GRAMMARS & PARSING

Lecture 7
CS2110 – Spring 2013
Java Tips

- Declare fields and methods **public** if they are to be visible outside the class; helper methods and private data should be declared **private**.

- Constants that will never be changed should be declared **final**.

- Public classes should appear in a file of the same name.

- Two kinds of boolean operators:
  - **e1 & e2**: evaluate both and compute their conjunction.
  - **e1 && e2**: evaluate **e1**; don’t evaluate **e2** unless necessary.

  ```java
  instead of
  if (s.equals("")) {
      f = true;
  } else {
      f = false;
  }
  write
  f = s.equals("");
  ```

  ```java
  instead of
  if (s.equals("")) {
      f = a;
  } else {
      f = b;
  }
  write
  f = s.equals("")? a : b;
  ```
Application of Recursion

- So far, we have discussed recursion on integers
  - Factorial, fibonacci, $a^n$, combinatorials

- Let us now consider a new application that shows off the full power of recursion: parsing

- Parsing has numerous applications: compilers, data retrieval, data mining,...
Motivation

- The cat ate the rat.
- The cat ate the rat slowly.
- The small cat ate the big rat slowly.
- The small cat ate the big rat on the mat slowly.
- The small cat that sat in the hat ate the big rat on the mat slowly.
- The small cat that sat in the hat ate the big rat on the mat slowly, then got sick.
- [...]

- Not all sequences of words are legal sentences
  - The ate cat rat the
- How many legal sentences are there?
- How many legal programs are there?
- Are all Java programs that compile legal programs?
- How do we know what programs are legal?

A Grammar

- Grammar: set of rules for generating sentences in a language
- Examples of Sentence:
  - boys see bunnies
  - bunnies like girls
  - ...
- White space between words does not matter
- The words boys, girls, bunnies, like, see are called *tokens* or *terminals*
- The words Sentence, Noun, Verb are called *nonterminals*
- This is a very boring grammar because the set of Sentences is finite (exactly 18 sentences)

- Our sample grammar has these rules:
  - A Sentence can be a Noun followed by a Verb followed by a Noun
  - A Noun can be ‘boys’ or ‘girls’ or ‘bunnies’
  - A Verb can be ‘like’ or ‘see’

- Sentence $\rightarrow$ Noun Verb Noun
- Noun $\rightarrow$ boys
- Noun $\rightarrow$ girls
- Noun $\rightarrow$ bunnies
- Verb $\rightarrow$ like
- Verb $\rightarrow$ see
A Recursive Grammar

- **Sentence → Sentence and Sentence**
- **Sentence → Sentence or Sentence**
- **Sentence → Noun Verb Noun**
- **Noun → boys**
- **Noun → girls**
- **Noun → bunnies**
- **Verb → like**
- **Verb → see**

- This grammar is more interesting than the last one because the set of Sentences is infinite

- **Examples of Sentences in this language:**
  - boys like girls
  - boys like girls and girls like bunnies
  - boys like girls and girls like bunnies and girls like bunnies
  - boys like girls and girls like bunnies and girls like bunnies and girls like bunnies
  - ........

- **What makes this set infinite?**
  - Recursive definition of Sentence
Detour

- What if we want to add a period at the end of every sentence?
  - Sentence → Sentence and Sentence .
  - Sentence → Sentence or Sentence .
  - Sentence → Noun Verb Noun .
  - Noun → ...

- Does this work?
- No! This produces sentences like: girls like boys . and boys like bunnies . .

