The exam is closed book and closed notes. Do not begin until instructed.

You have 90 minutes. Good luck!

Write your name and Cornell NetID at the top of every page! There are 5 questions on 10 numbered pages, front and back. Check that you have all the pages. When you hand in your exam, make sure your pages are still stapled together. If not, please use our stapler to reattach all your pages!

We have scrap paper available. If you do a lot of crossing out and rewriting, you might want to write code on scrap paper first and then copy it to the exam so that we can make sense of what you handed in.

Write your answers in the space provided. Ambiguous answers will be considered incorrect. You should be able to fit your answers easily into the space provided.

In some places, we have abbreviated or condensed code to reduce the number of pages that must be printed for the exam. In others, code has been obfuscated to make the problem more difficult. This does not mean that it’s good style.

**Academic Integrity Statement:** I pledge that I have neither given nor received any unauthorized aid on this exam.

______________________________

(signature)

1 of 10
1. True / False (10 points)

Circle T or F in the table below.

<p>| | | | |</p>
<table>
<thead>
<tr>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>a)</td>
<td>T</td>
<td>F</td>
<td>A method that computes a result in $O(n^3)$ time will always run slower than a method that computes a result in $O(n^2)$ time.</td>
</tr>
<tr>
<td>b)</td>
<td>T</td>
<td>F</td>
<td>If a graph with $n$ nodes has at least $n - 1$ edges, it is connected.</td>
</tr>
<tr>
<td>c)</td>
<td>T</td>
<td>F</td>
<td>The vertices of a finite graph can be topologically sorted if and only if the graph is acyclic.</td>
</tr>
<tr>
<td>d)</td>
<td>T</td>
<td>F</td>
<td>If a graph has thousands of vertices, each having at most 5 neighbors, you should implement the graph using an adjacency matrix instead of an adjacency list.</td>
</tr>
<tr>
<td>e)</td>
<td>T</td>
<td>F</td>
<td>LinkedList&lt;String&gt; is <strong>not</strong> a subtype of LinkedList&lt;Object&gt;, even though String[] is a subtype of Object[].</td>
</tr>
<tr>
<td>f)</td>
<td>T</td>
<td>F</td>
<td>Method m() processes a list of size $n$ using nested for-loops. Therefore, the runtime of m() can’t be $O(n^3)$.</td>
</tr>
<tr>
<td>g)</td>
<td>T</td>
<td>F</td>
<td>In the worst case, heap-sort and selection sort have the same run time.</td>
</tr>
<tr>
<td>h)</td>
<td>T</td>
<td>F</td>
<td>If a graph is bipartite, it is planar.</td>
</tr>
<tr>
<td>i)</td>
<td>T</td>
<td>F</td>
<td>Suppose a and b are objects. If a.equals(b) evaluates to <strong>true</strong>, then a.hashCode() == b.hashCode() must evaluate to <strong>true</strong>.</td>
</tr>
<tr>
<td>j)</td>
<td>T</td>
<td>F</td>
<td>Suppose a and b are objects. If a.equals(b) evaluates to <strong>false</strong>, then a.hashCode() == b.hashCode() must evaluate to <strong>false</strong>.</td>
</tr>
</tbody>
</table>
2. Complexity (30 points)

(a) 4 points  For each of the functions $f$ below, state the function $g(n)$ such that $f(n)$ is $O(g(n))$. $g(n)$ should be as simple and tight as possible. For example, one could say that $f(n) = 2n^2$ is $O(n^3)$ or $O(2n^2)$, but the best answer, the answer we want, is $O(n^2)$.

(i)  $f(n) = 2n \ast (\log n + n)$  \hspace*{1cm}  $g(n) = n^2$

(ii) $f(n) = 2^n + 3000n + 4$  \hspace*{1cm}  $g(n) = 2^n$

(iii) $f(n) = 9n \log n + \frac{n}{2}$  \hspace*{1cm}  $g(n) = n \log n$

(iv) $f(n) = 100 + 5n$  \hspace*{1cm}  $g(n) = n$

(b) 4 points  Recall that we proved in recitation that $f(n) = n + 6$ is $O(n)$. In a similar manner, prove that $f(n) = 17n^2 + 2n^3$ is $O(n^3)$.

\[
17n^2 + 2n^3 \leq n*n^2 + 2n^3 = 3n^3
\]

So, with $N = 17$ and $c = 3$, for all $n >= N$, $f(n) <= c g(n)$

Q.E.D. (meaning Quit, End, Done)

