

Async

Prof. Clarkson Fall 2015

Today's music: It's Gonna be Me by *NSYNC

Review

Previously in 3110

- Threads
- Async (futures): deferreds, scheduler, callbacks with upon and bind

Today:

- more on bind
- other sequencing operators
- ivars

Review: Async

- A third-party library for futures in OCaml
- Deferreds: values whose completed computation has been deferred until the future (and in fact is happening already)
- Scheduler runs **callbacks** that have been registered to consume the values of deferreds
 - Scheduler selects a callback whose input has become ready to consume,
 - runs the callback with that input,
 - Only ever one callback running at a time
 - Scheduler never interrupts the callback
 - and repeats.

Review: Deferred



An 'a **Deferred.t** is like a box:

- It starts out empty
- At some point in the future, it could be filled with a value of type 'a
- Once it's filled, the box's contents can never be changed ("write once")

Terminology:

- "box is filled" = "deferred is determined"
- "box is empty" = "deferred is undetermined"

Review: Registering a callback

upon:

```
'a Deferred.t
-> ('a -> unit)
-> unit
```

- use to register a callback (the function of type 'a -> unit) to run sometime after deferred is determined
- upon returns immediately with () no matter what
- sometime after box is filled (if ever), scheduler runs callback on contents of box
- callback produces () as return value, but never returned to anywhere

Suppose you create a deferred with **return 42**. When is that deferred determined?

- A. Immediately
- B. At some point in the future, but you don't know when.
- C. After the creator's callback returns control to the scheduler.
- D. Never
- E. None of the above

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A. Immediately

- B. At some point in the future, but you don't know when.
- C. After the creator's callback returns control to the scheduler.
- D. Never
- E. None of the above

Review: Creating deferreds

return : 'a -> 'a Deferred.t

- use to create a deferred that is already determined

after : Core.Std.Time.Span.t

- -> unit Deferred.t
- use to create a deferred that becomes determined sometime after a given length of time
- Core.Std.Time.Span.of_int_sec 10
 represents 10 seconds

BIND

Bind

```
bind :
    'a Deferred.t
    -> ('a -> 'b Deferred.t)
    -> 'b Deferred.t
```

- use to register a deferred computation after an existing one
- takes two inputs: a deferred d, and callback c
- bind d c immediately returns with a new deferred d'
- sometime after d is determined (if ever), scheduler runs c on contents of d
- c produces a new deferred, which if it ever becomes
 determined, also causes d' to be determined with same value

Bind

```
Deferred.bind
  (return 42)
  (fun n -> return (n+1))
```

- first argument is a deferred that is determined with value 42
- second argument is a callback that takes an integer n and returns a
 deferred that is determined with value n+1
- bind immediately returns with an undetermined deferred ud
- scheduler, when it next gets to run, can notice that first argument is determined, and run callback
- callback gets 42 out of box, binds it to n, and returns a new deferred that is determined with value 43
- scheduler can notice that output of callback has become determined, and make ud determined with same value

>>=

```
(>>=)

    infix operator version of bind

  - bind d c is the same as d >>= c
Deferred.bind
  (return 42)
  (fun n -> return (n+1))
(* equiv. *)
return 42 >>= fun n ->
return (n+1)
```

>>=

```
open Async.Std
let sec n = Core.Std.Time.Span.of int sec n
let return after v delay =
  after (sec delay) >>= fun () ->
  return v
let =
  (return after "First timer elapsed\n" 5) >>= fun s ->
 print string s;
  (return after "Second timer elapsed\n" 3) >>= fun s ->
 print string s;
  exit 0
let = print string "Hello\n"
let = Scheduler.go ()
```

```
let _ =
  (return_after "First timer elapsed\n" 5) >>= fun s ->
  print_string s;
  (return_after "Second timer elapsed\n" 3) >>= fun s ->
  print_string s;
  exit 0
let _ = print_string "Hello\n"
```

Which string will be printed first?

- A. "First timer elapsed"
- B. "Second timer elapsed"
- C. "Hello"

```
let _ =
  (return_after "First timer elapsed\n" 5) >>= fun s ->
  print_string s;
  (return_after "Second timer elapsed\n" 3) >>= fun s ->
  print_string s;
  exit 0
let _ = print_string "Hello\n"
```

Which string will be printed first?

- A. "First timer elapsed"
- B. "Second timer elapsed"
- C. "Hello"

```
let _ =
  (return_after "First timer elapsed\n" 5) >>= fun s ->
  print_string s;
  (return_after "Second timer elapsed\n" 3) >>= fun s ->
  print_string s;
  exit 0
let _ = print_string "Hello\n"
```

Which string will be printed second?

- A. "First timer elapsed"
- B. "Second timer elapsed"
- C. "Hello"

```
let _ =
  (return_after "First timer elapsed\n" 5) >>= fun s ->
  print_string s;
  (return_after "Second timer elapsed\n" 3) >>= fun s ->
  print_string s;
  exit 0
let _ = print_string "Hello\n"
```

Which string will be printed second?

