



CS 3110

More Variants

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Review

Previously in 3110:

- User-defined data types: records, variants
- Built-in type constructors: **list**, **option**

Today:

- * (i.e., Tuple Types)
- Type synonyms
- More about variants
- Exceptions

Tuples

- Several pieces of data glued together
- A **tuple** contains several **components**
- (Don't have to define tuple type before use)

e.g.,

- `(1, 2, 10)`
- `1, 2, 10`
- `(true, "Hello")`
- `([1; 2; 3], (0.5, 'X'))`

Tuple types

- `(1, 2, 10) : int*int*int`
- `1, 2, 10 : int*int*int`
- `(true, "Hello") : bool*string`
- `([1; 2; 3], (0.5, 'X'))
: int list * (float*char)`

Tuples

- 2-tuple: pair
- 3-tuple: triple
- beyond that: maybe better to use records

We need language constructs to *build* tuples and to *access* the components

- Building is easy: just write the tuple, as before
- Accessing uses pattern matching...

Accessing tuples

New kind of pattern, the **tuple pattern**: (p_1, \dots, p_n)

```
match (1,2,3) with  
| (x,y,z) -> x+y+z
```

```
(* ==> 6 *)
```

```
let thrd t =  
  match t with  
  | (x,y,z) -> z
```

```
(* thrd : 'a*'b*'c -> 'c *)
```

Note: we never needed more than one branch in the match expression...

Pattern matching without match

```
(* OK *)  
let thrd t =  
  match t with  
  | (x,y,z) -> z
```

```
(* good *)  
let thrd t =  
  let (x,y,z) = t in z
```

```
(* better *)  
let thrd t =  
  let (_,_,z) = t in z
```

```
(* best *)  
let thrd (_,_,z) = z
```

Variants vs. records vs. tuples

	Define	Build	Access
Variant	type	Constructor name	Pattern matching
Record	type	Record expression with {...}	Pattern matching OR field selection with dot operator .
Tuple	N/A	Tuple expression with (...)	Pattern matching OR fst or snd

- Variants: **one-of types** *aka* **sum types**
- Records, tuples: **each-of types** *aka* **product types**

Question

Which of the following would be better represented with records rather than variants?

- A. *Coins*, which can be pennies, nickels, dimes, or quarters
- B. *Students*, who have names and id numbers
- C. *A dessert*, which has a sauce, a creamy component, and a crunchy component
- D. A and C
- E. B and C

Question

Which of the following would be better represented with records rather than variants?

- A. *Coins*, which can be pennies, nickels, dimes, or quarters
- B. *Students*, who have names **and** NetIDs
- C. *A dessert*, which has a sauce, a creamy component, **and** a crunchy component
- D. A and C
- E. **B and C**

TYPE SYNONYMS

Type synonyms

Syntax: `type id = t`

- Anywhere you write `t`, you can also write `id`
- The two names are *synonymous*

e.g.

```
type point    = float * float
```

```
type vector  = float list
```

```
type matrix  = float list list
```

Type synonyms

```
type point = float * float
```

```
let getx : point -> float =  
  fun (x,_) -> x
```

```
let pt : point = (1.,2.)
```

```
let floatpair : float*float = (1.,3.)
```

```
let one = getx pt
```

```
let one' = getx floatpair
```

VARIANTS

Recall: Variants

```
type day = Sun | Mon | Tue | Wed  
         | Thu | Fri | Sat
```

```
type ptype = TNormal | TFire | TWater
```

```
type peff = ENormal | ENotVery | Esuper
```

So far, just enumerated sets of values

But they can do much more...

Variants that carry data

```
type shape =  
  | Point of point  
  | Circle of point * float (* center and radius *)  
  | Rect of point * point (* lower-left and  
                                upper-right corners *)
```

```
let area = function  
  | Point _ -> 0.0  
  | Circle (_,r) -> pi *. (r ** 2.0)  
  | Rect ((x1,y1),(x2,y2)) ->  
    let w = x2 -. x1 in  
    let h = y2 -. y1 in  
    w *. h
```


Variants that carry data

```
type shape =  
  | Point of point  
  | Circle of point * float  
  | Rect of point * point
```

```
let center = function  
  | Point p -> p  
  | Circle (p, _) -> p  
  | Rect ((x1, y1), (x2, y2)) ->  
    ((x2 -. x1) /. 2.0,  
     (y2 -. y1) /. 2.0)
```

Variants that carry data

```
type shape =  
  | Point of point  
  | Circle of point * float  
  | Rect of point * point
```

Every value of type **shape** is made from exactly one of the constructors and contains:

- a *tag* for which constructor it is from
- the data *carried* by that constructor

Called an **algebraic data type** because it contains product and sum types, aka **tagged union**

Tagged union

- *Union* because the set of all values of the type is the union of the set of all values of the individual constructors

```
type t = String of string  
        | Int of int
```

- *Tagged* because possible to determine which underlying set a value came from

```
type t = Left of int  
        | Right of int
```

