = []$

**Show:**

$$P([])$$

**Case:**  $\text{lst} = h::t$

**IH:**  $P(t) \equiv t @ [] = t$

**Show:**

$$P(h::t)$$

**QED**

# INDUCTION ON LISTS IN COQ

# **PROPOSITIONAL LOGIC**

# Logical connectives

- Implication:  $p \rightarrow p$
- Conjunction:  $p \wedge p$
- Disjunction:  $p \vee p$
- Negation:  $\sim p$

# Implication

Print `p_implies_p`.

*`p_implies_p =`*

*`fun (P : Prop) (P_assumed : P) => P_assumed`*  
*`: forall P : Prop, P -> P`*

`p_implies_p`  
is a function

first input is a  
proposition

second input is  
proof of first input

output is that proof

Coq proofs  
are  
functional programs



# Upcoming events

- [Last night] A8 out!
- [Today] Foster Office Hours @ 1:15pm

*This is formal.*

**THIS IS 3110**