

- Previous Lecture:
 - 2-d array—matrix
- Today's Lecture:
 - More examples on matrices
 - Optional reading: contour plot (7.2, 7.3 in *Insight*)
- Announcement:
 - Review your prelim and **re-do** the questions on which you didn't do well, perhaps with the help of a member of the course staff. Do not just read the solutions!

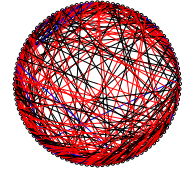
Storing and using data in tables

A company has 3 factories that make 5 products with these costs:

C

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

What is the best way to fill a given purchase order?



Connections between webpages

0	0	1	0	1	0	0
1	0	0	1	1	1	0
0	1	0	1	1	1	1
1	0	1	1	0	1	0
0	0	1	1	0	1	1
0	0	1	0	1	0	1
0	1	1	0	1	1	0

Lecture 13

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Pattern for traversing a matrix M

```

[nr, nc] = size(M)
for r= 1:nr
    % At row r
    for c= 1:nc
        % At column c (in row r)
        %
        % Do something with M(r,c) ...
    end
end
end

```

Lecture 14

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A Cost/Inventory Problem

- A company has 3 factories that make 5 different products
- The cost of making a product varies from factory to factory
- The inventory/capacity varies from factory to factory

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Problems

A customer submits a purchase order that is to be filled by a single factory.

1. How much would it cost a factory to fill the order?
2. Does a factory have enough inventory/capacity to fill the order?
3. Among the factories that can fill the order, who can do it most cheaply?

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Cost Array

C

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

The value of **C(i, j)** is what it costs factory **i** to make product **j**.

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Inventory (or Capacity) Array

Inv

38	5	99	34	42
82	19	83	12	42
51	29	21	56	87

The value of `Inv(i,j)` is the inventory in factory `i` of product `j`.

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Purchase Order

PO

1	0	12	29	5
---	---	----	----	---

The value of `PO(j)` is the number of product `j`'s that the customer wants

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C

10	36	22	15	62
12	35	20	12	66
13	37	21	16	59

PO

1	0	12	29	5
---	---	----	----	---

Cost for factory `i`:

```
s = 0; %Sum of cost
for j=1:5
    s = s + C(i,j)*PO(j)
end
```

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Encapsulate...

```
function TheBill = iCost(i,C,PO)
% The cost when factory i fills the
% purchase order

nProd = length(PO);
TheBill = 0;
for j=1:nProd
    TheBill = TheBill + C(i,j)*PO(j);
end
```

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Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
        % Found an Improvement
        iBest = i; minBill = iBill;
    end
end
```

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`inf` – a special value that can be regarded as positive infinity

```
x = 10/0    assigns inf to x
y = 1+x      assigns inf to y
z = 1/x      assigns 0 to z
w < inf      is always true if w is numeric
```

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Inventory/Capacity Considerations

What if a factory lacks the inventory/capacity to fill the purchase order?

Such a factory should be excluded from the find-the-cheapest computation.

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Who Can Fill the Order?

	38	5	99	34	42	Yes
Inv	82	19	83	12	42	No
	51	29	21	56	87	Yes
PO	1	0	12	29	5	

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Wanted: A True/False Function



DO is "true" if **factory i** can fill the order.
DO is "false" if **factory i** cannot fill the order.

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Example: Check inventory of factory 2

	38	5	99	34	42
Inv	82	19	83	12	42
	51	29	21	56	87
PO	1	0	12	29	5

Method 1: check the inventory for every product

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Still True...

	38	5	99	34	42
Inv	82	19	83	12	42
	51	29	21	56	87
PO	1	0	12	29	5

DO 1

DO = DO && (Inv(2,1) >= PO(1))

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Still True...

	38	5	99	34	42
Inv	82	19	83	12	42
	51	29	21	56	87
PO	1	0	12	29	5

DO 1

DO = DO && (Inv(2,2) >= PO(2))

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Encapsulate...

