Monoidal Categories

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Definition. Given a category \mathbf{C} , the category $\mathbb{L}\mathbf{C}$ is comprised of the following:

Objects An object is a list of objects of C.

Morphisms A morphism from $[A_1, \ldots, A_m]$ to $[B_1, \ldots, B_n]$ only exists when m equals n, in which case it is a list of morphisms $f_1 : A_1 \to B_1, \ldots, f_n : A_n \to B_n$ of **C**

Identities The identity of $[A_1, \ldots, A_n]$ is $[id_{A_1}, \ldots, id_{A_n}]$

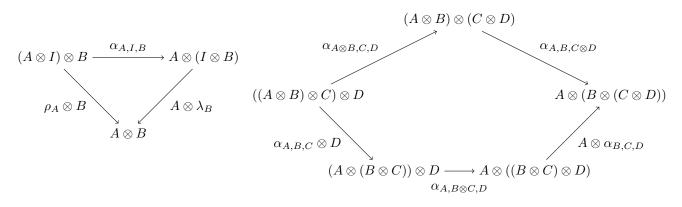
Composition The composition of $[f_1, \ldots, f_n]$ with $[g_1, \ldots, g_n]$ is $[f_1; g_1, \ldots, f_n; g_n]$

Remark. The above construction extends to a 2-monad on Cat, with $\eta_{\mathbf{C}}$ being the *functor* that maps each object and each morphism to the singleton list containing it, and $\mu_{\mathbf{C}}$ being the functor that maps each list of lists of objects and each list of lists of morphisms to its flattening.

Definition (Strict Monoidal Category). A strict monoidal category is equivalently defined as a (strict) monad algebra of \mathbb{L} on **Cat** or as an internal monoid of the multicategory **Cat**, i.e. a category **C** along with an object Iof **C** and a binary functor $\otimes : [\mathbf{C}, \mathbf{C}] \to \mathbf{C}$ satisfying identity and associativity. Similarly, a strict monoidal functor is equivalently defined as a (strict) morphism of monad algebras of \mathbb{L} on **Cat** or as an internal monoid homomorphism of the multicategory **Cat**, i.e. a functor F that preserves I and \otimes (strictly). A strict monoidal transformation is defined as a transformation of (strict) morphisms of monad algebras of \mathbb{L} on **Cat**, or equivalently as a natural transformation α from F to G such that α_I equals $id_{I'}$ and $\alpha_{A\otimes B}$ equals $\alpha_A \otimes' \alpha_B$.

Definition (Weak Monoidal Category). A weak monoidal category is equivalently defined as a weak monad algebra of \mathbb{L} on **Cat** or as the category of unary morphisms of a representable multicategory. Similarly, a weak monoidal functor is equivalently defined as a weak morphism of weak monad algebras of \mathbb{L} on **Cat** or as a functor corresponding to a tensor-preserving multifunctor between representable multicategories, i.e. multifunctors F with the property that FT with Ft is a tensor of FA_1, \ldots, FA_n whenever T with t is a tensor of A_1, \ldots, A_n . A weak monoidal transformation is defined as a transformation of weak morphisms of weak monad algebras of \mathbb{L} on **Cat**, or equivalently as a natural fransformation corresponding to a natural transformation α of tensor-preserving multifunctors such that $\alpha_T : FT \to GT$, where T with $t : \vec{A} \to T$ is a tensor of A_1, \ldots, A_n , has the property that $\Delta_{Ft} \alpha_T$ equals $\Delta_{\alpha_A, \ldots, \alpha_A_n} Gt$.

Remark. Yet another equivalent, and particularly common, definition of weak monoidal category is a category **C** along with an object I of **C**, a binary functor $\otimes : \mathbf{C} \times \mathbf{C} \to \mathbf{C}$, and natural isomorphisms $\{\rho_A : A \otimes I \to A\}_{A \in \mathbf{C}}$, $\{\lambda_A : I \otimes A \to A\}_{A \in \mathbf{C}}$, and $\{\alpha_{A,B,C} : (A \otimes B) \otimes C \to A \otimes (B \otimes C)\}_{A,B,C \in \mathbf{C}}$ making the following triangle and pentagon commute:



Definition (Lax Monoidal Functor). A lax monoidal functor between strict/weak monoidal categories is equivalently defined as a lax morphism of strict/weak monad algebras of \mathbb{L} on **Cat** or as a functor corresponding to a (not necessarily tensor-preserving) multifunctor between representable multicategories. A lax monoidal transformation is defined as a transformation of lax morphisms of weak monad algebras of \mathbb{L} on **Cat**, or equivalently as a natural fransformation corresponding to a natural transformation of multifunctors.

Remark. Yet another equivalent, and particularly common, definition of lax monoidal functor corresponding to the earlier common definition of weak monoidal categories is a functor F along with a morphism merge : $I' \to FI$ and a natural transformation $\{\text{merge}_{A,B} : FA \otimes' FB \to F(A \otimes B)\}_{A,B \in \mathbb{C}}$ such that the following diagrams commute:

