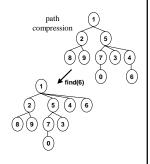
Finish Union/Find, Finish Graph Algorithms, Quick Overview

CS211 Fall 2000

Union/Find

- Operations
 - Union: Combine two sets
 - Find: Given an item, determine the "name" of the set that contains it
- Use reverse trees
 - Each item points at its parent
 The root is the "name" of the
 - The root is the "name" of the set
- Union-by-Size
 - Always make the larger tree be the root
- Path Compression
 - Every time we "find" something, we update every item we touch so that it points at the root

2



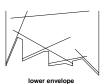
Union/Find Analysis

- Theorem (Tarjan)
 Using weighted union and path compression, a sequence of n union/find operations takes time O(n \u03c4(n))
- Note that α(n) ≤ 4 for any integer n that we are *ever* likely to encounter
- Is the α(n) factor really necessary?
 - Yes: Tarjan showed a lower bound of Ω(n α(n)) for union/find
 - Claim: the inverse Ackerman's function is not just an artifact of this one problem

Lower Envelope of Line Segments

Given n line segments in the plane, what is the worst-case complexity of their lower envelope?

 $\Theta(n \alpha(n))$



Two MST Algorithms (Both Greedy)

Kruskal's Algorithm

- Choose the shortest edge e such that
 - e is not yet processed
 - e does not make a cycle

Prim's Algorithm

- Choose the shortest edge e such that
 - e touches the tree
 - e touches a vertex not in the tree

Kruskal's MST Algorithm

KruskalMST(G):

E = edges of G; forest = empty;

do

<u,v> = least cost edge of E;
E = E - <u,v>;
if (u and v in different trees)
forest = forest ∪ <u,v>;

while (E is nonempty); return forest

- Can sort the edges initially (or can use a PQ)
- Use Union/Find to check for different trees and to combine trees
- Total worst-case time: O(e log e) when using adjacency lists
- Time is O(v² + e log e) for adjacency matrix

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Quick Review: Programming Topics

Object Oriented Programming

- Classes and Objects
 - ▲ Parameter Passing
 - ▲ Objects vs. References
 - ▲ Abstraction
- ▲ Encapsulation
- Inheritance ▲ Polymorphism
 - ▲ Dynamic Method Binding
 - ▲ Abstract Classes
 - ▲ Interfaces
 - ▲ Upcasting vs. Downcasting

- Other Topics
 - Access-Control Modifiers
 - Packages
 - Exceptions
 - Program Design
 - ▲ UML Diagrams
 - ▲ Pseudo-code
 - GUIs
- ▲ Javadoc ▲ Layout
 - ▲ Event Handling

Quick Review: Data-Structure Topics

- Searching and Sorting in Arrays
 - Binary Search
 - Quick Sort
 - Merge Sort
 - Insertion Sort
 - Heap Sort
 - Sorting Lower Bound
- Data Structures for Searching (Dictionaries)
 - Hash Table
 - Binary Search Tree
 - Balanced Trees

- Data Structures for Sequencing
 - Stack
 - Queue
 - Priority Queue
- Graphs (Adj List, Adj Matrix)
 - Shortest Paths
 - ▲ Breadth First Search
 - ▲ Dijkstra's Algorithm
 - Minimum Spanning Trees
 - ▲ Prim's Alg (single tree)
 - ▲ Kruskal's Alg (forest)

Quick Review: Additional Topics

- Recursion
 - Recursive Descent Parsing
 - Stack Frames
 - Induction
- Algorithm Analysis
 - Big-O
 - Worst- vs. Expected-Case
- Algorithm Design Methods
 - Divide and Conquer
- Greedy Method
- Union/Find
 - · union-by-size
 - path compression

- The Java Collections Framework
 - Interfaces: Set, SortedSet, List, Map, SortedMap, Iterator
 - Classes: HashSet, TreeSet, ArrayList, LinkedList, HashMap, TreeMap
 - Utilities: java.util.Arrays, java.util.Collections
 - Comparator vs.
 - Comparable

What I Do: Computational Geometry

- Using a computer to solve geometric problems
 - · Get to use lots of data structure ideas
 - Example
 - ▲ Given n line segments in the plane, report all intersections ▲ Uses both a PQ and a
 - Ralanced Tree
- Areas I work in
 - Motion Planning
 - Meshing
 - Shape Matching
 - ▲ computer vision
 - protein matching
 - · More theoretical questions