

April 18, 2019

Last time: network centrality

① PageRank ☆

$$\alpha P x + (1 - \alpha)v = x$$

② HITS (Hubs & Auth)

$$A = U \Sigma V^T$$

$$\text{auth} = v_1, \text{hub} = u_1$$

③ Eigenvector ☆

$$A x = \lambda_1 x$$

④ Katz

$$(\hat{I} - BA^T) R = BA^T \mathbf{1}$$

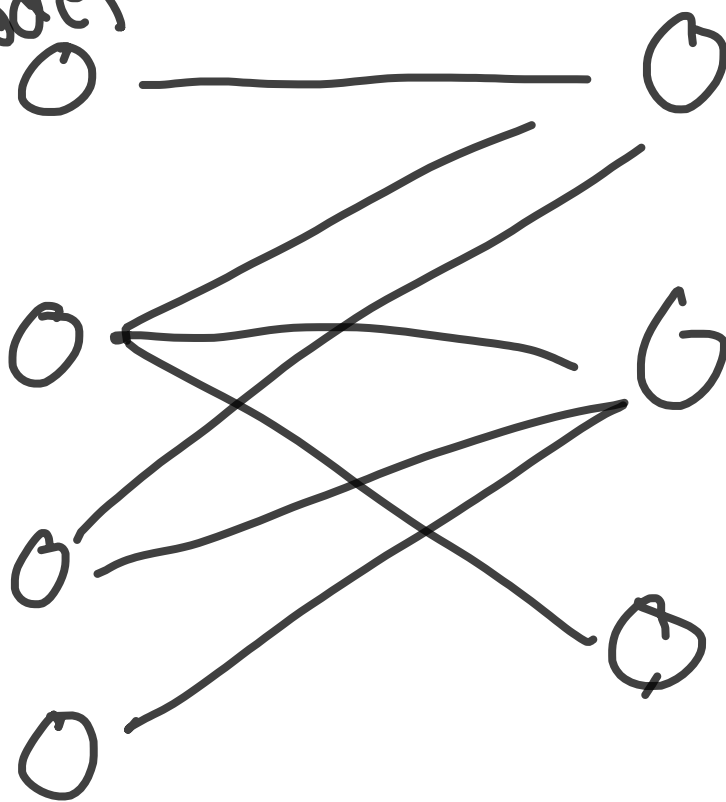
# Hypergraphs

instances of small sets

- multiple recipients on email
- authors on a paper
- multiple drugs/medications for patients

nodes

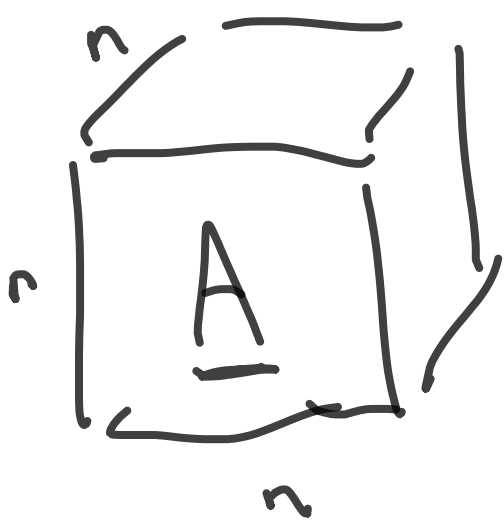
hyperedges



# Adjacency tensor

simple model:

- size-3 relationships
- no direction



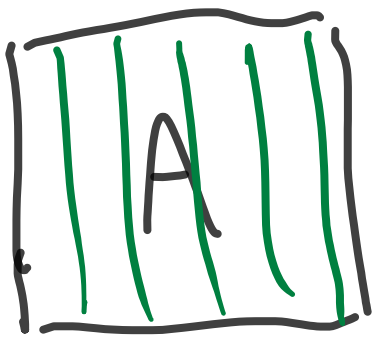
$$A_{ijk} = \begin{cases} 1 & \text{if } \{i, j, k\} \\ & \text{is in} \\ & \text{dataset} \\ 0 & \text{o/w} \end{cases}$$

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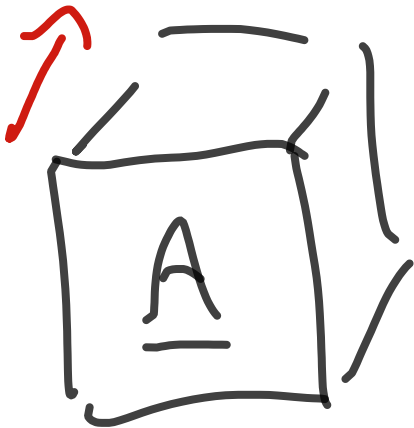
Tensor-based PageRank  
(Gleich, Lim, Yu 2015)

