

Expressions

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Today's music: Express Yourself (N.W.A.)

Recap

Last time:

- What is a functional language?
- Why study programming in a functional language

Today:

- Five aspects of a language
- Expressions, values, definitions

Question

Did you bring an iClicker today?

- A. Yes
- B. No
- C. I plead the 5th

No worries: Attendance point tracking starts in lecture on Tuesday; in section, on Monday

Five aspects of learning a PL

- 1. Syntax: How do you write language constructs?
- 2. Semantics: What do programs mean? (Type checking, evaluation rules)
- 3. Idioms: What are typical patterns for using language features to express your computation?
- **4. Libraries**: What facilities does the language (or a third-party project) provide as "standard"? (E.g., file access, data structures)
- **Tools**: What do language implementations provide to make your job easier? (E.g., top-level, debugger, GUI editor, ...)
- All are essential for good programmers to understand
- Breaking a new PL down into these pieces makes it easier to learn

Our Focus

We focus on semantics and idioms for OCaml

- Semantics is like a meta-tool: it will help you learn languages
- Idioms will make you a better programmer in those languages

Libraries and tools are a secondary focus: throughout your career you'll learn new ones on the job every year

Syntax is almost always boring

A fact to learn, like "Cornell was founded in 1865"

People obsess over subjective preferences {yawn}

Class rule: We don't complain about syntax

Phil Wadler



"In any language design, the total time spent discussing a feature in this list is proportional to two raised to the power of its position:

- 0. Semantics
- 1. Syntax
- 2. Lexical syntax
- 3. Lexical syntax of comments"

b. 1956

- CS Professor at Edinburgh
- One of the designers of the Haskell language
- The external examiner at my PhD defense @

Expressions

- Primary building block of OCaml programs
- Akin to statements or commands in imperative languages

Examples:

```
-21 + 21
```

- true
- -3.14159
- "Hello" ^ " " orld"

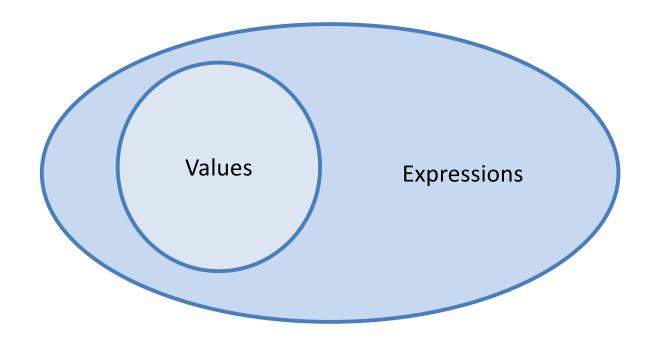
Expressions

Every kind of expression has:

- Syntax
- Semantics:
 - Type-checking rules (static semantics): produce a type or fail with an error message
 - Evaluation rules (dynamic semantics): produce a value
 - (or exception or infinite loop)
 - Used only on expressions that type-check

Values

A **value** is an expression that does not need any further evaluation



IF EXPRESSIONS

if expressions

Syntax:

if el

Write ==> to indicate evaluation Pronounce as "evaluates to"

Evaluation:

- if e1 evaluate to true, and if e2 evaluates to v,
 then if e1 then e2 else e3 evaluates to v
- if **e1** evaluates to **false**, and if **e3** evaluates to **v**, then **if e1 then e2**

Write colon to indicate type of expression Pronounce colon as "has type"

Type checking:

if e1 has type bool and e2 has bet and e3 has type t then if e1 then e2 else e3 has type t

if expressions

Syntax:

if e1 then e2 else e3

Evaluation:

- if e1 ==> true and e2 ==> v,
 then if e1 then e2 else e3 ==> v
- if e1 ==> false and e3 ==> v,
 then if e1 then e2 else e3 ==> v

Type checking:

```
if e1 : bool and e2 : t and e3 : t
then if e1 then e2 else e3 : t
```

if expressions

Syntax:

if e1 then e2 else e3

Evaluation:

```
• if e1 ==> true and e2 ==> v,
then (if e1 then e2 else e3) ==> v
```

```
• if e1 ==> false and e3 ==> v,
then (if e1 then e2 else e3) ==> v
```

Type checking:

```
if e1 : bool and e2 : t and e3 : t
then (if e1 then e2 else e3) : t
```

Type inference and annotation

- OCaml compiler infers types
 - Compilation fails with type error if it can't
 - Hard part of language design: guaranteeing compiler can infer types when program is correctly written

- You can manually annotate types anywhere
 - Replace e with (e : t)
 - Useful for resolving type errors

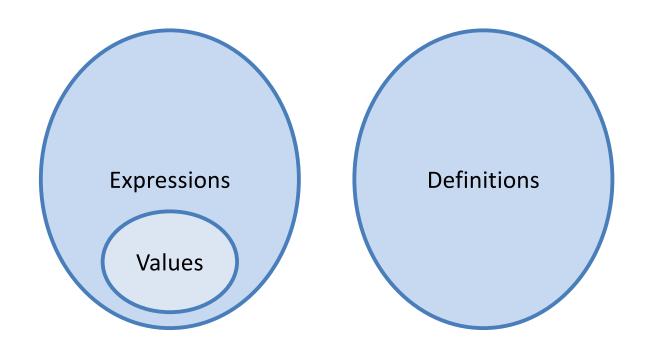
LET DEFINITIONS

Definitions

A **definition** gives a name to a value

Definitions are not expressions, or vice-versa

But definitions syntactically contain expressions



let definitions

Syntax:

let x = ewhere x is an *identifier*

Evaluation:

- Evaluate e to a value v
- Bind v to x: henceforth, x will evaluate to v
 (under the hood: there is a memory location named x that contains v)
- But the definition does not evaluate to a value

LET EXPRESSIONS

let expressions

Syntax:

```
let x = e1 in e2
```

x is an *identifier*

e1 is the binding expression

e2 is the body expression

let x = e1 in e2 is itself an expression

let expressions

let x = e1 in e2

Evaluation:

- Evaluate e1 to a value v1
- Substitute v1 for x in e2, yielding a new expression e2'
- Evaluate e2' to v2
- Result of evaluation is v2

let expressions

```
let x = e1 in e2
```

Type-checking:

```
If e1:t1 and x:t1 and e2:t2
then (let x = e1 in e2) : t2
```

This type-checking rule was stated incorrectly during lecture; it has been fixed

VARIABLE EXPRESSIONS

Variable expressions

How to evaluate just

X

?

let definitions in toplevel

```
let x = e
```

is implicitly, "in rest of what you type"

E.g., you type:

```
let a="big";;
let b="red";;
let c=a^b;;
```

Toplevel understands as

```
let a="big" in
let b="red" in
let c=a^b in...
```

Variable expressions

How to evaluate just

X

?

Answer: substitution from that giant nested **let** expression

Upcoming events

- [Today] Consulting hours in Rhodes 590 from 4:30-9:30pm
- [Tomorrow] Standby list closes
- [Monday/Tuesday] Discussion sections start
- [Tuesday] A0 released

This is expressive.

THIS IS 3110

WHAT ABOUT IMMUTABILITY?

Seems like variable can mutate...

```
let x = 1;;
let x = 2;;
x;;
```

But really it's just nested scopes

$$let x = 1 in$$

let x = 2 in

X

Allocate memory that will always be 1

Allocate memory that will always be 2

Which piece of memory does name mean? Innermost scope, as you would expect.