

Announcements:

- Prelim #1 Tuesday
 - Conflict exam: 5:45-7:15
 - Only for people with conflicts
 - Main exam: 7:30-9:00 in Thurston Hall 203/205
 - Graded (late) at night, back in section
- Practice prelim now on CMS, no answers
- Prelim review session: Sat/Sun afternoon or evening

Minimal correct induction proof

Example problem you might see on a prelim:

Recall that for any natural number n , we define $n!$ as $n(n-1)(n-2)\dots$, where $0! = 1$. Write a recursive definition `fact n` that computes $n!$, and prove your definition is correct using induction and the substitution model.

Solution:

```
let rec fact(n) = if n=0 then 1 else n*fact(n-1)
```

* Statement $P[n]$: the value of the OCaml expression `fact(n)` is $n!$

* Variable we are doing induction on: n , starting at 0

* Base case: we prove $P[0]$ as follows

`fact(0)`

b.s.m. (substitute) is

`if 0=0 then 1 else 0*fact(0-1)`

b.s.m. (primitives) is

`if true then 1 else 0*fact(0-1)`

b.s.m. (if) is

1

So the value of the expression `fact(0)` is 1 which is $0!$

* Induction step:

Pick an $n \geq 0$ and assume $P[n]$, then prove $P[n+1]$

`fact(n+1)`

b.s.m. (substitute) is

`if n+1=0 then 1 else n+1*fact(n+1-1)`

Since $n \geq 0$ the value of the expression `n+1=0` is false

b.s.m. (if) is

`n+1*fact(n+1-1)`

b.s.m. (primitives) is

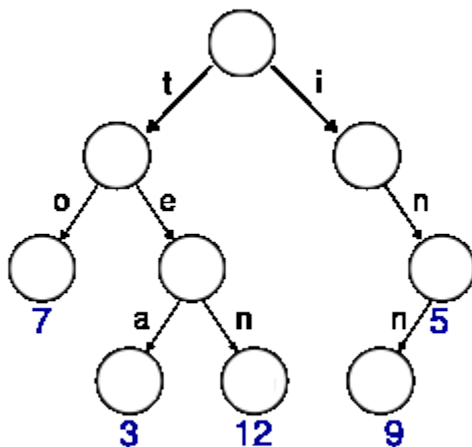
`n+1*fact(n)`

By the induction hypothesis $P[n]$ the value of `fact(n)` is $n!$ so this is

`n+1*n!`

which is $n+1!$

- You've seen binary trees in CS2110
- Let's look at a data structure called a "trie"
- A trie is a "finite map", like a dictionary. It maps keys to values. Typically for a trie the keys are strings and the values are numbers.
- A trie is sometimes called a "prefix tree". The basic idea is that a path through the tree represents a prefix, i.e. all strings that start with a particular substring.
 - Root is the empty string
- Example:



- This trie is the finite map {"to"->7, "tea"->3, "ten"->12, "in"->5, "inn"->9}
 - As you saw in CS2110, tree-like data structures of this form are very efficient when they are balanced
 - Note that a trie doesn't need to be binary, though this one is
 - In fact, 26 children or so (capitalization, punctuation)

- A trie is very efficient when there are lots of shared prefixes
 - Occurs in many situations (letters, genes, IP addresses)
- Lookup operation is obvious. Insert and delete are surprisingly similar. Everything takes time $O(L)$, which is the length of the longest entry.
- This is a huge advantage of a trie. Most data structures have very asymmetric costs for lookup/insert/delete, so you need to pick the right one for your application carefully.
- Also note that if you don't find what you are looking for you know something close to it. Useful for, e.g., spell checking.
 - Thought question: Google instant search?
- Important variant: radix tree (aka Patricia trie), where we ensure that every internal node has 2 or more children by merging nodes with 1 child
- Sub-variant: store at the end "black" or "white". Then you can use this to encode strings that are present and also strings that are absent. Application is for IP routing tables.

- We will go over the trie signature in section.
- An important idea, both in the trie and point example, is what is called a REP INVARIANT. This is a property of the representation that must be satisfied for the representation to be valid. For example, in our radix tree example, a node must have 2 or more children, and never 1 (could be 0 if it's a leaf).
- You will typically want to implement this with a function repOK that returns its argument or raises an exception.
- Check this on all inputs and on output.
 - This sanity check seems wasteful, and you can turn it off in production code (for example by making repOK into the identity function).
 - But it will catch a ton of subtle bugs
- Example: lists without duplicates, or in sorted order
 - In a certain sense these are types, but they can't be checked at compile time.
 - Another example: even numbers, or prime numbers, or even natural or whole numbers

- But let's now return to the idea of designing a proper specification.
- Deceptively simple example: square root function, float->float
- Spec: beyond the types, what is true before we call sqrt (precondition)
 - What is true after (postcondition)
- What is the actual spec?
 - Positive input
 - Returns "closest" positive float whose square is x
 - Sort of...
- What if the spec is violated?
 - Return something arbitrary? Rarely the right answer
 - Should raise an exception, in general
 - IEEE actually defines an "out of band" value, NaN
- Specs are interesting in large part for what they leave out (non-determinism)
- Refinement example: A) find(lst,x) index, versus B) smallest index
- Note that any implementation of B is also an implementation of A
- We say that B refines A, since it has more constraints