$$\alpha P x + (1-\alpha)v = x$$

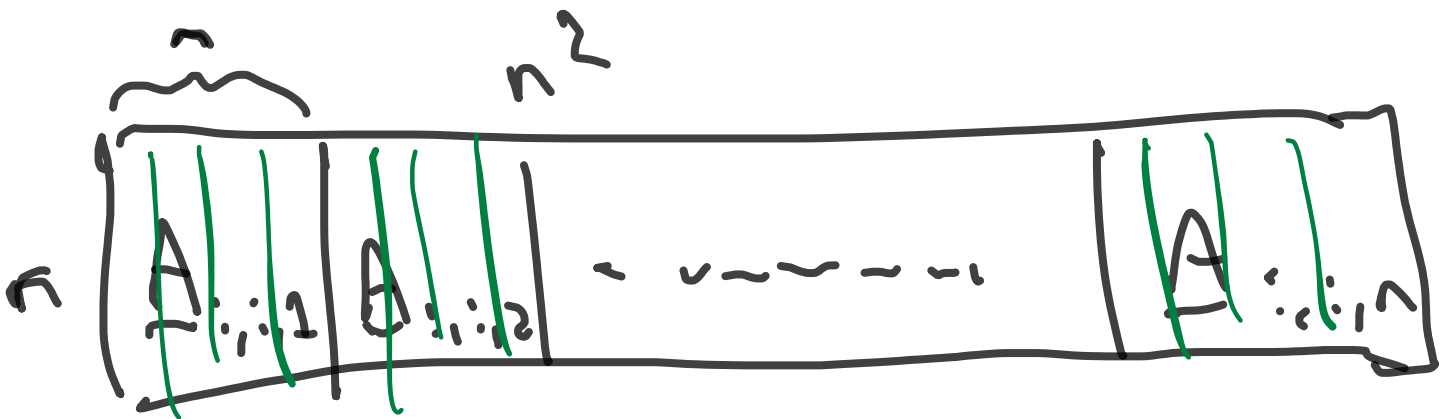
What is "P" matrix for tensors?



columns normalized  
to sum to 1



"unfold"  
↓



normalize to sum to 1  
refolding



Matrix case

$$\alpha P x + (1-\alpha)v = x$$

Tensor case

$$\alpha \underline{P} x^2 + (1-\alpha)v = x$$

$$[\underline{P} x^2]_i = \sum_{j,k} P_{ijk} x_j x_k$$

$x_1 x$	$x_2 x$	$\dots$	$x_n x$
$\underline{P}_{\dots i \dots 1}$	$\underline{P}_{\dots i \dots 2}$	$\dots$	$\underline{P}_{\dots i \dots n}$
+	+		+

$$= \underline{P} x^2$$

Does a solution even exist?

Yes!

Claim:  $f(x) = \alpha \underline{p}_x^2 + (1-\alpha)v$   
 is a stochastic vector

$$(\mathbf{1}^T f(x) = 1 \quad f(x) \geq 0)$$

Proof:  $\underline{p}_x^2 =$

$x_1 x$	$x_2 x$	$\dots$	$x_n x$
$\left[ \begin{array}{c} p_{\dots 11} \\ \vdots \\ p_{\dots 1n} \end{array} \right]$	$\left[ \begin{array}{c} p_{\dots 21} \\ \vdots \\ p_{\dots 2n} \end{array} \right]$	$\dots$	$\left[ \begin{array}{c} p_{\dots n1} \\ \vdots \\ p_{\dots nn} \end{array} \right]$

$\underline{p}_{i,j,k}$  is stochastic

$\gamma_j = \underline{p}_{\dots ij} x$  is stochastic

$\underline{p}_x^2 = \sum x_j \gamma_j$  is stochastic

$f(x) = \alpha \underline{p}_x^2 + (1-\alpha)v$  is stochastic

$$\alpha \frac{p}{x^2} + (1-\alpha)x = x$$

$f(x)$  is stochastic as long as  $x$  is stochastic

Theorem (Brouwer fixed point):

Let  $g: K \rightarrow K$  be a <sup>smooth</sup> function on a compact convex set  $K$

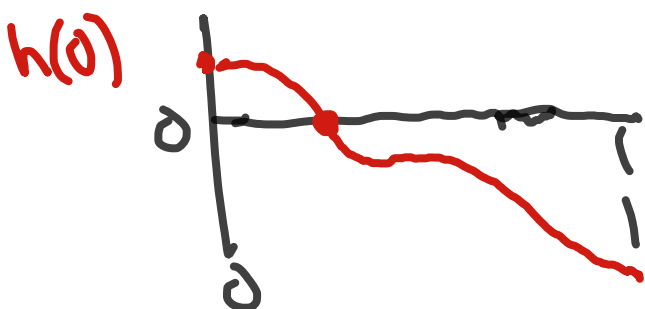
Then there exists  $x \in K$  such that  $g(x) = x$

Proof: for  $K = [0, 1]$

$$h(x) = g(x) - x$$

$$\exists \gamma \text{ s.t. } h(\gamma) = 0$$

$$\Leftrightarrow g(\gamma) = \gamma$$



$$K = \{w \in \mathbb{R}^n \mid \mathbb{1}^\top w = 1, w \geq 0\}$$

$$f: K \rightarrow K$$

$$f(x) = \alpha \underline{P} x^2 + (1-\alpha)v$$

BFP  $\Rightarrow \exists x$  such that

$$\alpha \underline{P} x^2 + (1-\alpha)v = x$$

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Can we actually compute a solution?

