Prelim 2

- 5:30 pm, Mallot 228 (even ids)
- 7:30pm, Kennedy Call Aud (odd ids)

- Those who need extra time should start at the 5:30 slot, and move into Kennedy Hall for the rest of their time when the 5:30 test ends.

- Those who need a quiet space should go directly to the small room Kennedy (it is 101) at 5:30?
Dijkstra: given a vertex v, finds shortest path from v to x for each vertex x in the graph

Key idea: maintain a 5-part invariant on three sets
1. Vertices already visited (“settled”). Distance known
2. Frontier nodes. One hop from the settled ones
3. Future nodes. > one hop from the settled ones

Algorithm: move the “closest” frontier node to settled, then adjust frontier and future sets to restore the invariant.
Graphs summary

- Minimum spanning tree: a tree that reaches every node while minimizing the summed weight of edges
  - Prim’s algorithm: repeatedly pick the lowest-weight edge that will connect some previously disconnected components. A “greedy” algorithm.
  - Kruskal’s algorithm: start with the whole graph, repeatedly remove the highest-weight edge that won’t disconnect the spanning tree. Also “greedy”.

- In all three cases, correctness is established using inductive proofs that focus on maintaining invariants!
Today: Start a new topic

- Modern computers have “multiple cores”
  - Instead of a single CPU on the chip
  - 5-10 common. Intel has prototypes with 80!

- And even with a single core your program may have more than one thing “to do” at a time
  - Argues for having a way to do many things at once

- Finally, we often run many programs all at once
Why Multicore?

- Moore’s Law: Computer speeds and memory densities nearly double each year
But a fast computer runs hot

- Power dissipation rises as the square of the CPU clock rate
- Chips were heading towards melting down!

- Multicore: with four CPUs (cores) on one chip, even if we run each at half speed we get more overall performance!
Keeping those cores busy

• The operating system provides support for multiple “processes”
• In reality there may be fewer processors than processes
• Processes are an illusion – at the hardware level, lots of multitasking
  – memory subsystem
  – video controller
  – buses
  – instruction prefetching
• Virtualization can even let one machine create the illusion of many machines (they share disks, etc)
What is a Thread?

- A separate “execution” that runs within a single program and can perform a computational task independently and concurrently with other threads.

- Many applications do their work in just a single thread: the one that called main() at startup.
  - But there may still be extra threads...
  - ... Garbage collection runs in a “background” thread.
  - GUIs have a separate thread that listens for events and “dispatches” upcalls.

- Today: learn to create new threads of our own.
What is a Thread?

- A thread is a kind of object that “independently computes”
  - Needs to be created, like any object
  - Then “started”. This causes some method (like main()) to be invoked. It runs side by side with other thread in the same program and they see the same global data

- The actual execution could occur on distinct CPU cores, but doesn’t need to
  - We can also simulate threads by *multiplexing* a smaller number of cores over a larger number of threads
Concurrency

- Concurrency refers to a single program in which several threads are running simultaneously
  - Special problems arise
  - They see the same data and hence can interfere with each other, e.g. if one thread is modifying a complex structure like a heap while another is trying to read it

- In cs2110 we focus on two main issues:
  - Race conditions
  - Deadlock
Thread class in Java

- Threads are instances of the class Thread
  - Can create many, but they do consume space & time

- The Java Virtual Machine created the thread that executes your main method.

- Threads have a priority
  - Higher priority threads are executed preferentially
  - A newly created Thread has initial priority equal to the thread that created it (but can change)
Creating a new Thread (Method 1)

```java
class PrimeThread extends Thread {
    long a, b;

    PrimeThread(long a, long b) {
        this.a = a; this.b = b;
    }

    public void run() {
        //compute primes between a and b
        ...
    }
}

PrimeThread p = new PrimeThread(143, 195);
p.start();
```

- **overrides** `Thread.run()`
- If you were to call `run()` directly, no new thread is used; the calling thread will run it.
- ...but if you create a new object and then call `start()`, Java invokes `run()` in a new thread.
Creating a new Thread (Method 2)

```java
class PrimeRun implements Runnable {
    long a, b;

    PrimeRun(long a, long b) {
        this.a = a; this.b = b;
    }

    public void run() {
        //compute primes between a and b
        ...
    }
}

PrimeRun p = new PrimeRun(143, 195);
new Thread(p).start();
```
```java
public class ThreadTest extends Thread {

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }

    public void run() {
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }
}
```
public class ThreadTest extends Thread {

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }

    public void run() {
        currThread().setPriority(4);
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n", Thread.currentThread(), i);
        }
    }
}
public class ThreadTest extends Thread {

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n",
                Thread.currentThread(), i);
        }
    }

    public void run() {
        currentThread().setPriority(6);
        for (int i = 0; i < 10; i++) {
            System.out.format("%s %d\n",
                Thread.currentThread(), i);
        }
    }

}
Example

```java
public class ThreadTest extends Thread {
    static boolean ok = true;

    public static void main(String[] args) {
        new ThreadTest().start();
        for (int i = 0; i < 10; i++) {
            System.out.println("waiting...");
            yield();
        }
        ok = false;
    }

    public void run() {
        while (ok) {
            System.out.println("running...");
            yield();
        }
        System.out.println("done");
    }
}
```

If threads happen to be sharing a CPU, yield allows other waiting threads to run. But if there are multiple cores, yield isn’t needed.
Terminating Threads is tricky

- Easily done... but only in certain ways
  - The safe way to terminate a thread is to have it return from its run method
  - If a thread throws an uncaught exception, whole program will be halted (but it can take a second or two...)