- A. "First timer elapsed"
- B. "Second timer elapsed"
- C. "Hello"

Concurrently

```
let t1 =
  return after "First timer elapsed\n" 5 >>= fun s ->
  print string s;
  return ()
let t2 =
  return_after "Second timer elapsed\n" 3 >>= fun s ->
  print string s;
  return ()
let =
  t1 >>= fun () ->
  t2 >>= fun () ->
  exit 0
```

Now the "second" timer string would be printed before the "first"

MORE SEQUENCING OPERATORS

Map

```
map :
    'a Deferred.t
    -> ('a -> 'b)
    -> 'b Deferred.t
```

- takes two inputs: a deferred **d**, and a function **f**
- map d f immediately returns with a new deferred d'
- sometime after d is determined (if ever), scheduler runs f on contents of d, immediately yielding a new value b, and d' is immediately determined with that value
- has its own infix operator (>>|)

Map

```
let return_after v delay =
   after (sec delay) >>= fun () ->
   return v

let return_after' v delay =
   after (sec delay)
   >>| fun () -> v
```

...how might you implement map?

Map

```
let map (d: 'a Deferred.t)
  (f: 'a -> 'b) : 'b Deferred.t
=

d >>= fun a ->
  return (f a)
```

Both

```
both :
    'a Deferred.t
    -> 'b Deferred.t
    -> ('a*'b) Deferred.t
```

- takes two inputs: a deferred **d1**, and a deferred **d2**
- both d1 d2 immediately returns with a new deferred d
- sometime after both d1 and d2 are determined (if ever), d is determined with the pair of values from inside d1 and d2

...how might you implement **both**?

Both

```
let both
  (d1: 'a Deferred.t)
  (d2: 'b Deferred.t)
  : ('a*'b) Deferred.t
  d1 >>= fun a ->
  d2 >>= fun b ->
  return (a,b)
```

Does this implementation force the contents of d1 to be computed before the contents of d2?

```
let both d1 d2 =
  d1 >>= fun a ->
  d2 >>= fun b ->
  return (a,b)
```

- A. Yes
- B. No

Does this implementation force the contents of d1 to be computed before the contents of d2?

```
let both d1 d2 =
  d1 >>= fun a ->
  d2 >>= fun b ->
  return (a,b)
```

A. Yes

B. No

Either

```
either :
    'a Deferred.t
    -> 'a Deferred.t
    -> 'a Deferred.t
```

- takes two inputs: a deferred d1, and a deferred d2
- either d1 d2 immediately returns with a new deferred d
- sometime after at least one of d1 and d2 is determined (if ever), d is determined with the same value
- no guarantee about timing of d1 vs d2: maybe d1 becomes determined first with value v1, then d2 with v2, then d with d2

...how might you implement **either**?

Either

```
let either
  (d1: 'a Deferred.t)
  (d2: 'a Deferred.t)
  : 'a Deferred.t
=
  failwith "You can't"
```

IVARS

Ivar



An 'a Ivar.t is like a box:

- It starts out empty
- At some point in the future, it could be filled with a value of type 'a
- Once it's filled, the box's contents can never be changed ("write once")
- You can fill the box

Ivar

- create : unit -> 'a Ivar.t
- is full : 'a Ivar.t -> bool
- fill : 'a Ivar.t -> 'a -> unit
 - Attempting to fill an already full ivar raises an exception
 - That's where the name comes from...

Digression on Cornell history

- i = incremental
- Originally [Arvind and Thomas 1981] Istructures were a kind of data structure for functional arrays in which each element could be assigned exactly once—hence the array was constructed incrementally
- Used for parallel computing in language called Id [Arvind, Nikhil, and Pingali 1986]
 - Keshav Pingali, Cornell CS prof 1986-2006?
- Implemented in *Concurrent ML* by John Reppy (Cornell PhD 1992)





Ivar

- create : unit -> 'a Ivar.t
- is_full : 'a Ivar.t -> bool
- fill : 'a Ivar.t -> 'a -> unit
 - Attempting to fill an already full ivar raises an exception
 - That's where the name comes from

...but how can you get a value out of the ivar?

Ivar

read : 'a Ivar.t -> 'a Deferred.t

- read i immediately returns a deferred that becomes determined after i is filled
- and to get a value out of that deferred, use any of the ways we've seen of registering callbacks

now we can implement **either**...

Either

```
let either d1 d2 =
  let result = Ivar.create () in
  let fill = fun x \rightarrow
    if Ivar.is empty result
    then Ivar.fill result x
    else () in
  upon d1 fill;
  upon d2 fill;
  Ivar.read result
```

Upcoming events

- [Thursday] A4 due
- [next Monday] project charter due

This is in sync.

THIS IS 3110