THREADS & CONCURRENCY
The CPU is the part of the computer that executes instructions.

**Java:** \( x = x + 2; \)

Suppose variable \( x \) is at Memory location 800,
Instructions at 10

**Machine language:**
10: load register 1, 800
11: Add register 1, 2
12: Store register 1, 800

Basic uniprocessor-CPU computer. Black lines indicate data flow, red lines indicate control flow

From wikipedia
>100 processes are competing for time. Here’s some of them:

<table>
<thead>
<tr>
<th>Process Name</th>
<th>% CPU</th>
<th>CPU Time</th>
<th>Threads</th>
</tr>
</thead>
<tbody>
<tr>
<td>Grab</td>
<td>4.1</td>
<td>3.33</td>
<td>7</td>
</tr>
<tr>
<td>ReportCrash</td>
<td>2.3</td>
<td>0.78</td>
<td>6</td>
</tr>
<tr>
<td>Eclipse</td>
<td>1.5</td>
<td>1:48:30.07</td>
<td>54</td>
</tr>
<tr>
<td>SuperTab</td>
<td>1.4</td>
<td>1:40:44.59</td>
<td>5</td>
</tr>
<tr>
<td>Activity Monitor</td>
<td>1.4</td>
<td>10.57</td>
<td>10</td>
</tr>
<tr>
<td><a href="https://www.wunderground.com">https://www.wunderground.com</a>...</td>
<td>1.1</td>
<td>1:34.19</td>
<td>23</td>
</tr>
<tr>
<td>Creative Cloud</td>
<td>0.8</td>
<td>58:32.81</td>
<td>27</td>
</tr>
<tr>
<td>Microsoft PowerPoint</td>
<td>0.6</td>
<td>3:24.02</td>
<td>9</td>
</tr>
<tr>
<td>Safari Networking</td>
<td>0.4</td>
<td>26:53.25</td>
<td>10</td>
</tr>
<tr>
<td>loginwindow</td>
<td>0.3</td>
<td>16:14.79</td>
<td>4</td>
</tr>
<tr>
<td>Google Drive</td>
<td>0.3</td>
<td>6.33</td>
<td>22</td>
</tr>
<tr>
<td>Safari</td>
<td>0.3</td>
<td>50:09.48</td>
<td>24</td>
</tr>
</tbody>
</table>
Clock rate

- Clock rate “frequency at which CPU is running”
  - Higher the clock rate, the faster instructions are executed.
- First CPUs: 5-10 Hz (cycles per second)
- Today MacBook Pro 3.5 GHz
- Your OS can control the clock rate, slow it down when idle, speed up when more work to do
Why multicore?

- Moore’s Law: Computer speeds and memory densities nearly double each year
But a fast computer runs hot

- Power dissipation rises as square of the clock rate
- Chips were heading toward melting down!
- Put more CPUs on a chip:
  with four CPUs on one chip, even if we run each at half speed we can perform more overall computations!
Today: Not one CPU but many

Processing Unit is called a core.

- Modern computers have “multiple cores” (processing units)
  - Instead of a single CPU (central processing unit) on the chip
    5-10 common. Intel has prototypes with 80!

- We often run many programs at the same time

- Even with a single core (processing unit), your program may
  have more than one thing “to do” at a time
  - Argues for having a way to do many things at once
Many programs. Each can have several “threads of execution”

We often run many programs at the same time
And each program may have several “threads of execution”

Example, in a Paint program, when you click the pencil tool, a new thread of execution is started to call the method to process it:

Main GUI thread          Process pencil click
Programming a Cluster...

- Sometimes you want to write a program that is executed on many machines!
- Atlas Cluster (at Cornell):
  - 768 cores
  - 1536 GB RAM
  - 24 TB Storage
  - 96 NICs (Network Interface Controller)
Many processes are executed simultaneously on your computer

- Operating system provides support for multiple “processes”
- Usually fewer processors than processes
- Processes are an abstraction: at hardware level, lots of multitasking
  - memory subsystem
  - video controller
  - buses
  - instruction prefetching
Concurrency

- *Concurrency* refers to a single program in which several processes, called threads, are running simultaneously
  - Special problems arise
  - They reference the same data and can interfere with each other, e.g. one process modifies a complex structure like a heap while another is trying to read it

- CS2110: we focus on two main issues:
  - Race conditions
  - Deadlock
Race conditions

- A “race condition” arises if two or more processes access the same variables or objects concurrently and at least one does updates

- Example: Processes t1 and t2

```
Process t1
...
x = x + 1;
```

```
Process t2
...
x = x + 1;
```

But `x = x+1;` is not an “atomic action”: it takes several steps
Race conditions

- Suppose $x$ is initially 5

**Thread t1**
- LOAD $x$
- ADD 1
- STORE $x$

**Thread t2**
- ...
- LOAD $x$
- ADD 1
- STORE $x$

- ... after finishing, $x = 6!$ We “lost” an update
Race conditions

- Typical race condition: two processes wanting to change a stack at the same time. Or make conflicting changes to a database at the same time.

- Race conditions are bad news

  - Race conditions can cause many kinds of bugs, not just the example we see here!
  
  - Common cause for “blue screens”: null pointer exceptions, damaged data structures
  
  - Concurrency makes proving programs correct much harder!
Deadlock

- To prevent race conditions, one often requires a process to “acquire” resources before accessing them, and only one process can “acquire” a given resource at a time.