- There are some old APIs but they have issues: stop(), interrupt(), suspend(), destroy(), etc.
  - Issue: they can easily leave the application in a “broken” internal state.
  - Many applications have some kind of variable telling the thread to stop itself.
Threads can pause

- When active, a thread is “runnable”.
  - It may not actually be “running”. For that, a CPU must schedule it. Higher priority threads could run first.

- A thread can also pause
  - It can call Thread.sleep(k) to sleep for k milliseconds
  - If it tries to do “I/O” (e.g. read a file, wait for mouse input, even open a file) this can cause it to pause
  - Java has a form of locks associated with objects. When threads lock an object, one succeeds at a time.
Background (daemon) Threads

- In many applications we have a notion of “foreground” and “background” (daemon) threads
  - Foreground threads are the ones doing visible work, like interacting with the user or updating the display
  - Background threads do things like maintaining data structures (rebalancing trees, garbage collection, etc)

- On your computer, the same notion of background workers explains why so many things are always running in the task manager.
Race Conditions

- A “race condition” arises if two or more threads access the same variables or objects concurrently and at least one does updates.

- Example: Suppose t1 and t2 simultaneously execute the statement \( x = x + 1 \); for some static global \( x \).
  - Internally, this involves loading \( x \), adding 1, storing \( x \).
  - If t1 and t2 do this concurrently, we execute the statement twice, but \( x \) may only be incremented once.
  - t1 and t2 “race” to do the update.
Suppose X is initially 5

... after finishing, X=6! We “lost” an update
Race Conditions

- Race conditions are bad news
  - Sometimes you can make code behave correctly despite race conditions, but more often they cause bugs
  - And they can cause many kinds of bugs, not just the example we see here!
  - A common cause for “blue screens”, null pointer exceptions, damaged data structures
Example – A Lucky Scenario

```java
private Stack<String> stack = new Stack<String>();

public void doSomething() {
    if (stack.isEmpty()) return;
    String s = stack.pop();
    //do something with s...
}
```

Suppose threads A and B want to call `doSomething()`, and there is one element on the stack

1. thread A tests `stack.isEmpty()` false
2. thread A pops ⇒ stack is now empty
3. thread B tests `stack.isEmpty()` ⇒ true
4. thread B just returns – nothing to do
Example — An Unlucky Scenario

```java
private Stack<String> stack = new Stack<String>();

public void doSomething() {
    if (stack.isEmpty()) return;
    String s = stack.pop();
    //do something with s...
}
```

Suppose threads A and B want to call `doSomething()`, and there is one element on the stack

1. thread A tests `stack.isEmpty() ⇒ false`
2. thread B tests `stack.isEmpty() ⇒ false`
3. thread A pops ⇒ stack is now empty
4. thread B pops ⇒ Exception!
Java has one “primary” tool for preventing these problems, and you must use it by carefully and explicitly — it isn’t automatic.

- Called a “synchronization barrier”

- We think of it as a kind of lock
  - Even if several threads try to acquire the lock at once, only one can succeed at a time, while others wait
  - When it releases the lock, the next thread can acquire it
  - You can’t predict the order in which contending threads will get the lock but it should be “fair” if priorities are the same
Solution – with synchronization

```java
private Stack<String> stack = new Stack<String>();

public void doSomething() {
    synchronized (stack) {
        if (stack.isEmpty()) return;
        String s = stack.pop();
    }
    //do something with s...
}
```

- Put critical operations in a `synchronized` block
- The `stack` object acts as a lock
- Only one thread can own the lock at a time
Solution – Locking

• You can lock on any object, including this

```java
public synchronized void doSomething() {
    ...
}
```

is equivalent to

```java
public void doSomething() {
    synchronized (this) {
        ...
    }
}
```
Combining mundane features can get you in trouble

Java has priorities... and synchronization

But they don’t “mix” nicely

High-priority runs before low priority

... The lower priority thread “starves”

Even worse...

With many threads, you could have a second high priority thread stuck waiting on that starving low priority thread! Now both are starving...
Fancier forms of locking

- Java developers have created various synchronization ADTs
  - Semaphores: a kind of synchronized counter
  - Event-driven synchronization

- The Windows and Linux and Apple O/S all have kernel locking features, like file locking

- But for Java, *synchronized* is the core mechanism
Deadlock

- The downside of locking – deadlock

- A deadlock occurs when two or more competing threads are waiting for one-another... forever

- Example:
  - Thread t1 calls synchronized b inside synchronized a
  - But thread t2 calls synchronized a inside synchronized b
  - t1 waits for t2... and t2 waits for t1...
Finer grained synchronization

- Java allows you to do fancier synchronization
  - But can only be used inside a synchronization block
  - Special primitives called wait/notify
Suppose we put this inside an object called `animator`:

```java
boolean isRunning = true;

public synchronized void run() {
    while (true) {
        while (isRunning) {
            //do one step of simulation
        }
        try {
            wait();
        } catch (InterruptedException ie) {
        }
        isRunning = true;
    }
}

public void stopAnimation() {
    animator.isRunning = false;
}

public void restartAnimation() {
    synchronized (animator) {
        animator.notify();
    }
}
```

- `wait/notify` must be synchronized!
- Relinquishes lock on `animator` – awaits notification
- Notifies processes waiting for `animator` lock
Use of multiple processes and multiple threads within each process can exploit concurrency
- Which may be real (multicore) or “virtual” (an illusion)

But when using threads, beware!
- Must lock (synchronize) any shared memory to avoid non-determinism and race conditions
- Yet synchronization also creates risk of deadlocks
- Even with proper locking concurrent programs can have other problems such as “livelock”

Serious treatment of concurrency is a complex topic (covered in more detail in cs3410 and cs4410)