13 Oct 2023 NP- Complete Pooblems in Graph Theory

To show a problem A is NP- Complete you need two poly-time algorithms Verticer V: takes an input of A (x) and a proposed solution (y) and verifies y is a valid solution to X. $A(x) = 1 \iff \exists y \quad \forall (x,y) = 1$ E.g. For 35AT V(x,y) takes byical formula x touth assignment y, outputs I ff y stisfies X. (z) Reduction R: takes an imput X of some other problem known to be hard, a.g. 3SAT, and fransforms Xo into an input of A. K. $3SAT(x_0) = 1 \iff A(x_1) = 1$

EX. INDEPENDENT SET. Japut: indirected graph G, psidle integer K.

adput: 1 If and only If I a subset of k vertices of G, sit. No edge how both of its endpoints in the subset. E.g. 2 => 1 (yes three is a 2-element the set)

 $\sum_{i=1}^{n} 3 = i p$ Why NP-Gamplete? O V(x,y) takes x = (G,k)and y = 16 inarry string checking S. subset of V(G), say S. It checks $\left[5 \right] = 2$, 0(n)For each edge (u,v) it checks u\$S or v\$S O(m)2 We will reduce 35AT to IND SET. A structure that looks like a Bodean variable, Monstated to IND SET: XI=1 XI=0 1-element independent subsets of this graph are in 1:1 correspondence with farth assignments of Xi. Selecting truth assignments for a variables: x,=1 x=0 0 0 en element indgendent



IND SET RESTRICTED TO GRAPHS OF MAX DEGREE 3. (d3-IND-SET). for a variable X, that lackings to S clauses. Geolget - - \ Make each variable X: Exto a gadget with 2n. vertices & edges forming on even cycle, where n; dentes # clauses Containing Xi or Xi. Set $p = (\Sigma_i n_i) + (\# clauses)$ GRAPH 3-COLORDBILITY: Given undweeted G

can we color its vertices with 3 colors

such that the endpoints of every edge are differently colored.