Sentence Sentence Sentence
Sentences with Periods

- **PunctuatedSentence** → **Sentence** .
- **Sentence** → **Sentence** and **Sentence**
- **Sentence** → **Sentence** or **Sentence**
- **Sentence** → **Noun** **Verb** **Noun**
- **Noun** → **boys**
- **Noun** → **girls**
- **Noun** → **bunnies**
- **Verb** → **like**
- **Verb** → **see**

- Add a new rule that adds a period only at the end of the sentence.
- The tokens here are the 7 words plus the period (.).
- This grammar is ambiguous:
  - boys like girls
  - and girls like boys
  - or girls like bunnies
Grammar for Simple Expressions

- $E \rightarrow \text{integer}$
- $E \rightarrow (E + E)$

**Simple expressions:**
- An $E$ can be an integer.
- An $E$ can be ‘(’ followed by an $E$ followed by ‘+’ followed by an $E$ followed by ‘)’

**Set of expressions defined by this grammar is a recursively-defined set**
- Is language finite or infinite?
- Do recursive grammars always yield infinite languages?

- Here are some legal expressions:
  - 2
  - $(3 + 34)$
  - $((4+23) + 89)$
  - $((89 + 23) + (23 + (34+12)))$

- Here are some illegal expressions:
  - $(3$
  - $3 + 4$

- The *tokens* in this grammar are (, +, ), and any integer
Grammars can be used in two ways:

- A grammar defines a language (i.e., the set of properly structured sentences)
- A grammar can be used to parse a sentence (thus, checking if the sentence is in the language)

To parse a sentence is to build a parse tree:

- This is much like diagramming a sentence:

Example: Show that \((4+23) + 89\) is a valid expression \(E\) by building a parse tree:

```
     E
   /   \
E +   E
 /   \
E +   89
 |     \
4 23
```
Recursive Descent Parsing

- Idea: Use the grammar to design a recursive program to check if a sentence is in the language
- To parse an expression E, for instance
  - We look for each terminal (i.e., each token)
  - Each nonterminal (e.g., E) can handle itself by using a recursive call
- The grammar tells how to write the program!

```java
boolean parseE() {
    if (first token is an integer) return true;
    if (first token is '(') {
        parseE();
        Make sure there is a '+' token;
        parseE();
        Make sure there is a ')' token;
        return true;
    }
    return false;
}
```
public static Node parseE(Scanner scanner) {
    if (scanner.hasNextInt()) {
        int data = scanner.nextInt();
        return new Node(data);
    }
    check(scanner, '(');
    left = parseE(scanner);
    check(scanner, '+');
    right = parseE(scanner);
    check(scanner, ')');
    return new Node(left, right);
}
Detour: Error Handling with Exceptions

- Parsing does two things:
  - It returns useful data (a parse tree)
  - It checks for validity (i.e., is the input a valid sentence?)

- How should we respond to invalid input?

- *Exceptions* allow us to do this without complicating our code unnecessarily
Exceptions

- Exceptions are usually thrown to indicate that something bad has happened
  - `IOException` on failure to open or read a file
  - `ClassCastException` if attempted to cast an object to a type that is not a supertype of the dynamic type of the object
  - `NullPointerException` if tried to dereference null
  - `ArrayIndexOutOfBoundsException` if tried to access an array element at index $i < 0$ or $i \geq \varepsilon$ the length of the array

- In our case (parsing), we should throw an exception when the input cannot be parsed
Handling Exceptions

- Exceptions can be caught by the program using a **try-catch** block

- **catch** clauses are called **exception handlers**

```java
Integer x = null;
try {
    x = (Integer)y;
    System.out.println(x.intValue());
} catch (ClassCastException e) {
    System.out.println("y was not an Integer");
} catch (NullPointerException e) {
    System.out.println("y was null");
}
```
Defining Your Own Exceptions

- An exception is an object (like everything else in Java)
- You can define your own exceptions and throw them

```java
class MyOwnException extends Exception {}

...if (input == null) {
    throw new MyOwnException();
}
```
Declaring Exceptions

- In general, any exception that could be thrown must be either declared in the method header or caught

```java
void foo(int input) throws MyOwnException {
    if (input == null) {
        throw new MyOwnException();
    }
    ...
}
```

- Note: `throws` means “can throw”, not “does throw”
- Subtypes of `RuntimeException` do not have to be declared (e.g., `NullPointerException`, `ClassCastException`)
  - These represent exceptions that can occur during “normal operation of the Java Virtual Machine”
How Exceptions are Handled

- If the exception is thrown from inside the `try` clause of a `try-catch` block with a handler for that exception (or a superclass of the exception), then that handler is executed.
  - Otherwise, the method terminates abruptly and control is passed back to the calling method.

