



Lecture 26: Sorting

CS 1110

Introduction to Computing Using Python



Cornell CIS
COMPUTING AND INFORMATION SCIENCE

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Plan of Attack

- Insertion Sort
- Partition
- Quick Sort

Insertion Sort

```

PRE: b[0..n] = [?(unknown values)]
k = 0
INV: b[0..k] = sorted, b[k+1..n] = [?(unknown)]

while k < n:
    # Push b[k] down into its
    # sorted position in b[0..k]
    k = k+1

POST: b[0..n] = sorted
    
```

Insertion Sort: Moving into Position

```

def push_down(b, k):
    while k > 0:
        if b[k-1] > b[k]:
            swap(b, k-1, k)
        k = k-1

k = 0
while k < n:
    push_down(b, k)
    k = k+1
    
```

Is this how you want to sort 500 exams?

Algorithm Complexity

```

def push_down(b, k):
    while k > 0:
        if b[k-1] > b[k]:
            swap(b, k-1, k)
        k = k-1

k = 0
while k < n:
    push_down(b, k)
    k = k+1
    
```

Nested loops multiply the number of operations required. We need to compare $b[k]$ to all n elements. n operations

Iterating through a sequence of length n requires n operations: `push_down` called n times

Q: Algorithm Complexity

```

def push_down(b, k):
    while k > 0:
        if b[k-1] > b[k]:
            swap(b, k-1, k)
        k = k-1

k = 0
while k < n:
    push_down(b, k)
    k = k+1
    
```

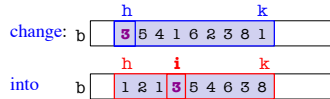
Approximately how many operations does this take?

A: ~ 1 operation
 B: ~ n operations
 C: ~ n² operations
 D: ~ n³ operations
 E: I don't know

Partition Algorithm

- Given a list segment $b[h..k]$ with some pivot value x in $b[h]$:
pre: b x | ?
- Swap elements of $b[h..k]$ and store in i to satisfy postcondition:
post: b $\leq x$ | x | $\geq x$

Example:



x

- Called the **pivot value**
- not a variable
- = whatever value is in $b[h]$

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Partition: What's the Invariant?

- Given a list segment $b[h..k]$ with some pivot value x in $b[h]$:
pre: b x | ?
- Swap elements of $b[h..k]$ and store in i to satisfy postcondition:
post: b $\leq x$ | x | $\geq x$



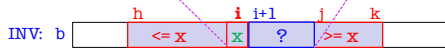
- Swap elements of $b[h..k]$ and store in i to satisfy postcondition:
post: b $\leq x$ | x | $\geq x$



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Partition: What's the Invariant?

- Given a list segment $b[h..k]$ with some pivot value x in $b[h]$:
pre: b x | ?
Initially $i = h, j = k+1$
- Swap elements of $b[h..k]$ and store in i to satisfy postcondition:
post: b $\leq x$ | x | $\geq x$



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Partition: What's the Invariant?

- Given a list segment $b[h..k]$ with some pivot value x in $b[h]$:
pre: b x | ?
- Swap elements of $b[h..k]$ and store in i to satisfy postcondition:
post: b $\leq x$ | x | $\geq x$



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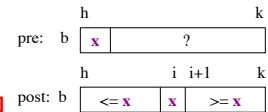
Sorting with Partitions

- Idea:** Pick a **pivot** element x
 - Partition sequence into $\leq x$ and $\geq x$
- Now Partition this and this, too
Keep recursing...
- b $\leq x$ | x | $\geq x$
- b Sorted!

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QuickSort

```
def quick_sort(b, h, k):
    """Sort the array fragment b[h..k]"""
    if k <= h:
        return
    i = partition(b, h, k)
    # INV: b[h..i-1] <= b[i] <= b[i+1..k]
    # Sort b[h..i-1] and b[i+1..k]
    quick_sort(b, h, i-1)
    quick_sort(b, i+1, k)
```



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