- Examples of resources are:
  - A file to be read
  - An object that maintains a stack, a linked list, a hash table, etc.

- But if processes have to acquire two or more resources at the same time in order to do their work, **deadlock** can occur. This is the subject of the next slides.
Dining philosopher problem

Five philosophers sitting at a table.

Each repeatedly does this:
1. think
2. eat

What do they eat?
spaghetti.

Need TWO forks to eat spaghetti!
Dining philosopher problem

Each does repeatedly:
1. think
2. eat (2 forks)

eat is then:
- pick up left fork
- pick up right fork
- pick up food, eat
- put down left fork
- put down right fork

At one point, they all pick up their left forks

DEADLOCK!
Dining philosopher problem

Simple solution to deadlock:
Number the forks. Pick up smaller one first
1. think
2. eat (2 forks)

eat is then:
pick up smaller fork
pick up bigger fork
pick up food, eat
put down bigger fork
put down smaller fork
Java: What is a Thread?

- A separate “execution” that runs within a single program and can perform a computational task independently and concurrently with other threads.

- Many applications do their work in just a single thread: the one that called main() at startup.
  - But there may still be extra threads...
  - ... Garbage collection runs in a “background” thread.
  - GUls have a separate thread that listens for events and “dispatches” calls to methods to process them.

- Today: learn to create new threads of our own in Java.
A thread is an object that “independently computes”
- Needs to be created, like any object
- Then “started” --causes some method to be called. It runs side by side with other threads in the same program; they see the same global data
- The actual executions could occur on different CPU cores, but but don’t have to
  - We can also simulate threads by *multiplexing* a smaller number of cores over a larger number of threads
Java class Thread

- threads are instances of class Thread
  - Can create many, but they do consume space & time
- The Java Virtual Machine creates the thread that executes your main method.
- Threads have a priority
  - Higher priority threads are executed preferentially
  - By default, newly created threads have initial priority equal to the thread that created it (but priority can be changed)
Creating a new Thread (Method 1)

```java
class PrimeThread extends Thread {
    long a, b;

    PrimeThread(long a, long b) {
        this.a = a;
        this.b = b;
    }

    @Override
    public void run() {
        //compute primes between a and b
        ...
    }
}
```

PrimeThread p = new PrimeThread(143, 195);
p.start();

Call `run()` directly? no new thread is used: Calling thread will run it
overrides `Thread.run()`

Do this and Java invokes `run()` in new thread
Creating a new Thread (Method 2)

class PrimeRun implements Runnable {
    long a, b;

    PrimeRun(long a, long b) {
        this.a = a; this.b = b;
    }

    public void run() {
        //compute primes between a and b
        ...
    }
}

PrimeRun p = new PrimeRun(143, 195);
new Thread(p).start();
Example

```java
public class ThreadTest extends Thread {

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }

    public void run() {
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }
}
```
```java
public class ThreadTest extends Thread {
    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }

    public void run() {
        currentThread().setPriority(4);
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }
}
```
Example

```java
public class ThreadTest extends Thread {

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d
", Thread.currentThread(), i);
        }
    }

    public void run() {
        currentThread().setPriority(6);
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d
", Thread.currentThread(), i);
        }
    }

}}
```
Example

```java
public class ThreadTest extends Thread {
    static boolean ok = true;

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.println("waiting...");
            yield();
        }
        ok = false;
    }

    public void run() {
        while (ok) {
            System.out.println("running...");
            yield();
        }
        System.out.println("done");
    }
}
```

If threads happen to be sharing a CPU, yield allows other waiting threads to run.
Terminating Threads is tricky

- Easily done... but only in certain ways
  - Safe way to terminate a thread: return from method run
  - Thread throws uncaught exception? whole program will be halted (but it can take a second or two ... )
- Some old APIs have issues: stop(), interrupt(), suspend(), destroy(), etc.
  - Issue: Can easily leave application in a “broken” internal state.
  - Many applications have some kind of variable telling the thread to stop itself.
Threads can pause

- When active, a thread is “runnable”.
  - It may not actually be “running”. For that, a CPU must schedule it. Higher priority threads could run first.

- A thread can pause
  - Call Thread.sleep(k) to sleep for k milliseconds
  - Doing I/O (e.g. read file, wait for mouse input, open file) can cause thread to pause
  - Java has a form of locks associated with objects. When threads lock an object, one succeeds at a time.
Background (daemon) Threads

- In many applications we have a notion of “foreground” and “background” (daemon) threads
  - Foreground threads are doing visible work, like interacting with the user or updating the display
  - Background threads do things like maintaining data structures (rebalancing trees, garbage collection, etc.)

- On your computer, the same notion of background workers explains why so many things are always running in the task manager.
Fancier forms of locking

- Java developers have created various synchronization abstract data types
  - Semaphores: a kind of synchronized counter (invented by Dijkstra)
  - Event-driven synchronization

- The Windows and Linux and Apple O/S have kernel locking features, like file locking

- But for Java, *synchronized* is the core mechanism
Summary

- Use of multiple processes and multiple threads within each process can exploit concurrency
  - Which may be real (multicore) or “virtual” (an illusion)
- When using threads, beware!
  - Synchronize any shared memory to avoid race conditions
  - Synchronize objects in certain order to avoid deadlocks
  - Even with proper synchronization, concurrent programs can have other problems such as “livelock”
- Serious treatment of concurrency is a complex topic (covered in more detail in cs3410 and cs4410)