CS212 Java Practicum

Introduction to SaM

What is SaM? Why SaM?

- . From last lecture:
 - computer stores data and instructions in memory
 - fetch-and-decode cycle:
 - JVM is ____ ____ of computers
 - bytecodes are _
- *SaM*:
 - stands for:
 - see SaM on CS212 for full instruction set
 - gives us legible instruction set

 - your compiler will generate
 BTW, what's a compiler? (last panel...)

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Samcode Instructions

- . Low-level instructions:
 - push and pop values in memory
 - *mnemonics* for bit patterns
- Structure:

opcode operand

. Areas (watch DIS play w/SaM)

Structure of Samcode File

- ASCII Text! (What's ASCII?)
- Write instructions on new lines
- // indicates single-line comments, which are ignored
- Program ends with
- Program must leave how many items on Stack?

Focus on Stack

- Call Stack (and other names):
 - function calls function calls ...
 - when last function done, go back, then back, then ...
 - how to picture this structure?
- - each function's portion of Stack
 - variables, data, administrative info
- Cells and addresses
 - start at 0!
- · Helpful picture?

Useful Registers

- Frame Based Register (FBR)
 - administrative information
 - keeps track of current frame (and thus, function)
- Stack Pointer (SP)
 - uses register
- store location of next free cell in stack
- · Helpful picture?

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Some Instructions

- . ALU:
 - arithmetic, boolean, comparisongenerally follows *below op top*

 - usually pops both values and pushes result
- Stack Manipulation:
 - pushing
 - swapping, duplicating

 - storing, retrieving
- . Register
- Control
- Descriptions:
 - see on-line documentation
 - see Chapter 1

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Some Examples

- Notation:
 - Infix: (1-2) 3
 - Postfix: 12-3-
- Logical: ~(4 <= 5)
 - Samcode rem: below op top
- Samcode?

Program Storage?

- Main memory model:
 - store programs as
 - so, instructions have patterns of _
- Where are they in SaM?
 - Samcode read into an array
 - array stores instruction objects
- Want more? See documentation and source code
 - SaM→Individual Files \rightarrow Core \rightarrow Instructions
 - See next page for example
- · How to load your own instructions?
 - recompile everything (a pain)
 - or...use SaM's instruction loader

Example package edu.cornell.cs.sam.core.instructions; import edu.cornell.cs.sam.core.*; public class SAM_ADD extends SamInstruction {
 public void exec() throws SystemException {
 int type1 = mem.getType(cpu.get(SP) - 2);
 int type2 = mem.getType(cpu.get(SP) - 1);
 mem.push(higherPrecedence(type1, type2), mem.pop() + mem.pop());
 cpu.inc(PC);

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Variable Scope

- Take an aside... is SaM really useful?
- · Example:

```
- is the following legal?
int x(int x) { return x++; }
int y(int x) { return x(x); }
```

- why? why not?
- · Scope of variable:
 - region of code in which variable represents something
 - how does Java indicate?
- Local and global variables:
 - each function has its own local variables
 - global variables shared
- . Does SaM help?

Variables and Frames

- . A way to picture variables in frames...
 - variable gets cell
 - Aside: SaM shows type of cell
- · Samcode program:
 - allocate cell
 - fill cell
 - later retrieve/change contents
 - finally deallocate cell (why?)

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```
Allocation and Deallocation
• Pushing:
- PUSHIMM... (see SaM website)
· Allocating:
    - Allocate v amount of vars: ADDSP v
    – Deallocate \mathbf{v} amount of vars: ADDSP -\mathbf{v}
• Example:
   ADDSP 3
   ADDSP -1
   ADDSP -1
   ADDSP -1
    STOP
    // error mesg (why?)
                                                         13
```

How to access a variable?

- . Addressing of variables:
 - absolute
 - relative
- Absolute:
 - don't worry about your current frame
 - figure out variable address on stack
 - eg) globals
- . Relative

 - do worry about your current framefigure out variable address with respect to FBR value
 - eg) locals

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Absolute Address

```
. Instructions:
```

```
- To store a value v at location i:
```

• PUSHIMM v: Stack[SP] $\leftarrow v$; SP++
• STOREABS i: Stack[i] \leftarrow Stack[SP-1]; SP--

- To **retrieve** a value **v** from location **k**:

• PUSHABS k; Stack[SP] \leftarrow Stack[k]; SP++

. Example:

```
int rv;
int x;
int y;
x = 10;
y = 20;
rv = x + y;
return rv;
```

Relative Address

• Instructions:

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- To store a value **v** at location **i**:

• PUSHIMM v: Stack[SP] ← v; SP++
• STOREOFF i: Stack[i+FBR] ← Stack[SP-1]; SP--

- To retrieve a value **v** from location **k**:

• PUSHOFF k: Stack[SP] ← Stack[k+FBR]; SP++

• Picture?

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```
Example
ADDSP 1 // rv of program
JSR add // new frame (jump to "add")
STOREOFF 0 // store rv of "add"
STOP // done
                                                                                                                                                              public int add()
    int x, y;
    x = 10;
    y = 20;
    return x+y;
}
                               // code for "add" function
// store old FBR (0) and set new FBR (2)
// allocate space for x, y, rv of add
// rv of add is at relative address 1
// push value 10
// store 10 in x's cell
// push value 20
// store 20 in y's cell
                                 // exchange rv of add for old FBR
// restore old FBR (0)
// exchange rv of add for return address
// return to Samcode just after "JSR add"
SWAP
UNLINK
SWAP
RST
NOTE: We will use a different frame structure later!
                                                                                                                                                                                                            17
```

Human Compiling

- Compiling:
 translate code (like Java) to machine code (like Samcode)
 - compiler (like javac) does the work for you
- Human Compiling (Part 1 of CS212):
 - you identify simple expressions and statements
 - you convert them into Samcode
 - you test your Samcode problems in SaM
 - we grade your correctness and style

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