## **Functors**

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**Definition** ((Biased) (**Set**-enriched) (Covariant) Functor from  $\langle O_{\mathbf{C}}, M_{\mathbf{C}}, ;, \iota, id_{\mathbf{C}}, \iota \rangle$  to  $\langle O_{\mathbf{D}}, M_{\mathbf{D}}, ;_{\mathbf{D}}, \iota, id_{\mathbf{D}}, \iota \rangle$ . A tuple  $\langle F_O, F_M, \mathfrak{d}, \mathfrak{i} \rangle$  where the components have the following types:

Function on Objects  $F_O: O_C \rightarrow O_D$ 

Function on Morphisms  $F_M: \forall \mathcal{C}_1, \mathcal{C}_2: O_{\mathbf{C}}. M_{\mathbf{C}}(\mathcal{C}_1, \mathcal{C}_2) \to M_{\mathbf{D}}(F_O(\mathcal{C}_1), F_O(\mathcal{C}_2))$  (objects implicit)

Distributivity  $\mathfrak{d}$ :  $\forall \mathcal{C}_1, \mathcal{C}_2, \mathcal{C}_3 : O_{\mathbf{C}}. \ \forall m_1 : M_{\mathbf{C}}(\mathcal{C}_1, \mathcal{C}_2), m_2 : M_{\mathbf{C}}(\mathcal{C}_2, \mathcal{C}_3). \ F_M(m_1) :_{\mathbf{D}} F_M(m_2) = F_M(m_1 :_{\mathbf{C}} m_2)$ 

Identity i:  $\forall C : O_{\mathbf{C}}$ .  $id_{\mathbf{D}} = F_M(id_{\mathbf{D}})$ 

*Notation.* We represent both  $F_O$  and  $F_M$  with just F.

Notation. When multiple categories (such as C and D) are present, we represent all their morphism types (e.g.  $M_C$  and  $M_D$ ) with just M or with the infix operator  $\rightarrow$ .

Notation. When multiple categories (such as C and D) are present, we represent all their composition operators (e.g. ;<sub>C</sub> and ;<sub>D</sub>) with just ;.

Notation. When multiple categories (such as  $\mathbf{C}$  and  $\mathbf{D}$ ) are present, we represent all their identity operators (e.g.  $id_{\mathbf{C}}$  and  $id_{\mathbf{D}}$ ) with just id.

*Notation.* We use  $C : \mathbf{C}$  to denote  $C : O_{\mathbf{C}}$ .

*Remark.* With these notations, the above definitions can be rephrased as: A functor from **C** to **D** is a tuple  $\langle F, F, \mathfrak{d}, \mathfrak{i} \rangle$  where the components have the following types:

Function on Objects  $F: O_{\mathbf{C}} \to O_{\mathbf{D}}$ 

Function on Morphisms  $F: \forall \mathcal{C}_1, \mathcal{C}_2 : \mathbf{C}. \ \mathcal{C}_1 \to \mathcal{C}_2 \to F(\mathcal{C}_1) \to F(\mathcal{C}_2)$  (objects implicit)

Distributivity  $\mathfrak{d}: \forall \mathcal{C}_1, \mathcal{C}_2, \mathcal{C}_3: \mathbf{C}. \forall \mathcal{C}_1 \xrightarrow{m_1} \mathcal{C}_2, \mathcal{C}_2 \xrightarrow{m_2} \mathcal{C}_3. F(m_1); F(m_2) = F(m_1; m_2)$ 

**Identity** i:  $\forall C : \mathbf{C}$ . id = F(id)

Example. Prost to Set:  $\langle \lambda \tau, \tau, \lambda \langle \tau_1, \tau_2 \rangle, \lambda \langle f, \cdot \rangle, f, \cdot, \cdot \rangle$ 

Mat to Set:  $\langle \lambda n. \mathbb{R}^n, \lambda \langle n_1, n_2 \rangle. \lambda M. \lambda \vec{x}. M \cdot \vec{x}, ., . \rangle$ 

 $\triangle$  to Prost:  $\langle \lambda n. \langle n, \leq, \bullet, \bullet \rangle, \lambda \langle n_1, n_2 \rangle. \lambda \sigma. \langle \sigma, \bullet \rangle, \bullet, \bullet \rangle$ 

 $\omega$  to  $\Delta$ :  $\langle \lambda n. n, \lambda \langle n_1, n_2 \rangle$ .  $\lambda$  Ite.  $\lambda n. n$  as  $m_2$  using Ite,  $\bullet$ ,  $\bullet$ 