```
function DO = iCanDo(i,Inv,PO)
% DO is true if factory i can fill
% the purchase order. Otherwise, false

nProd = length(PO);
DO = 1;
for j = 1:nProd
    DO = DO && ( Inv(i,j) >= PO(j) );
end
```

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Encapsulate...

```
function DO = iCanDo(i,Inv,PO)
% DO is true if factory i can fill
% the purchase order. Otherwise, false
nProd = length(PO);
j = 1;
while j<=nProd && Inv(i,j)>=PO(j)
    j = j+1;
end
DO = _____;
```

DO should be true when...

- A j < nProd
- B j == nProd
- C j > nProd

Lecture 14

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Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    iBill = iCost(i,C,PO);
    if iBill < minBill
        % Found an Improvement
        iBest = i; minBill = iBill;
    end
end
```

Don't bother with this unless there is sufficient inventory.

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Back To Finding the Cheapest

```
iBest = 0; minBill = inf;
for i=1:nFact
    if iCanDo(i,Inv,PO)
        iBill = iCost(i,C,PO);
        if iBill < minBill
            % Found an Improvement
            iBest = i; minBill = iBill;
        end
    end
end
```

See `Cheapest.m`
for alternative implementation

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Finding the Cheapest

C	10	36	22	15	62	1019	Yes
	12	35	20	12	66	930	No
	13	37	21	16	59	1040	Yes
PO	1	0	12	29	5		

As computed by `iCost` As computed by `iCanDo`

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Initialize vectors/matrices if dimensions are known
...instead of “building” the array one component at a time

```
% Initialize y
x=linspace(a,b,n);
y=zeros(1,n);
for k=1:n
    y(k)=myF(x(k));
end
```

```
% Build y on the fly
x=linspace(a,b,n);
for k=1:n
    y(k)=myF(x(k));
end
```

Much faster for large n!

Lecture 14

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Concatenating 2 vectors—copy 2 vectors into a new one

```
% given row vectors x and y
v= zeros(1,length(x)+length(y));
for k=1:length(x)
    v(k)= x(k);
end
for k=1:length(y)
    v(length(x)+k)= y(k);
end
```

This is **non-vectorized code**—operations are performed on one component (scalar) at a time

Lecture 13

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Concatenating 2 vectors—copy 2 vectors into a new one

```
% given row vectors x and y
v= zeros(1,length(x)+length(y));
for k=1:length(x)
    v(k)= x(k);
end
for k=1:length(y)
    v(length(x)+k)= y(k);
end
```

Below is **vectorized code**—ops are performed on multiple components (a vector) at the same time:
 $v = [x \ y];$

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Split a vector in 2—copy values into 2 vectors

```
% given row vector v
s= ceil(rand*length(v)); % split pt
x= zeros(1,s);
y= zeros(1,length(v)-s);
for k=1:s
    x(k)= v(k);
end
for k=1:length(y)
    y(k)= v(s+k);
end
```

This is **non-vectorized code**—operations are performed on one component (scalar) at a time

Lecture 13

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Accessing a submatrix

M	2	-1	.5	0	-3
	3	8	6	7	7
	5	-3	8.5	9	10
	52	81	.5	7	2

- **M** refers to the whole matrix
- **M(3,5)** refers to one component of **M**

Lecture 15

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Accessing a submatrix

M	2	-1	.5	0	-3
	3	8	6	7	7
	5	-3	8.5	9	10
	52	81	.5	7	2

- **M** refers to the whole matrix
- **M(3,5)** refers to one component of **M**
- **M(2:3,3:5)** refers to a submatrix of **M**

row indices

column indices

Lecture 15

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Split a vector in 2—copy values into 2 vectors

```
% given row vector v
s= ceil(rand*length(v)); % split pt
x= zeros(1,s);
y= zeros(1,length(v)-s);
for k=1:s
    x(k)= v(k);
end
for k=1:length(y)
    y(k)= v(s+k);
end
```

Below is **vectorized code**: multiple components (subvectors) are accessed at the same time:

$x = v(1:s);$
 $y = v(s+1:length(v));$

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