Matrix

$$x_{k+1} = \alpha P x_k + (1-\alpha)v$$

Tensor

$$\star x_{k+1} = f(x_k) = \alpha \underline{P} x_k^2 + (1-\alpha)v$$



# Theorem (Gleich, Lim, Yu)

if  $\alpha < 1/2$

- ★ converges (again, exponentially)

- solution  $\alpha \sum_{n=0}^{\infty} x^n + (1-\alpha)v = x$

is unique

if  $\alpha \geq 1/2$

- ★ doesn't necessarily converge

- solution may not be unique

# Hypergraph eigenvector centrality (Benson 2019)

graph evec centrality

$$Ax = \lambda_1 x \quad x > 0$$

*power*

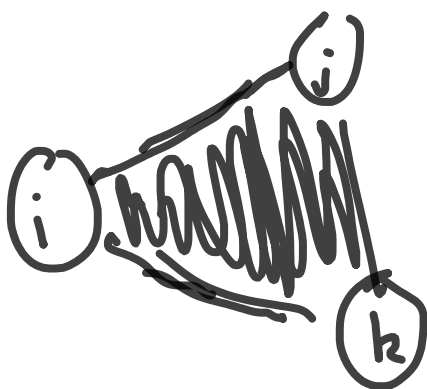
$$x_i = \gamma \sum_{(i,j) \in E} x_j \leftarrow \text{power of } j$$

$$\lambda_1 = \frac{1}{\gamma}$$

Theoretical tool:

Perron - Frobenius Theorem

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power of node  $i$   
prop. to powers of  
all such  $j, k$

$$\textcircled{1} \quad x_i = \gamma \sum_{(i,j,k)} x_j x_k$$

$$\Rightarrow \underline{A} x^2 = \frac{1}{\gamma} x$$

$x$ -eigenvector of a tensor

$$\textcircled{2} \quad x_i^2 = \gamma \sum_{(i,j,k)} x_j x_k$$

$$\Rightarrow \underline{A} x^2 = \frac{1}{\gamma} x^2$$

$$[x^2]_i = x_i^2$$

H-eigenvector of  $\underline{A}$

$$(3) \quad x_i = \gamma \sum_{(i,j,k)} x_j + x_k$$

$$\Rightarrow Wx = \frac{1}{2\gamma} x$$

$W_{ij} = \# k$  such that

$$A_{ijk} = 1$$

$$W = \sum_k A_{\dots, k}$$

Can apply P-F like last time

$\Rightarrow$  unique  $x > 0$  such that

$$Wx = \lambda_1 x$$

$$\lambda_1 > 0$$

$$\left. \begin{array}{l} \text{(Z)} \quad \underline{A} x^2 = \lambda x \\ \text{(H)} \quad \underline{A} x^2 = \lambda x^2 \end{array} \right\} \begin{array}{l} \text{also want} \\ x > 0 \end{array}$$


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Start with H

Definition:  $\underline{A}$  is irreducible  
 if  $W$  ( $W := \sum \underline{A}_{i,j}$ ) induces  
 a connected graph

Theorem: If  $\underline{A}$  is irreducible  
 and nonnegative, then exists  
<sup>unique</sup>  
 $\lambda > 0$ ,  $x > 0$  such that

$$\underline{A} x^2 = \lambda x^2$$

any other eigenvalue  $\lambda'$   
has  $|\lambda'| \leq 1$

$\Rightarrow$  this gives us the  
centrality vector

Can we compute it?

Yes! (Ng, Qi, Zhou 2009)

$$y_k = Ax_k^2$$

$$x_{k+1} = \frac{y_k}{\|y_k^{1/2}\|_2}$$

entry-wise  $\checkmark$

# Z-eigenvector centrality

Want:

$$\underline{A}x^2 = \lambda x$$

$$\lambda > 0, x > 0$$

Theorem: if  $\underline{A}$  is irreducible,  
then there exists  $\lambda > 0, x > 0$   
such that  $\underline{A}x^2 = \lambda x$

but...

- NP-hard to compute

- could be another  $\lambda' > 0, y > 0$

such  $\underline{A}y^2 = \lambda' y$

if we can compute some

$\lambda \geq 0, x \geq 0$  such that

$$\underline{A}x^2 = \lambda x$$

then we have tensor

$\mathbb{Z}$ -eigenvector centrality vector

Options:

- worst-case exponential time algorithms

- heuristics that tend to work very well in practice