Set to Rel:  $\langle \lambda \tau. \tau, \lambda \langle \tau_1, \tau_2 \rangle. \lambda f. \lambda \langle t_1, t_2 \rangle. f(t_1) = t_2, ., . \rangle$ 

Rel to Set:  $\langle \mathbb{P}, \lambda \langle \tau_1, \tau_2 \rangle$ .  $\lambda \phi$ .  $\lambda T_1$ .  $\{t_2 : \tau_2 \mid \exists t_1 : \tau_1 . t_1 \in T_1 \land \phi(t_1, t_2)\}, \bullet, \bullet \rangle$ 

**Mon**<sub>b</sub> to Alg(2,0):  $\langle \lambda \langle M, *, \bullet, e, \bullet \rangle$ .  $\langle M, *, e \rangle$ ,  $\lambda \langle M_1, M_2 \rangle$ .  $\lambda \langle f, \bullet, \bullet \rangle$ .  $\langle f, \bullet \rangle$ ,  $\bullet, \bullet \rangle$ 

Mon<sub>u</sub> to Alg( $\langle \mathbb{N}, \lambda n. \, \mathbb{n} \rangle$ ):  $\langle \lambda \langle M, \prod, \bullet \rangle . \, \langle M, \lambda n. \, \lambda(m_i)_{i \in \mathbb{n}}. \, \prod [m_1, \ldots, m_n], \bullet \rangle, \lambda \langle M_1, M_2 \rangle . \, \lambda \langle f, \bullet \rangle . \, \langle f, \bullet \rangle, \bullet, \bullet \rangle$ 

Prost to Rel(2):  $\langle \lambda \langle \tau, R, . \rangle$ .  $\langle \tau, R \rangle$ ,  $\lambda f. f. . . . \rangle$ 

**Exercise 1.** Prove that there are functors from **Set** to **Set** mapping a set S to  $\mathbb{L}S$ , to  $\mathbb{S}S$ , and to  $\mathbb{P}S$ .

**Exercise 2.** Prove that the set of functors from 1 to C is isomorphic to  $O_C$ .

**Exercise 3.** Prove that the set of functors from 2 to C is isomorphic to  $\sum_{C_1,C_2:C} C_1 \to C_2$ .

**Exercise 4.** Prove that there is a category **CAT** (of a higher universe) of categories whose objects are categories and whose morphisms are functors, such that there exists a functor from **CAT** to **SET** (i.e. **Set** for a higher universe) mapping categories to their set of objects and functors to their functions on objects.

**Definition** (Cat). The category of "small" categories (categories where O is an element of Type rather than Type<sub>1</sub>) and functors between them.

**Definition** ((Biased) (**Set**-enriched) Contravariant Functor from  $\langle O_{\mathbf{C}}, M_{\mathbf{C}}, ;, \iota, id_{\mathbf{C}}, \cdot \rangle$  to  $\langle O_{\mathbf{D}}, M_{\mathbf{D}}, ;_{\mathbf{D}}, \iota, id_{\mathbf{D}}, \cdot \rangle$ ). A tuple  $\langle F_O, F_M, \mathfrak{d}, \mathfrak{i} \rangle$  where the components have the following types:

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Function on Morphisms  $F: \forall \mathcal{C}_1, \mathcal{C}_2 : \mathbf{C}. \ \mathcal{C}_1 \to \mathcal{C}_2 \to F(\mathcal{C}_2) \to F(\mathcal{C}_1)$  (objects implicit)

**Distributivity 0:**  $\forall C_1, C_2, C_3 : \mathbf{C}. \ \forall C_1 \xrightarrow{m_1} C_2, C_2 \xrightarrow{m_2} C_3. \ F(m_2) ; F(m_1) = F(m_1; m_2)$ 

**Identity** i:  $\forall C : \mathbf{C}. id = F(id)$ 

**Exercise 5.** Prove that there is a contravariant functor from **Set** to **Set** mapping a set S to  $\mathbb{P}S$ .

Exercise 6. Prove that there is a contravariant functor from Rel to Rel whose object is component is the identity function.

Exercise 7. Prove that there is a contravariant functor from Sig to CAT mapping  $\Omega$ : Sig to Alg $(\Omega)$ : CAT.

**Exercise 8.** Prove that there is a contravariant functor from Sig to CAT mapping  $\Phi$ : Sig to Rel $(\Phi)$ : CAT.