(c) 4 points  What is the simplest and tightest time-complexity class of the following function? For example, if the function takes $2n^2$ time, write $O(n^2)$. Be careful! Consider the complexity of all operations and function calls below.

```java
/** Reverse list in place. */
public static void reverse(ArrayList<String> list) {
    int n = list.size();
    for (int i = 0; i < n; i++) {
        list.add(0, list.remove(i));
    }
}
```

The runtime of this function is $O(n^2)$. Note that $\text{list.add(0, list.remove(i))}$ is an $O(n)$ operation because $\text{add()}$ is $O(n)$ and $\text{remove()}$ is $O(i)$ for an ArrayList.
(d) 6 points For each of the following tasks, state the expected and worst-case time complexity. If the expected and worst-case time complexities are different, describe a situation in which the task will take the worst-case time.

(i) (2 points) Use merge-sort to sort an array of size $n$.

Expected run time: $O(n \log n)$ Worst-case run time: $O(n \log n)$

(ii) (2 points) Search for a value in a sorted array of size $n$.

Expected run time: $O(\log n)$ Worst-case run time: $O(\log n)$

(iii) (2 points) Check if a binary search tree of size $n$ contains a particular value.

Expected run time: $O(\log n)$ Worst case run time: $O(n)$

Suppose the binary search tree has all of it’s left pointers equal to null. Then in the worst case, we have to traverse all $n$ nodes to check if an element is present.

(e) 12 points Suppose we need to choose a data structure to store a changing collection of values. For each of the usage patterns listed below, mark an “X” in the following table to indicate which data structure serves that usage pattern best in terms of average-case time complexity. There is exactly one best data structure for each usage pattern.

<table>
<thead>
<tr>
<th>Usage Pattern</th>
<th>ArrayList</th>
<th>LinkedList</th>
<th>HashSet</th>
<th>Balanced BST</th>
<th>Heap</th>
</tr>
</thead>
<tbody>
<tr>
<td>(i)</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(ii)</td>
<td></td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(iii)</td>
<td></td>
<td></td>
<td>X</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(iv)</td>
<td></td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(v)</td>
<td></td>
<td></td>
<td></td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>(vi)</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td>X</td>
</tr>
</tbody>
</table>

The usage patterns:
(i) (2 points) Removing certain elements while iterating through all elements (and preserving the ordering of the remaining elements).
(ii) (2 points) Checking if the collection contains various values
(iii) (2 points) Processing the elements in sorted order
(iv) (2 points) Adding and removing elements from indices clustered near the end of the collection
(v) (2 points) Retrieving elements at several unpredictable indices
(vi) (2 points) Processing the elements in a priority order
3. Heaps (20 points)

Consider creating a max-heap of `ints` by adding the following elements in the order presented:

-5, 10, 3, -2, 27, 13, 84

(a) 8 points  Below, draw the resulting heap as a tree.

Note: If you provide a valid heap for these values but not the one that arises from adding the elements in order, you get 2 points and can still get full credit for the subsequent problems.

Solution

![Heap Diagram](image)

(b) 6 points  Draw the heap you provided as an array in the table below.

Solution

<table>
<thead>
<tr>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
</tr>
</thead>
<tbody>
<tr>
<td>84</td>
<td>10</td>
<td>27</td>
<td>-5</td>
<td>-2</td>
<td>3</td>
<td>13</td>
</tr>
</tbody>
</table>

(c) 4 points  Now we change the value of the element with value 27 to 92. Repair the heap you presented using the operations from class and draw the resulting heap as a tree below.

Note: If you provide a valid heap for these values but not the one that arises from performing this operation, you get 1 point and can still get full credit for the subsequent problem.

Solution

![Heap Diagram](image)

(d) 2 points  Draw the heap you provided as an array in the table below.

Solution

<table>
<thead>
<tr>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>5</th>
<th>6</th>
</tr>
</thead>
<tbody>
<tr>
<td>92</td>
<td>10</td>
<td>84</td>
<td>-5</td>
<td>-2</td>
<td>3</td>
<td>13</td>
</tr>
</tbody>
</table>
4. **Trees** (20 points)

(a) 2 points  Is the tree above balanced?

The tree is **not** balanced

(b) 6 points  Provide the pre-order, post-order, and in-order traversals of the tree above. Specifically, list the order in which the nodes of the tree are processed in each traversal. You should represent a node by the letter it contains, but please leave at least one space between letters for readability (e.g. a b c, *not* abc).