- "All for one and one for all":
 - all values of variant, regardless of constructor, have same type
 - any one value of variant built with exactly one constructor, all of which are specified in type definition



Variant types

Type definition syntax:

```
type t = C1 [of t1] | ... | Cn [of tn]
```

A constructor that carries data is *non-constant*

A constructor without data is *constant*

Semantics are straightforward; see notes

Question

Given our shape variant, which function would determine whether a shape is a circle centered at the origin?

```
type shape =  
  | Point of point  
  | Circle of point * float  
  | Rect of point * point
```

Possible answers on next slide...

```
type shape = Point of point | Circle of point * float | Rect of point * point
```

A

```
let cato = function  
| Point _ -> false  
| Circle _ -> true  
| Rect _ -> false
```

B

```
let cato = function  
| Point _ | Rect _ -> false  
| Circle (0.,0.), r -> true
```

E

```
let cato c =  
  c = Circle ((0.,0.),_)
```

```
let cato = function  
| Point p -> false  
| Circle (0.,0.) -> true  
| Rect (ll,ur) -> false
```

C

```
let cato = function  
| Circle ((0.,0.), _) -> true  
| _ -> false
```

D

```
type shape = Point of point | Circle of point * float | Rect of point * point
```

A

```
let cato = function  
| Point -> false  
| Circle -> true  
| Rect -> false
```

B

```
let cato = function  
| Point _ | Rect _ -> false  
| Circle (0.,0.), r -> true
```

E

```
let cato c =  
  c = Circle ((0.,0.),_)
```

```
let cato = function  
| Point p -> false  
| Circle (0.,0.) -> true  
| Rect (ll,ur) -> false
```

C

```
let cato = function  
| Circle ((0.,0.), _) -> true  
| _ -> false
```

D

RECURSIVE VARIANTS

Implement lists with variants

```
type intlist = Nil | Cons of int * intlist
```

```
let emp = Nil
```

```
let l3 = Cons (3, Nil) (* 3::[] or [3]*)
```

```
let l123 = Cons(1, Cons(2, l3)) (* [1;2;3] *)
```

```
let rec sum (l:intlist) =
```

```
  match l with
```

```
  | Nil -> 0
```

```
  | Cons(h,t) -> h + sum t
```

Implement lists with variants

```
let rec length = function  
  | Nil -> 0  
  | Cons (_, t) -> 1 + length t  
(* length : intlist -> int *)
```

```
let empty = function  
  | Nil -> true  
  | Cons _ -> false  
(* empty: intlist -> bool *)
```

PARAMETERIZED VARIANTS

Lists of any type

- **Have:** lists of ints
- **Want:** lists of ints, lists of strings, lists of pairs, lists of records that themselves contain lists of pairs, ...

Non-solution: copy code

```
type stringlist = SNil | SCons of string * stringlist
let empty = function
  | SNil -> true
  | SCons _ -> false
```

Lists of any type

Solution: **parameterize** types on other types

```
type 'a mylist = Nil | Cons of 'a * 'a mylist
```

```
let l3 = Cons (3, Nil) (* [3] *)
```

```
let lhi = Cons ("hi", Nil) (* ["hi"] *)
```

Lists of any type

```
type 'a mylist =  
  | Nil  
  | Cons of 'a * 'a mylist
```

mylist is not a type but a **type constructor**: takes a type as input and returns a type

- **int** mylist
- **string** mylist
- (**int*string**) mylist
- ...

Functions on parameterized variants

```
let rec length = function  
  | Nil -> 0  
  | Cons (_, t) -> 1 + length t  
(* length : 'a mylist -> int *)
```

```
let empty = function  
  | Nil -> true  
  | Cons _ -> false  
(* empty: 'a mylist -> bool *)
```

code stays the same; only the types change

Parametric polymorphism

- *poly* = many, *morph* = form
- write function that works for many arguments regardless of their type
- closely related to Java generics, related to C++ template instantiation, ...

THE POWER OF VARIANTS

Lists are just variants

OCaml effectively codes up lists as variants:

```
type 'a list = [] | :: of 'a * 'a list
```

- **list** is a type constructor parameterized on type variable **'a**
- **[]** and **::** are constructors
- Just a bit of syntactic magic in the compiler to use **[]** and **::** instead of alphabetic identifiers

Options are just variants

OCaml effectively codes up options as variants:

```
type 'a option = None | Some of 'a
```

- **option** is a type constructor parameterized on type variable 'a
- **None** and **Some** are constructors

Exceptions are (mostly) just variants

OCaml effectively codes up exceptions as slightly strange variants:

```
type exn
```

```
exception MyNewException of string
```

- Type **exn** is an *extensible* variant that may have new constructors added after its original definition
- Raise exceptions with **raise** **e**, where **e** is a value of type **exn**
- Handle exceptions with pattern matching, just like you would process any variant

Upcoming events

- [by Fri morning] A1 out
- [Feb 12 @ 5:30] WICC/ACSU/URMC G-Body on CIS Culture
- [Feb 12 @ 6:30] WICC Partner Finding Social