- If the calling method can handle the exception (i.e., if the call occurred within a `try-catch` block with a handler for that exception) then that handler is executed.
  - Otherwise, the calling method terminates abruptly, etc.

- If none of the calling methods handle the exception, the entire program terminates with an error message.
We can modify the parser so that it generates stack code to evaluate arithmetic expressions:

2
PUSH 2
STOP

(2 + 3)
PUSH 2
PUSH 3
ADD
STOP

Goal: Method parseE should return a string containing stack code for expression it has parsed

- Method parseE can generate code in a recursive way:
  - For integer i, it returns string “PUSH ” + i + “\n”
  - For (E1 + E2),
    - Recursive calls for E1 and E2 return code strings c1 and c2, respectively
    - Then to compile (E1 + E2), return c1 + c2 + “ADD\n”
  - Top-level method should tack on a STOP command after code received from parseE
Does Recursive Descent Always Work?

- There are some grammars that cannot be used as the basis for recursive descent
  - A trivial example (causes infinite recursion):
    - $S \rightarrow b$
    - $S \rightarrow Sa$

- For some constructs, recursive descent is hard to use
  - Can use a more powerful parsing technique (there are several, but not in this course)

- Can rewrite grammar
  - $S \rightarrow b$
  - $S \rightarrow bA$
  - $A \rightarrow a$
  - $A \rightarrow aA$
Sometimes a sentence has more than one parse tree
- $S \rightarrow A \mid axB$
- $A \rightarrow x \mid aAb$
- $B \rightarrow b \mid bB$
- The string $aaxbb$ can be parsed in two ways

This kind of ambiguity sometimes shows up in programming languages

if $E_1$ then if $E_2$ then $S_1$ else $S_2$

Which then does the else go with?

This ambiguity actually affects the program’s meaning

How do we resolve this?
- Provide an extra non-grammar rule (e.g., the else goes with the closest if)
- Modify the language (e.g., an if-statement must end with a ‘fi’)
- Operator precedence (e.g., $1 + 2 * 3$ should be parsed as $1 + (2 * 3)$, not $(1 + 2) * 3$)
- Other methods (e.g., Python uses amount of indentation)
Recursion is a very powerful technique for writing compact programs that do complex things.

Common mistakes:
- Incorrect or missing base cases
- Subproblems must be simpler than top-level problem

Try to write description of recursive algorithm and reason about base cases before writing code.

Why?
- Syntactic junk such as type declarations, etc. can create mental fog that obscures the underlying recursive algorithm.
- Best to separate the logic of the program from coding details.
Exercises

- Think about recursive calls made to parse and generate code for simple expressions
  - 2
  - (2 + 3)
  - ((2 + 45) + (34 + -9))

- Derive an expression for the total number of calls made to parseE for parsing an expression
  - **Hint:** think inductively

- Derive an expression for the maximum number of recursive calls that are active at any time during the parsing of an expression (i.e. max depth of call stack)
Exercises

- Write a grammar and recursive program for sentence palindromes that ignores white spaces & punctuation
  - Was it Eliot's toilet I saw?
  - No trace; not one carton
  - Go deliver a dare, vile dog!
  - Madam, in Eden I'm Adam

- Write a grammar and recursive program for strings $A^nB^n$
  - AB
  - AABB
  - AAAAAAABBBBBB

- Write a grammar and recursive program for Java identifiers
  - <letter> [<letter> or <digit>]\(^0...N\)
  - j27, but not 2j7