(i) (2 points) Pre-Order:

```
y e t b z w d a c x
```

(ii) (2 points) Post-Order:

```
b t z e d x c a w y
```

(iii) (2 points) In-Order:

```
t b e z y d w c x a
```
(c) 4 points  Draw the BST (binary search tree) resulting from adding the following values one by one to an empty tree:

\[ r, x, o, t, m, s, p, n \]

Note: If you provide a valid BST for these values but not the one that arises from adding the elements in order, you get 1 point.

Solution
(d) **8 points**  Now that you have practiced adding values to a BST, implement method `add` of class `BSTNode` below.

```java
/** An instance of this class represents a Binary Search Tree */
public class BSTNode {
    private int element; // the element of this BSTNode
    private BSTNode left; // left child of this BSTNode (null for an empty tree)
    private BSTNode right; // right child of this BSTNode (null for an empty tree)

    /** Constructor: a binary search tree containing only v */
    public BSTNode(int v) {
        element = v;
    }

    /** Add v to t and return t. If t is null return a new BSTNode containing v */
    // * If v is already in t, do nothing and return t. */
    public static BSTNode add(int v, BSTNode t) {
        if (t == null) return new BSTNode(v);
        if (v == t.element) return t;
        if (v < t.element) {
            t.left = add(v, t.left);
            return t;
        }
        t.right = add(v, t.right);
        return t;
    }
}
```
5. **Graphs** (20 points)

(a) **10 points** You are given a (weighted) directed graph, duplicated in the first column below for each subproblem so that you may mark it up with notes (which will be ignored by the graders). For each subproblem, write your answer in the *second* column, NOT THE FIRST COLUMN! When you need to choose among vertices to visit first, choose in alphabetical order.

(i) (3 points) In what order are the nodes visited by DFS, starting from vertex *a*?

(a c f d e g b h)

(ii) (3 points) In what order are the nodes visited by BFS, starting from vertex *a*?

(a c e f g h d b)

(iii) (4 points) In what order are the nodes removed from the frontier in Dijkstra’s algorithm, starting from vertex *a*?

(a e c g f h b d)
(b) 10 points  The graph algorithms we have covered have applications beyond the problems for which they were originally intended. For example, although DFS and BFS are search algorithms, they are just ways to traverse a graph. So besides reachability, they can be used for distance calculation and cycle detection. And, topological sort can be used to determine shortest paths in DAGs.

For each problem listed below, mark an “X” in the following table to indicate which graph algorithm solves the problem best in terms of worst-case time complexity. There is exactly one best graph algorithm for each problem.

<table>
<thead>
<tr>
<th>Problem</th>
<th>Topological sort</th>
<th>DFS</th>
<th>BFS</th>
<th>Dijkstra</th>
</tr>
</thead>
<tbody>
<tr>
<td>(i)</td>
<td></td>
<td></td>
<td>X</td>
<td></td>
</tr>
<tr>
<td>(ii)</td>
<td></td>
<td></td>
<td>X</td>
<td></td>
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<td></td>
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</tr>
<tr>
<td>(iv)</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>(v)</td>
<td>X</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

The problems:

(i) (2 points) You are on a bike trip in a park. Your map of the park shows you places to stop and rest, and you can calculate the distance between any two of them. You can bike only \( b \) miles without stopping at a resting place before you fall over, exhausted, and every mile takes you 10 minutes to bike. You are standing at the easternmost resting place. You want to determine how to most quickly reach the westernmost resting place without getting exhausted.

(ii) (2 points) You are intrigued by the concept of “six degrees of separation”: any two people on the planet are connected through at most five intermediate acquaintances. You want to test this theory at Cornell. You have a database that tells you whether any two Cornell students are friends. You want to find a student who is “furthest away” from you, i.e. the minimum number of intermediate friends needed to relate that student to you is greater than or equal to that of any other student.

(iii) (2 points) You are a maze designer, and you want to make sure your preliminary design isn’t too challenging. Write a program that, assuming there are no cycles, determines the length of the longest path from the maze’s entrance that doesn’t eventually lead to the maze’s exit.

(iv) (2 points) You have a list of courses you want to take at Cornell. Some are prerequisites of others. You want to find a sequence in which to take these classes without violating any prerequisite requirement.

(v) (2 points) Ithaca has bicycle paths, but not on every road. For each pair of intersections A and B, your map of Ithaca tells you whether there is a bicycle path from A to B; call such a path a direct bicycle path. You have to write a program to calculate the minimum number of direct bicycle paths needed to get from one place in Ithaca to another.